## The Test Ideals package for Macaulay2

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- Consider rings $R$ of characteristic $p>0$.
- No resolution of singularities (in general).
- Kunz proved:


## Theorem (Kunz)

$R$ is regular if and only if Frobenius is flat.

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- Because we are working with computers, domain finite type over $\mathbb{F}_{q}$.
- Kunz says Frobenius is flat if and only if $R^{1 / p^{e}}$ is locally free over R.
- We can weaken being locally free.


## Definition (Hochster-Roberts, Mehta-Ramanathan)

$R$ is $F$-pure if and only if $R \rightarrow R^{1 / p^{e}}$ splits.

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Checking F-purity can be pretty easy.

- Fedder's Criterion. $R=S / I, S$ is polynomial.


## Theorem (Fedder)

$R$ is $F$-nure at m if and only if $[p]: / \notin \mathrm{m}[p]$.

- If $I=(f)$, then $I^{[p]}: I=\left(f^{p-1}\right)$. $(B O A R D)$
- For example.

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```
i5 : \(S=Z Z / 7[x, y, z] ;\)
i6 : \(\mathrm{f}=\mathrm{x}^{\wedge} 3+\mathrm{y}^{\wedge} 3+\mathrm{z}^{\wedge} 3\);
i8 : isSubset (ideal (f^6), ideal ( \(\mathrm{x}^{\wedge} 7, \mathrm{y}^{\wedge} 7, \mathrm{z}^{\wedge} 7\) ))
o8 = false
```


## Macaulay2

We have written a package TestIdeals.m2 that computes whether a ring (or pair) is:

- F-regular
- F-rational
- F-injective
- Test ideals
- F-pure thresholds (with FThresholds.m2).


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- Fedder's criterion works because maps

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\phi_{R}: R^{1 / p^{e}} \rightarrow R
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come from maps

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\phi_{S}: S^{1 / p^{e}} \rightarrow S
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such that $\phi_{S}\left(I^{1 / p^{e}}\right) \subseteq I$.

- In fact,
$I^{\left[p^{e}\right]}: I \cong\left\{\phi \in \operatorname{Hom}_{S}\left(S^{1 / p^{e}}, S\right) \mid \phi\left(I^{1 / p^{e}}\right) \subseteq I\right\}$
- Translates questions on $P$ into questions in polynomial ring
- Note $\left\{\phi_{R} \neq 0\right\} \leftrightarrow\{\Delta \geq 0 \mathbb{Q}$-log boundary $\}$. (BOARD)
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## Frobenius trace

One more big tool.

- There exists $\Phi: S^{1 / p^{e}} \rightarrow S$.

- Other monomials to 0 .
- $\Phi$ generates $\mathrm{Hom}_{\mathrm{s}}\left(\mathrm{S}^{1 / p^{e}}, S\right)$.
- $\phi$ is Grothendieck dual to Frobenius.
- $\Phi\left(J^{1 / p^{e}}\right) \subseteq l$ if and only if

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f^{\left[p^{6}\right]} \subseteq J .
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## Theorem (Fedder restated)


defines locus where $R=S / I$ is not $F$-pure.

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$$

defines locus where $R=S / I$ is not $F$-pure.

We compute some Macaulay2 examples. $\Phi(J)$ is called the Frobenius root of $J$.

```
i12 : I =ideal (x^3 + \(\left.y^{\wedge} 3+z^{\wedge} 3\right)\);
i13 : frobeniusRoot(1, I^7 : I)
o13 = ideal 1
i14 : isFPure(S/I)
o14 = true
i15 : J = ideal (x^4+y^4+z^4);
i16 : frobeniusRoot (1, J^7 : J)
                                    2 2 2
o16 = ideal (z , y*z, x*z, y , x*y, x )
i19 : isFPure(S/J)
o19 = false
```

```
i20 : T = ZZ/5[a,b,c,d,e];
i21 : B = ZZ/5[x,y];
\(i 22: f=\operatorname{map}\left(B, T,\left\{x^{\wedge} 4, x^{\wedge} 3 * y, x^{\wedge} 2 * y^{\wedge} 2, x * y^{\wedge} 3, y^{\wedge}\right.\right.\)
                                432234
o22 = map (B, T, \(\{x, x y, x y, x * y, y\})\)
o22 : RingMap B <--- T
i23 : I = ker f
                                    2 2
o23 = ideal (d - c*e, c*d - b*e, b*d - a*e, c - a
o23 : Ideal of T
i24 : isFPure(T/I)
o24 = true
```


## $F$-regularity and test ideals

- Analog of KLT.


## Definition

$R$ is strongly $F$-regular if for every (interesting ${ }^{a}$ ) $c \in R$, there is some $e$ and $\phi: R^{1 / p^{e}} \rightarrow R$ so that $\phi\left(c^{1 / p^{e}}\right)=1$.
${ }^{a}$ In Jacobian ideal is good enough

- If translated by Fedder's methods,

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```
i3 : \(S=Z Z / 7[x, y, z]\);
i4 : R = S/ideal ( \(x^{\wedge} 2-y * z\) )
i5 : isFRegular(R);
o5 = true
\(i 20: A=Z Z / 7[x, y, z] /\left(y^{\wedge} 2 * z-x *(x-z) *(x+z)\right)\);
i21 : \(C=Z Z / 7[a, b, c, d, e, f]\);
\(i 22: g=\operatorname{map}\left(A, C,\left\{x^{\wedge} 2, x * y, x * z, y^{\wedge} 2, ~ y * z, ~ z \wedge 2\right\}\right)\)
i23 : I = ker g;
i26 : isFRegular (C/I);
o26 = false
```

- We can only show that $\mathbb{Q}$-Gorenstein rings are not $F$-regular.

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- The QGorensteinIndex=>infinity option can prove a non- $\mathbb{Q}$-Gorenstein ring is $F$-regular.

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i3 : S = ZZ/7[x,y,z];
i4 : R = S/ideal (x^2-y*z)
i6 : h = y;
i7 : isFRegular(1/2, y)
07 = false
i8 : isFRegular(1/3, y)
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- The pair $\left(R, h^{1 / 2}\right)$ is not $F$-regular but $\left(R, h^{1 / 3}\right)$ is.
- The FThresholds package can even compute F-pure thresholds.

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## F-rationality

- Analog of rational singularities.
- Implies (pseudo-)rational singularities in a fixed characteristic.
- Here's our definition:


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- $\left(c^{1 / p^{e}} \cdot \omega_{R^{1 / p^{e}}}\right) \xrightarrow{F^{e} \text {-dual }} \omega_{R}$ surjects.

The tricky part is writing the map:

$$
F-\text { dual }: \omega_{R^{1 / p^{e}}} \rightarrow \omega_{R}
$$

- Trick (Katzman) is to embed $\omega_{R}$ as an ideal in $R$.
- Extend $F$ - dual to $\phi_{R}: R^{1 / p^{e}} \rightarrow R$.
- Extend further to $\phi_{S}: S^{1 / p^{e}} \rightarrow S .(R=S / I)$
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Here is an example of an $F$-rational (but not $F$-regular) ring.

```
i8 : S = ZZ/3[a,b,c,d,t]; m = 4; n = 3;
i11 : M = matrix{ {a^2 + t^m, b, d},
                                {c, a^2, b^^n-d} };
    2 3
o11 : Matrix S <--- S
i12 : I = minors(2, M);
i13 : R = S/I;
i14 : isFRational(R)
o14 = true
```

Appeared in work of Anurag Singh (deform F-regularity)

## Characteristic zero applications

Characteristic $p>0$ conclusions imply results in characteristic zero.

## Theorem (Ma-•)

Suppose $R$ is a ring of mixed characteristic finite type over $\mathbb{Z}$. Suppose $p \in \mathbb{Z}$ is a regular element and $Q \subseteq R$ is a prime not containing any nonzero prime of $\mathbb{Z}$ so that $(p)+Q \neq R$.

If $R / p R$ is F-rational, then $R_{Q}=R_{Q} \otimes \mathbb{Q}$ has rational singularities.

- Analogous statement for log terminal/F-regular singularities, if the $\mathbb{Q}$-Gorenstein not divisible by $p$.
- Not known for log canonical/F-pure singularities (need mixed char inversion of adjunction).


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## $F$-injective

We can also study $F$-injective singularities (analog of Du Bois).
Definition
$R$ is $F$-injective if

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H^{-i} \omega_{R^{1 / p}}^{\bullet} \rightarrow H^{-i} \omega_{R}^{\bullet}
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surjects for all $i$.

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- Example

```
i10 : R = ZZ/[x,y,z]/ideal( (x^3+ y^3+z^3);
i11 : isFInjective(R)
o11 = true
```


## Test ideals

We can compute test ideals too. Including of pairs.

- In a $\mathbb{Q}$-Gorenstein ring.
- $\tau\left(R, f^{t}\right)$ equals sum of images of maps
$c$ as before.
- We use it to check F-regularity.
- Trick is stabilize image sums above.
- Can compute parameter test modules and parameter test ideals too.


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## Example

```
i2 : R = ZZ/5[x,y];
i3 : \(f=y^{\wedge} 2-x^{\wedge} 3\);
    32
o3 \(=-x+y\)
i4 : testIdeal(4/5, f);
○4 = ideal (y, x)
i5 : testIdeal(4/5-1/10000, f)
o5 = ideal 1
```

- We can compute $\tau\left(R, f^{t-\epsilon}\right)$, which is used to compute jumping numbers and $F$-pure thresholds.
- Needs HSLGModule function.

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We conclude with a discussion of the FThresholds package.

- If $R$ is $F$-regular, $F$-pure threshold (FPT) is the smallest $t \geq 0$ where $\tau\left(R, f^{t}\right) \neq R$.
- We do a binary-style search to a certain depth.
- However, if $f$ is a special form, we have other algorithms.
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## Example FPT

```
i2 : R = ZZ/5[x,y,z]
i3 : \(f=x^{\wedge} 5-y^{\wedge} 6+x^{\wedge} 3 * z^{\wedge} 5+2 * z^{\wedge} 8\)
    35665
03 = x \(\mathrm{z}+2 \mathrm{z}-\mathrm{y}+\mathrm{x}\)
i4 : fpt(f)
    1
○4 = -
    5
```

FPT of the cusp (in a nonstandard form).
i1 : $R=Z Z / 7[x, y]$
i4 : fpt ((x+y)^3- $\left.y^{\wedge} 2\right)$
5
$\circ 4=-$
6

## You can go to:

http://www.math.utah.edu/~schwede/M2.html to try it yourself!

