RATIONALMAPS, A PACKAGE FOR MACAULAY2

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Abstract. This paper describes the RationalMaps package for Macaulay2. This package provides functionality for computing several aspects of rational maps.

1. INTRODUCTION

This package aims to compute a number of things about rational maps between varieties. In particular, this package will compute

- The base locus of a rational map.
- Whether a rational map is birational.
- The inverse of a birational map.
- Whether a map is a closed embedding.
- And more!

Our functions have numerous options which allow them to run much more quickly in certain examples if configured correctly. The Verbuse option gives hints as to the best way to apply these.

Suppose \( \mathfrak{F} : X \to Y \) is given by homogeneous forms \( f \) of the same degree. The idea of looking at the syzygies of the forms \( f \) to detect the geometric properties of \( \mathfrak{F} \) goes back at least to [HKS92] in the case where \( X = \mathbb{P}^n \), \( Y = \mathbb{P}^m \) and \( m = n \) (see also [ST69]). In [RS01] this method was developed by Russo and Simis to handle the case \( X = \mathbb{P}^n \) and \( m \geq n \). Simis pushed the method further to the study of general rational maps between two integral projective schemes in arbitrary characteristic by an extended ideal-theoretic method emphasizing the role of the Rees algebra associated to the ideal generated by \( f \) [Sim04]. Recently, Doria, Hassanzadeh, and Simis applied these Rees algebra techniques to study the birationality of \( \mathfrak{F} \) [DHS12]. Our core functions, in particular the functions related to computing inverse maps, rely heavily on this work.

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2. BASE LOCI

We begin with the problem of computing the base locus of a map to projective space. Let \( X \) be a projective variety over any field \( k \) and let \( \mathfrak{F} : X \to \mathbb{P}^m_k \) be a rational map from \( X \) to projective space. Then we can choose some representative \( (f_0, \ldots, f_m) \) of \( \mathfrak{F} \), where each \( f_i \) is the \( i \)th coordinate of \( \mathfrak{F} \). A priori, each \( f_i \) is in \( K = \text{frac} R \), where \( R \) is the coordinate ring of \( X \). However, we can get another representative of \( \mathfrak{F} \) by clearing denominators. (Note this does not enlarge the base locus of

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vanish, and thus the base locus is the vanishing set of the ideal (f_i since \( \mathfrak{g} \) is undefined whenever the denominator of any of the \( f_i \) vanishes.) Thus we can assume that \( f_i \in R \) for all \( i \), and that all the \( f_i \) are homogenous of the same degree.

In this setting, one might naively think that the map \( \mathfrak{g} \) is undefined exactly when all of the \( f_i \) vanish, and thus the base locus is the vanishing set of the ideal \( (f_0, \ldots, f_m) \). However, this can yield a base locus that’s too big. Indeed, to find the base locus of a rational map, we must consider all possible representatives of the map and find where none of them are defined. To do this, we use the following result.

**Proposition 2.1.** \([\text{Sim04}, \text{Proposition } 1.1]\) Let \( \mathfrak{g} : X \dashrightarrow \mathbb{P}^m \) be a rational map and let \( f = \{f_0, \ldots, f_m\} \) be a representative of \( \mathfrak{g} \) with \( f_i \in R \) homogenous of degree \( d \) for all \( i \). Set \( I = (f_0, \ldots, f_m) \). Then the set of such representatives of \( \mathfrak{g} \) corresponds bijectively to the homogenous vectors in the rank 1 graded \( R \)-module \( \text{Hom}_R(I, R) \cong (R : \kappa I) \).

The bijection comes from multiplying our fixed representative \( f \) by \( h \in (R : \kappa I) \). Now, in the setting of [Proposition 2.1] let

\[
\bigoplus_s R(d_s) \xrightarrow{\varphi} R(-d)^m \xrightarrow{[f_0: \cdots : f_m]} I \rightarrow 0
\]

be a free resolution of \( I \). Then we get

\[
0 \to \text{Hom}_R(I, R) \to (R(-d)^{m+1})^\vee \xrightarrow{\varphi^t} \left( \bigoplus_s R(d_s) \right)^\vee
\]

where \( \varphi^t \) is the transpose of \( \varphi \) and \( R^\vee \) is the dual module of \( R \). Thus, we get that \( \text{Hom}_R(I, R) \cong \ker \varphi^t \), and so each representative of \( \mathfrak{g} \) corresponds to a vector in \( \ker \varphi^t \). The correspondence takes a representative \( (h_0f_0, \ldots, h_mf_m) \) to the map that multiplies vectors in \( R^{m+1} \) by \( [h_0f_0, \ldots, h_mf_m] \) on the left.

The base locus of \( \mathfrak{g} \) is the intersection of the sets \( V(f_0^i, \ldots, f_m^i) \) as \( f^i = (f_0^i, \ldots, f_m^i) \) ranges over all the representatives of \( \mathfrak{g} \). The above implies that this is the same as the intersection of the sets \( V(w_0^i, \ldots, w_m^i) \) as \( w^i = (w_0^i, \ldots, w_m^i) \) ranges over the vectors in \( \ker \varphi^t \). Now, given any \( a, f, g \in R \), we have \( V(af) \supseteq V(f) \) and \( V(f + g) \supseteq V(f) \cap V(g) \). Thus, it’s enough to take a generating set \( w_1^{\ast}, \ldots, w_n^{\ast} \) of \( \ker \varphi^t \) and take the intersection over this generating set.

The base locus of \( \mathfrak{g} \) is then the variety cut out by the ideal generated by all the entries of all of the \( w^i \). Our function \texttt{baseLocusOfMap} returns this ideal.

```plaintext
\texttt{i1 : loadPackage "RationalMaps";}
\texttt{i2 : R = QQ[x,y,z];}
\texttt{i3 : f = {x^2*y, x^2*z, x*y*z};}
\texttt{i4 : baseLocusOfMap(f);}
\texttt{o4 : Ideal of R}
```

If the \texttt{SaturateOutput} option is set \texttt{true}, our function will return the saturation of this ideal.

### 3. Birationality and Inverse Maps

Again, a rational map \( \mathfrak{g} : X \subseteq \mathbb{P}^n \dashrightarrow Y \subseteq \mathbb{P}^m \) between projective spaces is defined by \( m+1 \) forms \( f = \{f_0, \ldots, f_m\} \) of the same degree in \( n+1 \) variables, not all vanishing. Let \( f = \{f_0, \ldots, f_m\} \subseteq R \) be an ordered set of forms of the same degree, where \( R \), the coordinate ring of \( X \), denotes a standard graded \( k \)-algebra, with \( k \) a field and \( \dim R \geq 1 \). Thus, one has a presentation \( R \cong k[x_0, \ldots, x_n] = k[\mathbf{X}] / \mathfrak{a} \) with \( k[\mathbf{X}] = k[X_0, \ldots, X_n] \) standard graded and \( \mathfrak{a} \) a homogeneous ideal. Let \( R[\mathbf{Y}] = R[Y_0, \ldots, Y_m] \) be a polynomial ring over \( R \) with the standard bigrading where \( \deg(X_i) = (1, 0) \) and \( \deg(Y_j) = (0, 1) \). Since \( I = (f_0, \ldots, f_m) \) is generated in a fixed degree, the Rees algebra \( \mathcal{R}_R(I) := R \oplus I \oplus I^2 \oplus \cdots \cong R[It] \subset R[t] \)
is a standard bigraded $k$-algebra. Mapping $Y_j \mapsto f_j t$ yields a presentation $R[\mathbf{Y}]/\mathfrak{J} \simeq \mathcal{R}_R((\mathbf{f}))$, with $\mathfrak{J}$ a bihomogeneous presentation ideal. $\mathfrak{J}$ depends only on the rational map defined by $\mathbf{f}$ and not on this particular representative. Write

$$\mathfrak{J} = \bigoplus_{(p,q) \in \mathbb{N}^2} \mathfrak{J}_{(p,q)},$$

where $\mathfrak{J}_{(p,q)}$ denotes the $k$-vector space of forms of bidegree $(p, q)$.

For birationality, the following bihomogeneous piece is important:

$$\mathfrak{J}_{1,*} := \bigoplus_{r \in \mathbb{N}} \mathfrak{J}_{1,q}$$

with $\mathfrak{J}_{1,q}$ denoting the bigraded piece of $\mathfrak{J}$ spanned by the forms of bidegree $(1, q)$ for all $q \geq 0$. Now, a form of bidegree $(1, *)$ can be written as $\sum_{i=0}^n Q_i(\mathbf{Y}) x_i$, for suitable homogeneous $Q_i(\mathbf{Y}) \in k[\mathbf{Y}] \subset R[\mathbf{Y}]$ of the same degree. Since the $\mathbf{Y}$ are indeterminates over $R$, each pair of such representations of the same form gives a syzygy of $\{x_0, \ldots, x_n\}$ with coefficients in $k$. Thus every such representation is unique up to $k$-linear relations of $\{x_0, \ldots, x_n\}$, i.e., up to elements of $\mathfrak{a}_1$, the degree-1 part of $\mathfrak{a}$.

Next, one can pick a minimal set of generators of the ideal $(\mathfrak{J}_{1,*})$ consisting of a finite number of forms of bidegree $(1, q)$, for various $q$'s. Let $\{P_1, \ldots, P_3\} \subset k[\mathbf{X}, \mathbf{Y}]$ denote liftings of these biforms and let $\{\ell_1, \ldots, \ell_r\} \subset k[\mathbf{X}]$ be a $k$-vector space basis of $\mathfrak{a}_1$. Consider the Jacobian matrix of the polynomials $\{\ell_1, \ldots, \ell_r, P_1, \ldots, P_3\}$ with respect to $\mathbf{X}$. This is a matrix with entries in $k[\mathbf{Y}]$. Write $\psi$ for the corresponding matrix over $S = k[\mathbf{Y}]/b$, the coordinate ring of $\mathbf{Y}$. This matrix is called the weak Jacobian dual matrix associated to the given set of generators of $(\mathfrak{J}_{1,*})$. Note that a weak Jacobian matrix $\psi$ is not uniquely defined due to the lack of uniqueness in the expression of an individual form and to the choice of bihomogeneous generators. However, it is shown in [DHS12, Lemma 2.13] that if the weak Jacobian matrix associated to one set of bihomogeneous minimal generators of $(\mathfrak{J}_{1,*})$ has (finite) rank over $S$ then the weak Jacobian matrix associated to any other set of bihomogeneous minimal generators of $(\mathfrak{J}_{1,*})$ has (finite) rank over $S$ and the two ranks coincide.

The main theorem of loc.cit. is the following:

**Theorem 3.1.** Let $X \subseteq \mathbb{P}^n$ be non-degenerate. Then $\mathfrak{F}$ is birational if and only if $\text{rank}(\psi) = \text{edim}(R) - 1 (= n)$. Moreover

(i) We get a representative for the inverse of $\mathfrak{F}$ by taking the coordinates of any homogeneous vector of positive degree in the (rank one) null space of $\psi$ over $S$ for which these coordinates generate an ideal containing a regular element.

(ii) If, further, $R$ is a domain, the representative of $\mathfrak{F}$ in (i) can be taken to be the set of the (ordered, signed) $(\text{edim}(R) - 1)$-minors of an arbitrary $(\text{edim}(R) - 1) \times \text{edim}(R)$ submatrix of $\psi$ of rank $\text{edim}(R) - 1$.

As expected, the most expensive part of applying this theorem is computing the Rees ideal $\mathfrak{J}$. In the package RationalMaps we use ReesStrategy to compute the Rees equations. The algorithm is the standard elimination technique. However we do not use the ReesAlgebra package, since verifying birationality according to Theorem 3.1 only requires to computing a small part of the Rees ideal, namely elements of first-degree 1. This idea is applied in the SimisStrategy. More precisely, if the given map $\mathfrak{F}$ is birational, then the jacobian dual rank will attain its maximum value of $\text{edim}(R) - 1$ after computing the Rees equations up to degree $(1, N)$ for $N$ sufficiently large. This allows us to compute the inverse map. The down-side of SimisStrategy is that if $\mathfrak{F}$ is not birational, the desired number $N$ cannot be found and the process never terminates. To provide
a definitive answer for birationality, we use HybridStrategy, which is a hybrid of ReesStrategy and SimisStrategy. The default strategy is HybridStrategy.

HybridLimit is an option to switch SimisStrategy to ReesStrategy, if the computations up to degree \( (1, \text{HybridLimit}) \) do not lead to \( \text{rank}(\psi) = \text{edim}(R) - 1 \). The default value for HybridLimit is 15. The change from SimisStrategy to ReesStrategy is done in such a way that the generators of the Rees ideal computed in the SimisStrategy phase are not lost; the program computes other generators of the Rees ideal while keeping the generators it found before attaining HybridLimit.

There is yet another method for computing the Rees ideal called SaturationStrategy. In this option the whole Rees ideal is computed by saturating the defining ideal of the symmetric algebra with respect to a non-zero element in \( R \) (we assume \( R \) to be a domain). This Strategy appears to be slower in some examples, though one might be able improve this option in the future by stopping the computation of the saturation at a certain step.

Computing inverse maps is the most important functionality of this package, and is done by the function inverseOfMap. According to **Theorem 3.1**, there are two ways to compute the inverse of a map: (1) by finding any syzygy of the jacobian dual matrix, and (2) by finding a sub-matrix of \( \psi \) of rank \( \text{edim}(R) - 1 \). Each way has its own benefits. Method (1) is quite fast in many cases, however method (2) is very useful if the rank of the jacobian dual matrix \( \psi \) is relatively small compared to the degrees of the entries of \( \psi \). Our function inverseOfMap starts by using the second method and later switches to the first method if the second method didn’t work. The timing of this transition from the first method to the second method is controlled by the option MinorsCount. Setting MinorsCount to zero will mean that no minors are checked and the inverse map is computed just by looking at the syzygies of \( \psi \). If MinorsCount is left as null (the default value), the program will try to make an educated guess as to how big to set this option, depending on varieties the user is working with.

In addition, to improve the speed of the function inverseOfMap, we have two other options, AssumeDominant and CheckBirational. If AssumeDominant is set to be true, then inverseOfMap assumes that the map from \( X \) to \( Y \) is dominant and does not compute the image of the map; this is time consuming in certain cases. Similarly, if CheckBirational set false, inverseOfMap will not check birationality although it still computes the jacobian dual matrix.

In general, as long as **Verbose** is true, the function will make suggestions as to how to run it more quickly. For example.

```plaintext
i1 : loadPackage "RationalMaps";
i2 : Q=QQ[x,y,z,t,u];
i3 : phi=map(Q,Q,matrix{{x^5,y*x^4,z*x^4+y^5,t*x^4+z^5,u*x^4+t^5}});
o3 : RingMap Q <--- Q
i4 : time inverseOfMap(phi, AssumeDominant=>true,CheckBirational=>false, MinorsCount=>50000)
Starting inverseOfMapSimis(SimisStrategy or HybridStrategy)
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree \( \{1, 1\} \).
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree \( \{1, 2\} \).
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree \( \{1, 4\} \).
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree \( \{1, 7\} \).
inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree \( \{1, 11\} \).
inverseOfMapSimis: We give up. Using all of the previous computations,
we compute the whole Groebner basis of the rees ideal.
Increase HybridLimit and rerun to avoid this.
inverseOfMapSimis: Found Jacobian dual matrix (or a weak form of it), it has 5 columns
and about 20 rows.
inverseOfMapSimis: Looking for a nonzero minor
Starting nonZeroMinor, looking for rank: 4, we will run it 50000 times. If this is slow, rerun with MinorsCount=>0.
nonZeroMinor: Found a nonzero minor
inverseOfMapSimis: We found a nonzero minor.
-- used 3.51586 seconds
```
4. Embeddings

Our package also checks whether a rational map $\Phi : X \to Y$ is a closed embedding. The strategy is quite simple.

(a) We first check whether $\Phi$ is regular (by checking if its base locus is empty).
(b) We next invert the map (if possible).
(c) Finally, we check if the inverse map is also regular.

If all three conditions are met, then the map is a closed embedding and the function returns $\text{true}$. Otherwise $\text{isEmbedding}$ returns $\text{false}$.

```
i1 : needspackage "divisor"; --used to quickly define a map
i2 : c = QQ[x,y,z]/(x^4+x^2*y*z+y^4+z^3*x);
i3 : q = ideal(y,x+z); --a point on our curve
i6 : loadpackage "rationalmaps";
i8 : f2 = maptoprojectivespace(12*divisor(q));
o8 : RingMap C <--- QQ[YY , YY , YY , YY , YY , YY , YY , YY , YY ]
   1 2 3 4 5 6 7 8 9 10
i9 : time isEmbedding(f2)
   isEmbedding: About to find the image of the map.
   If you know the image, you may want to specify that and set AssumeDominant=>true option if this is slow.
   isEmbedding: Checking to see if the map is a regular map
   isEmbedding: computing the inverse map
   Starting inverseOfMapSimis(SimisStrategy or HybridStrategy)
   inverseOfMapSimis: About to compute partial Groebner basis of rees ideal up to degree {1, 1}.
   inverseOfMapSimis: We computed enough of the Groebner basis.
   inverseOfMapSimis: Found Jacobian dual matrix (or a weak form of it), it has 3 columns and about 17 rows.
   inverseOfMapSimis: Failed to find a nonzero minor. We now compute syzygies instead.
   If this doesn't terminate quickly, you may want to try increasing the option MinorsCount.
   isEmbedding: checking if the inverse map is a regular map
   -- used 3.01948 seconds
o9 = true
```

5. Functionality overlap with other packages

We note that our package has some overlaps in functionality with other packages.

While the Parametrization package [Boe10] focuses mostly on curves, it also includes a function called $\text{invertBirationalMap}$ which has the same functionality as $\text{inverseOfMap}$. On the other hand, these two functions were implemented somewhat differently and so sometimes one function can be substantially faster than the other.

The package Cremona [Sta16] focuses on very fast probabilistic computation in general cases and very fast deterministic computation for special kinds of maps from projective space. In particular, in Cremona,

- $\text{isBirational}$ gives a probabilisitic answer to the question of whether a map between varieties is birational. Furthermore, if the source is projective space, then $\text{degreeOfRationalMap}$ with $\text{MathMode=>true}$ can give a deterministic answer that is frequently faster than what our package can provide with $\text{isBirationalMap}$.
- $\text{invertBirMap}$ gives a very fast computation of the inverse of a birational map if the source is projective space and the map has maximal linear rank. If you pass this function a map not from projective space, then it calls a modified, faster version of $\text{invertBirationalMap}$ originally from Parametrization. Even in some cases with maximal linear rank, our $\text{inverseOfMap}$ function appears to be quite competitive however.
5.1. **Speed comparisons.** First we do comparisons using examples with maximal linear rank where *Cremona* excels. These examples were run using version 2.0 of *Cremona* and version 0.21 of *RationalMaps* and using version 1.9.1 of Macaulay2.

Indeed, in some of those examples with maximal linear rank, *Cremona* is substantially faster.

```plaintext
i1 : loadPackage "Cremona"; loadPackage "RationalMaps";
13: ringP20=QQ[t_0..t_20];
i4 : phi=map(ringP20,ringP20,{t_10*t_15-t_9*t_16+t_6*t_20,t_10*t_14-t_8*t_16+t_5*t_20,t_9*t_14-t_8*t_15+t_4*t_20,
t_6*t_14-t_5*t_15+t_4*t_16,t_11*t_13-t_16*t_17+t_15*t_18-t_14*t_19+t_7*t_20,t_10*t_12-t_13*t_16+t_8*t_19,
t_9*t_12-t_13*t_17+t_7*t_18-t_14*t_19+t_6*t_19+t_13*t_16,t_10*t_11-t_2*t_13-t_7*t_17+t_5*t_18-t_4*t_19,
t_9*t_11-t_3*t_16+t_2*t_17-t_1*t_18+t_0*t_19,t_8*t_11-t_3*t_15+t_1*t_16,t_7*t_11-t_3*t_14+t_0*t_16,
t_6*t_11-t_2*t_15+t_1*t_16,t_5*t_11-t_2*t_14+t_0*t_16,t_4*t_11-t_1*t_14+t_0*t_15,t_3*t_8-t_5*t_9+4*t_10,t_3*t_6-t_5*t_8+4*t_9+10,t_3*t_5-t_5*t_7+4*t_8+10,4*t_10,t_2*t_4-t_1*t_5+t_0*t_6});
i5 : time inverseOfMap(phi, AssumeDominant=>true, Verbose=>false)-- Function from "RationalMaps"
-- used 1.12522 seconds
i6 : time invertBirMap phi -- Function from "Cremona"
-- used 0.0789913 seconds
i7 : isSameMap(o5, o6) -- Function from "RationalMaps"
o7 = true

However, sometimes our function is substantially faster even in examples with maximal linear rank.

```plaintext
i1 : ZZ/33331[t_0..t_6];
i2 : loadPackage "Cremona"; loadPackage "RationalMaps";
i4 : phi=toMap minors(3,matrix{{t_0..t_4},{t_1..t_5},{t_2..t_6}});
i5 : J=kernelComponent(phi,2);
i6 : phi=toMap(phi,Dominant=>J);
i7 : time inverseOfMap(phi, AssumeDominant=>true, Verbose=>false) -- Function from "RationalMaps"
-- used 0.126314 seconds
i8 : time invertBirMap phi; -- Function from "Cremona"
-- used 0.375856 seconds
i9 : isSameMap(o7, o8)
o9 = true

Note both these examples were taken directly from the documentation of *Cremona*. On the other hand, when the source and target are not projective space, or the map does not have maximal linear rank, in our experience, our functions appear to be faster. Here is an example where the source and target are not projective space.

```plaintext
i1 : needsPackage "Divisor"; loadPackage "RationalMaps"; loadPackage("Cremona", Reload=>true);
i5 : C = QQ[x,y,z]/(x^4+x^2*y*z+y^4+z^3*x);
i6 : Q = ideal(y,x+z); --a point on our curve
i7 : (f3 = mapOntoImage(mapToProjectiveSpace(12*divisor(Q))));
i8 : (time g1 = inverseOfMap(f3, Verbose=>false)); -- Function from "RationalMaps"
-- used 1.1996 seconds
i9 : (time g2 = invertBirMap(f3)); -- Function from "Cremona"
-- used 13.7981 seconds
i10 : (g3 = map(target g1, source g2, sub(matrix g2, target g1)));
i11 : isSameMap(g3, g1)
o11 = true

We also include an example where the map does not have maximal linear rank.

```plaintext
i1 : loadPackage "RationalMaps"; loadPackage "Cremona";
i3 : Q=QQ[x,y,z,u,ul];
i4 : phi=map(Q,Q,matrix{{x^5,y*x^4+z^4+y^5,5*x*z^5,u*x^4*t^5}});
i6 : RingMap Q <-- Q
i5 : (time inverseOfMap(phi, AssumeDominant=>true,CheckBirational=>false, MinorsCount=>50000));
-- Function from "RationalMaps"
-- used 5.80638 seconds
i9 : (time h = invertBirMap(phi)); -- Function from "Cremona"
-- used 71.8998 seconds
i10 : isSameMap(o5, h)
o10 = true
```
REFERENCES


