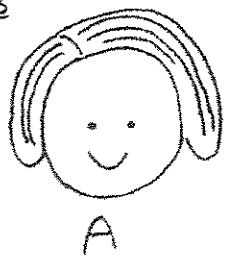
Math Access Noks
Friday June 19!
Public Key Cryptography with Secure Signature
numbering as in Davis' notes

Alice sets up to receive messages



1. Alice picks to large primes
   p, q
2. Alice computes her modulus
   $N := pq$

3a. She chooses the encryption power e relatively prime to $(p-1)(q-1)$
   [$\gcd(e, (p-1)(q-1)) = 1$]

3b. The modulus N and power e are Alice's public key
Anyone wishing to send Alice a "message", i.e a residue $0 \le x < N$, first encrypts it using the function

$$E(x) :\equiv x^e \mod N$$

7. Alice, because she knows number theory and p and q, can find her decryption power d. She solves the multiplicative inverse equation

$$ed \equiv 1 \mod (p-1)(q-1)$$

She will decrypt messages with the function

$$D(y) :\equiv y^d \mod N$$

since $D(E(x)) \equiv D(x^e) \equiv (x^e)^d = x^{ed} = x^{1 + k(p-1)(q-1)}$
$\equiv x \mod N$ by Euler's Thm Format-

the sender's original message!
(assuming x was one of the residue numbers, $0 \le x \le N-1$, and that a residue value is chosen for $D(E(x))$.)

Bob sends a message

4a. Bob wishes to send a message to Alice. He converts it into numbers x, using a conversion key like Dan's.

4b. Secure signature: Bob has created his own public key: e_B, N_B

and private key: d_B

He thinks of a sensible "signature", makes it numeric, s_B, and decrypts it using is private key, $D_B(s_B)$

5. Bob appends x to $D_B(s_B)$, creating

$$x \ast D_B(s_B)$$

(breaks this into blocks $< N_A$) and encrypts using Alice's public key

$$y = E_A\left(x \ast D_B(s_B)\right)$$

6. Bob sends y to Alice, e.g. over the internet

8. Alice wishes to decode the message y

She computes

$$D_A(y) = D_A\left(E_A\left(x \ast D_B(s_B)\right)\right)$$

$$= x \ast D_B(s_B)$$

↑ message ↑ gibberish

9. Alice uses Bob's public key to compute

$$E_B\left(D_B(s_B)\right) = s_B,$$

Bobs signature!

[Only Bob could make $D_B(s_B)$!]

Doesn't know D_A so can't get to $x \ast D_B(s_B)$;

so can't read the message and can't forge messages to Alice which look like they came from Bob.