

Getting back to GN method, why do we take steps

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$$s_k = -DR_k^T R_k \quad (*)$$

when DR_k is not full rank?

First order necessary optimality cond states:

$$\alpha^* \text{ minimizer of } \frac{1}{2} \|R(\alpha)\|^2 = f(\alpha) \Rightarrow Df(\alpha^*) = DR^T(\alpha^*) R(\alpha^*) =$$

Thus if our current iterate happens to be α^* ,

$$s_* = -DR_*^T R_* = - (DR_*^T DR_*)^+ \underbrace{DR_*^T R_*}_{=0} = 0$$

so we will not move from optimum if we find it.

(regardless of having small/large residual R_*)

Yet another way of seeing this:

s_* is the minimal norm sol to $\min_s \|R_* + DR_* s\|_2^2$

$$\Rightarrow DR_*^T DR_* s = -DR_*^T R_* = 0$$

And the smallest such s is clearly $s = 0$.

Large residual problems

Gauss Newton relies on approximating the Hessian by parts involving only first derivative information:

$$\text{if } f(x) = \frac{1}{2} \|R(x)\|^2 \text{ then } D^2 f(x) = DR^T(x) DR(x) + \sum_{i=1}^m R_i(x) D^2 R_i(x), \\ \approx DR^T(x) DR(x),$$

This approximation is not very good if $\|R(x)\|$ is large.

~ there are secant type methods (~ BFGS) that capture this missing second order information. (pp 262-265)

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§11 Non linear Equations

$F: \mathbb{R}^n \rightarrow \mathbb{R}^n$, want to find x^* s.t.

$$F(x^*) = 0$$

x^* is a root or zero of F .

Newton's method:

$$F(x_{k+1}) \approx m(x_{k+1}) = F(x_k) + DF(x_k)s \quad (\text{lin model})$$

where $DF(x_k)$ = Jacobian of F at x_k

$$= \begin{bmatrix} \nabla F_1(x_k)^T \\ \nabla F_2(x_k)^T \\ \vdots \\ \nabla F_n(x_k)^T \end{bmatrix}$$

compute s for which $m(x_k+s) = 0$ (we want a root!)

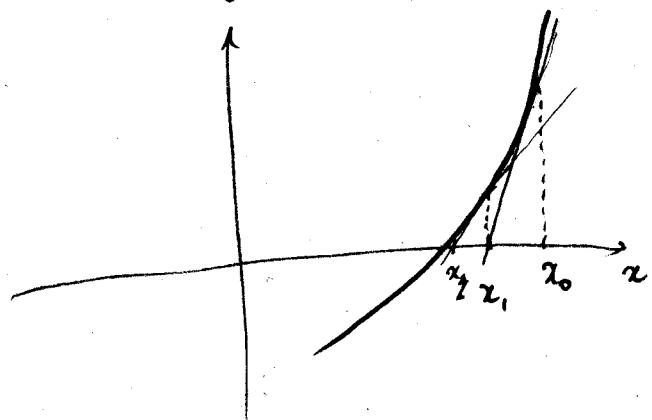
$$\Rightarrow DF_k s = -F_k \quad (2)$$

Newton's method for NL eq so : $x_{k+1} = x_k + s_k$, where s_k solves (2)

Note that:

$$T(x) = F(x_k) + F'(x_k)(x - x_k) \quad \text{to tangent to } F(x) \text{ at } x_k,$$

so by solving (2) we are finding root of tangent. For 1D function



Connection with unconstrained optimization

necessary optimality condition:

solving $\frac{DF(x)}{F(x)} = 0$ do solving a NL-type of eq.

$$DF(x) = \nabla^2 f(x) \quad (\text{guaranteed to be symmetric, but in general } DF(x) \text{ needs not be symmetric})$$

Non linear eq as an optimization problem

$$F(x) = 0 \quad (\Rightarrow \text{find } x \text{ that solves } \min_x \frac{1}{2} \|f(x)\|_2^2)$$

If we solve this problem using Gauss-Newton we find step as solution δ_k :

$$\begin{aligned} & \min_{\delta_k} \frac{1}{2} \|DF_k s + F_k\|_2^2 \quad \text{then assuming } DF_k \in \mathbb{R}^{n \times n}, \text{ non singular} \\ \Rightarrow \quad & \delta_k = - (DF_k^T DF_k)^{-1} DF_k^T F_k \\ & = - DF_k^{-1} F_k \quad \text{which is equivalent to what we derived before.} \end{aligned}$$

Convergence result for Newton's method for non-linear equations:

Let $F: \mathbb{R}^n \rightarrow \mathbb{R}^n$ be continuously diff'ble with $DF(x)$ Lipschitz continuous.

Let x^* be a point at which $DF(x^*)$ is invertible.

Then for x_0 close enough to x^* ,

$$\|x_{k+1} - x^*\| \leq C \|x_k - x^*\|^2 \quad (\text{q-quadratic convergence})$$

Proof: Banach fixed pt theorem + Banach lemma

It's possible to come up with Quasi-Newton methods that do not require the computation of $DF(x)$. Their derivation is similar to that of DFP, BFGS, and the most popular of such methods is

"Broyden's method" (P 279-283) which can achieve q-superlinear convergence.

Broyden's method:

choose initial point x_0

choose initial approx B_0 to $DR(x_0)$ (usually I)

B_0 non-singular.

for $k = 0, 1, \dots$

$$\text{Solve } B_k p_k = -R(x_k)$$

choose step size t_k (line search ...)

$$x_{k+1} = x_k + t_k p_k$$

$$s_k = x_{k+1} - x_k$$

$$y_k = R(x_{k+1}) - R(x_k)$$

$$B_{k+1} = B_k + \frac{(y_k - B_k s_k) s_k^T}{s_k^T s_k}$$

can also do update on approx to $DR(x_k)^{-1}$:

(as BFGS does)

Update so s.t. it satisfies secant equation:

$$y_k = B_{k+1} s_k$$

and

B_{k+1} solves $\min_B \|B - B_k\|$

$$\begin{matrix} B \\ \text{s.t.} \\ y_k = B s_k \end{matrix}$$

proof:

$$\|B_{k+1} - B_k\| = \left\| \frac{(y_k - B_k s_k) s_k^T}{s_k^T s_k} \right\|$$

$$= \left\| \frac{(B - B_k) s_k s_k^T}{s_k^T s_k} \right\| \leq \|B - B_k\| \left\| \frac{s_k s_k^T}{s_k^T s_k} \right\| = \|B - B_k\|$$

SVD $s_k = \begin{pmatrix} v \\ \sigma \\ u \end{pmatrix}$
 as largest singular value