

April 25, 2006

Lesson #28

(33) $K(A)$ and $D(A)$ as triangulated categories

Definition: An additive category \mathcal{K} is called triangulated if the following hold:

- (1) $\exists T: \mathcal{K} \rightarrow \mathcal{K}$ an equivalence (translation functor)
- (2) \exists class of distinguished triangles (exact triangles) satisfying the axioms:

(TR1) $\forall u: A \rightarrow B \exists$ an exact triangle $(u, v, w): A \xrightarrow{u} B \xrightarrow{v} C \xrightarrow{w} TA$

- $(id_A, 0, 0)$ is an exact triangle
- $(u, v, w) \cong (u', v', w')$ then (u, v, w) exact $\iff (u', v', w')$ exact.

(TR2) (Rotation) (u, v, w) exact $\iff (v, w, -Tu)$ exact

(TR3) (Morphisms) If (u, v, w) and (u', v', w') are exact triangles

$$\begin{array}{ccccccc} A & \xrightarrow{u} & B & \xrightarrow{v} & C & \xrightarrow{w} & TA \\ \downarrow f & & \downarrow g & & \downarrow h & & \downarrow Tf \\ A' & \xrightarrow{u'} & B' & \xrightarrow{v'} & C' & \xrightarrow{w'} & TA' \end{array}$$

$\exists h$ st we get a morphism of triangles.

(TR4) (The octahedral axiom)

$$\begin{array}{ccccccc} A & \xrightarrow{u} & B & \xrightarrow{f} & C' & \xrightarrow{\gamma} & TA \\ B & \xrightarrow{v} & C & \xrightarrow{\alpha} & A' & \xrightarrow{\beta} & TB \\ A & \xrightarrow{uv} & C & \xrightarrow{y} & B' & \xrightarrow{\delta} & TA \end{array} \quad \text{exact triangles}$$

\exists exact triangle $C' \xrightarrow{f} B' \xrightarrow{g} A' \xrightarrow{(\gamma_f) i} \Sigma C'$

And the following hold: $\delta f = \gamma, u \delta = \beta g, \alpha = g \gamma, \gamma v = f \beta$

to be isomorphic in $K(A)$ to a strict triangle.

(TR2): Let $A \xrightarrow{u} B \xrightarrow{v} Cone(u) \xrightarrow{z} A[-1]$ be a strict triangle. We want $B \rightarrow Cone(u) \rightarrow A[-1] \xrightarrow{-u[-1]} B[-1]$ to be isomorphic in $K(A)$ to a strict triangle.

$$\begin{array}{ccccccc}
 A & \rightarrow & A & \rightarrow & Cone(id_A) & \rightarrow & A[-1] \\
 \parallel & & \parallel & & \parallel & & \parallel \\
 A & \rightarrow & A & \rightarrow & Cone(id_A) & \rightarrow & A[-1] \\
 \parallel & & \parallel & & \parallel & & \parallel \\
 A & \rightarrow & A & \rightarrow & Cone(id_A) & \rightarrow & A[-1] \\
 \parallel & & \parallel & & \parallel & & \parallel \\
 A & \rightarrow & A & \rightarrow & Cone(id_A) & \rightarrow & A[-1]
 \end{array}$$

$u = id_A$

(TR3): any u stays in an exact (strict) triangle. Any triangle isomorphic to an exact triangle is exact.

Any triangle isomorphic to a strict triangle is called an exact triangle. $[in K(A)]$ is up to homotopy equivalence.

$$\begin{array}{ccc}
 A_{n-1} & \xrightarrow{u} & A_n \\
 \oplus & \searrow & \oplus \\
 B_{n-2} & \xrightarrow{v} & B_{n-1}
 \end{array}$$

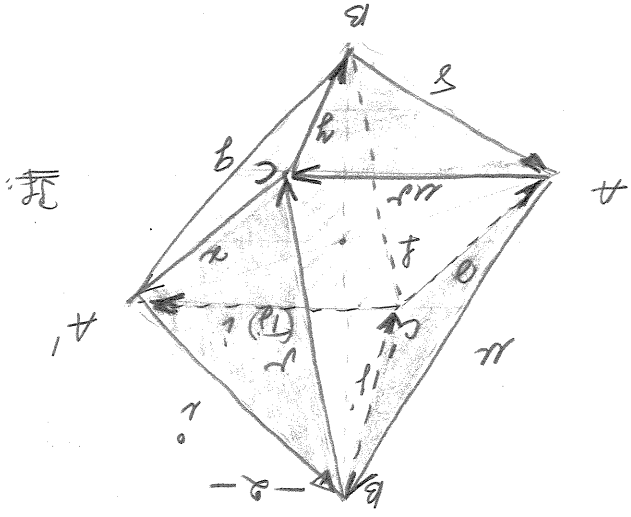
$Cone(u)_m = A_{m-1} \oplus B_m$

$$0 \rightarrow B \xrightarrow{v} Cone(u) \xrightarrow{z} A[-1] \rightarrow 0$$

strict triangles:

$\mathcal{K}(A)$ is an additive category. (see lesson #6). $TA := A[-1] \neq A \in \mathcal{K}(A)$ is an equivalence of categories.

Theorem: $\mathcal{K}(A)$ is a triangulated category. Proof:



Remark: $A \xrightarrow{u} B \xrightarrow{v} C \xrightarrow{w} TA$ exact triangle $\Rightarrow uv=0, vw=0$

$$\begin{array}{ccc}
 A & \xrightarrow{u} & B \\
 \uparrow & & \uparrow \\
 0 & \rightarrow & C = 0
 \end{array}$$

$$\begin{array}{ccccccc}
 B. & \xrightarrow{\psi} & \text{Cone}(u.) & \longrightarrow & \text{Cone}(y.) & \longrightarrow & B.[-1] \\
 \parallel & \parallel & \parallel & & \downarrow \theta & & \parallel \\
 B. & \xrightarrow{\varphi} & \text{Cone}(u.) & \xrightarrow{\delta} & A.[-1] & \xrightarrow{-u.[-1]} & B.[-1]
 \end{array}$$

$\exists \theta$. a homotopy equivalence. (Exercise Hw #5)

(TR3) We may assume that the two triangles are strict.

$$\begin{array}{ccccccc}
 A. & \xrightarrow{u.} & B. & \longrightarrow & \text{Cone}(u.) & \longrightarrow & A.[-1] \\
 f. \downarrow & \parallel & \downarrow g. & \parallel & \downarrow h. & \parallel & \downarrow \\
 A'. & \xrightarrow{u'.} & B'. & \longrightarrow & \text{Cone}(u'.) & \longrightarrow & A'.[-1]
 \end{array}$$

$h(a,b) := (f(a), g(b))$ the diagram is commutative.

(TRA) We may assume that all the triangles are strict.

$$\begin{array}{ll}
 A. \xrightarrow{u.} B. \xrightarrow{j} \text{Cone}(u.) \xrightarrow{\rho} A.[-1] & C' := \text{Cone}(u.) \\
 B. \xrightarrow{v.} C. \xrightarrow{x} \text{Cone}(v.) \xrightarrow{z} B.[-1] & A' := \text{Cone}(v.) \\
 A. \xrightarrow{v.u.} C. \xrightarrow{y} \text{Cone}(v.u.) \xrightarrow{\delta} A.[-1] & B' := \text{Cone}(v.u.) \\
 \text{Want: } C' \xrightarrow{f} B' \xrightarrow{g} A' \xrightarrow{(\rho) \circ i} C'[-1] &
 \end{array}$$

$$\begin{array}{ccccccc}
 \text{Cone}(u.) = C' & = \dots & \longrightarrow & A_n \oplus B_{n+1} & \longrightarrow & A_{n+1} \oplus B_n & \longrightarrow & A_{n-2} \oplus B_{n-2} & \longrightarrow \dots \\
 f. \downarrow & & & \text{id} \downarrow & & \downarrow \nu_{n+1} & \downarrow \text{id} & \downarrow \nu_{n+1} & \\
 \text{Cone}(v.u.) = B' & = \dots & \longrightarrow & A_n \oplus C_{n+1} & \longrightarrow & A_{n+1} \oplus C_n & \longrightarrow & \dots &
 \end{array}$$

$$\begin{array}{ccc}
 (a, \theta) & \longmapsto & (-d(a), d(\theta) - u.(a)) \quad d\nu = \nu d \\
 \downarrow & & \downarrow \nu \\
 (a, \nu(\theta)) & \longmapsto & (-d(a), d(\nu(\theta)) - \nu u.(a))
 \end{array}$$

$\boxed{\delta f = \rho}$ ✓

$\boxed{y \nu = f j}$

$y(\nu(\theta)) = (0, \nu(\theta))$
 $f j(\theta) = f((0, \theta)) = (0, \nu(\theta))$

$$\begin{array}{c}
 \text{cone}(v_0) = C' \rightarrow B_n \oplus C_{n+1} \rightarrow B_n \oplus C_n \rightarrow B_{n-2} \oplus C_{n-1} \rightarrow \dots \\
 \uparrow \quad \quad \quad \uparrow \quad \quad \quad \uparrow \quad \quad \quad \uparrow \\
 \text{cone}(v_1) = B'_1 \rightarrow A_n \oplus C_{n+1} \rightarrow A_{n-1} \oplus C_n \rightarrow A_{n-2} \oplus C_{n-1} \rightarrow \dots \\
 \uparrow \quad \quad \quad \uparrow \quad \quad \quad \uparrow \quad \quad \quad \uparrow \\
 \mu_n \downarrow \quad \quad \text{id} \downarrow \quad \quad \text{id} \downarrow \quad \quad \text{id} \downarrow \\
 \text{cone}(v_0) \rightarrow \text{cone}(v_1)
 \end{array}$$

$$\begin{array}{c}
 (a, c) \rightarrow (-d(a), d(c) - \gamma \cdot u(a)) \\
 \uparrow \\
 (u(a), c) \rightarrow (-d(u(a)), d(c) - \gamma \cdot u(a)) \\
 \uparrow \quad \quad \quad \uparrow \\
 \text{id} \downarrow \quad \quad \text{id} \downarrow \\
 (a, c) \rightarrow (u(a), c)
 \end{array}$$

$u d = d u$

$$u \gamma = \gamma u$$

$$u \gamma(a, c) = u(a)$$

$$\gamma u(a, c) = \gamma(u(a), c) = u(a)$$

$$x \gamma = \gamma x$$

$$\gamma(y(a, c)) = (0, c) = x(c)$$

$$\begin{array}{c}
 (T_1)^? : \text{cone}(f_0) \rightarrow \text{cone}(u) [-1] \\
 (b, c) \xrightarrow{?} (b, 0) \xrightarrow{?} (0, c) \\
 \uparrow \quad \quad \quad \uparrow \\
 (b, c) \xrightarrow{?} (b, 0) \xrightarrow{?} (0, c) \\
 \uparrow \quad \quad \quad \uparrow \\
 (b, c) \xrightarrow{?} (b, 0) \xrightarrow{?} (0, c)
 \end{array}$$

$$\begin{array}{c}
 C' \xrightarrow{f} B' \xrightarrow{g} A' \xrightarrow{h} C' [-1] \\
 \parallel \quad \quad \parallel \quad \quad \parallel \\
 C' \xrightarrow{f'} B' \xrightarrow{g'} A' \xrightarrow{h'} C' [-1] \\
 \uparrow \quad \quad \quad \uparrow \\
 (b, c) \xrightarrow{?} (b, 0) \xrightarrow{?} (0, c)
 \end{array}$$

γ is well defined

$$\begin{array}{c}
 A' = \dots \rightarrow B_m \oplus C_{m+1} \rightarrow B_m \oplus C_m \rightarrow B_{m-2} \oplus C_{m-1} \rightarrow \dots \\
 \uparrow \quad \quad \quad \uparrow \\
 (b, c) \xrightarrow{?} (b, 0) \xrightarrow{?} (0, c)
 \end{array}$$

$$\text{cone}(f_0) = A_{n+1} \oplus B_n \oplus A_n \oplus C_n \rightarrow A_{n-1} \oplus C_n \rightarrow A_{n-2} \oplus B_{n-2} \oplus A_{n-2} \oplus C_{n-1}$$

$$\begin{array}{c}
 (b, c) \rightarrow (-d(b), d(c) - \gamma(b)) \\
 \uparrow \\
 (0, b, 0, c) \rightarrow (0, d(b), 0, d(c)) \\
 \uparrow \\
 (0, b, 0, c) \rightarrow (0, d(b), 0, d(c)) \\
 \uparrow \\
 (0, b, 0, c) \rightarrow (0, d(b), 0, d(c))
 \end{array}$$

γ is a chain homotopy equivalence

$$\begin{array}{ccc} \varphi_* : \text{Cone}(f_*) = & \longrightarrow & (A_{n-2} \oplus B_{m-1}) \oplus (A_{n-1} \oplus C_n) \longrightarrow \\ \downarrow & & \downarrow \\ A' = & \longrightarrow & B_{m-1} \oplus C_n \longrightarrow \end{array}$$

$$(a_{n-2}, b, a_{n-1}, c)$$



$$(b + u(a_{n-1}), c)$$

$$\boxed{\varphi_* \delta = \text{id}_{A'}} \quad \varphi_* \delta(b, c) = \varphi(0, b, 0, c) = (b, c) \quad \checkmark$$

$$\boxed{\delta \varphi \sim \text{id}_{\text{Cone}(f_*)}} \quad \delta \varphi_*(a_{n-2}, b, a_{n-1}, c) = \delta_*(b + u(a_{n-1}), c) = (0, b + u(a_{n-1}), 0, c)$$

$$a_{n-1}, b, a_{n-1}, c$$

$$\begin{array}{ccccc} \text{Cone}(f_*) = & \longrightarrow & (A_{n-2} \oplus B_{m-1}) \oplus (A_{n-1} \oplus C_n) & \longrightarrow & \\ \downarrow & \swarrow s & \delta \varphi \downarrow \downarrow \text{id} & \swarrow s & \\ \text{Cone}(f_*) = & \longrightarrow & (A_{n-2} \oplus B_{m-1}) \oplus (A_{n-1} \oplus C_n) & \longrightarrow & \end{array}$$

Exercise*: Find s s.t. $ds + sd = \delta \varphi - \text{id}$.

Definition: Let \mathcal{K} be a triangulated category and \mathcal{A} an abelian category. An additive functor $H: \mathcal{K} \rightarrow \mathcal{A}$ is called

(*) homological functor if the following holds

$$(u, v, w) \text{ exact triangle} \implies \begin{array}{ccccccc} T^1 A. & \longrightarrow & A. & \longrightarrow & B. & \longrightarrow & C. \longrightarrow T A. \longrightarrow T B. \\ & & & & & & \longrightarrow T C. \longrightarrow T^2 A. \longrightarrow \dots \end{array}$$

induces an exact sequence

$$\dots \longrightarrow H(A.) \longrightarrow H(B.) \longrightarrow H(C.) \longrightarrow H(TA.) \longrightarrow H(TB.) \longrightarrow H(TC.) \longrightarrow H(T^2 A.) \longrightarrow \dots$$

$$H_i := H(T^i A) = H^i(A).$$

Example*: (1) $\text{Hom}_{\mathcal{K}}(X, -)$ is a cohomological functor.

(2) $\{H^0(-)\}$ is a cohomological functor.

Definition: If $H: \mathcal{K} \rightarrow \mathcal{A}$ is a (*) homological functor

$S = \{s \in \mathcal{K} \mid H(s) \text{ isom} \neq i\} \cup \{i\}$ called system arising from a homological functor. (i.e. $H(Ts)$ isom $\neq i$)

Theorem* If S is a system arising from a homological functor, then

- (1) S is a multiplicative system
- (2) \mathcal{K} triangulated category $\Rightarrow \mathcal{K}$ is a triangulated category and $\mathcal{K} \rightarrow \mathcal{K}$ is a morphism of triangulated categories.

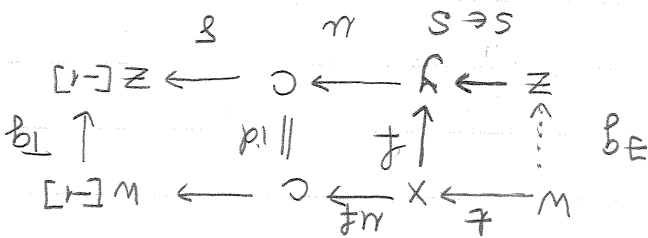
Corollary: $D(A)$ is a triangulated category over \mathcal{A} is locally small.

Definition: A morphism $\mathcal{K} \rightarrow \mathcal{K}'$ of triangulated categories is an additive functor that commutes with the translation functor T and sends exact triangles to exact triangles.

Proof of the theorem:

(1) S, τ are s.t. $H(S), H(\tau)$ are isomorphisms \Rightarrow

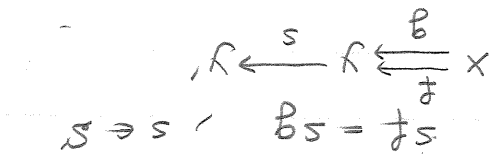
$$H(S\tau) = H(S)H(\tau) \text{ is an isomorphism.}$$



$H(S)$ isomorphism $\Rightarrow H^*(C) = 0 \Rightarrow H^*(\tau)$ isom $\Rightarrow \tau \in S$

Similarly, we show $\tau \in S$

Assume $s\tau = sg, s \in S, s\tau - g = 0 \Rightarrow s(\tau - g) = 0$



$$H^*(Z) = 0$$

$\text{Hom}_{\mathcal{K}}(X, -)$ is a homological functor \Rightarrow

$$\text{Hom}_{\mathcal{K}}(X, Z) \rightarrow \text{Hom}_{\mathcal{K}}(X, Y) \rightarrow \text{Hom}_{\mathcal{K}}(X, Y')$$

$f-g \longmapsto S(f-g) = 0$

$\Rightarrow \exists h: X \rightarrow Z$ s.t. $uh = f-g$ in \mathcal{K}

$$\exists X' \xrightarrow{t} X \xrightarrow{h} Z \xrightarrow{w} TX' \quad H^*(Z) = 0 \Rightarrow t \in S.$$

$$(f-g)t = u(ht) = 0 \Rightarrow ft = gt.$$

Similarly, $ft = gt \Rightarrow \exists s \in S$ s.t. $sf = sg$.

(2)

Assume $\bar{S}\mathcal{K}$ exists. We define a translation functor as follows:

$$T: \bar{S}\mathcal{K} \rightarrow \bar{S}\mathcal{K}, \quad T(f\bar{S}') = T(f)T(S')^{-1}$$

It is well defined since $s \in S \Rightarrow T(s) \in S'$

$$s \in S \Rightarrow H^i(s) = H^i(T'(s)) \text{ isomorphism } \forall i \in \mathbb{Z}.$$

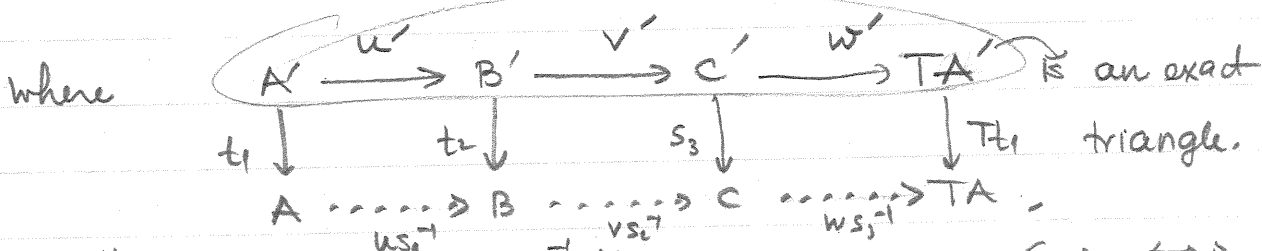
Exact triangles on $\bar{S}\mathcal{K}$

$$\begin{array}{ccc} A \xleftarrow{S_1} X \xrightarrow{u} B & uS_1^{-1} & A \xrightarrow{uS_1^{-1}} B \xrightarrow{vS_2^{-1}} C \xrightarrow{wS_3^{-1}} T(A) \\ B \xleftarrow{S_2} Y \xrightarrow{v} C & vS_2^{-1} & \\ C \xleftarrow{S_3} C' \xrightarrow{w} T(A) & wS_3^{-1} & \end{array}$$

exact triangle if:

$$wS_3^{-1} \rightsquigarrow S_3^{-1}(vS_2^{-1}) = v't_2^{-1} \rightsquigarrow t_2^{-1}(uS_1^{-1}) = u't_1^{-1}$$

$$\Rightarrow \boxed{uS_1^{-1} = t_2 u' t_1^{-1}} \text{ and } \boxed{vS_2^{-1} = S_3 v' t_2^{-1}}$$



Check*: Exact triangles on $\bar{S}\mathcal{K}$ satisfy the axioms (TR1) - (TR4).

\exists a functor of triangulated categories,

