

February 7, 2006

Lesson #9 (10) Derived functors (cont)

(11) Adjoint functors and left/right exactness

Theorem: Let  $F: A \rightarrow B$  be a right exact functor between two abelian categories. If  $A$  has enough projectives, then

- (1) Each  $L_n F$  is an additive functor.
- (2)  $L_* F = \{L_n F\}_{n \geq 0}$  form a homological  $\delta$ -functor.
- (3)  $L_* F$  is a universal  $\delta$ -functor.

Proof: (1) We show first that  $L_n F$  is a functor.

$$L_n F(\text{id}^A) := H_n F(\text{id}^B) = H_n(\text{id}^{F(P_\bullet)}) = \text{id}^{H_n F(P_\bullet)} = \text{id}^{H_n F(A)}$$

Let  $f: A' \rightarrow A$  and  $g: A \rightarrow A''$  be morphisms in  $A$ , and let  $P_\bullet \rightarrow A'$ ,  $P_\bullet \rightarrow A$  and  $P_\bullet \rightarrow A''$  be proj resolutions.

If  $f_\bullet: P'_\bullet \rightarrow P_\bullet$  and  $g_\bullet: P_\bullet \rightarrow P''_\bullet$  are extensions of  $f$  and  $g$  then  $g_\bullet f_\bullet: P'_\bullet \rightarrow P''_\bullet$  is an extension of  $gf: A' \rightarrow A''$ .

Therefore,  $L_n F(gf) := H_n F(g_\bullet f_\bullet) = H_n F(g_\bullet) H_n F(f_\bullet) = \dots$  since  $H_n$  and  $F$  are functors.  $\dots = L_n F(g) L_n F(f)$ .

$L_n F$  is additive // Let  $f_1, f_2: A' \rightarrow A$  be two morphisms and

Let  $(f_1)_\bullet: P'_\bullet \rightarrow P_\bullet$  and  $(f_2)_\bullet: P'_\bullet \rightarrow P_\bullet$  be two liftings; then

$(f_1)_\bullet + (f_2)_\bullet$  is a lifting of  $f_1 + f_2$ . Therefore,

$$L_n F(f_1 + f_2) := H_n F((f_1)_\bullet + (f_2)_\bullet) = H_n F((f_1)_\bullet) + H_n F((f_2)_\bullet) = L_n F(f_1) + L_n F(f_2)$$

(2) Let  $0 \rightarrow A' \rightarrow A \rightarrow A'' \rightarrow 0$  be an exact sequence in  $A$ .

By Horseshoe lemma there exist projective resolutions such that

$$\begin{array}{ccccccc} 0 & \rightarrow & A' & \rightarrow & A & \rightarrow & A'' \rightarrow 0 \\ & & \uparrow & & \uparrow & & \uparrow \\ 0 & \rightarrow & P'_\bullet & \rightarrow & P_\bullet & \rightarrow & P''_\bullet \rightarrow 0 \end{array}$$

exact rows.

Since each  $0 \rightarrow P'_n \rightarrow P_n \rightarrow P''_n \rightarrow 0$  split, we have an exact sequence  $0 \rightarrow F(P'_n) \rightarrow F(P_n) \rightarrow F(P''_n) \rightarrow 0$  for each  $n \in \mathbb{Z}$ . (it splits in  $B$ )

Therefore,  $0 \rightarrow F(P'_\bullet) \rightarrow F(P_\bullet) \rightarrow F(P''_\bullet) \rightarrow 0$  is an exact sequence of

complexes.

We consider the long exact sequence of homologies and obtain

$$\dots \rightarrow L_n F(A) \rightarrow L_n F(A') \rightarrow L_n F(A'') \xrightarrow{d_n} L_{n+1} F(A) \rightarrow \dots$$

The map  $d_n$  is natural  $\forall n$ .  
 Let  $0 \rightarrow A' \rightarrow A \rightarrow A'' \rightarrow 0$   
 $f' \downarrow \quad f \downarrow \quad f'' \downarrow$   
 $0 \rightarrow B' \rightarrow B \rightarrow B'' \rightarrow 0$

be a morphism of short exact

$$0 \rightarrow B' \rightarrow B \rightarrow B'' \rightarrow 0$$

sequences.

We consider liftings of the morphisms  $f'$  and  $f''$  to proj resolutions

$$0 \rightarrow P' \rightarrow P \rightarrow P'' \rightarrow 0$$

$$\uparrow \quad \uparrow \quad \uparrow$$

$$*f' \quad *f \quad *f''$$

$$0 \rightarrow Q' \rightarrow Q \rightarrow Q'' \rightarrow 0$$

ETS to construct  $f_0$  a lifting of  $f$  such that the diagram of projective resolutions is commutative. Indeed, it remains commutative after applying  $\mp$  and we that  $\mp$  is a homological  $\delta$ -functor.

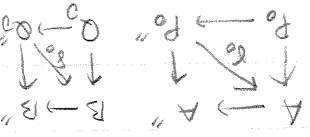
The construction of  $f_0$  is done by induction.

The differential of  $P$  is given by  $d_P = \begin{pmatrix} d_P \\ \theta_P \end{pmatrix}$ .

$Q$  is given by  $d_Q = \begin{pmatrix} d_Q \\ \theta_Q \end{pmatrix}$ .

Check that  $f_0 = \begin{pmatrix} f_0' \\ \Delta \end{pmatrix} \begin{pmatrix} f_0'' \\ \Delta \end{pmatrix}$  where  $\Delta: P'' \rightarrow Q' \rightarrow Q'' \rightarrow 0$

satisfy:  $\Delta_{i+1} d_i'' - d_i' \Delta_i = \theta_i' f_i'' - f_i' \theta_i''$  and  $P'' \eta'' \Delta_0 = f_0'' - \tau_0 f_0''$



$\Delta$  is constructed by induction.

(3) Remark: If  $P$  is projective, then  $L_n F(P) = 0 \forall n > 0$ .

Let  $T^*$  be another homological  $\delta$ -functor and let  $\varphi_0: T^* \rightarrow \mp$

be given. We construct natural transformations  $\varphi_n: T_n \rightarrow L_n F$  that commute with  $\delta_n$ , by induction.

Construction of  $\varphi_n$ , when  $\varphi_i$  is already defined (commuting with  $\delta_i$ ) for all  $i < n$ :

Let  $A$  be in  $\mathcal{A}$  and choose  $P \twoheadrightarrow A$ ,  $P$  projective and consider the short exact sequence  $0 \rightarrow K \rightarrow P \rightarrow A \rightarrow 0$ . We obtain:

$$\begin{array}{ccccccc}
 \dots & \rightarrow & T_n(K) & \rightarrow & T_n(P) & \rightarrow & T_n(A) \xrightarrow{\delta_n} T_{n-1}(K) \xrightarrow{\beta} \dots \\
 & & & & \exists! \varphi_n(A) \downarrow & & \downarrow \varphi_n(K) \cong \downarrow \delta \\
 & & & & \textcircled{*} & & \\
 & & \dots & \rightarrow & L_n F(P) & \rightarrow & L_n F(A) \xrightarrow{i} L_n F(K) \xrightarrow{\alpha} \dots
 \end{array}$$

$\parallel$   
 $0$

$\exists! \varphi_n(A)$  s.t.  $\textcircled{*}$  is comm, since  $\alpha \varphi_n(K) \delta_n = \delta \beta \delta_n = 0$

$\Rightarrow \text{Im}(\varphi_n(K) \delta_n) \subseteq \text{Ker } \alpha = \text{Im } i \cong L_n(F(A))$ .

It remains to check:

- $\varphi_n(A)$  does not depend on the choice of  $P$ .
- $\varphi_n$  is a natural transformation.
- $\varphi_n$  commutes with  $\delta_n$ . ■

### Right Derived Functors

Let  $F: \mathcal{A} \rightarrow \mathcal{B}$  be a left exact functor and assume that  $\mathcal{A}$  has enough injectives.

Let  $A \in \mathcal{A}$  and  $A \rightarrow I^\bullet$  an injective resolution.

Definition: The right derived functors are defined by:

$$R^n F(A) := H^n(F(I^\bullet))$$

Theorem: With the above hypothesis the following hold.

- (1) Each  $R^n F$  is an additive functor.
- (2)  $R^* F = \{R^n F\}_{n \geq 0}$  forms a cohomological  $\delta$ -functor.
- (3)  $R^* F$  is a universal cohomological  $\delta$ -functor.

Proof: Consider  $F^{op}: \mathcal{A}^{op} \rightarrow \mathcal{B}^{op}$ .  $F^{op}$  becomes a right derived functor. Since  $\mathcal{A}^{op}$  has enough projectives,

Proof: in the next lesson

$$\text{Ext}_R^i(A, B) = {}_R\text{Hom}_R(-, B)(A) \cong {}_R\text{Hom}_R(A, -)(B) \quad i \geq 0.$$

Theorem: If  $A$  and  $B$  are  $R$ -modules, then:

$${}_R\text{Hom}_R(-, B) \text{ exists.}$$

It is a covariant left exact functor, therefore

Example: Let  $B$  be an  $R$ -module, and define  $F(-) = \text{Hom}_R(-, B)$ ,

$$R^i F \text{ exists.}$$

$F: A^{\text{op}} \rightarrow B$ ; it is left exact and  $A^{\text{op}}$  has enough injectives  $\Rightarrow$   
 (1)  $R^i F$  (2) is done similarly. The functor  $F$  can be considered as:

They are universal  $\mathcal{F}$ -functors.

- the left derived functors  $\{L_n F\}_{n \geq 0}$  such that  $L_0 F = F$ .
- (2) If  $F$  is right exact and  $\mathcal{F}$  has enough injectives, then one can define the right derived functors  $\{R^n F\}_{n \geq 0}$  such that  $R^0 F = F$ .
- (1) If  $F$  is left exact and  $\mathcal{F}$  has enough projectives, then one can define

Let  $F: A \rightarrow B$  be a contravariant functor.

Remark:

Derived contravariant functors.

$I$  injective.

Remark:  $\text{Ext}_R^0(A, B) = \text{Hom}(A, B)$ , and  $\text{Ext}_R^n(A, I) = 0 \quad \forall n > 0$  and

$$\text{Ext}_R^n(A, B) := {}_R\text{Hom}_R(A, -)(B), \quad n \geq 0$$

$\exists$  right derived functors.

Definition:  $F(-) = \text{Hom}_R(A, -)$  is a left exact functor, then  
 Let  $A$  be an  $R$ -module, and define:

$$R^n F(A) := (L_n F^{\text{op}})^{\text{op}}(A) = H^n F(I)$$

⑪ Adjoint functors and left/right exactness.

Theorem

Let  $\mathcal{A}, \mathcal{B}$  be two abelian categories, and let  $(L, R) : \mathcal{A} \rightarrow \mathcal{B}$  be an adjoint pair of additive functors. Then:

- (1)  $L$  is right exact. (In particular, if  $\mathcal{A}$  has enough projectives,  $\exists L_n(L)$ .)
- (2)  $R$  is left exact. (In particular, if  $\mathcal{A}$  has enough injectives,  $\exists R^n(R)$ .)

Proof: (2)  $\implies$  (1)  $(R^{op}, L^{op}) : \mathcal{A}^{op} \rightarrow \mathcal{B}^{op}$  is also an adjoint pair. By (2)  $L^{op}$  is left exact, hence  $L$  is right exact.

(2) Let  $0 \rightarrow B' \rightarrow B \rightarrow B'' \rightarrow 0$  be an exact sequence in  $\mathcal{B}$ .

We have the commutative diagram:

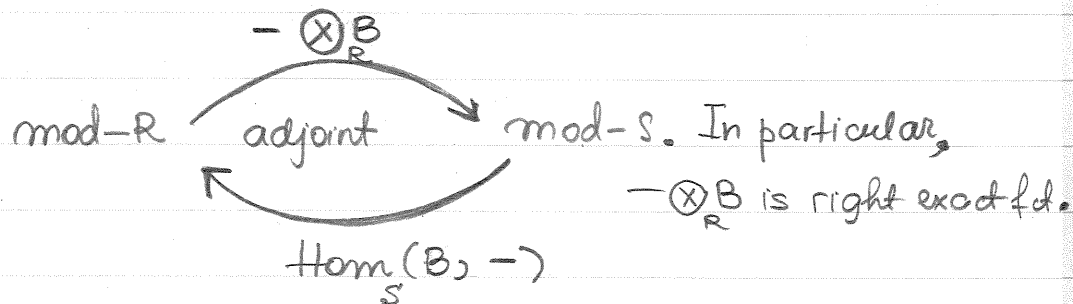
$$\begin{array}{ccccccc}
 0 & \longrightarrow & \text{Hom}_{\mathcal{B}}(L(A), B) & \longrightarrow & \text{Hom}_{\mathcal{B}}(L(A), B) & \longrightarrow & \text{Hom}_{\mathcal{B}}(L(A), B'') \\
 & & \downarrow \cong_{AB'} & \cong & \downarrow \cong_{AB} & & \downarrow \cong_{AB''} \\
 0 & \longrightarrow & \text{Hom}_{\mathcal{A}}(A, R(B)) & \longrightarrow & \text{Hom}_{\mathcal{A}}(A, R(B)) & \longrightarrow & \text{Hom}_{\mathcal{A}}(A, R(B''))
 \end{array}$$

The upper row is exact since  $\text{Hom}_{\mathcal{B}}(L(A), -)$  is left exact.

Next we apply the Yoneda lemma. ■

Proposition: Let  $R$  and  $S$  be two rings and let  ${}_R B_S$  be an  $R$ - $S$  bimodule.

Then:



(Sketch of)

- Proof: Check:
- $\text{Hom}_S(B, C)$  is a right  $R$ -module for all  $C$  right  $S$ -module
  - $A \otimes_R B$  is a right  $S$ -module for all  $A$  right  $R$ -mod

•  $\tau : \text{Hom}_S(A \otimes_R B, C) \xrightarrow{\cong} \text{Hom}_R(A, \text{Hom}_S(B, C))$

$$f \longleftarrow (a, \longleftarrow \tau f(a) : B \rightarrow C)$$

$$b \mapsto f(a \otimes b)$$

$$\tau^{-1} : a \otimes b \mapsto g(a)(b) \longrightarrow g$$

Definition: Let  $B$  be an  $R$ - $S$  bimodule (left  $R$ -module =

$R$ - $A$ - $B$ -bimodule). We define:

$$\text{Tor}_n^R(A, B) := L_n(- \otimes_R B)(A) \quad \text{they are } S\text{-modules.}$$

In particular,  $\text{Tor}_0^R(A, B) = A \otimes B$  and

$$\text{Tor}_n^R(P, B) = 0 \quad \forall n > 0 \text{ and } \forall P \text{ projective.}$$

since  $\text{Tor}_n^R(A, B) = H_n(P \otimes B)$ ,  $P \rightarrow A$  proj resolution.

Theorem:  $\text{Tor}_n^R(A, B) \cong L_n(- \otimes_R B)(A) \cong L_n(A \otimes_R -)(B)$ ,  $\forall n \geq 0$ .

Proof: in the next lesson.  $\square$