

Chap 5 Initial value problems for Ordinary Diff Eq (ODE)

Objective: develop methods to solve numerically problems of the kind

$$(1) \quad \begin{cases} \frac{dy}{dt} = f(t, y) & \text{for } t \in [a, b] \\ y(a) = \alpha \end{cases} \quad \left\{ \begin{array}{l} \text{often an exact solution to (1)} \\ \text{is too complicated or even} \\ \text{cannot be found.} \end{array} \right.$$

Extension to systems:

$$\begin{cases} \frac{dy_1}{dt} = f_1(t, y_1, y_2, \dots, y_n) \\ \frac{dy_2}{dt} = f_2(t, y_1, y_2, \dots, y_n) \\ \vdots \\ \frac{dy_n}{dt} = f_n(t, y_1, y_2, \dots, y_n) \end{cases}$$

for $t \in [a, b]$, s.t.

$$y_i(a) = \alpha_i, \quad i = 1, \dots, n$$

$$\left(\frac{dy}{dt} = f(t, y) \right)$$

And to n -th order IVP:

$$y^{(n)} = f(t, y, y', y'', \dots, y^{(n-1)}) \quad \text{s.t.} \quad \begin{cases} y(a) = \alpha_1 \\ y'(a) = \alpha_2 \\ \vdots \\ y^{(n-1)}(a) = \alpha_n \end{cases}$$

⚠ this is different from notion of stability for numerical methods

Fundamental questions to answer about the IVP (1)

- existence: does (1) admit a sol?
- uniqueness: is the sol to (1) unique?
- "stability": Do small changes in the statement of the problem introduce small changes in the solution?

A problem satisfying all 3 properties is said to be well posed
(in the sense of Hadamard)

The IVP (1) is well posed under mild assumptions on f . We need some terminology (maybe you've seen this many times before!) ②

Def (Lipschitz condition)

A function $f(t, y)$ satisfies a Lipschitz condition in the variable y on some $D \subset \mathbb{R}^2$ iff $\exists L > 0$

$$|f(t, y_1) - f(t, y_2)| \leq L |y_1 - y_2| \quad \forall (t, y_1) \in D, (t, y_2) \in D.$$

L = Lipschitz constant for f .

(means $f(t, \cdot)$ is Lipschitz continuous for all t).

Example:

$$D = \{(t, y) \mid 0 \leq t \leq 1, -1 \leq y \leq 1\}$$

$$f(t, y) = t|y|$$

$$|f(t, y_1) - f(t, y_2)| = |t|y_1| - t|y_2|| = |t| ||y_1| - |y_2|| \leq |y_1 - y_2|$$

$$L = 1$$

Def (Convex set)

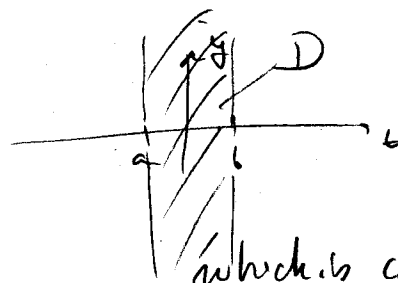
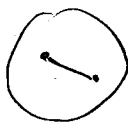
A set D is a convex set iff:

$$x, y \in D \Rightarrow \lambda x + (1-\lambda)y \in D \quad \forall \lambda \in [0, 1]$$

in our case:

$$(t_1, y_1) \in D, (t_2, y_2) \in D \rightarrow (\lambda t_1 + (1-\lambda)t_2, \lambda y_1 + (1-\lambda)y_2)$$

$$\forall \lambda \in [0, 1].$$



For IVP we usually have:

$$D = \{(t, y) \mid a \leq t \leq b, y \in \mathbb{R}\}$$

Here is a sufficient (but not necessary) condition for Lipschitz cdt^o: (3)

Theorem: Let $f(t, y)$ be defined on some convex set $D \subset \mathbb{R}^2$.

If $\exists L > 0$ s.t.

$$\left| \frac{\partial f}{\partial y}(t, y) \right| \leq L \quad \forall (t, y) \in D$$

then f satisfies a Lipschitz condition on D in y with Lip. const. L .

not: $f(t, y) = t|y|$ satisfies the Lipschitz condition and yet $\frac{\partial f}{\partial y}$ does not exist at $y=0$.

Existence and uniqueness for IVP (1) are taken care of by:

Theorem Let $D = \{(t, y) \mid a \leq t \leq b \text{ \& } y \in \mathbb{R}\}$ and that

- $f(t, y)$ is continuous on D
- f satisfies Lip. cdt^o on D in variable y , then the IVP (1) admits a unique solution.

"Stability": can be formulated as follows:

$$\text{Let } y(t) \text{ solve } \begin{cases} \frac{dy}{dt} = f(t, y), & a \leq t \leq b \\ y(a) = \alpha \end{cases}$$

$$\exists \epsilon_0 > 0, k > 0 \quad \forall \epsilon \text{ s.t. } \epsilon_0 > \epsilon > 0$$

$$\forall \delta(t) \text{ continuous s.t. } |\delta(t)| \leq \epsilon$$

$$\forall \delta_0 \in \mathbb{R} \quad \text{s.t. } |\delta_0| \leq \epsilon$$

$$\begin{cases} \frac{dz}{dt} = f(t, z) + \delta(t), & a \leq t \leq b, \\ z(a) = \alpha + \delta_0 \end{cases} \quad (\text{perturbed problem})$$

has a unique solution $z(t)$ and

$$|z(t) - y(t)| < k \epsilon$$

(Essentially the mapping data $(\alpha, f(t, y))$ to solution y is continuous).

This "stability" property is crucial to trust solutions given by a numerical method (sources of error can be from the method itself or from numerical roundoff).

$\leadsto$ All numerical methods assume IVP is well-posed.

Example: $D = \{(t, y) \mid t \in [0, 2], y \in \mathbb{R}\}$

$$\begin{cases} \frac{dy}{dt} = \underbrace{(y - t^2 + 1)}_{= f(t, y)}, & 0 \leq t \leq 2 \quad (\text{IVP}) \\ y(0) = \frac{1}{2} \end{cases}$$

$$\left| \frac{\partial f}{\partial y} \right| = |1| = 1 \Rightarrow f(t, y) \text{ satisfies Lip. cond on } D \text{ with variable } y.$$

& f constant $\Rightarrow$ problem is stable to perturbations in initial data.

We can verify this directly:

$$\begin{cases} \frac{dz}{dt} = z - t^2 + 1 + \delta & (\text{perturbed problem}) \\ z(0) = \frac{1}{2} + \delta_0 \end{cases}$$

$$\begin{aligned} \text{IVP has sol: } y(t) &= \frac{-1}{2} e^t + (t+1)^2 \\ z(t) &= \left(\frac{-1}{2} + \delta_0\right) e^t + (t+1)^2 + (e^t - 1)\delta \end{aligned}$$

$$|z - y| = |(\delta_0 + \delta)e^t - \delta| \leq |\delta_0 + \delta|e^2 + |\delta| \leq 2e^2 \epsilon + \epsilon = (2e^2 + 1)\epsilon$$

§5.2 Euler's method

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A simple numerical method to solve IVP:

$$\begin{cases} \frac{dy}{dt} = f(t, y), & a \leq t \leq b \\ y(a) = \alpha \end{cases} \quad (\text{assuming well-posedness})$$

Euler's method gives approx only on some mesh points:

$$t_j = a + hj \quad \text{where } h = \frac{b-a}{N} = \text{step size}$$

$$t_0 = a \quad t_1 \quad t_2 \quad \dots \quad t_{N-1} \quad t_N = b$$

(N+1) equally spaced points
(can be relaxed)

Idea for Euler's method: Taylor's theorem:

$$y(t_{i+1}) = y(t_i) + y'(t_i) \underbrace{(t_{i+1} - t_i)}_{=h} + \frac{1}{2} y''(\xi_i) \underbrace{(t_{i+1} - t_i)^2}_{=h^2}$$

for some $\xi_i \in (t_i, t_{i+1})$.

$$= y(t_i) + h y'(t_i) + \frac{h^2}{2} y''(\xi_i)$$

$$= y(t_i) + h f(t_i, y(t_i)) + \underbrace{\frac{h^2}{2} y''(\xi_i)}_{\text{neglect for Euler's method}} \quad (y \text{ satisfies DE in IVP})$$

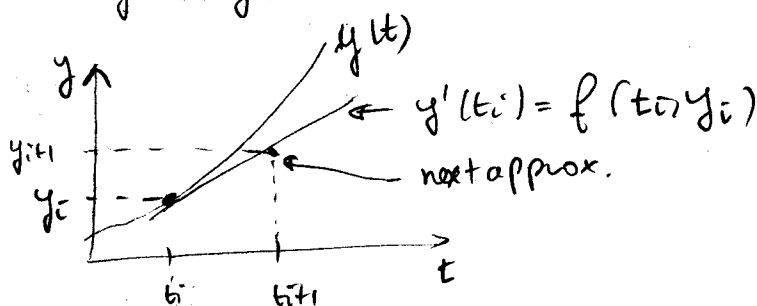
$$y_0 = \alpha$$

for $i = 0 \dots N-1$ known

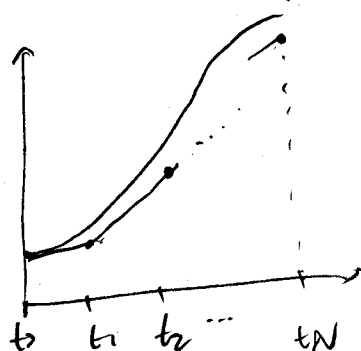
$$y_{i+1} = y_i + h f(t_i, y_i)$$

Geometrical interpretation

Assume $y_i = y(t_i)$



systematic error
is introduced at
many times: every step



Several types of errors in numerical methods for DE:

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- Local truncation error: error made in one step-

(Δ book defines this as $\frac{1}{h}$) For example in Euler's method: $O(h)$ since:
$$y(t_i + h) = y(t_i) + h f(t_i, y(t_i)) + \underline{O(h^2)}$$

- Local roundoff error: precision used for computation (10^{-16} double, 10^{-8} single)

- Global truncation error: accumulation of all the local truncation errors. If local truncation error was $O(h^{n+1})$ then the global truncation error must be $O(h^n)$ because the number of steps $O(1/h)$.

- Global roundoff error: accumulation of local roundoff errors of previous steps.

- Total error = global truncation error + global roundoff error

If the global truncation error is $O(h^n)$ then the method is of order n

Example: Euler's method is of order 1, and it is relatively simple to derive more precise bounds on the global truncation error:

then let f be continuous and satisfy Lipschitz condition with L -constant L on $D = \{(t, y) \mid t \in [a, b], y \in \mathbb{R}\}$ and that

$$|y''(t)| \leq M \quad \forall t \in [a, b] \text{ for some } M > 0.$$

let $y(t)$ be the sol to IVP
$$\begin{cases} \frac{dy}{dt} = f(t, y) \\ y(a) = \alpha \end{cases}$$

and $y_0, y_1, \dots, y_N$ the approximations given by Euler's method

then:

$$|y(t_i) - y_i| \leq \frac{hM}{2L} [e^{L(t_i - a)} - 1], \quad i = 0, 1, \dots, N.$$

Proof (sketch)

$$y(t_{i+1}) = y(t_i) + hf(t_i, y(t_i)) + \frac{h^2}{2} y''(\xi_i)$$

$$y_{i+1} = y_i + hf(t_i, y_i)$$

$$\Rightarrow |y(t_{i+1}) - y_{i+1}| \leq |y(t_i) - y_i| + h |f(t_i, y(t_i)) - f(t_i, y_i)| + \frac{h^2}{2} |y''(\xi_i)|$$

$$\leq (1 + hL) |y(t_i) - y_i| + \frac{h^2}{2} M$$

Using Lemma 5.8 in book: ($\sim$ geometric series + bounding exp)

$$|y_{i+1} - w_{i+1}| \leq e^{(i+1)hL} \left(\underbrace{|y_0 - w_0|}_{=0} + \frac{h^2 M}{2hL} \right) - \frac{h^2 M}{2hL}$$

$$\leq \frac{h^2 M}{2hL} (e^{\overbrace{(i+1)hL}^{t_{i+1} - t_0}} - 1)$$

$$t_{i+1} - t_0 = t_{i+1} - a.$$

Lemma 5.8 $\{a_i\}_{i=0}^k$ is a seq with $a_0 \geq -\frac{t}{s}$ and

$$a_{i+1} \leq a_i(1+s) + t \quad \text{for } i=0, \dots, k-1$$

$$\Rightarrow a_{i+1} \leq e^{(i+1)s} \left(a_0 + \frac{t}{s} \right) - \frac{t}{s}$$

Proof:

$$a_{i+1} \leq a_i(1+s) + t \leq ((1+s)a_{i-1} + t)(1+s) + t$$

$$\leq ((1+s)((1+s)a_{i-2} + t) + t)(1+s) + t$$

$$\vdots$$

$$\leq (1+s)^{i+1} a_0 + t \sum_{j=0}^i (1+s)^j$$

$$\Rightarrow a_{i+1} \leq (1+\delta)^{i+1} a_0 + t \left[\frac{1 - (1+\delta)^{i+1}}{1 - (1+\delta)} \right] \quad (8)$$

$$= (1+\delta)^{i+1} \left(a_0 + \frac{t}{\delta} \right) - \frac{t}{\delta}$$

$$\leq e^{(i+1)\delta} \left(a_0 + \frac{t}{\delta} \right) - \frac{t}{\delta}$$

$$\text{since } (1+x)^\alpha = e^{\alpha \ln(1+x)} \leq e^{\alpha x}$$

can be shown using Taylor's thm.

For Euler's method it's even possible to incorporate the round off errors in the analysis:

instead of having:

$$y_0 = \alpha$$

$$\text{for } i = 0 \dots N-1,$$

$$y_{i+1} = y_i + h f(t_i, y_i)$$

we commit a mistake at each step:

$$y_0 = \alpha + \delta_0$$

$$\text{for } i = 0 \dots N-1$$

$$y_{i+1} = y_i + h f(t_i, y_i) + \delta_{i+1}$$

If $|\delta_i| < \delta$ it is possible to show:

$$|y(t_i) - y_0| \leq \frac{1}{L} \left(\frac{hM}{2} + \frac{\delta}{h} \right) [e^{L(t_i-a)} - 1] + |\delta_0| e^{L(t_i-a)}$$

performance degrades as $h \rightarrow 0$!!

(similar to problems with numerical differentiation)

The h giving smallest error is

$$h = \sqrt{\frac{2\delta}{M}} \quad (\text{simple calculation})$$

§3 Higher order Taylor methods:

⑨

Same derivation as Euler method but carry Taylor series further:

$$y(t_{i+1}) = y(t_i) + h y'(t_i) + \frac{h^2}{2} y''(t_i) + \dots + \frac{h^n}{n!} y^{(n)}(t_i) + \frac{h^{n+1}}{(n+1)!} y^{(n+1)}(\xi_i)$$

for some $\xi_i \in (t_i, t_{i+1})$.

$$y'(t) = f(t, y(t))$$

$$y''(t) = \frac{d}{dt} [f(t, y(t))]$$

$$\vdots$$
$$y^{(k)}(t) = \frac{d^{k-1}}{dt^{k-1}} [f(t, y(t))]$$

Taylor method of order n

$$y_0 = \alpha$$

for $i = 0 \dots N-1$,

$$y_{i+1} = y_i + h f(t_i, y_i) + \frac{h^2}{2} \frac{d}{dt} [f(t_i, y_i)] + \dots + \frac{h^n}{n!} \frac{d^{n-1}}{dt^{n-1}} [f(t_i, y_i)]$$

Local truncation error is $O(h^{n+1})$

In general evaluating $\frac{d^k}{dt^k} [f(t, y(t))]$ can be quite involved because of the repeated application of the chain rule.

For example:

$$\begin{cases} y' = \cos t - \sin y + t^2 = f(t, y) \\ y(-1) = 3 \end{cases}$$

$$y'' = \frac{d}{dt} [f(t, y(t))] = -\sin t - y' \cos y + 2t$$

$$y''' = \frac{d^2}{dt^2} [f(t, y(t))] = -\cos t - y'' \cos y + (y')^2 \sin y + 2$$

$$y^{(4)} = \frac{d^3}{dt^3} [f(t, y(t))] = \sin t - y^{(3)} \cos y + 3y' y'' \sin y + (y')^3 \cos y$$

etc...

Note: • Assumes f is smooth ($n-1$ times differentiable if method of order n is desired)

(10)

- The higher the accuracy the more complicated the steps become
- differentiation can be carried out symbolically if possible!

§5.4 Runge-Kutta methods

Taylor series method for $\begin{cases} \frac{dy}{dt} = f(t, y), & t \in [a, b] \\ y(a) = A \end{cases}$

require us to compute formulas for $y'' = \frac{d}{dt}[f(t, y)]$
 $y^{(3)} = \frac{d^2}{dt^2}[f(t, y)]$
etc...

which can be quite involved.

Runge-Kutta methods avoid this difficulty by carefully chosen combinations of values of $f(t, y)$.

Second order Runge Kutta methods:

Start with Taylor series:

$$y(t+h) = y(t) + h y'(t) + \frac{h^2}{2} y''(t) + \frac{h^3}{3!} y^{(3)}(t) + \dots$$

From the DE we get:

$$y'(t) = f(t, y) = \boxed{f}$$

$$y''(t) = \frac{d}{dt} f(t, y) = \frac{\partial f}{\partial t}(t, y) + y' \frac{\partial f}{\partial y}(t, y) = \boxed{f_t + f f_y}$$

$$y'''(t) = \frac{d^2}{dt^2} f(t, y) = f_{tt} + f f_{ty} + (f_t + f f_y) f_y + f (f_{ty} + f f_{yy})$$

etc...

$$\Rightarrow y(t+h) = y(t) + hf + \frac{h^2}{2} (f_t + f f_y) + O(h^3)$$

The idea is to eliminate the partial deriv of f by using the first few terms of the two variable Taylor series for $f(t, y)$: (1)

$$\begin{aligned} f(t+h, y+h f) &= f + \nabla f \cdot \begin{pmatrix} h \\ h f \end{pmatrix} + \mathcal{O}(h^2) \\ &= f + h f_t + h f f_y + \mathcal{O}(h^2) \end{aligned}$$

Rewriting the Taylor series of y :

$$\begin{aligned} (*) \quad y(t+h) &= y(t) + \frac{1}{2} h f + \frac{1}{2} h \left[\underbrace{f + h f_t + h f f_y}_{\substack{\text{first 2 terms of 2 variable} \\ \text{Taylor series of } f}} \right] + \mathcal{O}(h^3) \\ &= y(t) + \frac{1}{2} h f + \frac{1}{2} h \left[f(t+h, y+h f) + \mathcal{O}(h^2) \right] + \mathcal{O}(h^3) \end{aligned}$$

Thus it's possible to construct an update which has the same $\mathcal{O}(h^3)$ local truncation error as 2nd-order Taylor method:

Modified Euler method

$$y_0 = A$$

for $i = 0, 1, \dots, N-1$

$$F_1 = h f(t_i, y_i)$$

$$F_2 = h f(t_i + h, y_i + F_1)$$

$$y_{i+1} = y_i + \frac{1}{2} F_1 + \frac{1}{2} F_2$$

The general form for Second order Runge Kutta updates is:

$$y(t+h) = y + w_1 h f + w_2 h f(t + \alpha h, y + \beta h f)$$

where w_1, w_2, α, β are parameters that we can adjust.

Using two variable Taylor expansion:

$$y(t+h) = y + w_1 h f + w_2 h \left[f + \alpha h f_t + \beta h f f_y \right]$$

Matching comparable terms with (*) we get:

(12)

$$\left\{ \begin{array}{l} w_1 + w_2 = 1 \\ w_2 \alpha = \frac{1}{2} \\ w_2 \beta = \frac{1}{2} \end{array} \right.$$

→ so there is a family of RK.
order 2 methods

Modified Euler: $w_1 = w_2 = \frac{1}{2}$, $\alpha = \beta = 1$

Midpoint method: $w_1 = 0$, $w_2 = 1$, $\alpha = \beta = \frac{1}{2}$

$$y_0 = A$$

for $i = 0, 1, \dots, N-1$

$$F_1 = h f(t_i, y_i)$$

$$F_2 = h f\left(t_i + \frac{h}{2}, y_i + \frac{1}{2} F_1\right)$$

$$y_{i+1} = y_i + F_2$$

(also $O(h^3)$ LTE)

Heun's method: $w_1 = \frac{1}{4}$, $w_2 = \frac{3}{4}$, $\alpha = \beta = \frac{2}{3}$

$$y_0 = A$$

for $i = 0, 1, \dots, N-1$

$$F_1 = h f(t_i, y_i)$$

$$F_2 = h f\left(t_i + \frac{2}{3}h, y_i + \frac{2}{3}F_1\right)$$

$$y_{i+1} = y_i + \frac{1}{4}F_1 + \frac{3}{4}F_2$$

Runge Kutta methods of order 3 are obtained by matching terms between Taylor expansion of order 3 of $y(t)$ and $f(t + \alpha h, y + w_1 f(t + \beta h, y + w_2 f(t, y)))$

The derivation is quite tedious, but here is one possible RK order 3 method:

$$y_0 = A$$

for $i = 0, \dots, N-1$

$$F_1 = h f(t_i, y_i)$$

$$F_2 = h f(t_i + \frac{1}{2}h, y_i + \frac{1}{2}F_1)$$

$$F_3 = h f(t_i + \frac{3}{4}h, y_i + \frac{3}{4}F_2)$$

$$y_{i+1} = y_i + \frac{1}{9}(2F_1 + 3F_2 + 4F_3)$$

$$\left(\begin{array}{l} \alpha = w_1 = \frac{3}{4} \\ \beta = w_2 = \frac{1}{2} \end{array} \right)$$

However it is not commonly used in practice.

The most popular RK method is that of order 4; again the derivation is tedious but the implementation straight forward.

Runge Kutta method of order 4 (LTE $\mathcal{O}(h^5)$)

$$y_0 = A$$

for $i = 0, \dots, N-1$

$$F_1 = h f(t_i, y_i)$$

$$F_2 = h f(t_i + \frac{1}{2}h, y_i + \frac{1}{2}F_1)$$

$$F_3 = h f(t_i + \frac{1}{2}h, y_i + \frac{1}{2}F_2)$$

$$F_4 = h f(t_i + h, y_i + F_3)$$

$$y_{i+1} = y_i + \frac{1}{6}(F_1 + 2F_2 + 2F_3 + F_4)$$

of function eval.

1 2 3 4 5 6 7 8 9 10 11 12

max order of RK method

1 2 3 4 4 5 6 6 7 7 8 9 ...

So # of function eval increases more rapidly than the max order of RK method. $\Rightarrow$ higher order RK methods are less attractive than the classical RK4. (it makes more sense to use smaller time steps for RK4 than using a higher order method with bigger steps).

§ 5.5 Adaptive Runge Kutta Fehlberg method

Idea: Estimate local truncation error and adjust the step length accordingly.

Local truncation error estimation

Assume we are given two update formulas for solving

$$\begin{cases} y' = f(t, y) & , t \in [a, b] \\ y(a) = \alpha \end{cases}$$

with local truncation errors differing by 1.

$$\textcircled{1} \quad y(t_{i+1}) = y(t_i) + h \phi(t_i, y(t_i), h) + \mathcal{O}(h^{n+1}) \quad (\text{order } n)$$

$$y_0 = \alpha$$

for $i = 0, \dots, N-1$

$$| y_{i+1} = y_i + h \phi(t_i, y_i, h)$$

$$\textcircled{2} \quad y(t_{i+1}) = y(t_i) + h \tilde{\phi}(t_i, y(t_i), h) + \mathcal{O}(h^{n+2}) \quad (\text{order } n+1)$$

$$\tilde{y}_0 = \alpha$$

for $i = 0, \dots, N-1$

$$| \tilde{y}_{i+1} = \tilde{y}_i + h \tilde{\phi}(t_i, y_i, h)$$

These update formulas could come from Taylor's method, RK etc... (5)

Assuming $y_i = \tilde{y}_i = y(t_i)$:

$$\underbrace{y_{i+1} - y(t_{i+1})}_{\substack{\text{local trunc err for ①} \\ \text{"} \\ \mathcal{O}(h^{n+1})}} = \underbrace{\tilde{y}_{i+1} - y(t_{i+1})}_{\substack{\text{local trunc error for ②} \\ \text{"} \\ \mathcal{O}(h^{n+2})}} + \underbrace{y_{i+1} - \tilde{y}_{i+1}}_{\text{"} \\ \mathcal{O}(h^{n+1})} \Rightarrow \mathcal{O}(h^{n+1})$$

Thus:

$$\underbrace{y_{i+1} - y(t_{i+1})}_{\text{local trunc err for ①}} = y_{i+1} - \tilde{y}_{i+1} + \mathcal{O}(h^{n+2})$$

How to use this estimate to predict step size needed to get a local truncation error $< \delta$ (where $\delta > 0$ is prescribed) ?

Simplifying assumption: K constant indep of h .

$$\tau_{i+1}(h) = y_{i+1} - y(t_{i+1}) \approx K h^{n+1} \quad (\text{① is } n \text{ order method})$$

If instead of h the step were qh ($q > 0$) :

$$\tau_{i+1}(qh) \approx K (qh)^{n+1} = q^{n+1} K h^{n+1} \approx q^{n+1} \tau_{i+1}(h) \approx q^{n+1} (y_{i+1} - \tilde{y}_{i+1})$$

$$\Rightarrow q^{n+1} |y_{i+1} - \tilde{y}_{i+1}| \approx |\tau_{i+1}(qh)| \leq \delta$$

$$\Rightarrow \text{for } q \leq \left(\frac{\delta}{|y_{i+1} - \tilde{y}_{i+1}|} \right)^{\frac{1}{n+1}} \text{ we can expect local trunc err to be less than } \delta.$$

Thus if we want to keep the local trunc error $\leq \delta$, we can

$$\text{set } h_{\text{new}} = h \left(\frac{\delta}{|y_{i+1} - \tilde{y}_{i+1}|} \right)^{\frac{1}{n+1}}$$

However it is recommended (based on empirical observations) to relax this to:

$$h_{\text{new}} = h \left(\frac{\delta}{2|y_{i+1} - \tilde{y}_{i+1}|} \right)^{\frac{1}{n}} \approx 0.9 h \left(\frac{\delta}{|y_{i+1} - \tilde{y}_{i+1}|} \right)^{\frac{1}{n}}$$

Fehlberg (1969) used an order 4 RK method with 5 fun evals
a $\xrightarrow{5}$ $\xrightarrow{6}$

total 11 fun. evals / iteration?? Does not seem very practical!
However the coefficients in Fehlberg's method are carefully chosen so that only 6 fun evals are needed / iter.
(by making the eval points for the order 4 method match those of the order 5 method).

The updates are of the form:

$$y_{i+1} = y_i + \sum_{j=1}^5 a_j F_j$$

$$\tilde{y}_{i+1} = y_i + \sum_{j=1}^6 b_j F_j$$

where the F_i are of the form:

$$F_i = h f(t + c_i h, x + \sum_{j=1}^{i-1} d_{ij} F_j)$$

(see book for the coeff. Sometimes only $a_i - b_i$ is specified)

The complete algo is:

(17)

$y_0 = \alpha$, $h = h_{\max}$, $\text{flag} = 1$
while ($\text{flag} = 1$)

- compute $F_1, \dots, F_6$
- compute $e = |y_{i+1} - \bar{y}_{i+1}|$ (a linear combination of the F_i , w/coeff given in tables)
- if $e \leq \delta$
 accept step:
 $t = t + h$
 $y_{i+1} = y_i + \sum_{j=1}^6 b_j F_j$ (use higher order method)
- Set $q = 0.9 (\delta/e)^{1/4}$
- Set $h = \begin{cases} 0.1 h & \text{if } q \leq 0.1 \\ 4 h & \text{if } q \geq 4 \\ q h & \text{otherwise} \end{cases}$
- Set $h = \min(h, h_{\max})$
- if $t \geq b$
 | $\text{flag} = 0$
 else
 | if $t + h > b$
 | | $h = b - t$ (adjust last step)
- if $h < h_{\min}$
 | $\text{flag} = 0$ (Step size is too small abnormal termination)

(normal termination)