

## §2.4: Cardinality "size" of a set = Cardinality

How can we compare the sizes of two sets? If  $S = \{x \in \mathbb{R} \mid x^2 = 9\}$ , then  $S = \{-3, 3\}$  and we say  $S$  has 2 elements. If  $T = \{1, 7, 11\}$  then  $T$  has 3 elements, and  $T$  is "larger" than  $S$ . How do we compare infinite sets like  $\mathbb{N}$  and  $\mathbb{R}$ ?

Def: Two sets  $S$  and  $T$  are equinumerous,  $S \sim T$ , if there exists a bijective function from  $S$  onto  $T$ .

Practice: If  $\mathcal{F}$  is a family of sets, then equinumerous is a relation on  $\mathcal{F}$ . Show it is an equivalence relation.

Def: A set  $S$  is finite if  $S = \emptyset$  or if  $\exists n \in \mathbb{N}$  and a bijection  $f: \{1, \dots, n\} \rightarrow S$ . If a set is not finite, it is infinite.

Notation:  $I_n = \{1, 2, \dots, n\}$ . If  $S \sim I_n$ ,  $S$  has  $n$  elements. The "size" of  $\emptyset$  is 0, and if the "size" of a set is not finite it is transfinite.

Practice: Suppose  $S$  and  $T$  are sets having  $n$  elements. Then from the definition  $\exists g, h$  bijections,  $g: I_n \rightarrow S$ ,  $h: I_n \rightarrow T$ . Show  $S$  and  $T$  are equinumerous directly by finding a bijection  $f: S \rightarrow T$ .  
 $(h \circ g^{-1}): S \rightarrow T$

Def: A set  $S$  is denumerable if  $\exists$  bijection  $f: \mathbb{N} \rightarrow S$ . If a set is finite or denumerable, it is countable. If a set is not countable it is uncountable.

The "size" of a countable set is  $\aleph_0$  (aleph)

Note: Not all infinite sets are  $\aleph_0$ . There are different "sizes" of infinite sets.

The "size" of a set  $S$  is <sup>the cardinality of a set,</sup> denoted  $|S|$ .

Ex: Let  $E = \{\text{even natural numbers}\}$ . How does  $|E|$  compare to  $|\mathbb{N}| = \aleph_0$ ? Since  $E \subset \mathbb{N}$ , we would think  $|E| < |\mathbb{N}|$ , but actually  $f: \mathbb{N} \rightarrow E$  given by  $f(n) = 2n$  is a bijection of  $\mathbb{N}$  onto  $E$ , so  $|E| = |\mathbb{N}| = \aleph_0$ .

Practice: Find a bijection  $f: \mathbb{N} \rightarrow \mathbb{Z} \Rightarrow \mathbb{Z}$  is denumerable.

$$f(n) = \begin{cases} \frac{n}{2} & \text{if } n \text{ is even} \\ -\frac{(n-1)}{2} & \text{if } n \text{ is odd} \end{cases} \quad \begin{array}{l} f(1) = s_1 = 0 \\ f(2) = s_2 = 1 \\ f(3) = s_3 = -1 \end{array} \quad \begin{array}{l} f(4) = s_4 = 2 \\ \vdots \end{array}$$

Let  $S$  be a nonempty set which is finite. Then there is a bijection  $f: I_n \rightarrow S$  which we can use to "count" the members of  $S$ ;

$f(1), f(2), \dots, f(n)$ . If we let  $s(k) = s_k$ , for  $1 \leq k \leq n$ , we have  $S = \{s_1, s_2, \dots, s_k\}$ . This process works for any denumerable set.

Axiom 1: (Wellordering Property of  $\mathbb{N}$ ) If  $S \subseteq \mathbb{N}$ ,  $S \neq \emptyset$ , then  $S$  has a least element. i.e.  $\exists m \in S$  st.  $m \leq k \forall k \in S$ .

Theorem 15: Let  $S$  be a countable set and let  $T \subseteq S$ . Then  $T$  is countable.

Pf: If  $T$  is finite, we are done. So suppose  $T$  is infinite  $\Rightarrow S$  is infinite.  $\exists f: \mathbb{N} \rightarrow S$  which is bijective, so we can write  $S = \{s_1, s_2, \dots\}$ . Let  $A = \{n \in \mathbb{N} \mid s_n \in T\}$ . Since  $A$  is a nonempty subset of  $\mathbb{N}$ , it has a least member, say  $a_1$ . Similarly  $A \setminus \{a_1\}$  has a least member,  $a_2$ . In general, having chosen  $\{a_1, a_2, \dots, a_k\}$ , let  $a_{k+1}$  be the least member of  $A \setminus \{a_1, a_2, \dots, a_k\}$ .

Now define  $g: \mathbb{N} \rightarrow \mathbb{N}$  by  $g(n) = a_n$ . Since  $T$  is infinite,  $g$  is defined for all  $n \in \mathbb{N}$ . Since  $a_{n+1} \notin \{a_1, \dots, a_n\}$ ,  $g$  is injective. Thus the composition  $f \circ g$  is injective. Since each element of  $T$  is somewhere in the list  $s_1, s_2, \dots$  of  $S$ ,  $g(\mathbb{N})$  includes all the subscripts of terms in  $T$ . Thus  $f \circ g$  is a bijection from  $\mathbb{N}$  onto  $T \Rightarrow T$  is denumerable.  $\square$

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Theorem 16: Let  $S$  be a nonempty set. TFAE:

- 1)  $S$  is countable
- 2)  $\exists f: S \rightarrow \mathbb{N}$  injective
- 3)  $\exists g: \mathbb{N} \rightarrow S$  surjective

Pf:

(1  $\Rightarrow$  2) Let  $S$  be countable. Then  $\exists h: I_n \rightarrow S$  for some  $n \in \mathbb{N}$  if  $S$  is finite or  $h: \mathbb{N} \rightarrow S$  if  $S$  is denumerable. In either case,  $h^{-1}: S \rightarrow \mathbb{N}$  is an injection.

(2  $\Rightarrow$  3) Suppose  $\exists f: S \rightarrow \mathbb{N}$  which is injective. Then  $f$  is a bijection from  $S$  to  $f(S)$ , and  $f^{-1}$  is a bijection from  $f(S)$  to  $S$ . Let  $p \in S$ . Define  $g: \mathbb{N} \rightarrow S$  by

$$g(n) = \begin{cases} f^{-1}(n) & \text{if } n \in f(S) \\ p & \text{if } n \notin f(S) \end{cases}$$

Then  $g[f(S)] = f^{-1}(f(S)) = S$  and  $g[\mathbb{N} \setminus f(S)] = \{p\}$  so  $g$  is a surjection from  $\mathbb{N}$  onto  $S$ .

(3  $\Rightarrow$  1) Let  $g: \mathbb{N} \rightarrow S$  be a surjection. Define  $h: S \rightarrow \mathbb{N}$  by  $h(s) =$  smallest  $n \in \mathbb{N}$  s.t.  $g(n) = s$ . Then  $h: S \rightarrow \mathbb{N}$  is injective, and bijective from  $S$  to  $h(S) \subseteq \mathbb{N}$ . Since  $\mathbb{N}$  is countable, our previous theorem tells us  $h(S)$  is countable. Since  $S$  and  $h(S)$  are equinumerous,  $S$  is also countable.  $\square$

Ex: a) Let  $S$  and  $T$  be nonempty countable sets. Then  $S \cup T$  is countable. We have (by the theorem)  $f: \mathbb{N} \rightarrow S$ ,  $g: \mathbb{N} \rightarrow T$  which are surjective. Define  $h: \mathbb{N} \rightarrow S \cup T$  by

$$h(n) = \begin{cases} f\left(\frac{n+1}{2}\right) & \text{if } n \text{ is odd} \\ g\left(\frac{n}{2}\right) & \text{if } n \text{ is even} \end{cases}$$

$h$  is surjective, so  $S \cup T$  is countable.

b) Let  $S$  and  $T$  be nonempty countable sets. We will show  $S \times T$  is countable, by using the factorization of numbers is unique, up to order,

$\exists f: S \rightarrow \mathbb{N}, \exists g: T \rightarrow \mathbb{N}$  injections.

Define  $h: S \times T \rightarrow \mathbb{N}$  by

$$h(s, t) = 2^{f(s)} 3^{g(t)} \quad \text{where } s \in S, t \in T.$$

Then  $h$  is injective, for if  $h(s, t) = h(u, v)$

$$\Rightarrow 2^{f(s)} 3^{g(t)} = 2^{f(u)} 3^{g(v)} \Rightarrow f(s) = f(u), g(t) = g(v)$$

$$\Rightarrow s = u, t = v.$$

$\therefore S \times T$  is countable.

c)  $\mathbb{Q}$  is countable.

First we will show  $\mathbb{Q}^+ = \{x \in \mathbb{Q} \mid x > 0\}$  is countable.

Every  $x \in \mathbb{Q}^+$  can be written as  $\frac{m}{n}$  with  $m, n \in \mathbb{N}, n \neq 0$  and  $m$  and  $n$  are relatively prime.

Define  $f: \mathbb{Q}^+ \rightarrow \mathbb{N}$  by  $f\left(\frac{m}{n}\right) = 2^m 3^n$ .

$f$  is injective, so  $\mathbb{Q}^+$  is countable. The mapping  $g: \mathbb{Q}^+ \rightarrow \mathbb{Q}^-$  given by  $g(r) = -r$  is clearly bijective, so  $\mathbb{Q}^+$  and  $\mathbb{Q}^-$  are equinumerous, so  $\mathbb{Q}^-$  is countable. Since  $\mathbb{Q} = \mathbb{Q}^+ \cup \{0\} \cup \mathbb{Q}^-$ , by applying (a) twice,  $\mathbb{Q}$  is countable.

d) The union of a countable family of countable sets is countable. Let  $\mathcal{C} = \{S_1, S_2, \dots\}$ .

$$S_1: s_{11} \quad s_{12} \quad s_{13} \quad \dots$$

$$S_2: s_{21} \quad s_{22} \quad s_{23} \quad \dots$$

$$S_3: s_{31} \quad s_{32} \quad s_{33} \quad \dots$$

Theorem 17:  $\mathbb{R}$  is uncountable.

Pf: Uses the "diagonal process" of Cantor.

It suffices to show  $J = (0, 1)$  is uncountable.

Suppose  $J$  is countable,  $J = \{x_1, x_2, x_3, \dots\} = \{x_n \mid n \in \mathbb{N}\}$ .

We shall construct a number  $y \in (0, 1)$  but  $y \notin \{x_1, x_2, \dots\}$ .

Each element in  $J$  has an infinite decimal expansion, so we write

$$x_1 = 0.a_{11} a_{12} a_{13} \dots$$

$$x_2 = 0.a_{21} a_{22} a_{23} \dots$$

$$x_3 = 0.a_{31} a_{32} a_{33} \dots$$

$\vdots$

$$x_n = 0.a_{n1} a_{n2} a_{n3} \dots$$

where each  $a_{ij} \in \{0, 1, 2, \dots, 9\}$ . (Some numbers have more than one representation such as  $0.5000\dots$  and  $0.4999\dots$ , but that's not a problem).

Let  $y = 0.b_1 b_2 b_3 \dots$  where

$$b_n = \begin{cases} 2 & \text{if } a_{nn} \neq 2 \\ 3 & \text{if } a_{nn} = 2 \end{cases}$$

So  $y \in J$ , but  $y \neq x_i$  for any  $i$ . This contradicts our assumption that  $J$  is countable, so  $J$  is uncountable.

Let's return to our question of comparing the "size" of sets.  $|S| = |T|$  iff  $S$  and  $T$  are equinumerous iff  $\exists$  bijection

$$f: S \rightarrow T.$$

We define  $|S| \leq |T|$  if  $\exists$  injection  $f: S \rightarrow T$ .

Theorem 18: Let  $S, T$  and  $U$  be sets.

a) If  $S \subseteq T$ , then  $|S| \leq |T|$

b)  $|S| \leq |S|$

c) If  $|S| \leq |T|$  and  $|T| \leq |U|$ , then  $|S| \leq |U|$ .

d) If  $m, n \in \mathbb{N}$ ,  $m \leq n$ , then  $|\{1, 2, \dots, m\}| \leq |\{1, 2, \dots, n\}|$ .

e) If  $S$  is finite, then  $|S| \leq \aleph_0$

It is customary to denote  $|\mathbb{R}|$  by  $c$ , for continuum. Since  $\aleph_1 \in \mathbb{R}$ , we have  $\aleph_1 \leq c$ . Therefore  $\aleph_1$  and  $c$  are unequal transfinite cardinals. There are others.

Notation: Given any set  $S$ , let  $\mathcal{P}(S)$  denote the collection of all subsets of  $S$ . This is called the power set of  $S$ .

Theorem 19: For any set  $S$ , we have  $|S| < |\mathcal{P}(S)|$ .

Pf: The function  $g: S \rightarrow \mathcal{P}(S)$  given by  $g(s) = \{s\}$  is clearly injective, so  $|S| \leq |\mathcal{P}(S)|$ . To prove  $|S| \neq |\mathcal{P}(S)|$ , we show that no function from  $S$  to  $\mathcal{P}(S)$  can be surjective. Suppose that  $f: S \rightarrow \mathcal{P}(S)$ . Then for each  $x \in S$ ,  $f(x) \subseteq S$ . Now for some  $x \in S$  it may be that  $x \in f(x)$  and for others it may not be. Let

$$T = \{x \in S \mid x \notin f(x)\}$$

Now  $T \subseteq S$ , so  $T \in \mathcal{P}(S)$ . If  $f$  were surjective, then  $T = f(y)$  for some  $y \in S$ . Now either  $y \in T$  or  $y \notin T$ , but both possibilities lead to contradictions: If  $y \in T$ , then  $y \notin f(y)$ , by how  $T$  is defined. But  $f(y) = T$ , so  $y \in T$ . On the other hand if  $y \notin T$ , then  $y \in f(y)$   $\Rightarrow y \in T$ .

Thus no function from  $S$  to  $\mathcal{P}(S)$  can be surjective, so  $|S| \neq |\mathcal{P}(S)|$ .  $\square$

By applying Theorem 19 over and over again, we get

$$\aleph_0 = |\mathbb{N}| < |\mathcal{P}(\mathbb{N})| < |\mathcal{P}(\mathcal{P}(\mathbb{N}))| < \dots$$

It can be shown that  $|\mathcal{P}(\mathbb{N})| = c$

Is there any set "in between"  $\mathbb{N}$  and  $\mathbb{R}$  in size? i.e.  $\exists$  a cardinal number  $\lambda$  s.t.  $\aleph_0 < \lambda < c$ ?

This is the continuum hypothesis. Famous unsolved problem.