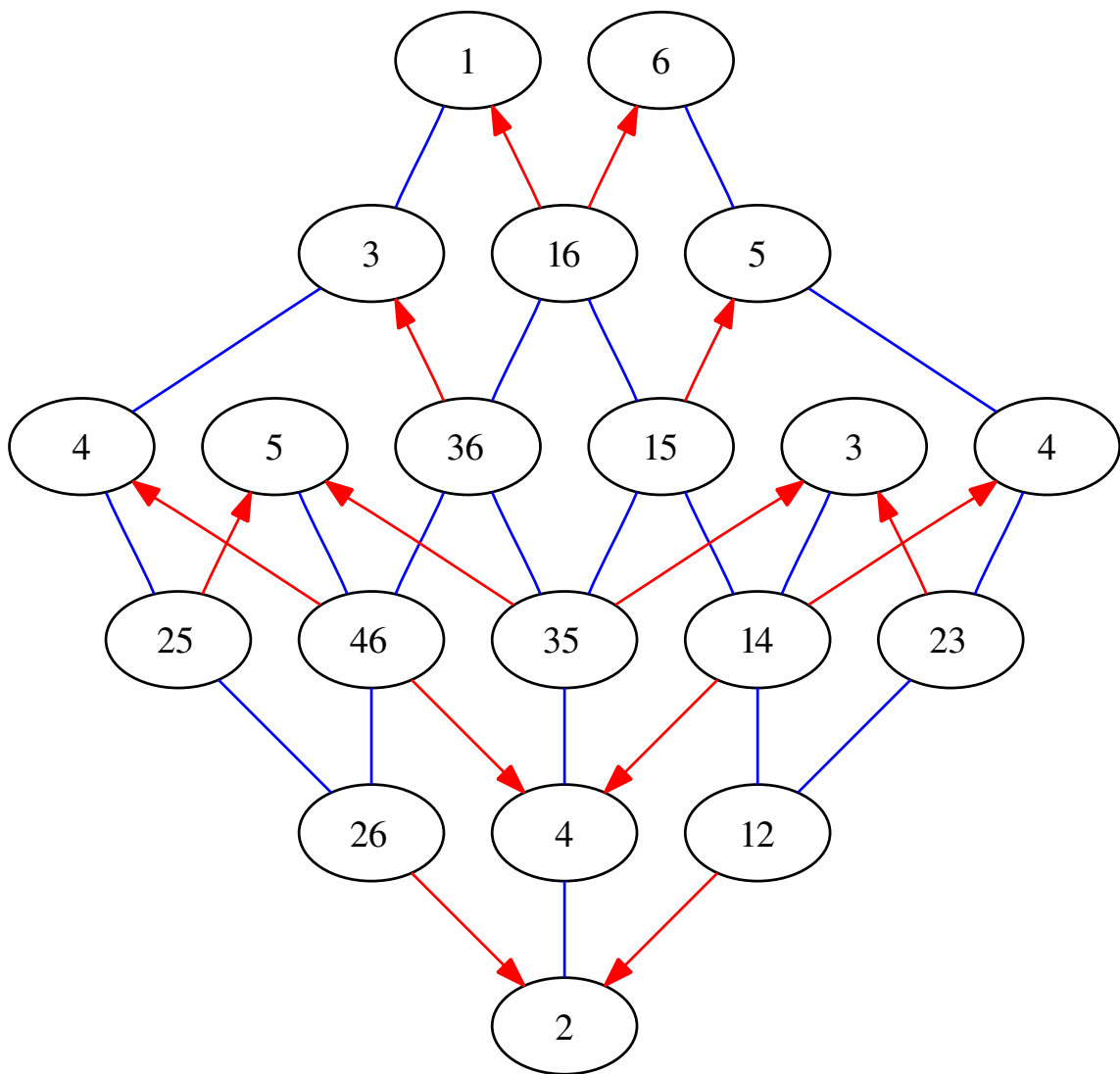


The Combinatorics of  $W$ -Graphs  
Computational Theory of Real Reductive Groups Workshop  
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JOHN STEMBRIDGE  $\langle$ jrs@umich.edu $\rangle$



## 1. What is a $W$ -Graph?

Let  $(W, S)$  be a Coxeter system,  $S = \{s_1, \dots, s_n\}$ .

For us,  $W$  will always be a finite Weyl group.

Let  $\mathcal{H} = \mathcal{H}(W, S)$  = the associated Iwahori-Hecke algebra over  $\mathbb{Z}[q^{\pm 1/2}]$ .  
 $= \langle T_1, \dots, T_n \mid (T_i - q)(T_i + 1) = 0, \text{ braid relations} \rangle$ .

DEFINITION. An  $S$ -labeled graph is a triple  $\Gamma = (V, m, \tau)$ , where

- $V$  is a (finite) vertex set,
- $m : V \times V \rightarrow \mathbb{Z}[q^{\pm 1/2}]$  (i.e., a matrix of edge-weights),
- $\tau : V \rightarrow 2^S = 2^{[n]}$ .

NOTATION. Write  $m(u \rightarrow v)$  for the  $(u, v)$ -entry of  $m$ .

Let  $M(\Gamma) = \text{free } \mathbb{Z}[q^{\pm 1/2}]\text{-module with basis } V$ .

Introduce operators  $T_i$  on  $M(\Gamma)$ :

$$T_i(v) = \begin{cases} qv & \text{if } i \notin \tau(v), \\ -v + q^{1/2} \sum_{u: i \notin \tau(u)} m(v \rightarrow u)u & \text{if } i \in \tau(v). \end{cases}$$

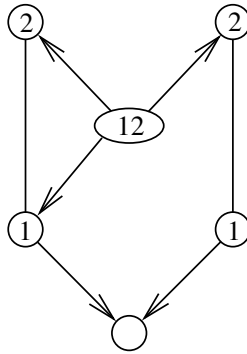
DEFINITION (K-L).  $\Gamma$  is a  $W$ -graph if this yields an  $\mathcal{H}$ -module.

NOTE:  $(T_i - q)(T_i + 1) = 0$  (always), so  $W$ -graph  $\Leftrightarrow$  braid relations.

$$T_i(v) = \begin{cases} qv & \text{if } i \notin \tau(v), \\ -v + q^{1/2} \sum_{u:i \notin \tau(u)} m(v \rightarrow u)u & \text{if } i \in \tau(v). \end{cases} \quad (1)$$

REMARKS.

- Kazhdan-Lusztig use  $T_i^t$ , not  $T_i$ .
- Restriction: for  $J \subset S$ ,  $\Gamma|_J := (V, m, \tau|_J)$  is a  $W_J$ -graph.
- At  $q = 1$ , we get a  $W$ -representation.
- However, braid relations at  $q = 1 \not\Rightarrow W$ -graph:



- If  $\tau(v) \subseteq \tau(u)$ , then (1) does not depend on  $m(v \rightarrow u)$ .

CONVENTION. WLOG, all  $W$ -graphs we consider will be **reduced**:

$$m(v \rightarrow u) = 0 \text{ whenever } \tau(v) \subseteq \tau(u).$$

DEFINITION. A  **$W$ -cell** is a strongly connected  $W$ -graph.

For every  $W$ -graph  $\Gamma$ ,  $M(\Gamma)$  has a filtration whose subquotients are cells.

Typically, cells are not irreducible as  $\mathcal{H}$ -reps or  $W$ -reps.

However (Gyoja, 1984): every irrep of  $W$  may be realized as a  $W$ -cell.

## 2. The Kazhdan-Lusztig $W$ -Graph

$\mathcal{H}$  has a distinguished basis  $\{C_w : w \in W\}$  (the Kazhdan-Lusztig basis).

The left and right action of  $T_i$  on  $C_w$  is encoded by a  $W \times W$ -graph

$$\Gamma_{LR} = (W, m, \tau_{LR}):$$

- $\tau_{LR}(v) = \tau_L(v) \cup \tau_R(v)$ , where
 
$$\tau_L(v) = \{i_L : \ell(s_i v) < \ell(v)\}, \quad \tau_R(v) = \{i_R : \ell(v s_i) < \ell(v)\}$$
- $m$  is determined by the Kazhdan-Lusztig polynomials:

$$m(u \rightarrow v) = \begin{cases} \mu(u, v) + \mu(v, u) & \text{if } \tau_{LR}(u) \not\subseteq \tau_{LR}(v), \\ 0 & \text{if } \tau_{LR}(u) \subseteq \tau_{LR}(v), \end{cases}$$

where  $\mu(u, v) = \text{coeff. of } q^{(\ell(v) - \ell(u) - 1)/2} \text{ in } P_{u,v}(q) (= 0 \text{ unless } u \leq v)$ .

REMARKS.

- Hard to compute  $\mu(x, y)$  without first computing  $P_{x,y}(q)$ .
- Restricting  $\Gamma_{LR}$  to the left action (say) yields a  $W$ -graph  $\Gamma_L$ .
- The cells of  $\Gamma_L$  decompose the regular representation of  $\mathcal{H}$ .
- Every two-sided K-L cell  $C$  has a “special”  $W$ -irrep associated to it that occurs with positive multiplicity in each left K-L cell  $\subset C$ .
- In type  $A$ , every left cell is irreducible, and the partition of  $W$  into left and right cells is given by the Robinson-Schensted correspondence.

The representation theory connection (complex groups):

- K-L “Conjecture”:  $P_{w_0 x, w_0 y}(1) = \text{multiplicity of } L_y \text{ in } M_x$ ,
- Vogan:  $\mu(x, y) = \dim \text{Ext}^1(M_x, L_y)$ ,

where  $M_w = \text{Verma module with h.w. } -w\rho - \rho$ ,  $L_w = \text{simple quotient}$ .

### 3. $W$ -Graphs for Real Groups

There is a similar story for real groups:

Let  $K$  = complexification of the maximal compact subgroup of  $G_{\mathbb{R}}$ .

Irreps can be assigned to  $K$ -orbits on  $G/B$  (complex case:  $W \approx B \backslash G/B$ ).

There are K-L-V polynomials  $P_{x,y}(q)$  generalizing K-L polynomials.

The top coefficients  $\mu(x,y)$  encode a  $W$ -graph structure  $\Gamma_K$  on  $K \backslash G/B$ .

Usually  $\Gamma_K$  will break into more than one component (block).

EXAMPLE. In the split real form of  $E_8$ , the  $W$ -graph has 6 blocks, the largest of which has 453,060 vertices and 104 cells.

Cells for real groups often appear as cells of  $\Gamma_L$ . Not always.

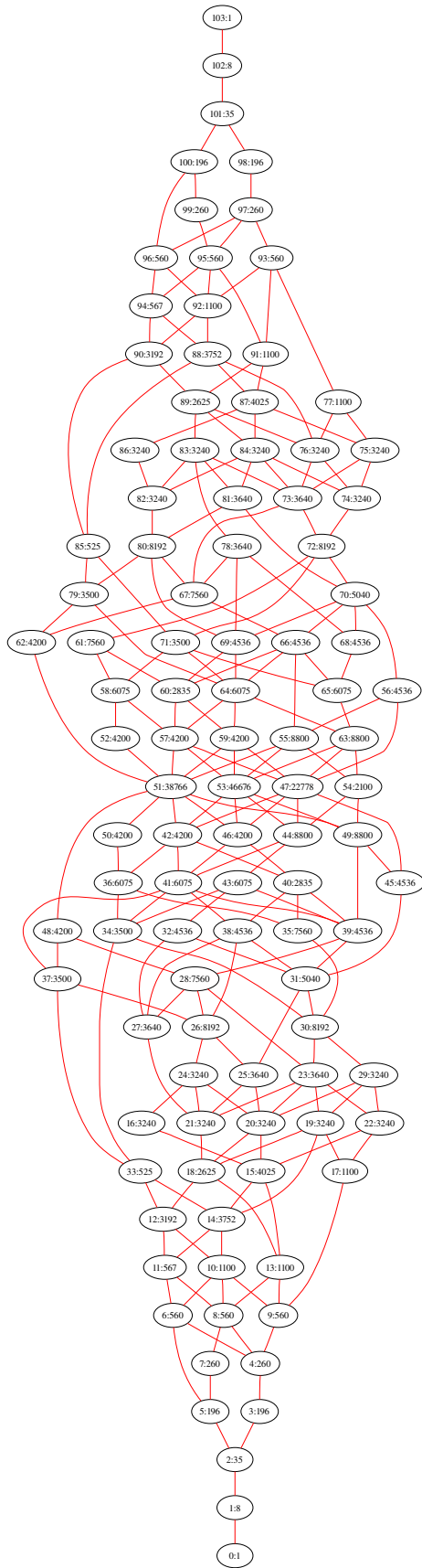
EXAMPLE.  $G_{\mathbb{C}}$  as a real group.

It has Weyl group  $W \times W$ ; its  $W \times W$ -graph is  $\Gamma_{LR}$ .

MAIN POINTS.

- The most basic constraints on these  $W$ -graphs are sufficiently strong that combinatorics alone can lend considerable insight into the structure of  $W$ -graphs and cells for real and complex groups.

- Sufficiently deep understanding of the combinatorics can yield constructions of  $W$ -cells without needing to compute K-L(-V) polynomials.



## 4. Admissible $W$ -Graphs

Three observations about the  $W$ -graphs for real and complex groups:

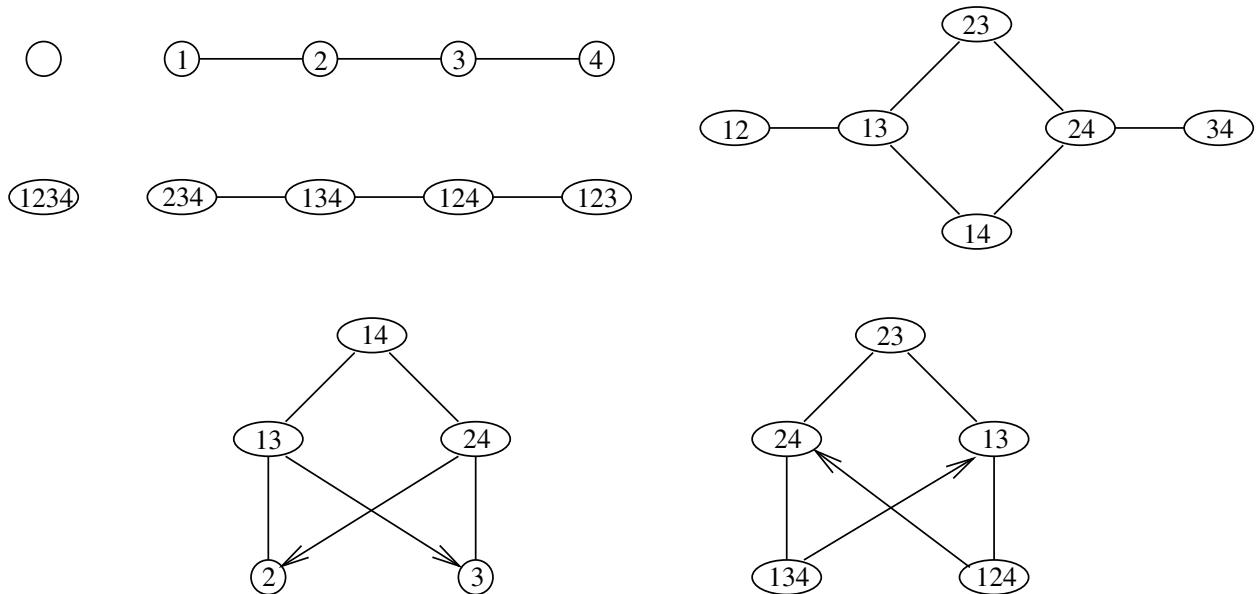
- (1) They have nonnegative integer edge weights.
- (2) They are **edge-symmetric**; i.e.,

$$m(u \rightarrow v) = m(v \rightarrow u) \quad \text{if } \tau(u) \not\subseteq \tau(v) \text{ and } \tau(v) \not\subseteq \tau(u).$$

- (3) They are bipartite. (If  $\mu(u, v) \neq 0$ , then  $\ell(u) \neq \ell(v) \pmod{2}$ .)

DEFINITION. A  $W$ -graph is **admissible** if it satisfies (1)–(3).

EXAMPLE. The admissible  $A_4$ -cells:

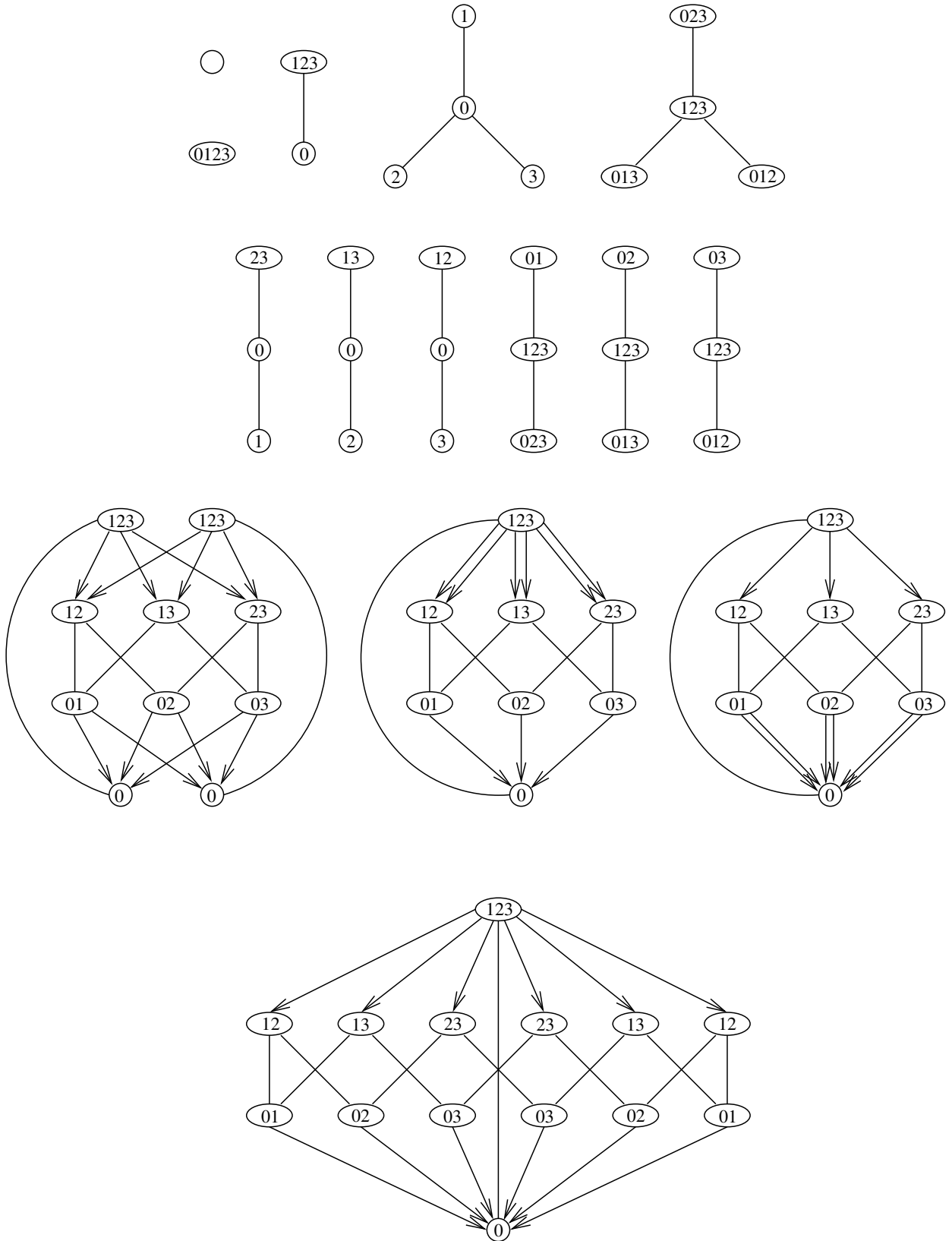


All of these are K-L cells; none are synthetic.

QUESTION. *Is every admissible  $A_n$ -cell a K-L cell?* (Confirmed for  $n \leq 9$ .)

CAUTION. McLarnan-Warrington: Interesting things happen in  $A_{15}$ .

The admissible  $D_4$ -cells (three are synthetic):





## 5. Some Interesting Questions

PROBLEM 1. *Are there finitely many admissible  $W$ -cells?*

- Confirmed for  $A_1, \dots, A_9, B_2, B_3, D_4, D_5, D_6, E_6, G_2$ .
- What about  $W_1 \times W_2$ -cells? More about this in Part II.

PROBLEM 2. *Classify/generate all admissible  $W$ -cells.*

PROBLEM 3. *How can we identify which admissible cells are synthetic?*

- Example: If  $\Gamma$  contains no “special”  $W$ -rep, then  $\Gamma$  is synthetic.
- Regard non-synthetics as closed under Levi restriction.

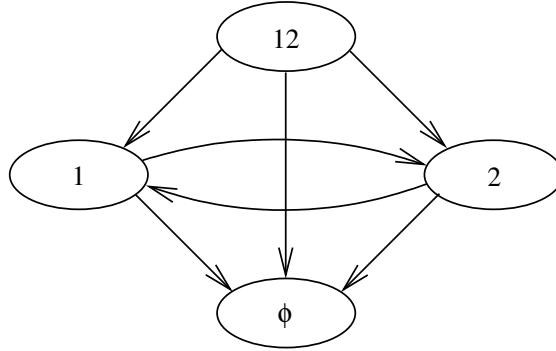
PROBLEM 4. *Understand “compressibility” of  $W$ -cells and  $W$ -graphs.*

- A given  $W$ -cell or  $W$ -graph should be reconstructible from a small amount of data. (Possible approaches: binding and branching rules.)

## 6. The Admissible Cells in Rank 2

Consider  $W = I_2(p)$  (dihedral group),  $2 \leq p < \infty$ .

Given an  $I_2(p)$ -graph, partition the vertices according to  $\tau$ :



Focus on non-trivial cells:  $\tau(v) = \{1\}$  or  $\{2\}$  for all  $v \in V$ .

The edge weight matrix will then have a block structure:  $m = \begin{bmatrix} 0 & B \\ A & 0 \end{bmatrix}$ .

The conditions on  $m$  are as follows:

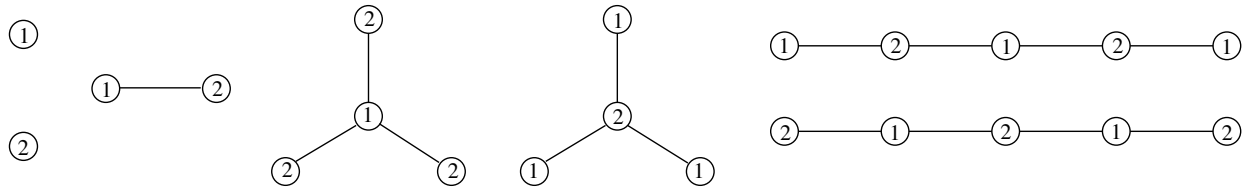
- $p = 2$ :  $m = 0$ .
- $p = 3$ :  $m^2 = 1$  (i.e.,  $AB = BA = 1$ ).
- $p = 4$ :  $m^3 = 2m$ .
- $p = 5$ :  $m^4 - 3m^2 + 1 = 0$ .
- $\vdots$

REMARKS.

- If we assume only  $\mathbb{Z}$ -weights, no classification is possible (cf.  $p = 3$ ).
- Edge symmetry  $\Leftrightarrow m = m^t$ .
- When  $p = 3$ , edge weights  $\in \mathbb{Z}^{\geq 0} \Rightarrow$  edge symmetry, but not in general.

THEOREM 1. A 2-colored graph is an admissible  $I_2(p)$ -cell iff it is a properly 2-colored A-D-E Dynkin diagram whose Coxeter number divides  $p$ .

EXAMPLE. The Dynkin diagrams with Coxeter number dividing 6 are  $A_1$ ,  $A_2$ ,  $D_4$ , and  $A_5$ . Therefore, the (nontrivial) admissible  $G_2$ -cells are



REMARK. The nontrivial K-L cells for  $I_2(p)$  are paths of length  $p - 2$ .

FACT (Vogan; cf. Problem 3). In a Levi restriction of type  $B_2 = I_2(4)$ , all nontrivial  $B_2$ -cells in  $\Gamma_K$  are paths of length 2.

*Proof Sketch.* Let  $\Gamma$  be any properly 2-colored graph.

Let  $\phi_p(t)$  be the Chebyshev polynomial such that  $\phi_p(2 \cos \theta) = \frac{\sin p\theta}{\sin \theta}$ .

Then  $\Gamma$  is an  $I_2(p)$ -cell  $\Leftrightarrow \phi_p(m) = 0$

$\Leftrightarrow m$  is diagonalizable with eigenvalues  $\subset \{2 \cos(\pi j/p) : 1 \leq j < p\}$ .

Now assume  $\Gamma$  is admissible ( $m = m^t$ ,  $\mathbb{Z}^{\geq 0}$ -entries).

If  $\Gamma$  is an  $I_2(p)$ -cell, then  $2 - m$  is positive definite.

Hence,  $2 - m$  is a (symmetric) Cartan matrix of finite type.

Conversely, let  $A$  be any Cartan matrix of finite type (symmetric or not).

Then the eigenvalues of  $A$  are  $2 - 2 \cos(\pi e_j/h)$ , where  $e_1, e_2, \dots$  are the exponents and  $h$  is the Coxeter number.  $\square$

## 7. Combinatorial Characterization

What are the graph-theoretic implications of the braid relations?

THEOREM 2. *An admissible  $S$ -labeled graph is a  $W$ -graph if and only if the following properties are satisfied:*

- *the Compatibility Rule,*
- *the Simplicity Rule,*
- *the Bonding Rule, and*
- *the Polygon Rule.*

THE COMPATIBILITY RULE (applies to all  $W$ -graphs for all  $W$ ):

*If  $m(u \rightarrow v) \neq 0$ , then*

*every  $i \in \tau(u) - \tau(v)$  is bonded to every  $j \in \tau(v) - \tau(u)$ .*

Necessity follows from analyzing commuting braid relations.

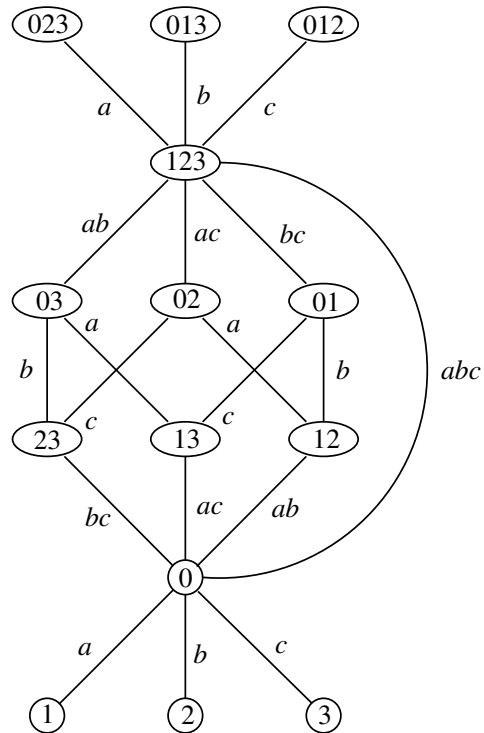
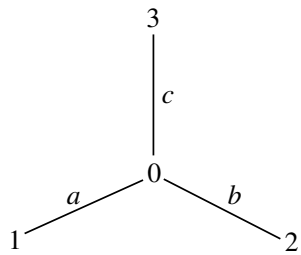
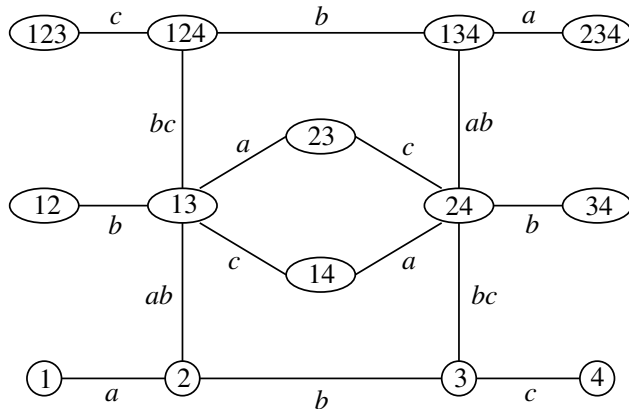
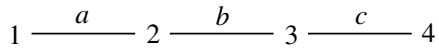
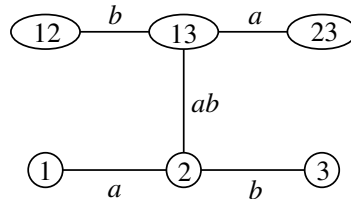
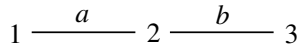
REFORMULATION: Define the **compatibility graph**  $\text{Comp}(W, S)$ :

- vertex set  $2^S = 2^{[n]}$ ,
- edges  $I \rightarrow J$  when

$I \not\subseteq J$  and every  $i \in I - J$  is bonded to every  $j \in J - I$ .

Compatibility means that  $\tau : \Gamma \rightarrow \text{Comp}(W, S)$  is a graph morphism.

# Compatibility graphs for $A_3$ , $A_4$ , and $D_4$



THE SIMPLICITY RULE:

Every edge  $u \rightarrow v$  is either

- an **arc**:  $\tau(u) \supsetneq \tau(v)$  (and there is no edge  $v \rightarrow u$ ), or
- a **simple edge**:  $m(u \rightarrow v) = m(v \rightarrow u) = 1$

Necessity follows from Theorem 1.

THE BONDING RULE:

If  $s_i s_j$  has order  $p_{ij} \geq 3$ , then the cells of  $\Gamma|_{\{i,j\}}$  must be

- singletons with  $\tau = \emptyset$  or  $\tau = \{i, j\}$ , and
- *A-D-E* Dynkin diagrams with Coxeter number dividing  $p_{ij}$ .

Necessity again follows from Theorem 1.

EXAMPLE. If  $p_{ij} = 3$ , then the nontrivial cells in  $\Gamma|_{\{i,j\}}$  are  $\{i\} - \{j\}$ .

Equivalently (for bonds with  $p_{ij} = 3$ ): if  $i \in \tau(u)$ ,  $j \notin \tau(u)$  then there is a unique vertex  $v$  adjacent to  $u$  such that  $i \notin \tau(v)$ ,  $j \in \tau(v)$ .

REMARK. The Compatibility, Simplicity, and Bonding Rules suffice to determine all admissible  $A_3$ -cells.

THE POLYGON RULE:

[Compare with G. Lusztig, *Represent. Theory* **1** (1997), Prop. A.4.]

Define

$$V^{ij} := \{v \in V : i \in \tau(v), j \in \tau(v)\},$$

$$V_j^i := \{v \in V : i \in \tau(v), j \notin \tau(v)\},$$

$$V_{ij} := \{v \in V : i \notin \tau(v), j \notin \tau(v)\}.$$

A path  $u \rightarrow v_1 \rightarrow \cdots \rightarrow v_{r-1} \rightarrow v$  is *alternating of type  $(i, j)$*  if

$$u \in V^{ij}, v_1 \in V_j^i, v_2 \in V_i^j, v_3 \in V_j^i, v_4 \in V_i^j, \dots, v \in V_{ij}.$$

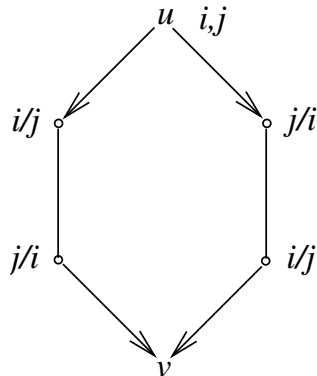
Set  $N_{ij}^r(u, v) := \sum m(u \rightarrow v_1)m(v_1 \rightarrow v_2) \cdots m(v_{r-1} \rightarrow v)$

(sum over all  $r$ -step alternating paths of type  $(i, j)$ ).

Then:

$$N_{ij}^r(u, v) = N_{ji}^r(u, v) \quad \text{for } 2 \leq r \leq p_{ij}.$$

EXAMPLE. 3-step alternating paths



REMARK. The Polygon Rule is **quadratic** in the arc weights.

## 8. Direct Products

Does the classification of admissible  $W_1 \times W_2$ -cells reduce to  $W_1$  and  $W_2$ ?

Not obviously. Not all cells are direct products.

Let  $\Gamma = (V, m, \tau_1 \cup \tau_2)$  be an admissible  $W_1 \times W_2$ -graph.

FACT. Every edge  $u \rightarrow v$  has one of three flavors:

- Type 1:  $\tau_1(u) \not\subseteq \tau_1(v)$ ,  $\tau_2(u) = \tau_2(v)$
- Type 2:  $\tau_1(u) = \tau_1(v)$ ,  $\tau_2(u) \not\subseteq \tau_2(v)$
- Type 12:  $\tau_1(u) \supsetneq \tau_1(v)$ ,  $\tau_2(u) \supsetneq \tau_2(v)$

Type 2 edges (and no others) are deleted when restricting  $\Gamma$  to  $W_1$ .

Hence,  $\tau_2$  is constant on  $W_1$ -cells.

KEY QUESTION. *Are there no arcs **between** cells in the  $W_1$ -restriction of a  $W_1 \times W_2$ -cell  $\Gamma$ ?*

True for two-sided K-L cells. If true for a general  $W_1 \times W_2$ -cell  $\Gamma$ , then

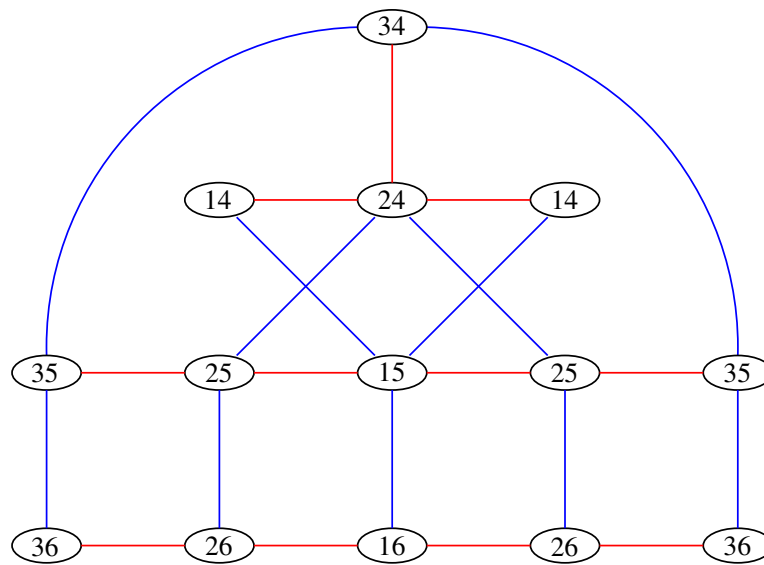
- Type 12 edges cannot exist within  $\Gamma$ .
- Every  $W_1$ -cell in  $\Gamma$  meets every  $W_2$ -cell.
- Bounds the number admissible cells for  $W_1 \times W_2$  in terms of  $W_1, W_2$ .
- Every  $W_1$ -cell in  $\Gamma$  has the same  $\tau_1$ -support.

Even if the answer is negative, something weaker is true.

FACT. The  $\tau_1$ -support of  $\Gamma$  equals the  $\tau_1$ -support of an admissible  $W_1$ -cell.



An admissible (K-L)  $B_3 \times B_3$ -cell



## 9. A Strategy for Resolving the Key Question

Consider two properties of an arbitrary admissible  $W$ -graph  $\Gamma = (V, m, \tau)$ :

PROPERTY A. *If  $\Gamma_1$  and  $\Gamma_2$  are cells of  $\Gamma$  such that  $\Gamma_1 < \Gamma_2$  in the induced partial order, then  $\tau(\Gamma_1) \neq \tau(\Gamma_2)$ .*

PROPERTY B. *If  $\Gamma_1$  and  $\Gamma_2$  are cells of  $\Gamma$  such that  $\Gamma_1 < \Gamma_2$  in the induced partial order and  $\tau(\Gamma_1) = \tau(\Gamma_2)$ , then there is a third cell  $\Gamma_3$  such that  $\Gamma_1 < \Gamma_3 < \Gamma_2$  and  $\tau(\Gamma_3) \not\subseteq \tau(\Gamma_1) = \tau(\Gamma_2)$ .*

- (Easy) Property A implies Property B.
- Property B affirmatively resolves the Key Question.
- Property A holds for the left K-L graph  $\Gamma_L$ . False in general.
- Property B has been confirmed for all low-rank admissible cells.

N.B. If Property B holds for  $W_1$ , then the Key Question has an affirmative answer for all  $W_1 \times W_2$ -cells, for all choices of  $W_2$ .

## 10. Support Families

It is natural to partition  $W$ -cells into families according to their  $\tau$ -support.

Any two left K-L cells either

- belong to the same two-sided cell, and
- have the same  $\tau$ -support, and
- contain the same “special”  $W$ -irrep,

or

- belong to distinct two-sided cells, and
- have unequal  $\tau$ -support, and
- have no  $W$ -irreducibles in common.

NOTE. The  $\tau$ -support of an admissible  $W$ -cell

- need not match the  $\tau$ -support of a left K-L cell, and
- need not contain a special  $W$ -irrep (a synthetic marker).

QUESTION. *For each  $\tau$ -support  $\mathcal{T} \subset 2^S$ , is there a  $W$ -irrep  $\sigma = \sigma(\mathcal{T})$  such that every admissible  $W$ -cell with  $\tau$ -support  $\mathcal{T}$  contains a copy of  $\sigma$ ?*

Assuming the Key Question has an affirmative answer, if  $\Gamma_1, \dots, \Gamma_l$  are  $W$ -cells that appear in some admissible  $W \times W'$ -cell for some  $W'$ , then they must have a  $W$ -irrep in common.

## 11. Molecular Components of $W$ -Graphs

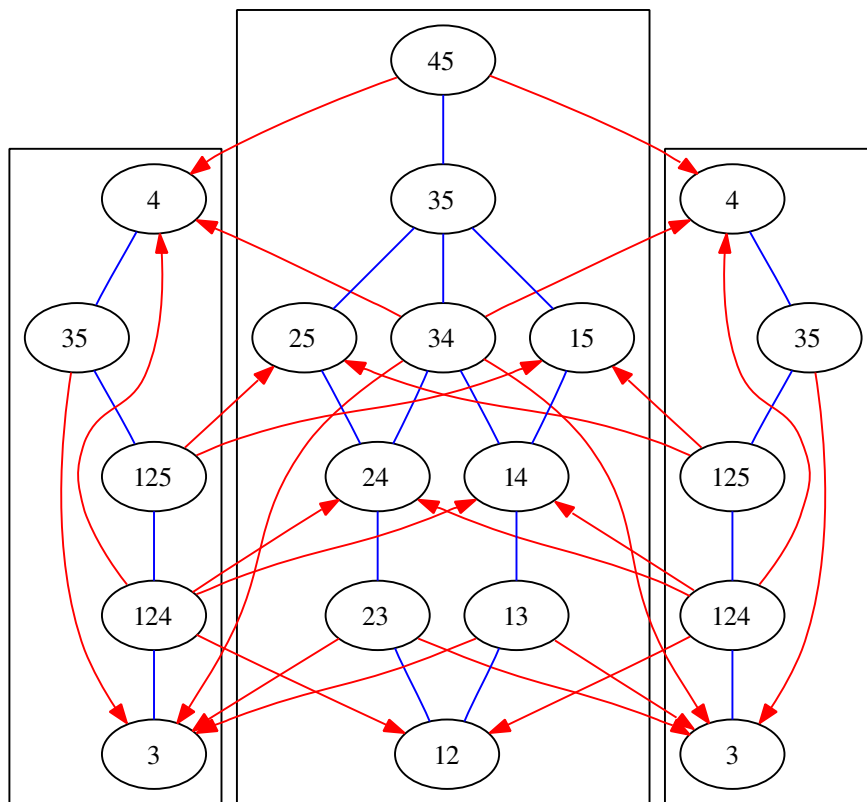
Recall the Simplicity Rule: every edge  $u \rightarrow v$  is either

- an **arc**:  $\tau(u) \supsetneq \tau(v)$  (and there is no edge  $v \rightarrow u$ ), or
- a **simple edge**:  $m(u \rightarrow v) = m(v \rightarrow u) = 1$

DEFINITION. A **molecular component** of an admissible  $W$ -graph  $\Gamma$  is a subgraph whose simple edges form a single connected component.

REMARK. All K-L cells in type  $A$  have only one molecular component.

A  $D_5$ -cell with three molecular components:



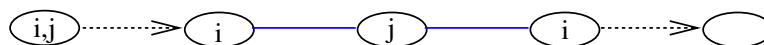
Classification strategy: first classify molecules, then classify all of the ways they may be glued together into (admissible) cells.

## 12. Synthesizing Molecules

IDEA #1: We can “easily” generate  $S$ -labeled graphs that satisfy the Compatibility, Simplicity, and Bonding Rules. No arc worries.

ISSUE: There are too many.

Need the Polygon Rule. Recall that it involves alternating  $(i, j)$ -paths:



FACT. Let  $(u, v, r, i, j)$  be an instance of the Polygon Rule

(initial point  $u$ , terminal point  $v$ , path length  $r$ ). Then

- if  $r = 2$  and there is  $k \in \tau(v) - \tau(u)$ , or
- if  $r = 3$  and there is  $k, l \in \tau(v) - \tau(u)$  such that  $k$  is not bonded to  $i$  and  $l$  is not bonded to  $j$ , or
- if  $r \geq 3$  and there is  $k \in \tau(v) - \tau(u)$  such that  $k$  is not bonded to  $i$  or  $j$ ,

then the resulting constraint is **linear** in weights of arcs.

An alternating path with only one arc can only involve the molecular components containing the two endpoints.

CONCLUSION: These instances of the Polygon Rule can be imposed **locally**.

So: add the **Local** Polygon Rule as a constraint on molecular components.

### 13. Stable Molecules

DEFINITION. An  $S$ -labeled graph that satisfies the Compatibility, Simplicity, Bonding, and Local Polygon Rules is **molecular**.

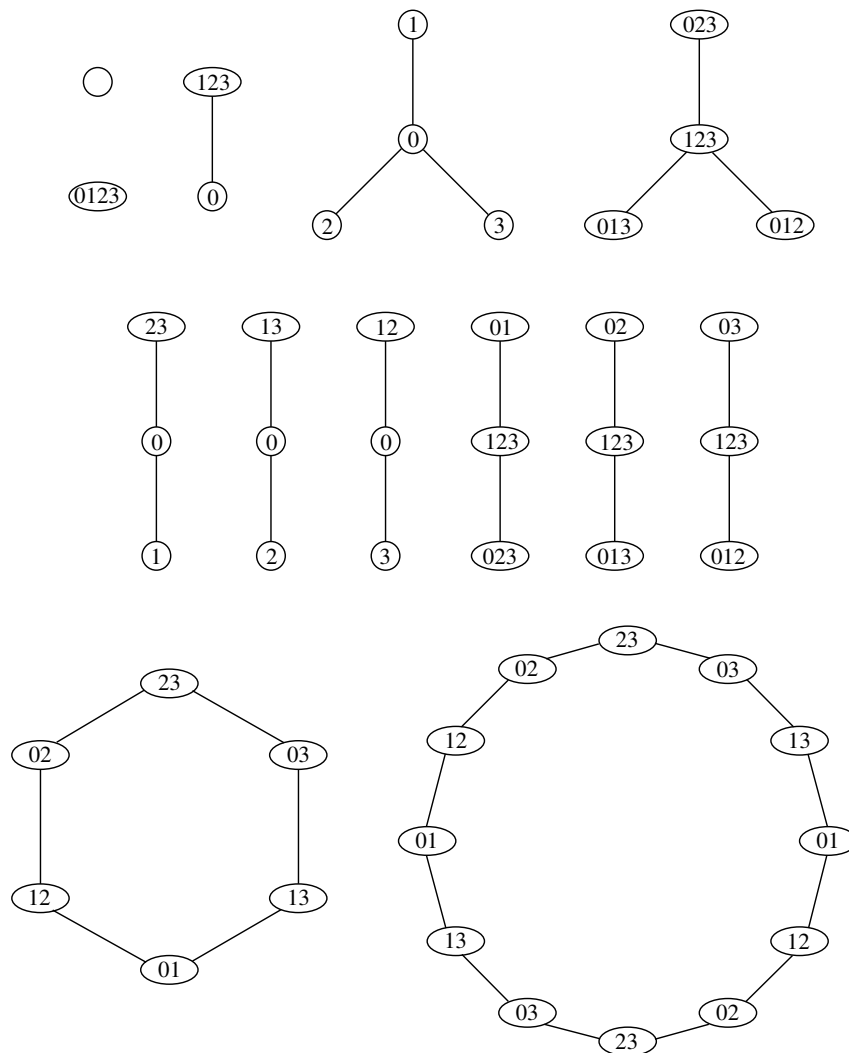
- If it has only one molecular component, it is a **molecule**.
- If it occurs in some admissible  $W$ -graph, it is **stable**.

For  $n \leq 9$ , the  $A_n$ -molecules are precisely the K-L cells!

There do exist unstable molecules. Sometimes infinitely many.

But in all cases so far, they have manageable structure.

The stable  $D_4$ -molecules:



## 14. Binding Spaces

Given a list of (stable)  $W$ -molecules, what are all of the (stable) molecular graphs that can be obtained by binding them together?

Focus on pairs of molecules, say  $\Gamma_1$  and  $\Gamma_2$ .

Regard every inclusion  $\tau(v_1) \supsetneq \tau(v_2)$  as a potential arc  $v_1 \rightarrow v_2$ .

DANGER: Admissible graphs must be bipartite!

Work in a category of **molecules-with-parity**:

every vertex has a parity, edges connect vertices of opposite parity.

Molecules are connected, so each affords two parity choices.

NOTATION:  $\Gamma \mapsto -\Gamma$  (parity-reversing operator).

DEFINITION. A **binding space** is the vector space  $B(\Gamma_1 \rightarrow \Gamma_2)$  of weight assignments for arcs  $\Gamma_1 \rightarrow \Gamma_2$  that satisfy the Local Polygon Rule.

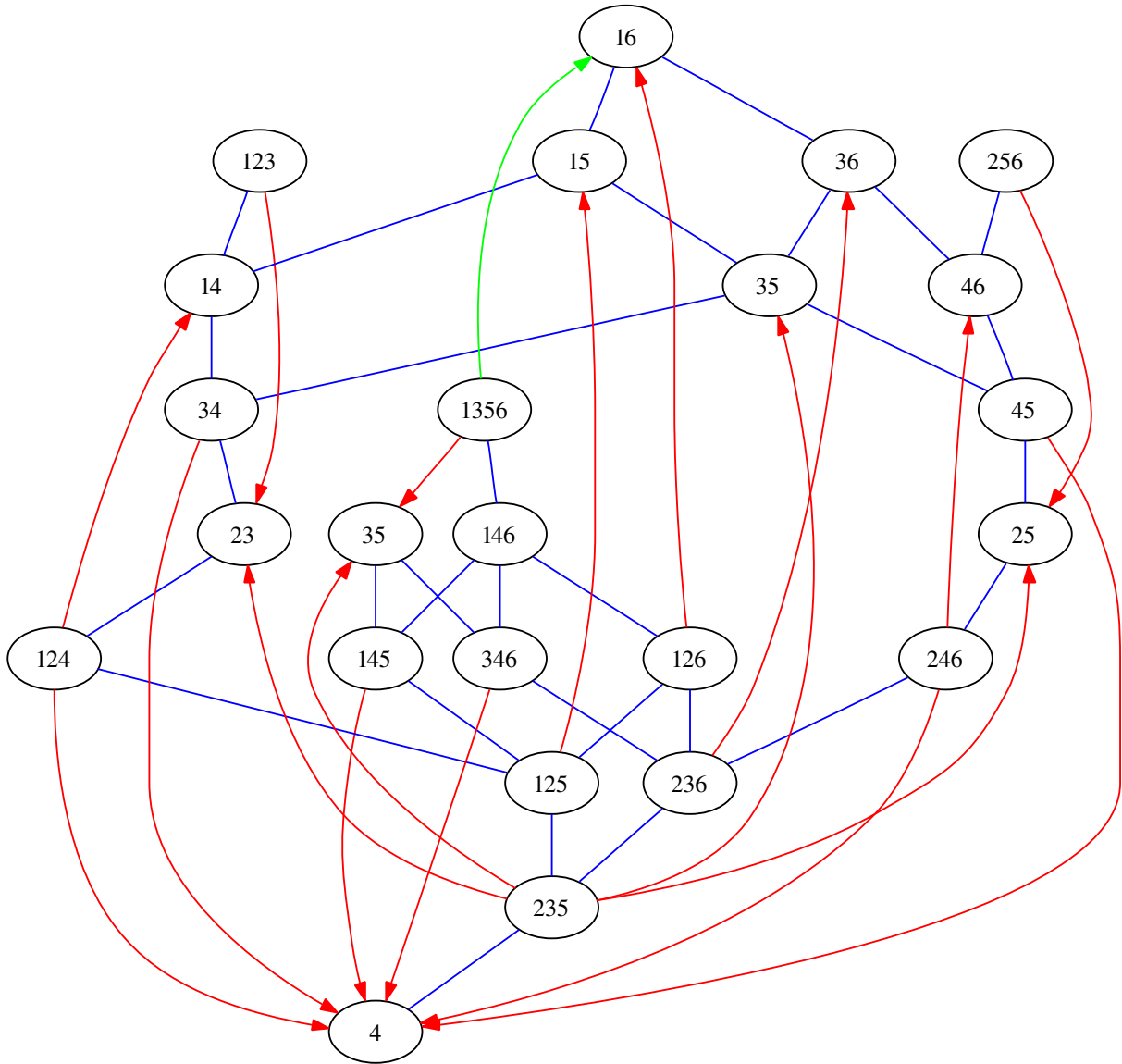
- Depends only on the simple edges of  $\Gamma_1$  and  $\Gamma_2$ .
- In simply-laced cases (at least), there is no torsion.
- Often,  $\dim B(\Gamma_1 \rightarrow \Gamma_2) = 0$  or  $1$ .
- Self-binding:  $B(\Gamma \rightarrow \Gamma)$  (even),  $B(\Gamma \rightarrow -\Gamma)$  (odd).

DEFINITION. A binding is **stable** if it occurs in some admissible  $W$ -graph.

NOTE. Each  $W$ -molecule  $\Gamma$  also has an **internal** binding space  $B(\Gamma)$ .

- $B(\Gamma)$  may be identified with an affine translate of  $B(\Gamma \rightarrow \Gamma)$ .

EXAMPLE. An  $E_6$ -molecule with  $\dim B(\cdot) = 1$ :





## 15. Binding Families

DEFINITION. The **bindability graph**  $BG(W)$  is the directed graph with

- vertices corresponding to  $W$ -molecules
- edges  $\Gamma \rightarrow \Gamma'$  whenever  $\dim B(\pm\Gamma \rightarrow \pm\Gamma') > 0$ .

Similarly, there is a **stable** bindability graph  $BG_{\text{st}}(W)$ .

Break  $BG(W)$  or  $BG_{\text{st}}(W)$  into strongly connected components.

NOTE. Every admissible  $W$ -cell is obtained by binding together one or more  $W$ -molecules from some strongly connected component of  $BG(W)$ .

- The same holds for  $BG_{\text{st}}(W)$ .
- This provides another natural way to partition  $W$ -cells into families.
- The resulting binding families of  $W$ -cells are partially ordered.
- For every admissible  $W$ -graph  $\Gamma$ , there is an order-preserving map

$$\phi(\Gamma) : \{\text{cells of } \Gamma\} \rightarrow \{\text{binding families of } W\text{-cells}\}.$$

QUESTIONS.

- Is  $\phi(\Gamma_L)$  surjective (i.e., does every binding family contain a K-L cell)?
- Are the fibers of  $\phi(\Gamma_L)$  unions of 2-sided cells?
- Is every binding family a union of support families?
- Are the binding families mutually orthogonal (as  $W$ -modules)?
- Is there a “special” molecule that occurs in every  $W$ -cell in a family?

Binding families of  $W$ -cells for  $W = D_5, D_6,$  and  $E_6$ .

