# Programming languages - C 

## ABSTRACT

(Cover sheet to be provided by ISO Secretariat.)
This International Standard specifies the form and establishes the interpretation of programs expressed in the programming language C. Its purpose is to promote portability, reliability, maintainability, and efficient execution of C language programs on a variety of computing systems.

Clauses are included that detail the C language itself and the contents of the C language execution library. Annexes summarize aspects of both of them, and enumerate factors that influence the portability of C programs.

Although this International Standard is intended to guide knowledgeable C language programmers as well as implementors of C language translation systems, the document itself is not designed to serve as a tutorial.

## Introduction

1 With the introduction of new devices and extended character sets, new features may be added to this International Standard. Subclauses in the language and library clauses warn implementors and programmers of usages which, though valid in themselves, may conflict with future additions.
2 Certain features are obsolescent, which means that they may be considered for withdrawal in future revisions of this International Standard. They are retained because of their widespread use, but their use in new implementations (for implementation features) or new programs (for language [6.9] or library features [7.20]) is discouraged.
3 This International Standard is divided into four major subdivisions:

- the introduction and preliminary elements;
- the characteristics of environments that translate and execute C programs;
- the language syntax, constraints, and semantics;
- the library facilities.

4 Examples are provided to illustrate possible forms of the constructions described. Footnotes are provided to emphasize consequences of the rules described in that subclause or elsewhere in this International Standard. References are used to refer to other related subclauses. A set of annexes summarizes information contained in this International Standard. The introduction, the examples, the footnotes, the references, and the annexes are not part of this International Standard.

5 The language clause (Clause 6) is derived from 'The C Reference Manual'" (see annex A).
6 The library clause (Clause 7) is based on the 1984 /usr/group Standard (see annex A).

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## Programming languages -C

## 1. Scope

1 This International Standard specifies the form and establishes the interpretation of programs written in the C programming language. ${ }^{1}$ It specifies

- the representation of C programs;
- the syntax and constraints of the C language;
- the semantic rules for interpreting C programs;
- the representation of input data to be processed by C programs;
- the representation of output data produced by C programs;
- the restrictions and limits imposed by a conforming implementation of C.

This International Standard does not specify

- the mechanism by which C programs are transformed for use by a data-processing system;
- the mechanism by which C programs are invoked for use by a data-processing system;
— the mechanism by which input data are transformed for use by a C program;
- the mechanism by which output data are transformed after being produced by a C program;
- the size or complexity of a program and its data that will exceed the capacity of any specific data-processing system or the capacity of a particular processor;
- all minimal requirements of a data-processing system that is capable of supporting a conforming implementation.

[^0]
## 2. Normative references

The following standards contain provisions which, through reference in this text, constitute provisions of this International Standard. At the time of publication, the editions indicated were valid. All standards are subject to revision, and parties to agreements based on this International Standard are encouraged to investigate the possibility of applying the most recent editions of the standards indicated below. Members of IEC and ISO maintain registers of currently valid International Standards.

IEC 559:1993, Binary floating-point arithmetic for microprocessor systems, second edition.

ISO 646:1983, Information processing - ISO 7-bit coded character set for information interchange.

ISO/IEC 2382-1:1993, Information technology - Vocabulary - Part 1: Fundamental terms.

ISO 4217:1987, Codes for the representation of currencies and funds.
ISO 8601:1988, Data elements and interchange formats - Information interchange - Representation of dates and times.

ISO/IEC TR 10176, Information technology - Guidelines for the preparation of programming language standards.

ISO/IEC 10646-1:1993, Information technology - Universal Multiple-Octet Coded Character Set (UCS) - Part 1: Architecture and Basic Multilingual Plane.

## 3. Definitions and conventions

1 In this International Standard, "shall" is to be interpreted as a requirement on an implementation or on a program; conversely, "shall not" is to be interpreted as a prohibition.

2 For the purposes of this International Standard, the following definitions apply. Other terms are defined where they appear in italic type or being on the left side of a syntax rule. Terms explicitly defined in this International Standard are not to be presumed to refer implicitly to similar terms defined elsewhere. Terms not defined in this International Standard are to be interpreted according to ISO 2382-1.

### 3.1 Alignment

1 A requirement that objects of a particular type be located on storage boundaries with addresses that are particular multiples of a byte address.

### 3.2 Argument

1 An expression in the comma-separated list bounded by the parentheses in a function call expression, or a sequence of preprocessing tokens in the comma-separated list bounded by the parentheses in a function-like macro invocation. Also known as "actual argument" or "actual parameter."

### 3.3 Bit

1 The unit of data storage in the execution environment large enough to hold an object that may have one of two values. It need not be possible to express the address of each individual bit of an object.

### 3.4 Byte

1 The unit of data storage large enough to hold any member of the basic character set of the execution environment. It shall be possible to express the address of each individual byte of an object uniquely. A byte is composed of a contiguous sequence of bits, the number of which is implementation-defined. The least significant bit is called the low-order bit; the most significant bit is called the high-order bit.

### 3.5 Character

1 A bit representation that fits in a byte. The representation of each member of the basic character set in both the source and execution environments shall fit in a byte.

### 3.6 Constraints

1 Restrictions, both syntactic and semantic, by which the exposition of language elements is to be interpreted.

### 3.7 Correctly rounded result

1 A representation in the result format that is nearest in value, subject to the effective rounding mode, to what the result would be given unlimited range and precision.

### 3.8 Diagnostic message

1 A message belonging to an implementation-defined subset of the implementation's message output.

### 3.9 Forward references

1 References to later subclauses of this International Standard that contain additional information relevant to this subclause.

### 3.10 Implementation

1 A particular set of software, running in a particular translation environment under particular control options, that performs translation of programs for, and supports execution of functions in, a particular execution environment.

### 3.11 Implementation-defined behavior

1 Unspecified behavior where each implementation shall document how the choice is made.

### 3.12 Implementation limits

1 Restrictions imposed upon programs by the implementation.

### 3.13 Locale-specific behavior

1 Behavior that depends on local conventions of nationality, culture, and language that each implementation shall document.

### 3.14 Multibyte character

1 A sequence of one or more bytes representing a member of the extended character set of either the source or the execution environment. The extended character set is a superset of the basic character set.

### 3.15 Object

1 A region of data storage in the execution environment, the contents of which can represent values. Except for bit-fields, objects are composed of contiguous sequences of one or more bytes, the number, order, and encoding of which are either explicitly specified or implementation-defined. When referenced, an object may be interpreted as having a particular type; see 6.2.2.1.

### 3.16 Parameter

1 An object declared as part of a function declaration or definition that acquires a value on entry to the function, or an identifier from the comma-separated list bounded by the parentheses immediately following the macro name in a function-like macro definition. Also known as "formal argument'" or 'formal parameter."

### 3.17 Recommended practice

1 Sections so entitled contain specification that is strongly recommended as being in keeping with the intent of the standard, but that may be impractical for some implementations.

### 3.18 Undefined behavior

1 Behavior, upon use of a nonportable or erroneous program construct, of erroneous data, or of indeterminately valued objects, for which this International Standard imposes no requirements. Permissible undefined behavior ranges from ignoring the situation completely with unpredictable results, to behaving during translation or program execution in a documented manner characteristic of the environment (with or without the issuance of a diagnostic message), to terminating a translation or execution (with the issuance of a diagnostic message).

2 If a "shall" or "shall not" requirement that appears outside of a constraint is violated, the behavior is undefined. Undefined behavior is otherwise indicated in this International Standard by the words "undefined behavior" or by the omission of any explicit definition of behavior. There is no difference in emphasis among these three; they all describe "behavior that is undefined."

3 The implementation must successfully translate a given program unless a syntax error is detected, a constraint is violated, or it can determine that every possible execution of that program would result in undefined behavior.

### 3.19 Unspecified behavior

1 Behavior where this International Standard provides two or more possibilities and imposes no requirements on which is chosen in any instance. A program that is correct in all other aspects, operating on correct data, containing unspecified behavior shall be a correct program and act in accordance with subclause 5.1.2.3.

## Examples

1. An example of unspecified behavior is the order in which the arguments to a function are evaluated.
2. An example of undefined behavior is the behavior on integer overflow.
3. An example of implementation-defined behavior is the propagation of the highorder bit when a signed integer is shifted right.
4. An example of locale-specific behavior is whether the islower function returns true for characters other than the 26 lowercase Latin letters.

Forward references: bitwise shift operators (6.3.7), expressions (6.3), function calls (6.3.2.2), the islower function (7.3.1.7), localization (7.5).

## 4. Compliance

1 A strictly conforming program shall use only those features of the language and library specified in this International Standard. ${ }^{2}$ It shall not produce output dependent on any unspecified, undefined, or implementation-defined behavior, and shall not exceed any minimum implementation limit.

2 The two forms of conforming implementation are hosted and freestanding. A conforming hosted implementation shall accept any strictly conforming program. A conforming freestanding implementation shall accept any strictly conforming program in which the use of the features specified in the library clause (clause 7) is confined to the contents of the standard headers <float.h>, <limits.h>, <stdarg.h>, <stddef.h>, and <iso646.h>. A conforming implementation may have extensions (including additional library functions), provided they do not alter the behavior of any strictly conforming program. ${ }^{3}$

3 A conforming program is one that is acceptable to a conforming implementation. ${ }^{4}$
4 An implementation shall be accompanied by a document that defines all implementation-defined characteristics and all extensions.

Forward references: limits <float.h> and <limits.h> (7.1.5), variable arguments <stdarg. $\mathrm{h}>$ (7.12), common definitions <stddef. $\mathrm{h}>$ (7.1.6), alternate spellings <iso646.h> (7.17).

```
2. This implies that a strictly conforming program can use features in a conditionally normative annex provided the use is conditioned by a \#ifdef directive with the conformance macro for the annex, as in
```

```
#ifdef __STDC_IEC_559__ /* FE_UPWARD defined */
```

\#ifdef __STDC_IEC_559__ /* FE_UPWARD defined */
/* ... */
/* ... */
fesetround(FE_UPWARD);
fesetround(FE_UPWARD);
/*... */
/*... */
\#endif

```
    #endif
```

3. This implies that a conforming implementation reserves no identifiers other than those explicitly reserved in this International Standard.
4. Strictly conforming programs are intended to be maximally portable among conforming implementations. Conforming programs may depend upon nonportable features of a conforming implementation.

## 5. Environment

1 An implementation translates C source files and executes C programs in two data-processing-system environments, which will be called the translation environment and the execution environment in this International Standard. Their characteristics define and constrain the results of executing conforming C programs constructed according to the syntactic and semantic rules for conforming implementations.

Forward references: In the environment clause (clause 5), only a few of many possible forward references have been noted.

### 5.1 Conceptual models

### 5.1.1 Translation environment

### 5.1.1.1 Program structure

1 A C program need not all be translated at the same time. The text of the program is kept in units called source files, also known as preprocessing files, in this International Standard. A source file together with all the headers and source files included via the preprocessing directive \#include is known as a preprocessing translation unit. After preprocessing, a preprocessing translation unit is called a translation unit. Previously translated translation units may be preserved individually or in libraries. The separate translation units of a program communicate by (for example) calls to functions whose identifiers have external linkage, manipulation of objects whose identifiers have external linkage, or manipulation of data files. Translation units may be separately translated and then later linked to produce an executable program.

Forward references: conditional inclusion (6.8.1), linkages of identifiers (6.1.2.2), source file inclusion (6.8.2), external definitions (6.7), preprocessing directives (6.8).

### 5.1.1.2 Translation phases

1 The precedence among the syntax rules of translation is specified by the following phases. ${ }^{5}$

1. Physical source file multibyte characters are mapped to the source character set (introducing new-line characters for end-of-line indicators) if necessary. Any multibyte source file character not in the basic source character set is replaced

[^1]by the universal-character-name that designates that multibyte character. ${ }^{6}$ Then, trigraph sequences are replaced by corresponding single-character internal representations.
2. Each instance of a backslash character immediately followed by a newline character is deleted, splicing physical source lines to form logical source lines. Only the last backslash on any physical source line shall be eligible for being part of such a splice. A source file that is not empty shall end in a new-line character, which shall not be immediately preceded by a backslash character before any such splicing takes place.
3. The source file is decomposed into preprocessing tokens ${ }^{7}$ and sequences of white-space characters (including comments). A source file shall not end in a partial preprocessing token or comment. Each comment is replaced by one space character. New-line characters are retained. Whether each nonempty sequence of white-space characters other than new-line is retained or replaced by one space character is implementation-defined.
4. Preprocessing directives are executed, macro invocations are expanded, and pragma unary operator expressions are executed. If a character sequence that matches the syntax of a universal-character-name is produced by token concatenation (6.8.3.3), the behavior is undefined. A \#include preprocessing directive causes the named header or source file to be processed from phase 1 through phase 4, recursively. All preprocessing directives are then deleted.
5. Each source character set member, escape sequence, and universal-charactername in character constants and string literals is converted to a member of the execution character set.
6. Adjacent character string literal tokens are concatenated and adjacent wide string literal tokens are concatenated.
7. White-space characters separating tokens are no longer significant. Each preprocessing token is converted into a token. The resulting tokens are
6. The process of handling extended characters is specified in terms of mapping to an encoding that uses only the basic source character set, and, in the case of character literals and strings, further mapping to the execution character set. In practical terms, however, any internal encoding may be used, so long as an actual extended character encountered in the input, and the same extended character expressed in the input as a universal-character-name (i.e., using the $\backslash \boldsymbol{U}$ or $\backslash \mathbf{u}$ notation), are handled equivalently.
7. As described in 6.1, the process of dividing a source file's characters into preprocessing tokens is context-dependent. For example, see the handling of < within a \#include preprocessing directive.
syntactically and semantically analyzed and translated as a translation unit.
8. All external object and function references are resolved. Library components are linked to satisfy external references to functions and objects not defined in the current translation. All such translator output is collected into a program image which contains information needed for execution in its execution environment.

## Constraints

2 A universal-character-name shall not specify a character short identifier in the range 0000 through 0020 or 007 F through 009 F inclusive. A universal-character-name shall not designate a character in the basic source character set.

Forward references: lexical elements (6.1), preprocessing directives (6.8), trigraph sequences (5.2.1.1), external definitions (6.7).

### 5.1.1.3 Diagnostics

1 A conforming implementation shall produce at least one diagnostic message (identified in an implementation-defined manner) if a preprocessing translation unit or translation unit contains a violation of any syntax rule or constraint, even if the behavior is also explicitly specified as undefined or implementation-defined. Diagnostic messages need not be produced in other circumstances. ${ }^{8}$

## Examples

2 An implementation shall issue a diagnostic for the translation unit:

```
char i;
int i;
```

because in those cases where wording in this International Standard describes the behavior for a construct as being both a constraint error and resulting in undefined behavior, the constraint error shall be diagnosed.

[^2]
### 5.1.2 Execution environments

1 Two execution environments are defined: freestanding and hosted. In both cases, program startup occurs when a designated C function is called by the execution environment. All objects in static storage shall be initialized (set to their initial values) before program startup. The manner and timing of such initialization are otherwise unspecified. Program termination returns control to the execution environment.

Forward references: initialization (6.5.8).

### 5.1.2.1 Freestanding environment

1 In a freestanding environment (in which C program execution may take place without any benefit of an operating system), the name and type of the function called at program startup are implementation-defined. Any library facilities available to a freestanding program are implementation-defined.

2 The effect of program termination in a freestanding environment is implementationdefined.

### 5.1.2.2 Hosted environment

1 A hosted environment need not be provided, but shall conform to the following specifications if present.

### 5.1.2.2.1 Program startup

1 The function called at program startup is named main. The implementation declares no prototype for this function. It shall be defined with no parameters:
int main(void) \{ /* ... */ \}
or with two parameters (referred to here as argc and argv, though any names may be used, as they are local to the function in which they are declared):
int main(int argc, char *argv[]) \{ /* ... */ \} or equivalent, ${ }^{9}$ or in some other implementation-defined manner.
2 If they are defined, the parameters to the main function shall obey the following constraints:

[^3]- The value of argc shall be nonnegative.
- argv[argc] shall be a null pointer.
- If the value of argc is greater than zero, the array members argv[0] through argv[argc-1] inclusive shall contain pointers to strings, which are given implementation-defined values by the host environment prior to program startup. The intent is to supply to the program information determined prior to program startup from elsewhere in the hosted environment. If the host environment is not capable of supplying strings with letters in both uppercase and lowercase, the implementation shall ensure that the strings are received in lowercase.
- If the value of argc is greater than zero, the string pointed to by argv[0] represents the program name; argv[0][0] shall be the null character if the program name is not available from the host environment. If the value of argc is greater than one, the strings pointed to by argv[1] through argv[argc-1] represent the program parameters.
- The parameters argc and argv and the strings pointed to by the argv array shall be modifiable by the program, and retain their last-stored values between program startup and program termination.


### 5.1.2.2.2 Program execution

1 In a hosted environment, a program may use all the functions, macros, type definitions, and objects described in the library clause (clause 7).

### 5.1.2.2.3 Program termination

1 A return from the initial call to the main function is equivalent to calling the exit function with the value returned by the main function as its argument. ${ }^{10}$ If the \} that terminates the main function is reached, the termination status returned to the host environment is unspecified.

Forward references: definition of terms (7.1.1), the exit function (7.14.4.3).

[^4]
### 5.1.2.3 Program execution

1 The semantic descriptions in this International Standard describe the behavior of an abstract machine in which issues of optimization are irrelevant.

2 Accessing a volatile object, modifying an object, modifying a file, or calling a function that does any of those operations are all side effects, ${ }^{11}$ which are changes in the state of the execution environment. Evaluation of an expression may produce side effects. At certain specified points in the execution sequence called sequence points, all side effects of previous evaluations shall be complete and no side effects of subsequent evaluations shall have taken place.

3 In the abstract machine, all expressions are evaluated as specified by the semantics. An actual implementation need not evaluate part of an expression if it can deduce that its value is not used and that no needed side effects are produced (including any caused by calling a function or accessing a volatile object).

4 When the processing of the abstract machine is interrupted by receipt of a signal, only the values of objects as of the previous sequence point may be relied on. Objects that may be modified between the previous sequence point and the next sequence point need not have received their correct values yet.

5 An instance of each object with automatic storage duration is associated with each entry into its block. Such an object exists and retains its last-stored value during the execution of the block and while the block is suspended (by a call of a function or receipt of a signal).
6 The least requirements on a conforming implementation are:

- At sequence points, volatile objects are stable in the sense that previous evaluations are complete and subsequent evaluations have not yet occurred.
- At program termination, all data written into files shall be identical to the result that execution of the program according to the abstract semantics would have produced.
- The input and output dynamics of interactive devices shall take place as specified in 7.13.3. The intent of these requirements is that unbuffered or line-buffered

11. The IEC 559 standard for binary floating-point arithmetic requires certain status flags and control modes, with user access. Floating-point operations implicitly set the status flags; modes affect result values of floating-point operations. Implementations that support such floating-point state will need to regard changes to it as side effects - see Annex F for details. The floating-point environment library <fenv. h > provides a programming facility for indicating when these side effects matter, freeing the implementations in other cases.
output appear as soon as possible, to ensure that prompting messages actually appear prior to a program waiting for input.

7 What constitutes an interactive device is implementation-defined.
8 More stringent correspondences between abstract and actual semantics may be defined by each implementation.

## Examples

1. An implementation might define a one-to-one correspondence between abstract and actual semantics: at every sequence point, the values of the actual objects would agree with those specified by the abstract semantics. The keyword volatile would then be redundant.

Alternatively, an implementation might perform various optimizations within each translation unit, such that the actual semantics would agree with the abstract semantics only when making function calls across translation unit boundaries. In such an implementation, at the time of each function entry and function return where the calling function and the called function are in different translation units, the values of all externally linked objects and of all objects accessible via pointers therein would agree with the abstract semantics. Furthermore, at the time of each such function entry the values of the parameters of the called function and of all objects accessible via pointers therein would agree with the abstract semantics. In this type of implementation, objects referred to by interrupt service routines activated by the signal function would require explicit specification of volatile storage, as well as other implementation-defined restrictions.
2. In executing the fragment

```
char c1, c2;
/* ... */
c1 = c1 + c2;
```

the "integer promotions'" require that the abstract machine promote the value of each variable to int size and then add the two ints and truncate the sum. Provided the addition of two chars can be done without overflow, or with overflow wrapping silently to produce the correct result, the actual execution need only produce the same result, possibly omitting the promotions.
3. Similarly, in the fragment

```
float f1, f2;
double d;
/* ... */
f1 = f2 * d;
```

the multiplication may be executed using single-precision arithmetic if the implementation can ascertain that the result would be the same as if it were executed using double-precision arithmetic (for example, if $\mathbf{d}$ were replaced by the constant 2.0 , which has type double).
4. Implementations employing wide registers must take care to honor appropriate semantics. Values must be independent of whether they are represented in a register or in memory. For example, an implicit spilling of a register must not alter the value. Also, an explicit store and load must round to the precision of the storage type. In particular, casts and assignments must perform their specified conversion: for the fragment

```
double d1, d2;
float f;
d1 = f = expression;
d2 = (float) expressions;
```

the values assigned to $\mathbf{d} 1$ and $\mathbf{d} 2$ must have been converted to float.
5. Rearrangement for floating-point expressions is often restricted because of limitations in precision as well as range. The implementation cannot generally apply the mathematical associative rules for addition or multiplication, nor the distributive rule, because of roundoff error, even in the absence of overflow and underflow. Likewise, implementations cannot generally replace decimal constants in order to rearrange expressions. In the following fragment, rearrangements suggested by mathematical rules for real numbers are often not valid. See Annex F.8.

```
double x, y, z;
/* ... */
x = (x * y) * z; // not equivalent to x *= y * z;
z = (x - y) + y ; // not equivalent to z = x;
z = x + x * y; // not equivalent to z = x * (1.0 + y);
y = x / 5.0; // not equivalent of y = x * 0.2;
```

6. To illustrate the grouping behavior of expressions, in the following fragment
```
int a, b;
/* ... */
a = a + 32760 + b + 5;
```

the expression statement behaves exactly the same as

$$
a=(((a+32760)+b)+5) ;
$$

due to the associativity and precedence of these operators. Thus, the result of the sum $(\mathbf{a}+32760)$ is next added to $\mathbf{b}$, and that result is then added to $\mathbf{5}$ which results in the value assigned to $\mathbf{a}$. On a machine in which overflows produce an explicit trap and in which the range of values representable by an int is $[-32768,+32767]$, the implementation cannot rewrite this expression as

```
a = ((a + b) + 32765);
```

since if the values for $\mathbf{a}$ and $\mathbf{b}$ were, respectively, -32754 and -15 , the sum $\mathbf{a}+\mathbf{b}$ would produce a trap while the original expression would not; nor can the expression be rewritten either as

$$
a=((a+32765)+b) ;
$$

or

$$
a=(a+(b+32765)) ;
$$

since the values for $\mathbf{a}$ and $\mathbf{b}$ might have been, respectively, 4 and -8 or -17 and 12. However on a machine in which overflow silently generates some value and where positive and negative overflows cancel, the above expression statement can be rewritten by the implementation in any of the above ways because the same result will occur.
7. The grouping of an expression does not completely determine its evaluation. In the following fragment

```
#include <stdio.h>
int sum;
char *p;
/* ... */
sum = sum * 10 - '0' + (*p++ = getchar());
```

the expression statement is grouped as if it were written as

```
sum = (((sum * 10) - '0') + ((*(p++)) = (getchar())));
```

but the actual increment of $\mathbf{p}$ can occur at any time between the previous sequence point and the next sequence point (the ; ), and the call to getchar can occur at any point prior to the need of its returned value.

Forward references: compound statement, or block (6.6.2), expressions (6.3), files (7.13.3), sequence points (6.3, 6.6), the signal function (7.11), type qualifiers (6.5.3).

### 5.2 Environmental considerations

### 5.2.1 Character sets

1 Two sets of characters and their associated collating sequences shall be defined: the set in which source files are written, and the set interpreted in the execution environment. The values of the members of the execution character set are implementation-defined; any additional members beyond those required by this subclause are locale-specific.

2 In a character constant or string literal, members of the execution character set shall be represented by corresponding members of the source character set or by escape sequences consisting of the backslash $\backslash$ followed by one or more characters. A byte with all bits set to 0 , called the null character, shall exist in the basic execution character set; it is used to terminate a character string.

3 Both the basic source and basic execution character sets shall have at least the following members: the 26 uppercase letters of the Latin alphabet

| A | $\mathbf{B}$ | $\mathbf{C}$ | $\mathbf{D}$ | $\mathbf{E}$ | $\mathbf{F}$ | $\mathbf{G}$ | $\mathbf{H}$ | $\mathbf{I}$ | $\mathbf{J}$ | $\mathbf{K}$ | $\mathbf{L}$ | $\mathbf{M}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{N}$ | $\mathbf{O}$ | $\mathbf{P}$ | $\mathbf{Q}$ | $\mathbf{R}$ | $\mathbf{S}$ | $\mathbf{T}$ | $\mathbf{U}$ | $\mathbf{V}$ | $\mathbf{W}$ | $\mathbf{X}$ | $\mathbf{Y}$ | $\mathbf{Z}$ |

the 26 lowercase letters of the Latin alphabet

the 10 decimal digits

```
0
```

the following 29 graphic characters

the space character, and control characters representing horizontal tab, vertical tab, and form feed. In both the source and execution basic character sets, the value of each character after 0 in the above list of decimal digits shall be one greater than the value of the previous. In source files, there shall be some way of indicating the end of each line of text; this International Standard treats such an end-of-line indicator as if it were a single new-line character. In the execution character set, there shall be control characters representing alert, backspace, carriage return, and new line. If any other characters are encountered in a source file (except in a character constant, a string literal, a header name, a comment, or a preprocessing token that is never converted to a token), the behavior is undefined.

4 The universal-character-name construct provides a way to name other characters.

```
hex-quad:
            hexadecimal-digit hexadecimal-digit
            hexadecimal-digit hexadecimal-digit
universal-character-name:
    \u hex-quad
    \U hex-quad hex-quad
```

5 The character designated by the universal-character-name \Unnnnnnnn is that character whose character short identifier is nnnnnnnn specified by ISO/IEC 10646-1; the character designated by the universal-character-name \unnnn is that character whose character short identifier is 0000nnnn specified by ISO/IEC 10646-1.
Forward references: identifiers (6.1.2), character constants (6.1.3.4), preprocessing directives (6.8), string literals (6.1.4), comments (6.1.9), string (7.1.1).

### 5.2.1.1 Trigraph sequences

1 All occurrences in a source file of the following sequences of three characters (called trigraph sequences ${ }^{12}$ ) are replaced with the corresponding single character.


No other trigraph sequences exist. Each ? that does not begin one of the trigraphs listed above is not changed.

[^5]
## Examples

2 The following source line
printf("Eh???/n");
becomes (after replacement of the trigraph sequence ??/)
printf("Eh?\n");

### 5.2.1.2 Multibyte characters

1 The source character set may contain multibyte characters, used to represent members of the extended character set. The execution character set may also contain multibyte characters, which need not have the same encoding as for the source character set. For both character sets, the following shall hold:

- The single-byte characters defined in 5.2.1 shall be present.
- The presence, meaning, and representation of any additional members is localespecific.
- A multibyte character may have a state-dependent encoding, wherein each sequence of multibyte characters begins in an initial shift state and enters other locale-specific shift states when specific multibyte characters are encountered in the sequence. While in the initial shift state, all single-byte characters retain their usual interpretation and do not alter the shift state. The interpretation for subsequent bytes in the sequence is a function of the current shift state.
- A byte with all bits zero shall be interpreted as a null character independent of shift state.
- A byte with all bits zero shall not occur in the second or subsequent bytes of a multibyte character.

For the source character set, the following shall hold:

- A comment, string literal, character constant, or header name shall begin and end in the initial shift state.
- A comment, string literal, character constant, or header name shall consist of a sequence of valid multibyte characters.


### 5.2.2 Character display semantics

1 The active position is that location on a display device where the next character output by the fputc function would appear. The intent of writing a printable character (as defined by the isprint function) to a display device is to display a graphic representation of that character at the active position and then advance the active position to the next position on the current line. The direction of writing is localespecific. If the active position is at the final position of a line (if there is one), the behavior is unspecified.

2 Alphabetic escape sequences representing nongraphic characters in the execution character set are intended to produce actions on display devices as follows:
\a (alert) Produces an audible or visible alert. The active position shall not be changed.
$\backslash \mathbf{b}$ (backspace) Moves the active position to the previous position on the current line. If the active position is at the initial position of a line, the behavior is unspecified.
$\backslash \mathbf{f}($ form feed $)$ Moves the active position to the initial position at the start of the next logical page.

In (new line) Moves the active position to the initial position of the next line.
Ir (carriage return) Moves the active position to the initial position of the current line.
\t (horizontal tab) Moves the active position to the next horizontal tabulation position on the current line. If the active position is at or past the last defined horizontal tabulation position, the behavior is unspecified.

Iv (vertical tab) Moves the active position to the initial position of the next vertical tabulation position. If the active position is at or past the last defined vertical tabulation position, the behavior is unspecified.

3 Each of these escape sequences shall produce a unique implementation-defined value which can be stored in a single char object. The external representations in a text file need not be identical to the internal representations, and are outside the scope of this International Standard.

Forward references: the isprint function (7.3.1.8), the fputc function (7.13.7.3).

### 5.2.3 Signals and interrupts

1 Functions shall be implemented such that they may be interrupted at any time by a signal, or may be called by a signal handler, or both, with no alteration to earlier, but still active, invocations' control flow (after the interruption), function return values, or objects with automatic storage duration. All such objects shall be maintained outside the function image (the instructions that comprise the executable representation of a function) on a per-invocation basis.

### 5.2.4 Environmental limits

1 Both the translation and execution environments constrain the implementation of language translators and libraries. The following summarizes the environmental limits on a conforming implementation.

### 5.2.4.1 Translation limits

1 The implementation shall be able to translate and execute at least one program that contains at least one instance of every one of the following limits: ${ }^{13}$

- 127 nesting levels of compound statements, iteration statements, and selection statements
- 63 nesting levels of conditional inclusion
- 12 pointer, array, and function declarators (in any combinations) modifying an arithmetic, structure, union, or incomplete type in a declaration
- 63 nesting levels of parenthesized declarators within a full declarator
- 63 nesting levels of parenthesized expressions within a full expression
- 63 significant initial characters in an internal identifier or a macro name
- 31 significant initial characters in an external identifier
- 4095 external identifiers in one translation unit
- 511 identifiers with block scope declared in one block
- 4095 macro identifiers simultaneously defined in one preprocessing translation unit
- 127 parameters in one function definition

13. Implementations should avoid imposing fixed translation limits whenever possible.

- 127 arguments in one function call
- 127 parameters in one macro definition
- 127 arguments in one macro invocation
- 4095 characters in a logical source line
- 4095 characters in a character string literal or wide string literal (after concatenation)
- 65535 bytes in an object (in a hosted environment only)
- 15 nesting levels for \#included files
- 1023 case labels for a switch statement (excluding those for any nested switch statements)
- 1023 members in a single structure or union
- 1023 enumeration constants in a single enumeration
- 63 levels of nested structure or union definitions in a single struct-declaration-list


### 5.2.4.2 Numerical limits

1 A conforming implementation shall document all the limits specified in this subclause, which shall be specified in the headers <limits.h> and <float.h>.

### 5.2.4.2.1 Sizes of integer types <limits.h>

1 The values given below shall be replaced by constant expressions suitable for use in \#if preprocessing directives. Moreover, except for CHAR_BIT and MB_LEN_MAX, the following shall be replaced by expressions that have the same type as would an expression that is an object of the corresponding type converted according to the integer promotions. Their implementation-defined values shall be equal or greater in magnitude (absolute value) to those shown, with the same sign.

- number of bits for smallest object that is not a bit-field (byte)

CHAR_BIT

- minimum value for an object of type signed char

SCHAR_MIN -127

- maximum value for an object of type signed char SCHAR_MAX +127
- maximum value for an object of type unsigned char

UCHAR_MAX 255

- minimum value for an object of type char

CHAR_MIN
see below

- maximum value for an object of type char

CHAR_MAX
see below

- maximum number of bytes in a multibyte character, for any supported locale

MB_LEN_MAX 1

- minimum value for an object of type short int

SHRT_MIN -32767

- maximum value for an object of type short int

SHRT_MAX +32767

- maximum value for an object of type unsigned short int USHRT_MAX

65535

- minimum value for an object of type int

INT_MIN
-32767

- maximum value for an object of type int

INT_MAX +32767

- maximum value for an object of type unsigned int UINT_MAX 65535
- minimum value for an object of type long int

LONG_MIN $\mathbf{- 2 1 4 7 4 8 3 6 4 7}$

- maximum value for an object of type long int

LONG_MAX +2147483647

- maximum value for an object of type unsigned long int

ULONG_MAX
4294967295

- minimum value for an object of type long long int

LLONG_MIN -9223372036854775807

- maximum value for an object of type long long int

LLONG_MAX +9223372036854775807

- maximum value for an object of type unsigned long long int

ULLONG_MAX 18446744073709551615
If the value of an object of type char is treated as a signed integer when used in an expression, the value of CHAR_MIN shall be the same as that of SCHAR_MIN and the value of CHAR_MAX shall be the same as that of SCHAR_MAX. Otherwise, the value of CHAR_MIN shall be 0 and the value of CHAR_MAX shall be the same as that of

UCHAR_MAX. ${ }^{14}$ The value UCHAR_MAX+1 shall equal 2 raised to the power CHAR_BIT.

### 5.2.4.2.2 Characteristics of floating types <float.h>

1 The characteristics of floating types are defined in terms of a model that describes a representation of floating-point numbers and values that provide information about an implementation's floating-point arithmetic. ${ }^{15}$ The following parameters are used to define the model for each floating-point type:
$s \quad \operatorname{sign}( \pm 1)$
$b \quad$ base or radix of exponent representation (an integer $>1$ )
$e \quad$ exponent (an integer between a minimum $e_{\min }$ and a maximum $e_{\max }$ )
$p \quad$ precision (the number of base- $b$ digits in the significand)
$f_{k} \quad$ nonnegative integers less than $b$ (the significand digits)
2 A normalized floating-point number $x\left(f_{1}>0\right.$ if $\left.x \neq 0\right)$ is defined by the following model:

$$
x=s \times b^{e} \times \sum_{k=1}^{p} f_{k} \times b^{-k}, \quad e_{\min } \leq e \leq e_{\max }
$$

3 Floating types might include values that are not normalized floating-point numbers, for example subnormal numbers $\left(x \neq 0, e=e_{\min }, f_{1}=0\right)$, infinities, and NaNs. A $N a N$ is an encoding signifying Not-a-Number. A quiet NaN propagates through almost every arithmetic operation without raising an exception; a signaling NaN generally raises an exception when occurring as an arithmetic operand. ${ }^{16}$

4 All integer values in the <float.h> header, except FLT_ROUNDS, shall be constant expressions suitable for use in \#if preprocessing directives; all floating values shall be constant expressions. All except FLT_EVAL_METHOD, FLT_RADIX, and FLT_ROUNDS have separate names for all three floating-point types. The floatingpoint model representation is provided for all values except FLT_EVAL_METHOD and FLT_ROUNDS.

5 The rounding mode for floating-point addition is characterized by the value of FLT_ROUNDS: ${ }^{17}$

[^6]-1 indeterminable
0 toward zero
1 to nearest
2 toward positive infinity
3 toward negative infinity
All other values for FLT_ROUNDS characterize implementation-defined rounding behavior.

The values of operations with floating operands and values subject to the usual arithmetic conversions and of floating constants are evaluated to a format whose range and precision may be greater than required by the type. The use of evaluation formats is characterized by the value of FLT_EVAL_METHOD: ${ }^{18}$
-1 indeterminable;
0 evaluate all operations and constants just to the range and precision of the type;
1 evaluate operations and constants of type float and double to the range and precision of the double type, evaluate long double operations and constants to the range and precision of the long double type;

2 evaluate all operations and constants to the range and precision of the long double type.

All other negative values for FLT_EVAL_METHOD characterize implementationdefined behavior.

7 The values given in the following list shall be replaced by implementation-defined expressions that shall be equal or greater in magnitude (absolute value) to those shown, with the same sign:

- radix of exponent representation, $b$

FLT_RADIX

## 2

- number of base-FLT_RADIX digits in the floating-point significand, $p$

18. The evaluation method determines evaluation formats of expressions involving all floating types, not just real types. For example, if FLT_EVAL_METHOD is 1, then the product of two float complex operands is represented in the double complex format, and its parts are evaluated to double.
```
FLT_MANT_DIG
DBL_MANT_DIG
LDBL_MANT_DIG
```

- number of decimal digits, $q$, such that any floating-point number with $q$ decimal digits can be rounded into a floating-point number with $p$ radix $b$ digits and back again without change to the $q$ decimal digits,
$\left\lfloor(p-1) \times \log _{10} b\right\rfloor+ \begin{cases}1 & \text { if } b \text { is a power of } 10 \\ 0 & \text { otherwise }\end{cases}$

| FLT_DIG | 6 |
| :--- | ---: |
| DBL_DIG | 10 |

- minimum negative integer such that $\mathbf{F L T} \_$RADIX raised to that power minus 1 is a normalized floating-point number, $e_{\text {min }}$

FLT_MIN_EXP
DBL_MIN_EXP
LDBL_MIN_EXP

- minimum negative integer such that 10 raised to that power is in the range of normalized floating-point numbers, $\left\lceil\log _{10} b^{e_{\min }-1}\right\rceil$
FLT_MIN_10_EXP -37
DBL_MIN_10_EXP -37
LDBL_MIN_10_EXP -37
- maximum integer such that FLT_RADIX raised to that power minus 1 is a representable finite floating-point number, $e_{\max }$

FLT_MAX_EXP
DBL_MAX_EXP
LDBL_MAX_EXP

- maximum integer such that 10 raised to that power is in the range of representable finite floating-point numbers, $\left\lfloor\log _{10}\left(\left(1-b^{-p}\right) \times b^{e_{\text {max }}}\right)\right\rfloor$
FLT_MAX_10_EXP
+37
DBL_MAX_10_EXP +37
LDBL_MAX_10_EXP +37
The values given in the following list shall be replaced by implementation-defined expressions with values that shall be equal to or greater than those shown:
- maximum representable finite floating-point number, $\left(1-b^{-p}\right) \times b^{e_{\text {max }}}$

```
FLT_MAX
1E+37
DBL_MAX
1E+37
LDBL_MAX
1E+37
```

The values given in the following list shall be replaced by implementation-defined expressions with values that shall be equal to or less than those shown:

- the difference between 1 and the least value greater than 1 that is representable in the given floating point type, $b^{1-p}$

```
FLT_EPSILON 1E-5
DBL_EPSILON 1E-9
LDBL_EPSILON 1E-9
```

- minimum normalized positive floating-point number, $b^{e_{\min }-1}$

| FLT_MIN | $1 \mathrm{E}-37$ |
| :--- | :--- |
| DBL_MIN | $1 \mathrm{E}-37$ |
| LDBL_MIN | $1 \mathrm{E}-37$ |

Examples

1. The following describes an artificial floating-point representation that meets the minimum requirements of this International Standard, and the appropriate values in a <float. $\mathrm{h}>$ header for type float:

$$
x=s \times 16^{e} \times \sum_{k=1}^{6} f_{k} \times 16^{-k}, \quad-31 \leq e \leq+32
$$

| FLT_RADIX | 16 |
| :--- | ---: |
| FLT_MANT_DIG | 6 |
| FLT_EPSILON | $9.53674316 \mathrm{E}-07 \mathrm{~F}$ |
| FLT_DIG | 6 |
| FLT_MIN_EXP | -31 |
| FLT_MIN | $2.93873588 \mathrm{E}-39 \mathrm{~F}$ |
| FLT_MIN_10_EXP | -38 |
| FLT_MAX_EXP | +32 |
| FLT_MAX | $3.40282347 \mathrm{E}+38 \mathrm{~F}$ |
| FLT_MAX_10_EXP | +38 |

2. The following describes floating-point representations that also meet the requirements for single-precision and double-precision normalized numbers in IEC $559,{ }^{19}$ and the appropriate values in a <float. $\mathrm{h}>$ header for types float and double:

$$
\begin{array}{ll}
x_{f}=s \times 2^{e} \times \sum_{k=1}^{24} f_{k} \times 2^{-k}, & -125 \leq e \leq+128 \\
x_{d}=s \times 2^{e} \times \sum_{k=1}^{53} f_{k} \times 2^{-k}, & -1021 \leq e \leq+1024
\end{array}
$$

| FLT_RADIX | 2 |  |
| :---: | :---: | :---: |
| FLT_MANT_DIG | 24 |  |
| FLT_EPSILON | 1.19209290E-07F | // decimal constant |
| FLT_EPSILON | 0X1P-23F | // hex constant |
| FLT_DIG | 6 |  |
| FLT_MIN_EXP | -125 |  |
| FLT_MIN | 1.17549435E-38F | // decimal constant |
| FLT_MIN | 0X1P-126F | // hex constant |
| FLT_MIN_10_EXP | -37 |  |
| FLT_MAX_EXP | +128 |  |
| FLT_MAX | 3.40282347E+38F | // decimal constant |
| FLT_MAX | 0X1.fffffepl27F | // hex constant |
| FLT_MAX_10_EXP | +38 |  |
| DBL_MANT_DIG | 53 |  |
| DBL_EPSILON 2. | $460492503131 \mathrm{E}-16$ | // decimal constant |
| DBL_EPSILON | 0X1P-52 | // hex constant |
| DBL_DIG | 15 |  |
| DBL_MIN_EXP | -1021 |  |
| DBL_MIN 2.2 | 38585072014E-308 | // decimal constant |
| DBL_MIN | 0X1P-1022 | // hex constant |
| DBL_MIN_10_EXP | -307 |  |
| DBL_MAX_EXP | +1024 |  |
| DBL_MAX 1.7 | $31348623157 \mathrm{E}+308$ | // decimal constant |
| DBL_MAX OX | ffffffffffeP1023 | // hex constant |
| DBL_MAX_10_EXP | +308 |  |

19. The floating-point model in that standard sums powers of $b$ from zero, so the values of the exponent limits are one less than shown here.

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Forward references: conditional inclusion (6.8.1).

## 6. Language

1 In the syntax notation used in the language clause (clause 6), syntactic categories (nonterminals) are indicated by italic type, and literal words and character set members (terminals) by bold type. A colon (:) following a nonterminal introduces its definition. Alternative definitions are listed on separate lines, except when prefaced by the words '"one of.' An optional symbol is indicated by the subscript 'opt,'" so that

$$
\left\{\text { expression }_{\text {opt }}\right\}
$$

indicates an optional expression enclosed in braces.

### 6.1 Lexical elements

Syntax
1 token:
keyword
identifier
constant
string-literal
operator
punctuator
preprocessing-token:
header-name
identifier
pp-number
character-constant
string-literal
operator
punctuator
each non-white-space character that cannot be one of the above

## Constraints

2 Each preprocessing token that is converted to a token shall have the lexical form of a keyword, an identifier, a constant, a string literal, an operator, or a punctuator.

## Semantics

3 A token is the minimal lexical element of the language in translation phases 7 and 8. The categories of tokens are: keywords, identifiers, constants, string literals, operators, and punctuators. A preprocessing token is the minimal lexical element of the language in translation phases 3 through 6 . The categories of preprocessing token are: header names, identifiers, preprocessing numbers, character constants, string
literals, operators, punctuators, and single non-white-space characters that do not lexically match the other preprocessing token categories. If a ' or a " character matches the last category, the behavior is undefined. Preprocessing tokens can be separated by white space; this consists of comments (described later), or white-space characters (space, horizontal tab, new-line, vertical tab, and form-feed), or both. As described in 6.8 , in certain circumstances during translation phase 4 , white space (or the absence thereof) serves as more than preprocessing token separation. White space may appear within a preprocessing token only as part of a header name or between the quotation characters in a character constant or string literal.

4 If the input stream has been parsed into preprocessing tokens up to a given character, the next preprocessing token is the longest sequence of characters that could constitute a preprocessing token.

5 A header name preprocessing token is only recognized within a \#include preprocessing directive, and within such a directive, a sequence of characters that could be either a header name or a string literal is recognized as the former.

## Examples

1. The program fragment 1Ex is parsed as a preprocessing number token (one that is not a valid floating or integer constant token), even though a parse as the pair of preprocessing tokens $\mathbf{1}$ and $\mathbf{E x}$ might produce a valid expression (for example, if Ex were a macro defined as $\mathbf{+ 1}$ ). Similarly, the program fragment 1E1 is parsed as a preprocessing number (one that is a valid floating constant token), whether or not $\mathbf{E}$ is a macro name.
2. The program fragment $\mathbf{x + + + + + \mathbf { y }}$ is parsed as $\mathbf{x}+\boldsymbol{+}+\boldsymbol{+} \mathbf{y}$, which violates a constraint on increment operators, even though the parse $\mathbf{x}+\boldsymbol{+}+\boldsymbol{+} \mathbf{y}$ might yield a correct expression.

Forward references: character constants (6.1.3.4), comments (6.1.9), expressions (6.3), floating constants (6.1.3.1), header names (6.1.7), macro replacement (6.8.3), postfix increment and decrement operators (6.3.2.4), prefix increment and decrement operators (6.3.3.1), preprocessing directives (6.8), preprocessing numbers (6.1.8), string literals (6.1.4).

### 6.1.1 Keywords

Syntax
1

| keyword: one of |  |  |  |
| :---: | :--- | :--- | :--- |
| auto | break | case | char |
| complex | const | continue | default |
| do | double | else | enum |
| extern | float | for | goto |
| if | imaginary | inline | int |
| long | register | restrict | return |
| short | signed | sizeof | static |
| struct | switch | typedef | union |
| unsigned | void | volatile | while |

## Semantics

2 The token complex is reserved in translation units where the header <complex.h> is included; the token imaginary is reserved in translation units where both the header <complex.h> is included and the macro _Imaginary_I is defined; all other keyword tokens are reserved in all translation units. When reserved, the above tokens (entirely in lowercase) are keywords (in translation phases 7 and 8), and shall not be used otherwise. When the token complex or imaginary is reserved, its use prior to the first inclusion of the header <complex.h> results in undefined behavior.

### 6.1.2 Identifiers

Syntax
identifier:
nondigit
identifier nondigit
identifier digit
nondigit: one of
universal-character-name

digit: one of
$\begin{array}{llllllllll}0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9\end{array}$

## Description

2 An identifier is a sequence of nondigit characters (including the underscore _ and the lowercase and uppercase letters) and digits. Each universal-character-name in an identifier shall designate a character whose encoding in ISO 10646-1 falls into one of the ranges specified in Annex H. ${ }^{20}$ The first character shall be a nondigit character.

## Semantics

3 An identifier can denote an object, a function, or one of the following entities that will be described later: a tag or a member of a structure, union, or enumeration; a typedef name; a label name; a macro name; or a macro parameter. The same identifier can denote different entities at different points in the program. A member of an enumeration is called an enumeration constant. Macro names and macro parameters are not considered further here, because prior to the semantic phase of program translation any occurrences of macro names in the source file are replaced by the preprocessing token sequences that constitute their macro definitions.

4 There is no specific limit on the maximum length of an identifier.
5 When preprocessing tokens are converted to tokens during translation phase 7, if a preprocessing token could be converted to either a keyword or an identifier, it is converted to a keyword.

## Implementation limits

6 The implementation shall treat at least the first 63 characters of an internal name (a macro name or an identifier that does not have external linkage) as significant. The implementation may further restrict the significance of an external name (an identifier that has external linkage) to 31 characters. In both external and internal names, lower-case and upper-case letters are different. The number of significant characters in an identifier is implementation-defined.

7 Any identifiers that differ in a significant character are different identifiers. If two identifiers differ in a nonsignificant character, the behavior is undefined.

[^7]Forward references: linkages of identifiers (6.1.2.2), macro replacement (6.8.3).

### 6.1.2.1 Scopes of identifiers

1 For each different entity that an identifier designates, the identifier is visible (i.e., can be used) only within a region of program text called its scope. Different entities designated by the same identifier either have non-overlapping scopes, or are in different name spaces. There are four kinds of scopes: function, file, block, and function prototype. (A function prototype is a declaration of a function that declares the types of its parameters.)

2 A label name is the only kind of identifier that has function scope. It can be used (in a goto statement) anywhere in the function in which it appears, and is declared implicitly by its syntactic appearance (followed by a : and a statement). Label names shall be unique within a function.

3 Every other identifier has scope determined by the placement of its declaration (in a declarator or type specifier). If the declarator or type specifier that declares the identifier appears outside of any block or list of parameters, the identifier has file scope, which terminates at the end of the translation unit. If the declarator or type specifier that declares the identifier appears inside a block or within the list of parameter declarations in a function definition, the identifier has block scope, which terminates at the \} that closes the associated block. If the declarator or type specifier that declares the identifier appears within the list of parameter declarations in a function prototype (not part of a function definition), the identifier has function prototype scope, which terminates at the end of the function declarator. If an identifier designates two different entities in the same name space, the scopes might overlap. If so, the scope of one entity (the inner scope) will be a strict subset of the scope of the other entity (the outer scope). Within the inner scope, the identifier designates the entity declared in the inner scope; the entity declared in the outer scope is hidden (and not visible) within the inner scope.

4 Unless explicitly stated otherwise, where this International Standard uses the term identifier to refer to some entity (as opposed to the syntactic construct), it refers to the entity in the relevant name space whose declaration is visible at the point the identifier occurs.

5 Two identifiers have the same scope if and only if their scopes terminate at the same point.

Structure, union, and enumeration tags have scope that begins just after the appearance of the tag in a type specifier that declares the tag. Each enumeration constant has scope that begins just after the appearance of its defining enumerator in an enumerator list. Any other identifier has scope that begins just after the completion of its declarator.

Forward references: compound statement, or block (6.6.2), declarations (6.5), enumeration specifiers (6.5.2.2), function calls (6.3.2.2), function declarators (including prototypes) (6.5.5.3), function definitions (6.7.1), the goto statement (6.6.6.1), labeled statements (6.6.1), name spaces of identifiers (6.1.2.3), scope of macro definitions (6.8.3.5), source file inclusion (6.8.2), tags (6.5.2.3), type specifiers (6.5.2).

### 6.1.2.2 Linkages of identifiers

1 An identifier declared in different scopes or in the same scope more than once can be made to refer to the same object or function by a process called linkage. There are three kinds of linkage: external, internal, and none.

2 In the set of translation units and libraries that constitutes an entire program, each instance of a particular identifier with external linkage denotes the same object or function. Within one translation unit, each instance of an identifier with internal linkage denotes the same object or function. Identifiers with no linkage denote unique entities.

3 If the declaration of a file scope identifier for an object or a function contains the storage-class specifier static, the identifier has internal linkage. ${ }^{21}$

4 For an identifier declared with the storage-class specifier extern in a scope in which a prior declaration of that identifier is visible, ${ }^{22}$ if the prior declaration specifies internal or external linkage, the linkage of the identifier at the later declaration becomes the linkage specified at the prior declaration. If no prior declaration is visible, or if the prior declaration specifies no linkage, then the identifier has external linkage.

5 If the declaration of an identifier for a function has no storage-class specifier, its linkage is determined exactly as if it were declared with the storage-class specifier extern. If the declaration of an identifier for an object has file scope and no storage-class specifier, its linkage is external.

6 The following identifiers have no linkage: an identifier declared to be anything other than an object or a function; an identifier declared to be a function parameter; a block scope identifier for an object declared without the storage-class specifier extern.

7 If, within a translation unit, the same identifier appears with both internal and external linkage, the behavior is undefined.

[^8]Forward references: compound statement, or block (6.6.2), declarations (6.5), expressions (6.3), external definitions (6.7).

### 6.1.2.3 Name spaces of identifiers

1 If more than one declaration of a particular identifier is visible at any point in a translation unit, the syntactic context disambiguates uses that refer to different entities. Thus, there are separate name spaces for various categories of identifiers, as follows:

- label names (disambiguated by the syntax of the label declaration and use);
— the tags of structures, unions, and enumerations (disambiguated by following any ${ }^{23}$ of the keywords struct, union, or enum);
- the members of structures or unions; each structure or union has a separate name space for its members (disambiguated by the type of the expression used to access the member via the . or $\rightarrow>$ operator);
- all other identifiers, called ordinary identifiers (declared in ordinary declarators or as enumeration constants).

Forward references: enumeration specifiers (6.5.2.2), labeled statements (6.6.1), structure and union specifiers (6.5.2.1), structure and union members (6.3.2.3), tags (6.5.2.3).

### 6.1.2.4 Storage durations of objects

1 An object has a storage duration that determines its lifetime. There are three storage durations: static, automatic, and allocated. Allocated storage is described in 7.14.3.

2 An object whose identifier is declared with external or internal linkage, or with the storage-class specifier static has static storage duration. For such an object, storage is reserved and its stored value is initialized only once, prior to program startup. The object exists, has a constant address, and retains its last-stored value throughout the execution of the entire program. ${ }^{24}$

An object whose identifier is declared with no linkage and without the storage-class specifier static has automatic storage duration. Storage is guaranteed to be reserved for a new instance of such an object on each normal entry into the block with which it is associated. If the block with which the object is associated is entered by a

[^9]jump from outside the block to a labeled statement in the block or in an enclosed block, then storage is guaranteed to be reserved provided the object does not have a variable length array type. If the object is variably modified and the block is entered by a jump to a labeled statement, then the behavior is undefined. If an initialization is specified for the value stored in the object, it is performed on each normal entry, but not if the block is entered by a jump to a labeled statement beyond the declaration. A backwards jump might cause the initializer to be evaluated more than once; if so, a new value will be stored each time. Storage for the object is no longer guaranteed to be reserved when execution of the block ends in any way. (Entering an enclosed block suspends but does not end execution of the enclosing block. Calling a function suspends but does not end execution of the block containing the call.) The value of a pointer that referred to an object with automatic storage duration that is no longer guaranteed to be reserved is indeterminate. During execution of the associated block, the object has a constant address.

Forward references: compound statement, or block (6.6.2), function calls (6.3.2.2), declarators (6.5.5), array declarators (6.5.5.2), initialization (6.5.8).

### 6.1.2.5 Types

1 The meaning of a value stored in an object or returned by a function is determined by the type of the expression used to access it. (An identifier declared to be an object is the simplest such expression; the type is specified in the declaration of the identifier.) Types are partitioned into object types (types that describe objects), function types (types that describe functions), and incomplete types (types that describe objects but lack information needed to determine their sizes).

An object declared as type char is large enough to store any member of the basic execution character set. If a member of the required source character set enumerated in 5.2.1 is stored in a char object, its value is guaranteed to be positive. If any other character is stored in a char object, the resulting value is implementation-defined but shall be within the range of values that can be represented in that type.

There are five standard signed integer types, designated as signed char, short int, int, long int, and long long int. (These and other types may be designated in several additional ways, as described in 6.5.2.) There may also be implementation-defined extended signed integer types. ${ }^{25}$ The standard and extended signed integer types are collectively called just signed integer types. ${ }^{26}$

[^10]4 An object declared as type signed char occupies the same amount of storage as a 'plain' char object. A 'plain’" int object has the natural size suggested by the architecture of the execution environment (large enough to contain any value in the range INT_MIN to INT_MAX as defined in the header <limits.h>).

5 For each of the signed integer types, there is a corresponding (but different) unsigned integer type (designated with the keyword unsigned) that uses the same amount of storage (including sign information) and has the same alignment requirements. The unsigned integer types that correspond to the standard signed integer types are the standard unsigned integer types. The unsigned integer types that correspond to the extended signed integer types are the extended unsigned integer types.

6 The extended signed integer types and extended unsigned integer types are collectively called the extended integer types.

7 For any two types with the same signedness and different integer conversion rank, the range of values of the type with smaller integer conversion rank is a subrange of the values of the other type.

8 The range of nonnegative values of a signed integer type is a subrange of the corresponding unsigned integer type, and the representation of the same value in each type is the same. ${ }^{27}$ A computation involving unsigned operands can never overflow, because a result that cannot be represented by the resulting unsigned integer type is reduced modulo the number that is one greater than the largest value that can be represented by the resulting unsigned integer type.

9 There are three real floating types, designated as float, double, and long double. The set of values of the type float is a subset of the set of values of the type double; the set of values of the type double is a subset of the set of values of the type long double.

10 There are three complex types, designated as float complex, double complex, and long double complex. ${ }^{28}$ The real floating and complex types are collectively called the floating types.

11 For each floating type there is a corresponding real type, which is always a real floating type. For real floating types, it is the same type. For complex types, it is the type given by deleting the keyword complex from the type name.

[^11]12 Each complex type has the same representation and alignment requirements as an array type containing exactly two elements of the corresponding real type; the first element is equal to the real part, and the second element to the imaginary part, of the complex number.

13 The type char, the signed and unsigned integer types, and the floating types are collectively called the basic types. Even if the implementation defines two or more basic types to have the same representation, they are nevertheless different types. ${ }^{29}$

14 The three types char, signed char, and unsigned char are collectively called the character types. The implementation shall define char to have the same range, representation, and behavior as one of signed char and unsigned char. ${ }^{30}$

15 An enumeration comprises a set of named integer constant values. Each distinct enumeration constitutes a different enumerated type.

16 The void type comprises an empty set of values; it is an incomplete type that cannot be completed.

17 Any number of derived types can be constructed from the object, function, and incomplete types, as follows:

- An array type describes a contiguously allocated nonempty set of objects with a particular member object type, called the element type. ${ }^{31}$ Array types are characterized by their element type and by the number of elements in the array. An array type is said to be derived from its element type, and if its element type is $T$, the array type is sometimes called 'array of $T$.'" The construction of an array type from an element type is called 'array type derivation.'
- A structure type describes a sequentially allocated nonempty set of member objects, each of which has an optionally specified name and possibly distinct type.
- A union type describes an overlapping nonempty set of member objects, each of which has an optionally specified name and possibly distinct type.

29. An implementation may define new keywords that provide alternative ways to designate a basic (or any other) type. An alternate way to designed a basic type does not violate the requirement that all basic types be different. Implementation-defined keywords must have the form of an identifier reserved for any use as described in 7.1.3.
30. CHAR_MIN, defined in <limits.h>, will have one of the values 0 or SCHAR_MIN, and this can be used to distinguish the two options. Irrespective of the choice made, char is a separate type from the other two, and it not compatible with either.
31. Since object types do not include incomplete types, an array of incomplete type cannot be constructed.

- A function type describes a function with specified return type. A function type is characterized by its return type and the number and types of its parameters. A function type is said to be derived from its return type, and if its return type is $T$, the function type is sometimes called "function returning $T$." The construction of a function type from a return type is called "function type derivation."
- A pointer type may be derived from a function type, an object type, or an incomplete type, called the referenced type. A pointer type describes an object whose value provides a reference to an entity of the referenced type. A pointer type derived from the referenced type $T$ is sometimes called "pointer to $T$." The construction of a pointer type from a referenced type is called "pointer type derivation."

The type char, the signed and unsigned integer types, and the enumerated types are collectively called integer types. The integer and real floating types are collectively called real types.

Integer and floating types are collectively called arithmetic types. Arithmetic types and pointer types are collectively called scalar types. Array and structure types are collectively called aggregate types. ${ }^{32}$

21 Each arithmetic type belongs to one type-domain. The real type-domain comprises the real types. The complex type-domain comprises the complex types.

An array type of unknown size is an incomplete type. It is completed, for an identifier of that type, by specifying the size in a later declaration (with internal or external linkage). A structure or union type of unknown content (as described in 6.5.2.3) is an incomplete type. It is completed, for all declarations of that type, by declaring the same structure or union tag with its defining content later in the same scope.

Array, function, and pointer types are collectively called derived declarator types. A declarator type derivation from a type $T$ is the construction of a derived declarator type from $T$ by the application of an array-type, a function-type, or a pointer-type derivation to $T$.

24 A type is characterized by its type category, which is either the outermost derivation of a derived type (as noted above in the construction of derived types), or the type itself if the type consists of no derived types.

[^12]Any type so far mentioned is an unqualified type. Each unqualified type has several qualified versions of its type, ${ }^{33}$ corresponding to the combinations of one, two, or all three of the const, volatile, and restrict qualifiers. The qualified or unqualified versions of a type are distinct types that belong to the same type category and have the same representation and alignment requirements. ${ }^{27}$ A derived type is not qualified by the qualifiers (if any) of the type from which it is derived.

A pointer to void shall have the same representation and alignment requirements as a pointer to a character type. Similarly, pointers to qualified or unqualified versions of compatible types shall have the same representation and alignment requirements. ${ }^{27}$ All pointers to structure types shall have the same representation and alignment requirements as each other. All pointers to union types shall have the same representation and alignment requirements as each other. Pointers to other types need not have the same representation or alignment requirements.

## Examples

27 1. The type designated as "float *'" has type 'pointer to float." Its type category is pointer, not a floating type. The const-qualified version of this type is designated as "float * const" whereas the type designated as "const float *'" is not a qualified type - its type is "pointer to const-qualified float" and is a pointer to a qualified type.
2. The type designated as "struct tag (*[5]) (float)" has type "array of pointer to function returning struct tag." The array has length five and the function has a single parameter of type float. Its type category is array.

Forward references: character constants (6.1.3.4), compatible type and composite type (6.1.2.6), integer conversion rank (6.2.1.1), declarations (6.5), tags (6.5.2.3), type qualifiers (6.5.3).

### 6.1.2.6 Compatible type and composite type

1 Two types have compatible type if their types are the same. Additional rules for determining whether two types are compatible are described in 6.5 .2 for type specifiers, in 6.5 .3 for type qualifiers, and in 6.5 .5 for declarators. ${ }^{34}$ Moreover, two structure, union, or enumerated types declared in separate translation units are compatible if their tags and members satisfy the following requirements. If one is declared with a tag, the other shall be declared with the same tag. If both are
33. See 6.5 .3 regarding qualified array and function types.
34. Two types need not be identical to be compatible.
completed types, then the following additional requirements apply: there shall be a one-to-one correspondence between their members such that each pair of corresponding members are declared with compatible types, and such that if one member of a corresponding pair is declared with a name, the other member is declared with the same name. For two structures, corresponding members shall be declared in the same order. For two structures or unions, corresponding bit-fields shall have the same widths. For two enumerations, corresponding members shall have the same values.

2 All declarations that refer to the same object or function shall have compatible type; otherwise, the behavior is undefined.

3 A composite type can be constructed from two types that are compatible; it is a type that is compatible with both of the two types and satisfies the following conditions:

- If one type is an array of known constant size, the composite type is an array of that size; otherwise, if one type is a variable length array, the composite type is that type.
- If only one type is a function type with a parameter type list (a function prototype), the composite type is a function prototype with the parameter type list.
- If both types are function types with parameter type lists, the type of each parameter in the composite parameter type list is the composite type of the corresponding parameters.

These rules apply recursively to the types from which the two types are derived.
For an identifier with internal or external linkage declared in a scope in which a prior declaration of that identifier is visible, ${ }^{35}$ if the prior declaration specifies internal or external linkage, the type of the identifier at the later declaration becomes the composite type.

## Examples

5 Given the following two file scope declarations:

```
int f(int (*)(), double (*) [3]);
int f(int (*) (char *), double (*) []);
```

The resulting composite type for the function is:
35. As specified in 6.1.2.1, the later declaration might hide the prior declaration.

```
int f(int (*)(char *), double (*) [3]);
```

Forward references: declarators (6.5.5), enumeration specifiers (6.5.2.2), structure and union specifiers (6.5.2.1), type definitions (6.5.7), type qualifiers (6.5.3), type specifiers (6.5.2).

### 6.1.2.7 Predefined identifiers

1 The following identifier shall be defined by the implementation:
__func__ The name of the lexically-enclosing function.
Forward references: the identifier __func__ (6.3.1.1).

### 6.1.2.8 Representations of types

1 The representations of all types are unspecified except as stated in this subclause.

### 6.1.2.8.1 General

1 Values of type unsigned char shall be represented using a pure binary notation. ${ }^{36}$
2 When stored in objects of any other object type, values of that type consist of $n *$ CHAR_BIT bits, where $n$ is the size of an object of that type, in bytes. The value may be copied into an object of type unsigned char [ $n$ ] (e.g., by memcpy); the resulting set of bytes is called the object representation of the value. Two values with the same object representation shall compare equal, but values that compare equal might have different object representations.

3 Certain object representations might not represent a value of that type. If the stored value of an object has such a representation and is accessed by an lvalue expression that does not have character type, the behavior is undefined. If such a representation is produced by a side effect that modifies all or any part of the object by an lvalue expression that does not have character type, the behavior is undefined. ${ }^{37}$ Such representations are called trap representations.

4 When a value is stored in an object of structure or union type, including in a member object, the bytes of the object representation that correspond to any padding bytes take

[^13]unspecified values. ${ }^{38}$ The values of padding bytes shall not affect whether the value of such an object is a trap representation. Those bits of a structure or union object that are in the same byte as a bit-field member, but are not part of that member, shall similarly not affect whether the value of such an object is a trap representation.

When a value is stored in a member of an object of union type, the bytes of the object representation that do not correspond to that member but do correspond to other members take unspecified values, but the value of the union object shall not thereby become a trap representation.

6 Where an operator is applied to a value which has more than one object representation, which object representation is used shall not affect the value of the result. Where a value is stored in an object using a type that has more than one object representation for that value, it is unspecified which representation is used, but a trap representation shall not be generated.

### 6.1.2.8.2 Integer types

1 For unsigned integer types other than unsigned char, the bits of the object representation shall be divided into two groups: value bits and padding bits (there need not be any of the latter). If there are $N$ value bits, each bit shall represent a different power of 2 between 1 and $2^{N-1}$, so that objects of that type shall be capable of representing values from 0 to $2^{N}-1$ using a pure binary representation; this shall be known as the value representation. The values of any padding bits are unspecified. ${ }^{39}$

2 For signed integer types, the bits of the object representation shall be divided into three groups: value bits, padding bits, and the sign bit. There need not be any padding bits; there shall be exactly one sign bit. Each bit that is a value bit shall have the same value as the same bit in the object representation of the corresponding unsigned type (if there are $M$ value bits in the signed type and $N$ in the unsigned type, then $M \leq N$ ). If the sign bit is zero, it shall not affect the resulting value. If the sign bit is one, then the value shall be modified in one of the following ways:

- the corresponding value with sign bit 0 is negated;

[^14]- the sign bit has the value $-2^{N}$;
- the sign bit has the value $1-2^{N}$.

3 The values of any padding bits are unspecified. ${ }^{39}$ A valid (non-trap) object representation of a signed integer type where the sign bit is zero is a valid object representation of the corresponding unsigned type, and shall represent the same value.

4 The precision of an integer type is the number of bits it uses to represent values, excluding any sign and padding bits. The width of an integer type is the same but including any sign bit; thus for unsigned integer types the two values are the same, while for signed integer types the width is one greater than the precision.

### 6.1.3 Constants

## Syntax

constant:
floating-constant
integer-constant
enumeration-constant
character-constant

## Constraints

2 The value of a constant shall be in the range of representable values for its type.

## Semantics

3 Each constant has a type, determined by its form and value, as detailed later.

### 6.1.3.1 Floating constants

## Syntax

```
floating-constant:
    decimal-floating-constant
    hexadecimal-floating-constant
decimal-floating-constant:
    fractional-constant exponent-part opt floating-suffix opt
    digit-sequence exponent-part floating-suffix
```

hexadecimal-floating-constant:
0x hexadecimal-fractional-constant binary-exponent-part floating-suffix ${ }_{\text {opt }}$
0x hexadecimal-fractional-constant binary-exponent-part floating-suffix opt $0 \times$ hexadecimal-digit-sequence binary-exponent-part floating-suffix ${ }_{o p t}$ 0x hexadecimal-digit-sequence binary-exponent-part floating-suffix opt
fractional-constant:
digit-sequence ${ }_{\text {opt }}$. digit-sequence
digit-sequence.
exponent-part:
e sign ${ }_{\text {opt }}$ digit-sequence
E sign opt digit-sequence
sign: one of
$+\quad-$
digit-sequence:
digit
digit-sequence digit
hexadecimal-fractional-constant:
hexadecimal-digit-sequence ${ }_{\text {opt }}$.
hexadecimal-digit-sequence
hexadecimal-digit-sequence .
binary-exponent-part:
$\mathbf{P}$ sign opt digit-sequence
$\mathbf{P}$ sign digit-sequence
hexadecimal-digit-sequence:
hexadecimal-digit
hexadecimal-digit-sequence hexadecimal-digit
hexadecimal-digit: one of
$\begin{array}{llllllllll}0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9\end{array}$

A $\quad$ B $\quad$ C $\quad$ D $\quad$ E $\quad$ F
floating-suffix: one of
$\mathbf{f} \quad \mathbf{F} \quad \mathbf{L}$

## Description

2 A floating constant has a significand part that may be followed by an exponent part and a suffix that specifies its type. The components of the significand part may include a digit sequence representing the whole-number part, followed by a period (.), followed by a digit sequence representing the fraction part. The components of the exponent part are an $\mathbf{e}, \mathbf{E}, \mathbf{p}$, or $\mathbf{P}$ followed by an exponent consisting of an optionally signed digit sequence. Either the whole-number part or the fraction part shall be present; for decimal floating constants, either the period or the exponent part shall be present.

## Semantics

3 The significand part is interpreted as a (decimal or hexadecimal) rational number; the digit sequence in the exponent part is interpreted as a decimal integer. For decimal floating constants, the exponent indicates the power of 10 by which the significand part is to be scaled. For hexadecimal floating constants, the exponent indicates the power of 2 by which the significand part is to be scaled. For decimal floating constants, and also for hexadecimal floating constants when FLT_RADIX is not a power of 2 , if the scaled value is in the range of representable values (for its type) the result is either the nearest representable value, or the larger or smaller representable value immediately adjacent to the nearest representable value, chosen in an implementation-defined manner. For hexadecimal floating constants, if FLT_RADIX is a power of 2 and the scaled value is in the range of representable values (for its type), then the result of a hexadecimal floating constant is correctly rounded.

4 An unsuffixed floating constant has type double. If suffixed by the letter $\mathbf{f}$ or $\mathbf{F}$, it has type float. If suffixed by the letter 1 or $\mathbf{L}$, it has type long double.

## Recommended practice

5 The implementation produces a diagnostic message if a hexadecimal constant cannot be represented exactly in its evaluation format; the implementation then proceeds with the translation of the program.

6 The translation-time conversion of floating constants matches the execution-time conversion of character strings by library functions, such as strtod, given matching inputs suitable for both conversions, the same result format, and default execution-time rounding. ${ }^{40}$

[^15]
### 6.1.3.2 Integer constants

## Syntax

integer-constant:
decimal-constant integer-suffix ${ }_{\text {opt }}$
octal-constant integer-suffix opt
hexadecimal-constant integer-suffix ${ }_{\text {opt }}$
decimal-constant:
nonzero-digit
decimal-constant digit
octal-constant:
0
octal-constant octal-digit
hexadecimal-constant:
0x hexadecimal-digit
0x hexadecimal-digit
hexadecimal-constant hexadecimal-digit
nonzero-digit: one of
$\begin{array}{lllllllll}1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9\end{array}$
octal-digit: one of
$\begin{array}{llllllll}0 & 1 & 2 & 3 & 4 & 5 & 6 & 7\end{array}$
integer-suffix:
unsigned-suffix long-suffix opt
long-suffix unsigned-suffix opt
unsigned-suffix long-long-suffix opt
long-long-suffix unsigned-suffix opt
unsigned-suffix: one of
u U
long-suffix: one of
1 L
long-long-suffix: one of
11 LL

## Description

2 An integer constant begins with a digit, but has no period or exponent part. It may have a prefix that specifies its base and a suffix that specifies its type.

3 A decimal constant begins with a nonzero digit and consists of a sequence of decimal digits. An octal constant consists of the prefix 0 optionally followed by a sequence of the digits $\mathbf{0}$ through $\mathbf{7}$ only. A hexadecimal constant consists of the prefix $\mathbf{0 x}$ or $\mathbf{0 x}$ followed by a sequence of the decimal digits and the letters a (or $\mathbf{A}$ ) through $\mathbf{f}$ (or $\mathbf{F}$ ) with values 10 through 15 respectively.

## Semantics

4 The value of a decimal constant is computed base 10; that of an octal constant, base 8; that of a hexadecimal constant, base 16 . The lexically first digit is the most significant.

The type of an integer constant is the first of the corresponding list in which its value can be represented. Unsuffixed decimal: int, long int, long long int; unsuffixed octal or hexadecimal: int, unsigned int, long int, unsigned long int, long long int, unsigned long long int; suffixed by the letter $u$ or U : unsigned int, unsigned long int, unsigned long long int; decimal suffixed by the letter 1 or L : long int, long long int; octal or hexadecimal suffixed by the letter 1 or L : long int, unsigned long int, long long int, unsigned long long int; suffixed by both the letters $u$ or $U$ and 1 or L: unsigned long int, unsigned long long int; decimal suffixed by 11 or LL: long long int; octal or hexadecimal suffixed by the letter 11 or LL: long long int, unsigned long long int; suffixed by both $u$ or $U$ and $l l$ or LL: unsigned long long int. If an integer constant can not be represented by any type in its list, it may have an extended integer type, if the extended integer type can represent its value. If all of the types in the list for the constant are signed, the extended integer type shall be signed. If all of the types in the list for the constant are unsigned, the extended integer type shall be unsigned. If the list contains both signed and unsigned types, the extended integer type may be signed or unsigned.

### 6.1.3.3 Enumeration constants

## Syntax

enumeration-constant: identifier

## Semantics

2 An identifier declared as an enumeration constant has type int.
Forward references: enumeration specifiers (6.5.2.2).

### 6.1.3.4 Character constants

Syntax

```
character-constant:
    'c-char-sequence'
    L'c-char-sequence'
    c-char-sequence:
    c-char
    c-char-sequence c-char
c-char:
    any member of the source character set except
            the single-quote ', backslash \, or new-line character
    escape-sequence
    universal-character-name
escape-sequence:
    simple-escape-sequence
    octal-escape-sequence
    hexadecimal-escape-sequence
simple-escape-sequence: one of
    \' \" \? \\
    \a \b \f \n \r \t \v
octal-escape-sequence:
            \octal-digit
    \ octal-digit octal-digit
    \ octal-digit octal-digit octal-digit
hexadecimal-escape-sequence:
    \x hexadecimal-digit
    hexadecimal-escape-sequence hexadecimal-digit
```


## Description

2 An integer character constant is a sequence of one or more multibyte characters enclosed in single-quotes, as in ' $\mathbf{x}^{\prime}$ or ' $\mathbf{a b}^{\prime}$. A wide character constant is the same, except prefixed by the letter $\mathbf{L}$. With a few exceptions detailed later, the elements of the sequence are any members of the source character set; they are mapped in an
implementation-defined manner to members of the execution character set.
3 The single-quote ', the double-quote ", the question-mark ?, the backslash <br>, and arbitrary integer values, are representable according to the following table of escape sequences:

| single-quote ' | $\backslash \prime$ |
| :--- | :--- |
| double-quote " | $\backslash "$ |
| question-mark ? | $\backslash ?$ |
| backslash \} $&{\backslash \backslash} \\ {\text { octal integer }} &{\text { \octal digits }} \\ {\text { hexadecimal integer }} &{\backslash \mathbf{x} \text { hexadecimal digits }}$ |  |

4 The double-quote " and question-mark ? are representable either by themselves or by the escape sequences $\backslash "$ and $\backslash$ ?, respectively, but the single-quote ' and the backslash $\backslash$ shall be represented, respectively, by the escape sequences $\backslash^{\prime}$ and $\backslash \backslash$.

5 The octal digits that follow the backslash in an octal escape sequence are taken to be part of the construction of a single character for an integer character constant or of a single wide character for a wide character constant. The numerical value of the octal integer so formed specifies the value of the desired character or wide character.

6 The hexadecimal digits that follow the backslash and the letter $\mathbf{x}$ in a hexadecimal escape sequence are taken to be part of the construction of a single character for an integer character constant or of a single wide character for a wide character constant. The numerical value of the hexadecimal integer so formed specifies the value of the desired character or wide character.

7 Each octal or hexadecimal escape sequence is the longest sequence of characters that can constitute the escape sequence.

8 In addition, certain nongraphic characters are representable by escape sequences consisting of the backslash $\backslash$ followed by a lowercase letter: $\backslash \mathbf{a}, \backslash \mathbf{b}, \backslash \mathbf{f}, \backslash \mathbf{n}, \backslash \mathbf{r}, \backslash \mathbf{t}$, and $\backslash \mathbf{v} .{ }^{41}$ If any other escape sequence is encountered, the behavior is undefined. ${ }^{42}$

[^16]
## Constraints

The value of an octal or hexadecimal escape sequence shall be in the range of representable values for the type unsigned char for an integer character constant, or the unsigned type corresponding to wchar_t for a wide character constant.

## Semantics

An integer character constant has type int. The value of an integer character constant containing a single character that maps to a member of the basic execution character set is the numerical value of the representation of the mapped character interpreted as an integer. The value of an integer character constant containing more than one character, or containing a character or escape sequence not represented in the basic execution character set, is implementation-defined. If an integer character constant contains a single character or escape sequence, its value is the one that results when an object with type char whose value is that of the single character or escape sequence is converted to type int.

11 A wide character constant has type wchar_t, an integer type defined in the <stddef. h > header. The value of a wide character constant containing a single multibyte character that maps to a member of the extended execution character set is the wide character (code) corresponding to that multibyte character, as defined by the mbtowc function, with an implementation-defined current locale. The value of a wide character constant containing more than one multibyte character, or containing a multibyte character or escape sequence not represented in the extended execution character set, is implementation-defined.

## Examples

1. The construction ' $\backslash 0^{\prime}$ ' is commonly used to represent the null character.
2. Consider implementations that use two's-complement representation for integers and eight bits for objects that have type char. In an implementation in which type char has the same range of values as signed char, the integer character constant ' $\backslash \mathbf{x F F}$ ' has the value -1 ; if type char has the same range of values as unsigned char, the character constant ' $\backslash \mathbf{x F F} \mathbf{I}^{\prime}$ has the value +255 .
3. Even if eight bits are used for objects that have type char, the construction ' $\backslash \times 123^{\prime}$ specifies an integer character constant containing only one character, since a hexadecimal escape sequence is terminated only by a non-hexadecimal character. To specify an integer character constant containing the two characters whose values are ' $\backslash \mathbf{x 1 2 '}$ and ' $\mathbf{3 '}^{\prime}$, the construction ' $\backslash 0223^{\prime}$ may be used, since an octal escape sequence is terminated after three octal digits. (The value of this two-character integer character constant is implementation-defined.)
4. Even if 12 or more bits are used for objects that have type wchar_t, the construction $\mathbf{L}^{\prime} \backslash 123 \mathbf{4}^{\prime}$ specifies the implementation-defined value that results from the combination of the values 0123 and ${ }^{\prime} 4^{\prime}$.

Forward references: characters and integers (6.2.1.1), common definitions <stddef.h> (7.1.6), the mbtowc function (7.14.7.2).

### 6.1.4 String literals

Syntax

## string-literal:

"s-char-sequence ${ }_{\text {opt }}$ "
L "s-char-sequence ${ }_{\text {opt }} "$
$s$-char-sequence:
s-char
s-char-sequence s-char
$s$-char:
any member of the source character set except
the double-quote ", backslash <br>, or new-line character
escape-sequence
universal-character-name

## Description

2 A character string literal is a sequence of zero or more multibyte characters enclosed in double-quotes, as in "xyz". A wide string literal is the same, except prefixed by the letter $\mathbf{L}$.

3 The same considerations apply to each element of the sequence in a character string literal or a wide string literal as if it were in an integer character constant or a wide character constant, except that the single-quote ' is representable either by itself or by the escape sequence $\backslash^{\prime}$, but the double-quote " shall be represented by the escape sequence \".

## Semantics

4 In translation phase 6, the multibyte character sequences specified by any sequence of adjacent character and wide string literal tokens are concatenated into a single multibyte character sequence. If any of the tokens are wide string literal tokens, the resulting multibyte character sequence is treated as a wide string literal; otherwise, it is treated as a character string literal.

In translation phase 7, a byte or code of value zero is appended to each multibyte character sequence that results from a string literal or literals. ${ }^{43}$ The multibyte character sequence is then used to initialize an array of static storage duration and length just sufficient to contain the sequence. For character string literals, the array elements have type char, and are initialized with the individual bytes of the multibyte character sequence; for wide string literals, the array elements have type wchar_t, and are initialized with the sequence of wide characters corresponding to the multibyte character sequence.

6 These arrays need not be distinct provided their elements have the appropriate values. If the program attempts to modify such an array, the behavior is undefined.

## Examples

7 This pair of adjacent character string literals

```
"\x12" "3"
```

produces a single character string literal containing the two characters whose values are ' $\backslash \mathbf{x 1 2 '}$ and ' $\mathbf{3 '}^{\prime}$, because escape sequences are converted into single members of the execution character set just prior to adjacent string literal concatenation.

Forward references: common definitions <stddef. $\mathrm{h}>$ (7.1.6).

### 6.1.5 Operators

## Syntax

operator: one of


[^17]
## Constraints

2 The operators [ ], ( ), and ? : (independent of spelling) shall occur in pairs, possibly separated by expressions. The operators \# and \#\# (also spelled \%: and $\%: \%$, respectively) shall occur in macro-defining preprocessing directives only.

## Semantics

3 An operator specifies an operation to be performed (an evaluation) that yields a value, or yields a designator, or produces a side effect, or a combination thereof. An operand is an entity on which an operator acts.

4 In all aspects of the language, these six tokens

```
<: :> <% %> %: %:%:
```

behave, respectively, the same as these six tokens
[ ] \{ \} \# \#\#
except for their spelling. ${ }^{44}$
Forward references: expressions (6.3), macro replacement (6.8.3).

### 6.1.6 Punctuators

Syntax
punctuator: one of

## Constraints

2 The punctuators [ ], ( ), and \{ \} (independent of spelling) shall occur (after translation phase 4) in pairs, possibly separated by expressions, declarations, or statements. The punctuator \# (also spelled \%:) shall occur in preprocessing directives only.
44. Thus [ and < : behave differently when "stringized" (see subclause 6.8.3.2), but can otherwise be
freely interchanged. freely interchanged.

## Semantics

3 A punctuator is a symbol that has independent syntactic and semantic significance but does not specify an operation to be performed that yields a value. Depending on context, the same symbol may also represent an operator or part of an operator.

Forward references: expressions (6.3), declarations (6.5), preprocessing directives (6.8), statements (6.6).

### 6.1.7 Header names

Syntax

```
header-name:
    <h-char-sequence>
    "q-char-sequence"
h-char-sequence:
    h-char
    h-char-sequence h-char
```

$h$-char:
any member of the source character set except
the new-line character and >
q-char-sequence:
q-char
q-char-sequence q-char
$q$-char:
any member of the source character set except
the new-line character and "

## Semantics

2 The sequences in both forms of header names are mapped in an implementationdefined manner to headers or external source file names as specified in 6.8.2.

3 If the characters ', $\backslash$, ", //, or /* occur in the sequence between the $<$ and $>$ delimiters, the behavior is undefined. Similarly, if the characters ', <br>, //, or /* occur in the sequence between the " delimiters, the behavior is undefined. ${ }^{45}$ A header name preprocessing token is recognized only within a \#include preprocessing directive.

[^18]
## Examples

4 The following sequence of characters:

```
0x3<1/a.h>1e2
#include <1/a.h>
#define const.member@$
```

forms the following sequence of preprocessing tokens (with each individual preprocessing token delimited by a $\{$ on the left and a $\}$ on the right).
$\{0 \times 3\}\{<\}\{1\}\{/\}\{a\}\{\}.\{h\}\{>\}\{1 e 2\}$
\{\#\}\{include\} \{<1/a.h>\}
\{\#\}\{define\} \{const\}\{.\}\{member\}\{@\}\{\$\}
Forward references: source file inclusion (6.8.2).

### 6.1.8 Preprocessing numbers

Syntax
digit
. digit
pp-number digit
pp-number nondigit
pp-number e sign
pp-number $\mathbf{E}$ sign
pp-number $\mathbf{p}$ sign
pp-number $\mathbf{P}$ sign
pp-number .

## Description

2 A preprocessing number begins with a digit optionally preceded by a period (.) and may be followed by letters, underscores, digits, periods, and $\mathbf{e}+$, $\mathbf{e}-\mathbf{E} \mathbf{E}, \mathbf{E}-, \mathbf{p}+, \mathbf{p}-$, $\mathbf{P}+$, or $\mathbf{P}$ - character sequences.

3 Preprocessing number tokens lexically include all floating and integer constant tokens.

## Semantics

4 A preprocessing number does not have type or a value; it acquires both after a successful conversion (as part of translation phase 7) to a floating constant token or an integer constant token.

### 6.1.9 Comments

1 Except within a character constant, a string literal, or a comment, the characters /* introduce a comment. The contents of a comment are examined only to identify multibyte characters and to find the characters */ that terminate it. ${ }^{46}$

2 Except within a character constant, a string literal, or a comment, the characters // introduce a comment that includes all multibyte characters up to, but not including, the next new-line character. The contents of such a comment are examined only to identify multibyte characters and to find the terminating new-line character.

## Examples

3

```
    "a//b" // four-character string literal
    \#include "//e" // undefined behavior
    // */ // comment, not syntax error
    \(\mathbf{f}=\mathbf{g} / * * / / \mathbf{h} ; \quad / /\) equivalent to \(\mathbf{f}=\mathbf{g} / \mathbf{h}\);
    /ハ
    i(); // part of a two-line comment
    ハ
    / j(); // part of a two-line comment
    \#define glue \((x, y)\) x\#\#y
    glue(/,/) \(\mathbf{k}() ; / /\) syntax error, not comment
    /*//*/ l(); // equivalent to l();
    \(\mathrm{m}=\mathrm{n} / / * * / 0\)
    \(+\mathrm{p} ; \quad / /\) equivalent to \(\mathrm{m}=\mathrm{n}+\mathrm{p}\);
```

[^19]
### 6.2 Conversions

1 Several operators convert operand values from one type to another automatically. This subclause specifies the result required from such an implicit conversion, as well as those that result from a cast operation (an explicit conversion). The list in 6.2.1.7 summarizes the conversions performed by most ordinary operators; it is supplemented as required by the discussion of each operator in 6.3.

2 Conversion of an operand value to a compatible type causes no change to the value or the representation.

Forward references: cast operators (6.3.4).

### 6.2.1 Arithmetic operands

### 6.2.1.1 Characters and integers

1 Every integer type has an integer conversion rank defined as follows:

- No two signed integer types shall have the same rank, even if they have the same representation.
- The rank of a signed integer type shall be greater than the rank of any signed integer type with less precision.
- The rank of any standard signed integer type shall be greater than the rank of any extended signed integer type with the same precision.
- The rank of long long int shall be greater than the rank of long int, which shall be greater than the rank of int, which shall be greater than the rank of short int, which shall be greater than the rank of signed char.
- The rank of any unsigned integer type shall equal the rank of the corresponding signed integer type.
- The rank of char shall equal the rank of signed char and unsigned char.
- The rank of any enumerated type shall equal the rank of the compatible integer type.
- The rank of any extended signed integer type relative to another extended signed integer type with the same precision is implementation-defined, but still subject to the other rules for determining the integer conversion rank.
- For all integer types $\mathbf{T 1}, \mathbf{T 2}$, and $\mathbf{T 3}$, if $\mathbf{T 1}$ has greater rank than $\mathbf{T 2}$ and $\mathbf{T 2}$ has greater rank than $\mathbf{T 3}$, then $\mathbf{T 1}$ has greater rank than $\mathbf{T 3}$.

The following may be used in an expression wherever an int or unsigned int may be used.

- An object or expression with an integer type whose integer conversion rank is less than the rank of int and unsigned int.
- A bit-field of type int, signed int, or unsigned int.

3 If an int can represent all values of the original type, the value is converted to an int; otherwise, it is converted to an unsigned int. These are called the integer promotions. ${ }^{47}$ All other types are unchanged by the integer promotions.

4 The integer promotions preserve value including sign. As discussed earlier, whether a 'plain' char is treated as signed is implementation-defined.

Forward references: enumeration specifiers (6.5.2.2), structure and union specifiers (6.5.2.1).

### 6.2.1.2 Signed and unsigned integers

1 When a value with integer type is converted to another integer type, if the value can be represented by the new type, it is unchanged.

2 Otherwise, if the new type is unsigned, the value is converted by repeatedly adding or subtracting one more than the maximum value that can be represented in the new type until the value is in the range of the new type.

3 Otherwise, the new type is signed and the value cannot be represented in it; the result is implementation-defined.

### 6.2.1.3 Real floating and integer

1 When a value of real floating type is converted to integer type, the fractional part is discarded. If the value of the integral part cannot be represented by the integer type, the behavior is undefined. ${ }^{48}$

2 When a value of integer type is converted to real floating type, if the value being converted is in the range of values that can be represented but cannot be represented exactly, the result is either the nearest higher or nearest lower value, chosen in an implementation-defined manner. If the value being converted is outside the range of values that can be represented, the behavior is undefined.

[^20]
### 6.2.1.4 Real floating types

1 When a float is promoted to double or long double, or a double is promoted to long double, its value is unchanged.

2 When a double is demoted to float or a long double to double or float, if the value being converted is outside the range of values that can be represented, the behavior is undefined. If the value being converted is in the range of values that can be represented but cannot be represented exactly, the result is either the nearest higher or nearest lower value, chosen in an implementation-defined manner.

### 6.2.1.5 Complex types

1 When a value of complex type is converted to another complex type, both the real and imaginary parts follow the conversion rules for the corresponding real types.

### 6.2.1.6 Real and complex

1 When a value of real type is converted to a complex type, the real part of the complex result value is determined by the rules of conversion to the corresponding real type and the imaginary part of the complex result value is a positive zero or an unsigned zero.

When a value of complex type is converted to a real type, the imaginary part of the complex value is discarded and the value of the real part is converted according to the conversion rules for the corresponding real type.

### 6.2.1.7 Usual arithmetic conversions

1 Many operators that expect operands of arithmetic type cause conversions and yield result types in a similar way. The purpose is to determine a common real type for the operands and result. For the specified operands, each operand is converted, without change of type-domain, to a type whose corresponding real type is the common real type. Unless explicitly stated otherwise, the common real type is also the corresponding real type of the result, whose type-domain is determined by the operator. This pattern is called the usual arithmetic conversions:

First, if the corresponding real type of either operand is long double, the other operand is converted, without change of type-domain, to a type whose corresponding real type is long double.

Otherwise, if the corresponding real type of either operand is double, the other operand is converted, without change of type-domain, to a type whose corresponding real type is double.

Otherwise, if the corresponding real type of either operand is float, the other operand is converted, without change of type-domain, to a type whose corresponding real type is float. ${ }^{49}$

Otherwise, the integer promotions are performed on both operands. Then the following rules are applied to the promoted operands:

If both operands have the same type, then no further conversion is needed.
Otherwise, if both operands have signed integer types or both have unsigned integer types, the operand with the type of lesser integer conversion rank is converted to the type of the operand with greater rank.

Otherwise, if the operand that has unsigned integer type has rank greater or equal to the rank of the type of the other operand, then the operand with signed integer type is converted to the type of the operand with unsigned integer type.

Otherwise, if the type of the operand with signed integer type can represent all of the values of the type of the operand with unsigned integer type, then the operand with unsigned integer type is converted to the type of the operand with signed integer type.

Otherwise, both operands are converted to the unsigned integer type corresponding to the type of the operand with signed integer type.

2 The values of floating operands and of the results of floating expressions may be represented in greater precision and range than that required by the type; the types are not changed thereby. ${ }^{50}$

### 6.2.2 Other operands

### 6.2.2.1 Lvalues and function designators

1 An lvalue is an expression (with an object type or an incomplete type other than void) that designates an object. ${ }^{51}$ When an object is said to have a particular type, the type is specified by the lvalue used to designate the object. A modifiable lvalue is

[^21]an lvalue that does not have array type, does not have an incomplete type, does not have a const-qualified type, and if it is a structure or union, does not have any member (including, recursively, any member or element of all contained aggregates or unions) with a const-qualified type.

2 Except when it is the operand of the sizeof operator, the unary \& operator, the ++ operator, the -- operator, or the left operand of the . operator or an assignment operator, an lvalue that does not have array type is converted to the value stored in the designated object (and is no longer an lvalue). If the lvalue has qualified type, the value has the unqualified version of the type of the lvalue; otherwise, the value has the type of the lvalue. If the lvalue has an incomplete type and does not have array type, the behavior is undefined.

3 Except when it is the operand of the sizeof operator or the unary $\&$ operator, or is a character string literal used to initialize an array of character type, or is a wide string literal used to initialize an array with element type compatible with wchar_t, an lvalue that has type "array of type" is converted to an expression that has type "pointer to type" that points to the initial element of the array object and is not an lvalue. If the array object has register storage class, the behavior is undefined.

4 A function designator is an expression that has function type. Except when it is the operand of the sizeof operator ${ }^{52}$ or the unary $\&$ operator, a function designator with type "function returning type" is converted to an expression that has type "pointer to function returning type."

Forward references: address and indirection operators (6.3.3.2), assignment operators (6.3.16), common definitions <stddef. $\mathrm{h}>$ (7.1.6), initialization (6.5.8), postfix increment and decrement operators (6.3.2.4), prefix increment and decrement operators (6.3.3.1), the sizeof operator (6.3.3.4), structure and union members (6.3.2.3).

### 6.2.2.2 void

1 The (nonexistent) value of a void expression (an expression that has type void) shall not be used in any way, and implicit or explicit conversions (except to void) shall not be applied to such an expression. If an expression of any other type occurs in a context where a void expression is required, its value or designator is discarded. (A void expression is evaluated for its side effects.)

[^22]
### 6.2.2.3 Pointers

1 A pointer to void may be converted to or from a pointer to any incomplete or object type. A pointer to any incomplete or object type may be converted to a pointer to void and back again; the result shall compare equal to the original pointer.

2 For any qualifier $q$, a pointer to a non- $q$-qualified type may be converted to a pointer to the $q$-qualified version of the type; the values stored in the original and converted pointers shall compare equal.

3 An integer constant expression with the value 0 , or such an expression cast to type void *, is called a null pointer constant. ${ }^{53}$ If a null pointer constant is assigned to or compared for equality to a pointer, the constant is converted to a pointer of that type. Such a pointer, called a null pointer, is guaranteed to compare unequal to a pointer to any object or function.

4 Conversion of a null pointer to another pointer type yields a null pointer of that type. Any two null pointers shall compare equal.

5 An integer may be converted to any pointer type. The result is implementationdefined, might not be properly aligned, and might not point to an entity of the referenced type. ${ }^{54}$

6 Any pointer type may be converted to an integer type; the result is implementationdefined. If the result cannot be represented in the integer type, the behavior is undefined. The result need not be in the range of values of any integer type. ${ }^{55}$

7 A pointer to an object or incomplete type may be converted to a pointer to a different object or incomplete type. If the resulting pointer is not correctly aligned ${ }^{56}$ for the pointed to type, the behavior is undefined. Otherwise, when converted back again, the result shall compare equal to the original pointer.

8 A pointer to a function of one type may be converted to a pointer to a function of another type and back again; the result shall compare equal to the original pointer. If a converted pointer is used to call a function that has a type that is not compatible with the type of the called function, the behavior is undefined.
53. The macro NULL is defined in <stddef. $\mathrm{h}>$ as a null pointer constant; see 7.1.6.
54. The mapping functions for converting a pointer to an integer or an integer to a pointer are intended to be consistent with the addressing structure of the execution environment.
55. Thus, if the conversion is to unsigned int but yields a negative value, the behavior is undefined.
56. In general, the concept correctly aligned is transitive: if a pointer to type A is correctly aligned for a pointer to type $B$, which in turn is correctly aligned for a pointer to type $C$, then a pointer to type $A$ is correctly aligned for a pointer to type $C$.

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Forward references: cast operators (6.3.4), equality operators (6.3.9), simple assignment (6.3.16.1).

### 6.3 Expressions

1 An expression is a sequence of operators and operands that specifies computation of a value, or that designates an object or a function, or that generates side effects, or that performs a combination thereof.

2 Between the previous and next sequence point an object shall have its stored value modified at most once by the evaluation of an expression. Furthermore, the prior value shall be accessed only to determine the value to be stored. ${ }^{57}$
3 Except as indicated by the syntax ${ }^{58}$ or otherwise specified later (for the function-call operator (), \&\&, \|, ?:, and comma operators), the order of evaluation of subexpressions and the order in which side effects take place are both unspecified.

4 Some operators (the unary operator $\sim$, and the binary operators $\lll \gg, \&, \wedge$, and |, collectively described as bitwise operators) shall have operands that have integer type. These operators return values that depend on the internal representations of integers, and have implementation-defined and undefined aspects for signed types.

5 If an exception occurs during the evaluation of an expression (that is, if the result is not mathematically defined or not in the range of representable values for its type), the behavior is undefined.

6 The effective type of an object for an access to its stored value is the declared type of the object, if it has one. If a value is stored into an object having no declared type through an lvalue having a type that is not a character type, then the type of the lvalue becomes the effective type of the object for that access and for subsequent accesses that do not modify the stored value. If a value is copied into an object having no
57. This paragraph renders undefined statement expressions such as

## i $=++i+1 ;$

while allowing

$$
i=i+1 ;
$$

58. The syntax specifies the precedence of operators in the evaluation of an expression, which is the same as the order of the major subclauses of this subclause, highest precedence first. Thus, for example, the expressions allowed as the operands of the binary + operator (6.3.6) shall be those expressions defined in 6.3.1 through 6.3.6. The exceptions are cast expressions (6.3.4) as operands of unary operators (6.3.3), and an operand contained between any of the following pairs of operators: grouping parentheses () (6.3.1), subscripting brackets [] (6.3.2.1), function-call parentheses () (6.3.2.2), and the conditional operator ? : (6.3.15).

Within each major subclause, the operators have the same precedence. Left- or right-associativity is indicated in each subclause by the syntax for the expressions discussed therein.
declared type using memcpy or memmove, or is copied as an array of character type, then the effective type of the modified object for that access and for subsequent accesses that do not modify the value is the effective type of the object from which the value is copied, if it has one. For all other accesses to an object having no declared type, the effective type of the object is simply the type of the lvalue used for the access.

7 An object shall have its stored value accessed only by an lvalue expression that has one of the following types: ${ }^{59}$

- a type compatible with the effective type of the object,
- a qualified version of a type compatible with the effective type of the object,
- a type that is the signed or unsigned type corresponding to the effective type of the object,
- a type that is the signed or unsigned type corresponding to a qualified version of the effective type of the object,
- an aggregate or union type that includes one of the aforementioned types among its members (including, recursively, a member of a subaggregate or contained union), or
- a character type.

A floating expression may be contracted, that is, evaluated as though it were an atomic operation, thereby omitting rounding errors implied by the source code and the expression evaluation method. ${ }^{60}$ The FP_CONTRACT pragma in <math.h> provides a way to disallow contracted expressions. Otherwise, whether and how expressions are contracted is implementation-defined. ${ }^{61}$

[^23]
### 6.3.1 Primary expressions

Syntax

```
primary-expression:
    identifier
    constant
    string-literal
    ( expression )
```


## Semantics

2 An identifier is a primary expression, provided it has been declared as designating an object (in which case it is an lvalue) or a function (in which case it is a function designator). ${ }^{62}$

3 A constant is a primary expression. Its type depends on its form and value, as detailed in 6.1.3.

4 A string literal is a primary expression. It is an lvalue with type as detailed in 6.1.4.
5 A parenthesized expression is a primary expression. Its type and value are identical to those of the unparenthesized expression. It is an lvalue, a function designator, or a void expression if the unparenthesized expression is, respectively, an lvalue, a function designator, or a void expression.

Forward references: declarations (6.5).

### 6.3.1.1 The identifier __func__

## Semantics

1 The identifier __func__ is implicitly declared by the translator as if, immediately following the opening brace of each function definition, the declaration

```
static const char __func__[] = "function-name";
```

appeared, where function-name is the name of the lexically-enclosing function. ${ }^{63}$ This name is the unadorned name of the function.

2 This name is encoded as if the implicit declaration had been written in the source character set and then translated into the execution character set as indicated in
62. Thus, an undeclared identifier is a violation of the syntax.
63. Note that since the name __func__ is reserved for any use by the implementation (7.1.3), if any other identifier is explicitly declared using the name __func__, the behavior is undefined.
translation phase 5.

## Examples

Consider the code fragment:

```
#include <stdio.h>
void myfunc (void)
{
    printf("%s\n", __func__);
    /* ... */
}
```

Each time the function is called, it will print to the standard output stream:

```
myfunc
```


### 6.3.2 Postfix operators

Syntax
postfix-expression:
primary-expression
postfix-expression [ expression ]
postfix-expression ( argument-expression-list ${ }_{\text {opt }}$ )
postfix-expression . identifier
postfix-expression -> identifier
postfix-expression ++
postfix-expression --
( type-name ) \{ initializer-list \}
(type-name ) \{ initializer-list , \}
argument-expression-list:
assignment-expression
argument-expression-list, assignment-expression

### 6.3.2.1 Array subscripting

## Constraints

1 One of the expressions shall have type "pointer to object type," the other expression shall have integer type, and the result has type "type."

## Semantics

2 A postfix expression followed by an expression in square brackets [] is a subscripted designation of an element of an array object. The definition of the subscript operator [] is that $\mathbf{E 1}[\mathbf{E 2}$ ] is identical to (*(E1+(E2))). Because of the conversion rules that apply to the binary + operator, if $\mathbf{E} \mathbf{1}$ is an array object (equivalently, a pointer to the initial element of an array object) and $\mathbf{E} 2$ is an integer, $\mathbf{E 1}$ [ $\mathbf{E} 2]$ designates the E2-th element of E1 (counting from zero).

3 Successive subscript operators designate an element of a multidimensional array object. If $\mathbf{E}$ is an $n$-dimensional array ( $n \geq 2$ ) with dimensions $i \times j \times \ldots \times k$, then $\mathbf{E}$ (used as other than an lvalue) is converted to a pointer to an ( $n-1$ )-dimensional array with dimensions $j \times \ldots \times k$. If the unary * operator is applied to this pointer explicitly, or implicitly as a result of subscripting, the result is the pointed-to ( $n-1$ )-dimensional array, which itself is converted into a pointer if used as other than an lvalue. It follows from this that arrays are stored in row-major order (last subscript varies fastest).

## Examples

Consider the array object defined by the declaration

```
int x[3][5];
```

Here $\mathbf{x}$ is a $3 \times 5$ array of ints; more precisely, $\mathbf{x}$ is an array of three element objects, each of which is an array of five ints. In the expression $\mathbf{x}[\boldsymbol{i}]$, which is equivalent
 Then $\mathbf{i}$ is adjusted according to the type of $\mathbf{x}$, which conceptually entails multiplying i by the size of the object to which the pointer points, namely an array of five int objects. The results are added and indirection is applied to yield an array of five ints. When used in the expression $\mathbf{x}[\mathbf{i}][j]$, that in turn is converted to a pointer to the first of the ints, so $\mathbf{x}[i][j]$ yields an int.

Forward references: additive operators (6.3.6), address and indirection operators (6.3.3.2), array declarators (6.5.5.2).

### 6.3.2.2 Function calls

## Constraints

1 The expression that denotes the called function ${ }^{64}$ shall have type pointer to function returning void or returning an object type other than an array type.

[^24]2 If the expression that denotes the called function has a type that includes a prototype, the number of arguments shall agree with the number of parameters. Each argument shall have a type such that its value may be assigned to an object with the unqualified version of the type of its corresponding parameter.

## Semantics

3 A postfix expression followed by parentheses () containing a possibly empty, comma-separated list of expressions is a function call. The postfix expression denotes the called function. The list of expressions specifies the arguments to the function.

4 An argument may be an expression of any object type. In preparing for the call to a function, the arguments are evaluated, and each parameter is assigned the value of the corresponding argument. ${ }^{65}$ If the expression that denotes the called function has type pointer to function returning an object type, the function call expression has the same type as that object type, and has the value determined as specified in 6.6.6.4. Otherwise, the function call has type void.

5 If the expression that denotes the called function has a type that does not include a prototype, the integer promotions are performed on each argument, and arguments that have type float are promoted to double. These are called the default argument promotions. If the number of arguments does not agree with the number of parameters, the behavior is undefined. If the function is defined with a type that does not include a prototype, and the types of the arguments after promotion are not compatible with those of the parameters after promotion, the behavior is undefined, except for the following cases:

- one promoted type is a signed integer type, the other promoted type is the corresponding unsigned integer type, and the value is representable in both types;
- one type is pointer to void and the other is a pointer to a character type.

6 If the function is defined with a type that includes a prototype, and the types of the arguments after promotion are not compatible with the types of the parameters, or if the prototype ends with an ellipsis (, . . .), the behavior is undefined.

7 If the expression that denotes the called function has a type that includes a prototype, the arguments are implicitly converted, as if by assignment, to the types of the corresponding parameters, taking the type of each parameter to be the unqualified

[^25]version of its declared type. The ellipsis notation in a function prototype declarator causes argument type conversion to stop after the last declared parameter. The default argument promotions are performed on trailing arguments.

8 If the function is defined with a type that is not compatible with the type (of the expression) pointed to by the expression that denotes the called function, the behavior is undefined.

9 No other conversions are performed implicitly; in particular, the number and types of arguments are not compared with those of the parameters in a function definition that does not include a function prototype declarator.

10 The order of evaluation of the function designator, the arguments, and subexpressions within the arguments is unspecified, but there is a sequence point before the actual call.

11 Recursive function calls shall be permitted, both directly and indirectly through any chain of other functions.

## Examples

12 In the function call

$$
(* p f[f 1()]) \quad(f 2(), f 3()+f 4())
$$

the functions $\mathbf{f 1} \mathbf{f} \mathbf{5}, \mathbf{f 3}$, and $\mathbf{f 4}$ may be called in any order. All side effects shall be completed before the function pointed to by $\mathrm{pf}[\mathrm{f} 1()]$ is entered.

Forward references: function declarators (including prototypes) (6.5.5.3), function definitions (6.7.1), the return statement (6.6.6.4), simple assignment (6.3.16.1).

### 6.3.2.3 Structure and union members

## Constraints

1 The first operand of the . operator shall have a qualified or unqualified structure or union type, and the second operand shall name a member of that type.

2 The first operand of the $\rightarrow$ operator shall have type "pointer to qualified or unqualified structure" or 'pointer to qualified or unqualified union," and the second operand shall name a member of the type pointed to.

## Semantics

3 A postfix expression followed by the . operator and an identifier designates a member of a structure or union object. The value is that of the named member, and is an lvalue if the first expression is an lvalue. If the first expression has qualified type, the result has the so-qualified version of the type of the designated member.

A postfix expression followed by the $\rightarrow$ operator and an identifier designates a member of a structure or union object. The value is that of the named member of the object to which the first expression points, and is an lvalue. ${ }^{66}$ If the first expression is a pointer to a qualified type, the result has the so-qualified version of the type of the designated member.

5 With one exception, if the value of a member of a union object is used when the most recent store to the object was to a different member, the behavior is implementationdefined. ${ }^{67}$ One special guarantee is made in order to simplify the use of unions: If a union contains several structures that share a common initial sequence (see below), and if the union object currently contains one of these structures, it is permitted to inspect the common initial part of any of them anywhere that a declaration of the completed type of the union is visible. Two structures share a common initial sequence if corresponding members have compatible types (and, for bit-fields, the same widths) for a sequence of one or more initial members.

## Examples

1. If $\mathbf{f}$ is a function returning a structure or union, and $\mathbf{x}$ is a member of that structure or union, $\mathbf{f}() . \mathbf{x}$ is a valid postfix expression but is not an lvalue.
2. The following is a valid fragment:
[^26]```
union {
    struct {
        int alltypes;
        } n;
        struct {
        int type;
        int intnode;
    } ni;
    struct {
        int type;
        double doublenode;
    } nf;
} u;
u.nf.type = 1;
u.nf.doublenode = 3.14;
/* ... */
if (u.n.alltypes == 1)
    if (sin(u.nf.doublenode) == 0.0)
        /* ... */
```

3. The following is not a valid fragment (because the union type is not visible within function $\mathbf{f}$ ):
```
struct t1 { int m; };
struct t2 { int m; };
int f(struct t1 * p1, struct t2 * p2)
{
    if (p1->m < 0)
            p2->m = -p2->m;
    return p1->m;
}
int g()
{
        union {
            struct t1 s1;
            struct t2 s2;
        } u;
        /* ... */
        return f(&u.s1, &u.s2);
}
```

Forward references: address and indirection operators (6.3.3.2), structure and union specifiers (6.5.2.1).

### 6.3.2.4 Postfix increment and decrement operators

## Constraints

1 The operand of the postfix increment or decrement operator shall have qualified or unqualified real or pointer type and shall be a modifiable lvalue.

## Semantics

2 The result of the postfix ++ operator is the value of the operand. After the result is obtained, the value of the operand is incremented. (That is, the value 1 of the appropriate type is added to it.) See the discussions of additive operators and compound assignment for information on constraints, types, and conversions and the effects of operations on pointers. The side effect of updating the stored value of the operand shall occur between the previous and the next sequence point.

3 The postfix -- operator is analogous to the postfix ++ operator, except that the value of the operand is decremented (that is, the value 1 of the appropriate type is subtracted from it).

Forward references: additive operators (6.3.6), compound assignment (6.3.16.2).

### 6.3.2.5 Compound literals

## Constraints

1 The type name shall specify an object type or an array of unknown size.
2 No initializer shall attempt to provide a value for an object not contained within the entire unnamed object specified by the compound literal.

3 If the compound literal occurs outside the body of a function, the initializer list shall consist of constant expressions.

## Semantics

4 A postfix expression that consists of a parenthesized type name followed by a braceenclosed list of initializers is a compound literal. It provides an unnamed object whose value is given by the initializer list. ${ }^{68}$

[^27] types or void only, and the result of a cast expression is not an lvalue.

If the type name specifies an array of unknown size, the size is determined by the initializer list as specified in 6.5.7, and the type of the compound literal is that of the completed array type. Otherwise (when the type name specifies an object type), the type of the compound literal is that specified by the type name. In either case, the result is an lvalue.

6 The value of the compound literal is that of an unnamed object initialized by the initializer list. The object has static storage duration if and only if the compound literal occurs outside the body of a function; otherwise, it has automatic storage duration associated with the enclosing block.

7 All the semantic rules and constraints for initializer lists in 6.5.8 are applicable to compound literals. ${ }^{69}$

8 String literals, and compound literals with const-qualified types, need not designate distinct objects. ${ }^{70}$

## Examples

9 1. The file scope definition

```
int *p = (int []){2, 4};
```

initializes $\mathbf{p}$ to point to the first element of an array of two ints, the first having the value two and the second, four. The expressions in this compound literal must be constant. The unnamed object has static storage duration.
2. In contrast, in

```
void f(void)
{
        int *p;
        /*...*/
        p = (int [2]){*p};
        /*...*/
    }
```

$\mathbf{p}$ is assigned the address of the first element of an array of two ints, the first having the value previously pointed to by $\mathbf{p}$ and the second, zero. The expressions in this compound literal need not be constant. The unnamed object
69. For example, subobjects without explicit initializers are initialized to zero.
70. This allows implementations to share storage for string literals and constant compound literals with the same or overlapping representations.
has automatic storage duration.
3. Initializers with designations can be combined with compound literals. Structure objects created using compound literals can be passed to functions without depending on member order:

```
drawline((struct point) {.x=1, . y=1},
    (struct point) {.x=3, . y=4});
```

Or, if drawline instead expected pointers to struct point:

```
drawline(&(struct point) {.x=1, . y=1},
    &(struct point) {.x=3, . y=4});
```

4. A read-only compound literal can be specified through constructions like:
(const float []) \{1e0, 1e1, 1e2, 1e3, 1e4, 1e5, 1e6\}
5. The following three expressions have different meanings:
"/tmp/filexxxxxx"
(char []) \{"/tmp/fileXXXXXX"\}
(const char[]) \{"/tmp/fileXXXXXX"\}
The first always has static storage duration and has type array of char, but need not be modifiable; the last two have automatic storage duration when they occur within the body of a function, and the first of these two is modifiable.
6. Like string literals, const-qualified compound literals can be placed into readonly memory and can even be shared. For example,
```
(const char[]){"abc"} == "abc"
```

might yield 1 if the literals' storage is shared.
7. Since compound literals are unnamed, a single compound literal cannot specify a circularly linked object. For example, there is no way to write a self-referential compound literal that could be used as the function argument in place of the named object endless_zeros below:

```
struct int_list { int car; struct int_list *cdr; };
struct int_list endless_zeros = {0, &endless_zeros};
eval(endless_zeros);
```

8. Outside the body of a function, a compound literal is an initialization of a static object; however, inside a function body, it is an assignment to an automatic object. Therefore, the following two loops produce the same sequence of values for the objects associated with their respective compound literals.
```
for (int i = 0; i < 10; i++) {
    f((struct s){.a = i, .b = 42});
}
for (int i = 0; i < 10; i++)
    f((struct s){.a = i, .b = 42});
```

9. Each compound literal creates only a single object in a given scope:
```
struct s { int i; };
int f (void)
{
    struct s *p = 0, *q;
    int j;
        for (j = 0; j < 2; j++)
            q = p, p = &((struct s){ j });
        return p == q && q.i == 1;
}
```

The function $\mathbf{f}()$ always returns the value 1 .
Note that if the body of the for loop were enclosed in braces, the lifetime of the unnamed object would be the body of the loop only, and on entry next time around $\mathbf{p}$ would be pointing to an object which is no longer guaranteed to exist, which is undefined behavior.

### 6.3.3 Unary operators

## Syntax

```
unary-expression:
    postfix-expression
    ++ unary-expression
    -- unary-expression
    unary-operator cast-expression
    sizeof unary-expression
    sizeof ( type-name )
```

    unary-operator: one of
    \& * + - ~ !
    
### 6.3.3.1 Prefix increment and decrement operators

## Constraints

1 The operand of the prefix increment or decrement operator shall have qualified or unqualified real or pointer type and shall be a modifiable lvalue.

## Semantics

2 The value of the operand of the prefix ++ operator is incremented. The result is the new value of the operand after incrementation. The expression $\mathbf{+ + E}$ is equivalent to (E+=1). See the discussions of additive operators and compound assignment for information on constraints, types, side effects, and conversions and the effects of operations on pointers.

3 The prefix -- operator is analogous to the prefix ++ operator, except that the value of the operand is decremented.

Forward references: additive operators (6.3.6), compound assignment (6.3.16.2).

### 6.3.3.2 Address and indirection operators

## Constraints

1 The operand of the unary $\&$ operator shall be either a function designator, the result of a [] or unary * operator, or an lvalue that designates an object that is not a bit-field and is not declared with the register storage-class specifier.

2 The operand of the unary * operator shall have pointer type.

## Semantics

3 The result of the unary $\&$ (address-of) operator is a pointer to the object or function designated by its operand. If the operand has type "type", the result has type "pointer to type". If the operand is the result of a unary * operator, neither that operator nor the $\&$ operator are evaluated, and the result shall be as if both were omitted, even if the intermediate object does not exist, except that the constraints on the operators still apply and the result is not an lvalue. Similarly, if the operand is the result of a [] operator, neither the $\&$ operator nor the unary * that is implied by the [] are evaluated, and the result shall be as if the $\&$ operator was removed and the [] operator was changed to $\mathrm{a}+$ operator.

4 The unary * operator denotes indirection. If the operand points to a function, the result is a function designator; if it points to an object, the result is an lvalue designating the object. If the operand has type "pointer to type", the result has type "type". If an invalid value has been assigned to the pointer, the behavior of the unary * operator is undefined. ${ }^{71}$

Forward references: storage-class specifiers (6.5.1), structure and union specifiers (6.5.2.1).

### 6.3.3.3 Unary arithmetic operators

## Constraints

1 The operand of the unary + or - operator shall have arithmetic type; of the $\sim$ operator, integer type; of the ! operator, scalar type.

## Semantics

2 The result of the unary + operator is the value of its operand. The integer promotion is performed on the operand, and the result has the promoted type.

3 The result of the unary - operator is the negative of its operand. The integer promotion is performed on the operand, and the result has the promoted type.

4 The result of the $\sim$ operator is the bitwise complement of its operand (that is, each bit in the result is set if and only if the corresponding bit in the converted operand is not set). The integer promotion is performed on the operand, and the result has the promoted type. The expression $\sim \mathbf{E}$ is equivalent to (ULLONG_MAX-E) if $\mathbf{E}$ is promoted to type unsigned long long, to (ULONG_MAX-E) if $\mathbf{E}$ is promoted to type unsigned long, to (UINT_MAX-E) if $E$ is promoted to type unsigned int. (The constants ULLONG_MAX, ULONG_MAX, and UINT_MAX are defined in the header <limits.h>.)

5 The result of the logical negation operator ! is 0 if the value of its operand compares unequal to 0,1 if the value of its operand compares equal to 0 . The result has type int. The expression $!\mathbf{E}$ is equivalent to ( $0=\boldsymbol{=} \mathbf{E}$ ).

[^28]Forward references: limits <float. h> and <limits.h> (7.1.5).

### 6.3.3.4 The sizeof operator

## Constraints

1 The sizeof operator shall not be applied to an expression that has function type or an incomplete type, to the parenthesized name of such a type, or to an lvalue that designates a bit-field object.

## Semantics

2 The sizeof operator yields the size (in bytes) of its operand, which may be an expression or the parenthesized name of a type. The size is determined from the type of the operand. The result is an integer. If the type of the operand is a variable length array type, the operand is evaluated; otherwise, the operand is not evaluated and the result is an integer constant.

3 When applied to an operand that has type char, unsigned char, or signed char, (or a qualified version thereof) the result is 1 . When applied to an operand that has array type, the result is the total number of bytes in the array. ${ }^{72}$ When applied to an operand that has structure or union type, the result is the total number of bytes in such an object, including internal and trailing padding.

The value of the result is implementation-defined, and its type (an unsigned integer type) is size_t defined in the <stddef. $\mathbf{h} \boldsymbol{\text { s }}$ header.

## Examples

5 1. A principal use of the sizeof operator is in communication with routines such as storage allocators and I/O systems. A storage-allocation function might accept a size (in bytes) of an object to allocate and return a pointer to void. For example:

```
extern void *alloc(size_t);
double *dp = alloc(sizeof *dp);
```

The implementation of the alloc function should ensure that its return value is aligned suitably for conversion to a pointer to double.
2. Another use of the sizeof operator is to compute the number of elements in an array:

[^29]```
sizeof array / sizeof array[0]
```

3. In this example, the size of a variable-length array is computed and returned from a function:
```
size_t fsize3 (int n)
{
    char b[n+3]; // Variable length array.
    return sizeof b; // Execution time sizeof.
}
int main()
{
    size_t size;
    size = fsize3(10); // fsize3 returns l3.
    return 0;
}
```

Forward references: common definitions <stddef.h> (7.1.6), declarations (6.5), structure and union specifiers (6.5.2.1), type names (6.5.6), array declarators (6.5.5.2).

### 6.3.4 Cast operators

## Syntax

cast-expression:
unary-expression
( type-name ) cast-expression

## Constraints

2 Unless the type name specifies a void type, the type name shall specify qualified or unqualified scalar type and the operand shall have scalar type.

3 Conversions that involve pointers, other than where permitted by the constraints of 6.3.16.1, shall be specified by means of an explicit cast.

## Semantics

4 Preceding an expression by a parenthesized type name converts the value of the expression to the named type. This construction is called a cast. ${ }^{73}$ A cast that specifies no conversion has no effect on the type or value of an expression. ${ }^{74}$
73. A cast does not yield an lvalue. Thus, a cast to a qualified type has the same effect as a cast to the unqualified version of the type.
74. If the value of the expression is represented with greater precision or range than required by the type named by the cast (6.2.1.7), then the cast specifies a conversion even if the type of the expression is the same as the named type.

Forward references: equality operators (6.3.9), function declarators (including prototypes) (6.5.5.3), simple assignment (6.3.16.1), type names (6.5.6).

### 6.3.5 Multiplicative operators

## Syntax

        multiplicative-expression:
    cast-expression
multiplicative-expression * cast-expression
multiplicative-expression / cast-expression
multiplicative-expression \% cast-expression

## Constraints

2 Each of the operands shall have arithmetic type. The operands of the \% operator shall have integer type.

## Semantics

3 The usual arithmetic conversions are performed on the operands.
4 The result of the binary * operator is the product of the operands.
5 The result of the / operator is the quotient from the division of the first operand by the second; the result of the \% operator is the remainder. In both operations, if the value of the second operand is zero, the behavior is undefined.

6 When integers are divided, the result of the / operator is the algebraic quotient with any fractional part discarded. ${ }^{75}$ If the quotient $\mathbf{a} / \mathbf{b}$ is representable, the expression $(a / b) * b+a \% b$ shall equal $a$.

7 If either operand has complex type, the result has complex type.

[^30]
### 6.3.6 Additive operators

## Syntax

multiplicative-expression
additive-expression + multiplicative-expression
additive-expression - multiplicative-expression

## Constraints

2 For addition, either both operands shall have arithmetic type, or one operand shall be a pointer to an object type and the other shall have integer type. (Incrementing is equivalent to adding 1.)

3 For subtraction, one of the following shall hold:

- both operands have arithmetic type;
- both operands are pointers to qualified or unqualified versions of compatible object types; or
- the left operand is a pointer to an object type and the right operand has integer type. (Decrementing is equivalent to subtracting 1.)


## Semantics

4 If both operands have arithmetic type, the usual arithmetic conversions are performed on them.

5 The result of the binary + operator is the sum of the operands.
6 The result of the binary - operator is the difference resulting from the subtraction of the second operand from the first.

7 For the purposes of these operators, a pointer to a nonarray object behaves the same as a pointer to the first element of an array of length one with the type of the object as its element type.

When an expression that has integer type is added to or subtracted from a pointer, the result has the type of the pointer operand. If the pointer operand points to an element of an array object, and the array is large enough, the result points to an element offset from the original element such that the difference of the subscripts of the resulting and original array elements equals the integer expression. In other words, if the expression $\mathbf{P}$ points to the $i$-th element of an array object, the expressions $(\mathbf{P})+\mathbf{N}$ (equivalently, $\mathbf{N}+(\mathbf{P})$ ) and $(\mathbf{P})-\mathbf{N}$ (where $\mathbf{N}$ has the value $n$ ) point to, respectively, the $i+n$-th and $i-n$-th elements of the array object, provided they exist. Moreover, if the expression $\mathbf{P}$ points to the last element of an array object, the expression $(P)+1$ points one past the last element of the array object, and if the expression $\mathbf{Q}$ points one past the last
element of an array object, the expression (Q) $\mathbf{- 1}$ points to the last element of the array object. If both the pointer operand and the result point to elements of the same array object, or one past the last element of the array object, the evaluation shall not produce an overflow; otherwise, the behavior is undefined. Unless both the pointer operand and the result point to elements of the same array object, or the pointer operand points one past the last element of an array object and the result points to an element of the same array object, the behavior is undefined if the result is used as an operand of a unary * operator that is actually evaluated.

9 When two pointers to elements of the same array object are subtracted, the result is the difference of the subscripts of the two array elements. The size of the result is implementation-defined, and its type (a signed integer type) is ptrdiff_t defined in the <stddef. $\mathrm{h}>$ header. If the result is not representable in an object of that type, the behavior is undefined. In other words, if the expressions $\mathbf{P}$ and $\mathbf{Q}$ point to, respectively, the $i$-th and $j$-th elements of an array object, the expression ( $\mathbf{P}$ ) - (Q) has the value $i-j$ provided the value fits in an object of type ptrdiff_t. Moreover, if the expression $\mathbf{P}$ points either to an element of an array object or one past the last element of an array object, and the expression $\mathbf{Q}$ points to the last element of the same array object, the expression $((Q)+1)-(P)$ has the same value as $((Q)-(P))+1$ and as $-((P)-((Q)+1))$, and has the value zero if the expression $P$ points one past the last element of the array object, even though the expression $(Q)+1$ does not point to an element of the array object. Unless both pointers point to elements of the same array object, or one past the last element of the array object, the behavior is undefined. ${ }^{76}$

10 If either operand has complex type, the result has complex type.

## Examples

11 Pointer arithmetic is well defined with pointers to variable length array types.

[^31]```
{
```

```
    int n = 4, m = 3;
int a[n][m];
int (*p)[m] = a; // p == &a[0]
p += 1; // p == &a[1]
    (*p)[2] = 99; // a[1][2] == 99
    n = p - a; // n == 1
```

\}

12 If array a in the above example is declared to be an array of known constant size, and pointer $\mathbf{p}$ is declared to be a pointer to an array of the same know constant size that points to $\mathbf{a}$, the results are the same.

Forward references: array declarators (6.5.5.2), common definitions <stddef.h> (7.1.6).

### 6.3.7 Bitwise shift operators

## Syntax

shift-expression:
additive-expression
shift-expression $\ll$ additive-expression
shift-expression >> additive-expression

## Constraints

2 Each of the operands shall have integer type.

## Semantics

3 The integer promotions are performed on each of the operands. The type of the result is that of the promoted left operand. If the value of the right operand is negative or is greater than or equal to the number of value and sign bits in the object representation of the promoted left operand, the behavior is undefined.

4 The result of $\mathbf{E} 1 \ll \mathbf{E} 2$ is $\mathbf{E} 1$ left-shifted $\mathbf{E} 2$ bit positions; vacated bits are filled with zeros. If $\mathbf{E 1}$ has an unsigned type, the value of the result is $\mathbf{E} \mathbf{1} \times 2^{\mathbf{E} 2}$, reduced modulo ULLONG_MAX+1 if E1 has type unsigned long long, ULONG_MAX+1 if E1 has type unsigned long, UINT_MAX+1 otherwise. (The constants ULLONG_MAX, ULONG_MAX, and UINT_MAX are defined in the header <limits.h>.) If E1 has a signed type and nonnegative value, and $\mathbf{E} \mathbf{1} \times 2^{\mathbf{E 2}}$ is less than or equal to INT_MAX (if E1 has type int), LONG_MAX (if E1 has type long int), or LLONG_MAX (if E1 has type long long int), then that is the resulting value. Otherwise, the behavior is undefined.

5 The result of $\mathbf{E 1} \gg \mathbf{E} 2$ is $\mathbf{E} 1$ right-shifted $\mathbf{E} 2$ bit positions. If $\mathbf{E} 1$ has an unsigned type or if $\mathbf{E 1}$ has a signed type and a nonnegative value, the value of the result is the integral part of the quotient of $\mathbf{E 1}$ divided by the quantity, 2 raised to the power $\mathbf{E} \mathbf{2}$. If E1 has a signed type and a negative value, the resulting value is implementationdefined.

### 6.3.8 Relational operators

## Syntax

> relational-expression: shift-expression relational-expression < relational-expression $>$ shift-expression relational-expression $<=$ shift-expression relational-expression $>=$ shift-expression

## Constraints

2 One of the following shall hold:

- both operands have real type;
- both operands are pointers to qualified or unqualified versions of compatible object types; or
- both operands are pointers to qualified or unqualified versions of compatible incomplete types.


## Semantics

3 If both of the operands have arithmetic type, the usual arithmetic conversions are performed.

4 For the purposes of these operators, a pointer to a nonarray object behaves the same as a pointer to the first element of an array of length one with the type of the object as its element type.

5 When two pointers are compared, the result depends on the relative locations in the address space of the objects pointed to. If two pointers to object or incomplete types both point to the same object, or both point one past the last element of the same array object, they compare equal. If the objects pointed to are members of the same aggregate object, pointers to structure members declared later compare greater than pointers to members declared earlier in the structure, and pointers to array elements with larger subscript values compare greater than pointers to elements of the same array with lower subscript values. All pointers to members of the same union object compare equal. If the expression $\mathbf{P}$ points to an element of an array object and the expression $\mathbf{Q}$ points to the last element of the same array object, the pointer expression
$\mathbf{Q}+\mathbf{1}$ compares greater than $\mathbf{P}$. In all other cases, the behavior is undefined.
6 Each of the operators $<$ (less than), $>$ (greater than), $<=$ (less than or equal to), and $>=$ (greater than or equal to) shall yield 1 if the specified relation is true and 0 if it is false. ${ }^{77}$ The result has type int.

### 6.3.9 Equality operators

## Syntax

```
equality-expression:
    relational-expression
    equality-expression == relational-expression
    equality-expression != relational-expression
```


## Constraints

2 One of the following shall hold:

- both operands have arithmetic type;
- both operands are pointers to qualified or unqualified versions of compatible types;
- one operand is a pointer to an object or incomplete type and the other is a pointer to a qualified or unqualified version of void; or
- one operand is a pointer and the other is a null pointer constant.


## Semantics

3 The $==$ (equal to) and the $!=$ (not equal to) operators are analogous to the relational operators except for their lower precedence. ${ }^{78}$ Where the operands have types and values suitable for the relational operators, the semantics detailed in 6.3.8 apply.

4 If two pointers to object or incomplete types are both null pointers, they compare equal. If two pointers to object or incomplete types compare equal, they both are null pointers, or both point to the same object, or both point one past the last element of the same array object. ${ }^{79}$ If two pointers to function types are both null pointers or both point to the same function, they compare equal. If two pointers to function types compare equal, either both are null pointers, or both point to the same function. If one

[^32]of the operands is a pointer to an object or incomplete type and the other has type pointer to a qualified or unqualified version of void, the pointer to an object or incomplete type is converted to the type of the other operand.

5 Values of complex types are equal if and only if both their real parts are equal and also their imaginary parts are equal. Any two values of arithmetic types from different type-domains are equal if and only if the results of their conversion to the complex type corresponding to the common real type determined by the usual arithmetic conversions are equal.

### 6.3.10 Bitwise AND operator

Syntax
AND-expression:
equality-expression
AND-expression \& equality-expression

## Constraints

2 Each of the operands shall have integer type.

## Semantics

3 The usual arithmetic conversions are performed on the operands.
4 The result of the binary \& operator is the bitwise AND of the operands (that is, each bit in the result is set if and only if each of the corresponding bits in the converted operands is set).

### 6.3.11 Bitwise exclusive OR operator

Syntax

> exclusive-OR-expression:
> AND-expression
> exclusive-OR-expression ^ AND-expression

## Constraints

2 Each of the operands shall have integer type.

## Semantics

3 The usual arithmetic conversions are performed on the operands.
4 The result of the ${ }^{\wedge}$ operator is the bitwise exclusive OR of the operands (that is, each bit in the result is set if and only if exactly one of the corresponding bits in the converted operands is set).

### 6.3.12 Bitwise inclusive OR operator

## Syntax

```
inclusive-OR-expression:
    exclusive-OR-expression
    inclusive-OR-expression | exclusive-OR-expression
```


## Constraints

2 Each of the operands shall have integer type.

## Semantics

3 The usual arithmetic conversions are performed on the operands.
4 The result of the $\mid$ operator is the bitwise inclusive OR of the operands (that is, each bit in the result is set if and only if at least one of the corresponding bits in the converted operands is set).

### 6.3.13 Logical AND operator

Syntax

> logical-AND-expression:
> inclusive-OR-expression
> logical-AND-expression $\& \&$ inclusive-OR-expression

## Constraints

2 Each of the operands shall have scalar type.

## Semantics

3 The $\& \&$ operator shall yield 1 if both of its operands compare unequal to 0 ; otherwise, it yields 0 . The result has type int.

4 Unlike the bitwise binary $\&$ operator, the $\& \&$ operator guarantees left-to-right evaluation; there is a sequence point after the evaluation of the first operand. If the first operand compares equal to 0 , the second operand is not evaluated.

### 6.3.14 Logical OR operator

## Syntax

logical-AND-expression
logical-OR-expression || logical-AND-expression

## Constraints

2 Each of the operands shall have scalar type.

## Semantics

3 The || operator shall yield 1 if either of its operands compare unequal to 0 ; otherwise, it yields 0 . The result has type int.

4 Unlike the bitwise \| operator, the || operator guarantees left-to-right evaluation; there is a sequence point after the evaluation of the first operand. If the first operand compares unequal to 0 , the second operand is not evaluated.

### 6.3.15 Conditional operator

## Syntax

conditional-expression:
logical-OR-expression
logical-OR-expression ? expression : conditional-expression

## Constraints

2 The first operand shall have scalar type.
3 One of the following shall hold for the second and third operands:

- both operands have arithmetic type;
- both operands have compatible structure or union types;
- both operands have void type;
- both operands are pointers to qualified or unqualified versions of compatible types;
- one operand is a pointer and the other is a null pointer constant; or
- one operand is a pointer to an object or incomplete type and the other is a pointer to a qualified or unqualified version of void.


## Semantics

4 The first operand is evaluated; there is a sequence point after its evaluation. The second operand is evaluated only if the first compares unequal to 0 ; the third operand is evaluated only if the first compares equal to 0 ; the value of the second or third operand (whichever is evaluated) is the result. ${ }^{80}$

5 If both the second and third operands have arithmetic type, the usual arithmetic conversions are performed to bring them to a common type and the result has that type. If both the operands have structure or union type, the result has that type. If both operands have void type, the result has void type.
6 If both the second and third operands are pointers or one is a null pointer constant and the other is a pointer, the result type is a pointer to a type qualified with all the type qualifiers of the types pointed-to by both operands. Furthermore, if both operands are pointers to compatible types or differently qualified versions of a compatible type, the result has the composite type; if one operand is a null pointer constant, the result has the type of the other operand; otherwise, one operand is a pointer to void or a qualified version of void, in which case the other operand is converted to type pointer to void, and the result has that type.

## Examples

7 The common type that results when the second and third operands are pointers is determined in two independent stages. The appropriate qualifiers, for example, do not depend on whether the two pointers have compatible types.

Given the declarations

```
const void *c_vp;
void *vp;
const int *c_ip;
volatile int *v_ip;
int *ip;
const char *c_cp;
```

the third column in the following table is the common type that is the result of a conditional expression in which the first two columns are the second and third operands (in either order):

[^33]| c_vp c_ip | const void * |
| :--- | :--- |
| v_ip 0 | volatile int * |
| c_ip v_ip | const volatile int * |
| vp c_cp | const void * |
| ip $\quad$ c_ip | const int * |
| vp ip | void * |

### 6.3.16 Assignment operators

## Syntax

assignment-expression:
conditional-expression
unary-expression assignment-operator assignment-expression
assignment-operator: one of

$$
=*=1=\% \quad+=-=\lll>=\Leftrightarrow=\wedge=\mid=
$$

## Constraints

2 An assignment operator shall have a modifiable lvalue as its left operand.

## Semantics

3 An assignment operator stores a value in the object designated by the left operand. An assignment expression has the value of the left operand after the assignment, but is not an lvalue. The type of an assignment expression is the type of the left operand unless the left operand has qualified type, in which case it is the unqualified version of the type of the left operand. The side effect of updating the stored value of the left operand shall occur between the previous and the next sequence point.

4 The order of evaluation of the operands is unspecified.

### 6.3.16.1 Simple assignment

## Constraints

1 One of the following shall hold: ${ }^{81}$

- the left operand has qualified or unqualified arithmetic type and the right has arithmetic type;

[^34]- the left operand has a qualified or unqualified version of a structure or union type compatible with the type of the right;
- both operands are pointers to qualified or unqualified versions of compatible types, and the type pointed to by the left has all the qualifiers of the type pointed to by the right;
- one operand is a pointer to an object or incomplete type and the other is a pointer to a qualified or unqualified version of void, and the type pointed to by the left has all the qualifiers of the type pointed to by the right; or
- the left operand is a pointer and the right is a null pointer constant.


## Semantics

2 In simple assignment ( $=$ ), the value of the right operand is converted to the type of the assignment expression and replaces the value stored in the object designated by the left operand.

If the value being stored in an object is accessed from another object that overlaps in any way the storage of the first object, then the overlap shall be exact and the two objects shall have qualified or unqualified versions of a compatible type; otherwise, the behavior is undefined.

## Examples

1. In the program fragment
```
int f(void);
char c;
/* ... */
if ((c = f()) == -1)
    /* ... */
```

the int value returned by the function may be truncated when stored in the char, and then converted back to int width prior to the comparison. In an implementation in which 'plain'" char has the same range of values as unsigned char (and char is narrower than int), the result of the conversion cannot be negative, so the operands of the comparison can never compare equal. Therefore, for full portability, the variable c should be declared as int.
2. In the fragment:

```
char c;
int i;
long l;
```

$1=(c=i) ;$
the value of $\mathbf{i}$ is converted to the type of the assignment-expression $\mathbf{c}=\mathbf{i}$, that is, char type. The value of the expression enclosed in parentheses is then converted to the type of the outer assignment-expression, that is, long type.

### 6.3.16.2 Compound assignment

## Constraints

1 For the operators $+=$ and $-=$ only, either the left operand shall be a pointer to an object type and the right shall have integer type, or the left operand shall have qualified or unqualified arithmetic type and the right shall have arithmetic type.

2 For the other operators, each operand shall have arithmetic type consistent with those allowed by the corresponding binary operator.

## Semantics

3 A compound assignment of the form $\mathbf{E 1}$ op $=\mathbf{E} 2$ differs from the simple assignment expression $\mathbf{E} 1=\mathbf{E} 1$ op ( $\mathbf{E} \mathbf{2}$ ) only in that the lvalue $\mathbf{E} 1$ is evaluated only once.

### 6.3.17 Comma operator

## Syntax

assignment-expression
expression, assignment-expression

## Semantics

2 The left operand of a comma operator is evaluated as a void expression; there is a sequence point after its evaluation. Then the right operand is evaluated; the result has its type and value. ${ }^{82}$

[^35]
## Examples

As indicated by the syntax, in contexts where a comma is a punctuator (in lists of arguments to functions and lists of initializers) the comma operator as described in this subclause cannot appear. On the other hand, it can be used within a parenthesized expression or within the second expression of a conditional operator in such contexts. In the function call

```
f(a, (t=3, t+2), c)
```

the function has three arguments, the second of which has the value 5 .
Forward references: initialization (6.5.8).

### 6.4 Constant expressions

## Syntax

## constant-expression:

conditional-expression

## Description

2 A constant expression can be evaluated during translation rather than runtime, and accordingly may be used in any place that a constant may be.

## Constraints

3 Constant expressions shall not contain assignment, increment, decrement, function-call, or comma operators, except when they are contained within the operand of asizeof operator. ${ }^{83}$

4 Each constant expression shall evaluate to a constant that is in the range of representable values for its type.

## Semantics

5 An expression that evaluates to a constant is required in several contexts. ${ }^{84}$ If a floating expression is evaluated in the translation environment, the arithmetic precision and range shall be at least as great as if the expression were being evaluated in the execution environment.

6 An integer constant expression shall have integer type and shall only have operands that are integer constants, enumeration constants, character constants, sizeof expressions whose operand does not have variable length array type or a parenthesized name of such a type, and floating constants that are the immediate operands of casts. Cast operators in an integer constant expression shall only convert arithmetic types to integer types, except as part of an operand to the sizeof operator.

7 More latitude is permitted for constant expressions in initializers. Such a constant expression shall be, or evaluate to, one of the following:
83. The operand of a sizeof operator is not evaluated (6.3.3.4), and thus any operator in 6.3 may be used.
84. An integer constant expression must be used to specify the size of a bit-field member of a structure, the value of an enumeration constant, the size of an array, or the value of a case constant. Further constraints that apply to the integer constant expressions used in conditional-inclusion preprocessing directives are discussed in 6.8.1.

- an arithmetic constant expression,
- a null pointer constant,
- an address constant, or
- an address constant for an object type plus or minus an integer constant expression.

An arithmetic constant expression shall have arithmetic type and shall only have operands that are integer constants, floating constants, enumeration constants, character constants, and sizeof expressions. Cast operators in an arithmetic constant expression shall only convert arithmetic types to arithmetic types, except as part of an operand to the sizeof operator.

9 An address constant is a null pointer, a pointer to an lvalue designating an object of static storage duration, or to a function designator; it shall be created explicitly using the unary \& operator or an integer constant cast to pointer type, or implicitly by the use of an expression of array or function type. The array-subscript [] and memberaccess . and $\rightarrow$ operators, the address $\&$ and indirection * unary operators, and pointer casts may be used in the creation of an address constant, but the value of an object shall not be accessed by use of these operators.

10 An implementation may accept other forms of constant expressions.
11 The semantic rules for the evaluation of a constant expression are the same as for nonconstant expressions. ${ }^{85}$

Forward references: array declarators (6.5.5.2), initialization (6.5.8).

[^36]
### 6.5 Declarations

## Syntax

2 A declaration shall declare at least a declarator (excluding the parameters of a function or the members of a structure or union), a tag, or the members of an enumeration.

3 If an identifier has no linkage, there shall be no more than one declaration of the identifier (in a declarator or type specifier) with the same scope and in the same name space, except for tags as specified in 6.5.2.3.
4 All declarations in the same scope that refer to the same object or function shall specify compatible types.

## Semantics

5 A declaration specifies the interpretation and attributes of a set of identifiers. A definition of an identifier is a declaration for that identifier that:

- for an object, causes storage to be reserved for that object;
— for a function, includes the function body; ${ }^{86}$
- for an enumeration constant or typedef name, is the (only) declaration of the identifier.

86. Function definitions have a different syntax, described in 6.7.1.

6 The declaration specifiers consist of a sequence of specifiers that indicate the linkage, storage duration, and part of the type of the entities that the declarators denote. The init-declarator-list is a comma-separated sequence of declarators, each of which may have additional type information, or an initializer, or both. The declarators contain the identifiers (if any) being declared.

7 If an identifier for an object is declared with no linkage, the type for the object shall be complete by the end of its declarator, or by the end of its init-declarator if it has an initializer.

Forward references: declarators (6.5.5), enumeration specifiers (6.5.2.2), initialization (6.5.8), tags (6.5.2.3).

### 6.5.1 Storage-class specifiers

Syntax
typedef
extern
static
auto
register

## Constraints

2 At most, one storage-class specifier may be given in the declaration specifiers in a declaration. ${ }^{87}$

## Semantics

3 The typedef specifier is called a "storage-class specifier" for syntactic convenience only; it is discussed in 6.5.7. The meanings of the various linkages and storage durations were discussed in 6.1.2.2 and 6.1.2.4.

4 A declaration of an identifier for an object with storage-class specifier register suggests that access to the object be as fast as possible. The extent to which such suggestions are effective is implementation-defined. ${ }^{88}$
87. See 'future language directions" (6.9.2).
88. The implementation may treat any register declaration simply as an auto declaration. However, whether or not addressable storage is actually used, the address of any part of an object declared with storage-class specifier register may not be computed, either explicitly (by use of the unary \& operator as discussed in 6.3.3.2) or implicitly (by converting an array name to a pointer as discussed in 6.2.2.1). Thus the only operator that can be applied to an array declared with storage-class specifier register is sizeof.

5 The declaration of an identifier for a function that has block scope shall have no explicit storage-class specifier other than extern.

6 If an aggregate or union object is declared with a storage-class specifier other than typedef, the properties resulting from the storage-class specifier, except with respect to linkage, also apply to the members of the object, and so on recursively for any aggregate or union member objects.

Forward references: type definitions (6.5.7).

### 6.5.2 Type specifiers

Syntax
type-specifier:
void
char
short
int
long
float
double
complex
signed
unsigned
struct-or-union-specifier
enum-specifier
typedef-name

## Constraints

2 At least one type specifier shall be given in the declaration specifiers in a declaration. Each list of type specifiers shall be one of the following sets (delimited by commas, when there is more than one set on a line); the type specifiers may occur in any order, possibly intermixed with the other declaration specifiers.

- void
- char
- signed char
- unsigned char
- short, signed short, short int, or signed short int
- unsigned short, or unsigned short int
- int, signed, or signed int
- unsigned, or unsigned int
- long, signed long, long int, or signed long int
- unsigned long, or unsigned long int
- long long, signed long long, long long int, or signed long long int
- unsigned long long, or unsigned long long int
- float
- double
- long double
- float complex
- double complex
- long double complex
- struct-or-union specifier
- enum-specifier
- typedef-name

Semantics
3 Specifiers for structures, unions, and enumerations are discussed in 6.5.2.1 through 6.5.2.3. Declarations of typedef names are discussed in 6.5.7. The characteristics of the other types are discussed in 6.1.2.5.

4 Each of the comma-separated sets designates the same type, except that for bit-fields, it is implementation-defined whether the specified int is the same type as signed int or is the same type as unsigned int.

Forward references: enumeration specifiers (6.5.2.2), structure and union specifiers (6.5.2.1), tags (6.5.2.3), type definitions (6.5.7).

### 6.5.2.1 Structure and union specifiers

Syntax
struct-or-union-specifier:
struct-or-union identifier ${ }_{\text {opt }}\{$ struct-declaration-list \}
struct-or-union identifier

```
struct-or-union:
    struct
    union
struct-declaration-list:
    struct-declaration
    struct-declaration-list struct-declaration
struct-declaration:
    specifier-qualifier-list struct-declarator-list ;
specifier-qualifier-list:
    type-specifier specifier-qualifier-list \({ }_{\text {opt }}\)
    type-qualifier specifier-qualifier-list opt
struct-declarator-list:
    struct-declarator
    struct-declarator-list , struct-declarator
struct-declarator:
    declarator
    declarator \(_{\text {opt }}\) : constant-expression
```


## Constraints

2 A structure or union shall not contain a member with incomplete or function type, except that the last element of a structure may have incomplete array type. Hence it shall not contain an instance of itself (but may contain a pointer to an instance of itself).

3 The expression that specifies the width of a bit-field shall be an integer constant expression that has nonnegative value that shall not exceed the number of bits in an object of the type that is specified if the colon and expression are omitted. If the value is zero, the declaration shall have no declarator.

## Semantics

4 As discussed in 6.1.2.5, a structure is a type consisting of a sequence of members, whose storage is allocated in an ordered sequence, and a union is a type consisting of a sequence of members whose storage overlap.

5 Structure and union specifiers have the same form.
6 The presence of a struct-declaration-list in a struct-or-union-specifier declares a new type, within a translation unit. The struct-declaration-list is a sequence of declarations for the members of the structure or union. If the struct-declaration-list contains no named members, the behavior is undefined. The type is incomplete until after the \} that terminates the list.

7 A member of a structure or union may have any object type other than a variably modified type. ${ }^{89}$ In addition, a member may be declared to consist of a specified number of bits (including a sign bit, if any). Such a member is called a bit-field; ${ }^{90}$ its width is preceded by a colon.

8 A bit-field shall have a type that is a qualified or unqualified version of signed int or unsigned int. A bit-field is interpreted as a signed or unsigned integer type consisting of the specified number of bits. ${ }^{91}$

9 An implementation may allocate any addressable storage unit large enough to hold a bit-field. If enough space remains, a bit-field that immediately follows another bitfield in a structure shall be packed into adjacent bits of the same unit. If insufficient space remains, whether a bit-field that does not fit is put into the next unit or overlaps adjacent units is implementation-defined. The order of allocation of bit-fields within a unit (high-order to low-order or low-order to high-order) is implementation-defined. The alignment of the addressable storage unit is unspecified.

10 A bit-field declaration with no declarator, but only a colon and a width, indicates an unnamed bit-field. ${ }^{92}$ As a special case of this, a bit-field structure member with a width of 0 indicates that no further bit-field is to be packed into the unit in which the previous bit-field, if any, was placed.

11 Each non-bit-field member of a structure or union object is aligned in an implementation-defined manner appropriate to its type.

12 Within a structure object, the non-bit-field members and the units in which bit-fields reside have addresses that increase in the order in which they are declared. A pointer to a structure object, suitably converted, points to its initial member (or if that member is a bit-field, then to the unit in which it resides), and vice versa. There may be unnamed padding within a structure object, but not at its beginning.

13 The size of a union is sufficient to contain the largest of its members. The value of at most one of the members can be stored in a union object at any time. A pointer to a

[^37]union object, suitably converted, points to each of its members (or if a member is a bit-field, then to the unit in which it resides), and vice versa.

14 There may be unnamed padding at the end of a structure or union, were the structure or union to be an element of an array.

15 As a special case, the last element of a structure may be an incomplete array type. This is called a flexible array member, and the size of the structure shall be equal to the offset of the last element of an otherwise identical structure that replaces the flexible array member with an array of one element. When an lvalue whose type is a structure with a flexible array member is used to access an object, it behaves as if that member were replaced by the longest array that would not make the structure larger than the object being accessed. If this array would have no elements, then it behaves as if it has one element, but the behavior is undefined if any attempt is made to access that element.

## Examples

After the declarations:

```
struct s { int n; double d[]; };
struct ss { int n; double d[1]; };
```

the three expressions:

```
sizeof (struct s)
offsetof(struct s, d)
offsetof(struct ss, d)
```

have the same value. The structure struct $\mathbf{s}$ has a flexible array member d.
If sizeof (double) is 8 , then after the following code is executed:

```
struct s *s1;
struct s *s2;
s1 = malloc(sizeof (struct s) + 64);
s2 = malloc(sizeof (struct s) + 46);
```

and assuming that the calls to malloc succeed, s1 and s2 behave as if they had been declared as:
struct $\{$ int $n$; double d[8]; \} *s1;
struct \{ int $n$; double d[5]; \} *s2;
18
Following the further successful assignments:

```
s1 = malloc(sizeof (struct s) + 10);
s2 = malloc(sizeof (struct s) + 6);
```

they then behave as if they had been declared as:

```
struct { int n; double d[1]; } *s1, *s2;
```

and:

```
double *dp;
dp = &(s1->d[0]); // Permitted
*dp = 42; // Permitted
dp = &(s2->d[0]); // Permitted
*dp = 42;
// Undefined behavior
```

Forward references: tags (6.5.2.3).

### 6.5.2.2 Enumeration specifiers

## Syntax

enum-specifier:
enum identifier ${ }_{\text {opt }}\{$ enumerator-list \}
enum identifier opt $\{$ enumerator-list , \}
enum identifier
enumerator-list:
enumerator
enumerator-list , enumerator
enumerator:
enumeration-constant
enumeration-constant $=$ constant-expression

## Constraints

2 The expression that defines the value of an enumeration constant shall be an integer constant expression that has a value representable as an int.

## Semantics

3 The identifiers in an enumerator list are declared as constants that have type int and may appear wherever such are permitted. ${ }^{93}$ An enumerator with $=$ defines its enumeration constant as the value of the constant expression. If the first enumerator has no $=$, the value of its enumeration constant is 0 . Each subsequent enumerator with no $=$ defines its enumeration constant as the value of the constant expression obtained

[^38]by adding 1 to the value of the previous enumeration constant. (The use of enumerators with $=$ may produce enumeration constants with values that duplicate other values in the same enumeration.) The enumerators of an enumeration are also known as its members.

4 Each enumerated type shall be compatible with an integer type. The choice of type is implementation-defined, but shall be capable of representing the values of all the members of the enumeration.

5 The enumerated type is complete at the \} that terminates the list of enumerator declarations.

## Examples

```
enum hue \{ chartreuse, burgundy, claret=20, winedark \};
enum hue col, *cp;
col = claret;
cp = \&col;
if (*cp != burgundy)
/* ... */
```

makes hue the tag of an enumeration, and then declares col as an object that has that type and cp as a pointer to an object that has that type. The enumerated values are in the set $\{0,1,20,21\}$.

Forward references: tags (6.5.2.3).

### 6.5.2.3 Tags

## Constraints

1 A specific type shall have its content defined at most once.
2 A type specifier of the form
enum identifier
without an enumerator list shall only appear after the type it specifies is completed.

## Semantics

3 All declarations of structure, union, or enumerated types that have the same scope and use the same tag declare the same type. The type is complete ${ }^{94}$ until the closing

[^39]brace of the list defining the content, and complete thereafter.
4 Two declarations of structure, union, or enumerated types which are in different scopes or use different tags declare distinct types. Each declaration of a structure, union, or enumerated type which does not include a tag declares a distinct type.

5 A type specifier of the form
struct-or-union identifier $r_{o p t}\{$ struct-declaration-list \}
or
enum identifier \{ enumerator-list \}
or
enum identifier \{ enumerator-list , \}
declares a structure, union, or enumerated type. The list defines the structure content, union content, or enumeration content. If an identifier is provided, ${ }^{95}$ the type specifier also declares the identifier to be the tag of that type.

6 A declaration of the form
struct-or-union identifier
specifies a structure of union type and declares the identifier as a tag of that type. ${ }^{96}$
7 If a type specifier of the form
struct-or-union identifier
occurs other than as part of one of the above forms, and no other declaration of the identifier as a tag is visible, then it declares an incomplete structure or union type, and declares the identifier as the tag of that type. ${ }^{96}$

8 If a type specifier of the form

## struct-or-union identifier

or

```
enum identifier
```

occurs other than as part of one of the above forms, and a declaration of the identifier as a tag is visible, then it specifies the same type as that other declaration, and does

[^40]not redeclare the tag.

## Examples

1. This mechanism allows declaration of a self-referential structure.
```
struct tnode {
    int count;
    struct tnode *left, *right;
};
```

specifies a structure that contains an integer and two pointers to objects of the same type. Once this declaration has been given, the declaration
struct tnode $s, \quad * s p ;$
declares $\boldsymbol{s}$ to be an object of the given type and $\mathbf{s p}$ to be a pointer to an object of the given type. With these declarations, the expression sp->left refers to the left struct tnode pointer of the object to which sp points; the expression s.right->count designates the count member of the right struct tnode pointed to from s.

The following alternative formulation uses the typedef mechanism:

```
typedef struct tnode TNODE;
struct tnode {
    int count;
    TNODE *left, *right;
};
TNODE s, *sp;
```

2. To illustrate the use of prior declaration of a tag to specify a pair of mutually referential structures, the declarations
```
struct s1 { struct s2 *s2p; /* ... */ }; // D1
struct s2 { struct s1 *s1p; /* ... */ }; // D2
```

specify a pair of structures that contain pointers to each other. Note, however, that if s2 were already declared as a tag in an enclosing scope, the declaration D1 would refer to it, not to the tag s2 declared in D2. To eliminate this context sensitivity, the declaration

```
struct s2;
```

may be inserted ahead of D1. This declares a new tag s2 in the inner scope; the declaration D2 then completes the specification of the new type.
3. An enumeration type is compatible with some integer type. An implementation may delay the choice of which integer type until all enumeration constants have
been seen. Thus in:

```
enum f { C = sizeof (enum f) };
```

the behavior is undefined since the size of the respective enumeration type is not necessarily known when sizeof is encountered.

Forward references: declarators (6.5.5), array declarators (6.5.5.2), type definitions (6.5.7).

### 6.5.3 Type qualifiers

## Syntax

const
restrict
volatile

## Constraints

2 Types other than pointer types derived from object or incomplete types shall not be restrict-qualified.

## Semantics

3 The properties associated with qualified types are meaningful only for expressions that are lvalues. ${ }^{97}$

4 If the same qualifier appears more than once in the same specifier-qualifier-list, either directly or via one or more typedefs, the behavior is the same as if it appeared only once.

5 If an attempt is made to modify an object defined with a const-qualified type through use of an lvalue with non-const-qualified type, the behavior is undefined. If an attempt is made to refer to an object defined with a volatile-qualified type through use of an lvalue with non-volatile-qualified type, the behavior is undefined. ${ }^{98}$

[^41]6 An object that has volatile-qualified type may be modified in ways unknown to the implementation or have other unknown side effects. Therefore any expression referring to such an object shall be evaluated strictly according to the rules of the abstract machine, as described in 5.1.2.3. Furthermore, at every sequence point the value last stored in the object shall agree with that prescribed by the abstract machine, except as modified by the unknown factors mentioned previously. ${ }^{99}$ What constitutes an access to an object that has volatile-qualified type is implementation-defined.

7 An object that is referenced through a restrict-qualified pointer has a special association with that pointer. This association, defined in 6.5.3.1 below, requires that all references to that object shall use, directly or indirectly, the value of that pointer. For example, a statement that assigns a value returned by malloc to a single pointer establishes this association between the allocated object and the pointer. The intended use of the restrict qualifier (like the register storage class) is to promote optimization, and deleting all instances of the qualifier from a conforming program does not change its meaning (i.e., observable behavior).

8 If the specification of an array type includes any type qualifiers, the element type is so-qualified, not the array type. If the specification of a function type includes any type qualifiers, the behavior is undefined. ${ }^{100}$

For two qualified types to be compatible, both shall have the identically qualified version of a compatible type; the order of type qualifiers within a list of specifiers or qualifiers does not affect the specified type.

## Examples

1. An object declared
```
extern const volatile int real_time_clock;
```

may be modifiable by hardware, but cannot be assigned to, incremented, or decremented.
2. The following declarations and expressions illustrate the behavior when type qualifiers modify an aggregate type:

[^42]```
const struct s { int mem; } cs = { 1 };
struct s ncs; // the object ncs is modifiable
typedef int A[2][3];
const A a = {{4, 5, 6}, {7, 8, 9}}; // array of array of
                                    // const int
int *pi;
const int *pci;
ncs = cs; // valid
cs = ncs; // violates modifiable lvalue constraint for =
pi = &ncs.mem; // valid
pi = &cs.mem; // violates type constraints for =
pci = &cs.mem; // valid
pi =a[0]; // invalid: a[0] has type 'const int *'"
```


### 6.5.3.1 Formal definition of restrict

1 Let $\mathbf{D}$ be a declaration of an ordinary identifier that provides a means of designating an object $\mathbf{P}$ as a restrict-qualified pointer.

If $\mathbf{D}$ appears inside a block and does not have storage class extern, let $\mathbf{B}$ denote the block. If $\mathbf{D}$ appears in the list of parameter declarations of a function definition, let $\mathbf{B}$ denote the associated block. Otherwise, let B denote the block of main (or the block of whatever function is called at program startup in a freestanding environment).

In what follows, a pointer expression $\mathbf{E}$ is said to be based on object $\mathbf{P}$ if (at some sequence point in the execution of $\mathbf{B}$ prior to the evaluation of $\mathbf{E}$ ) modifying $\mathbf{P}$ to point to a copy of the array object into which it formerly pointed would change the value of $\mathbf{E}$. (In other words, $\mathbf{E}$ depends on the value of $\mathbf{P}$ itself rather than on the value of an object referenced indirectly through $\mathbf{P}$. For example, if identifier $\mathbf{p}$ has type (int **restrict), then the pointer expressions $\mathbf{p}$ and $\mathrm{p}+1$ are based on the restricted pointer object designated by p , but the pointer expressions *p and $\mathrm{p}[1]$ are not.) Note that "based" is defined only for expressions with pointer types.

4 During each execution of $\mathbf{B}$, let $\mathbf{A}$ be the array object that is determined dynamically by all references through pointer expressions based on $\mathbf{P}$. Then all references to values of $\mathbf{A}$ shall be through pointer expressions based on $\mathbf{P}$. Furthermore, if $\mathbf{P}$ is assigned the value of a pointer expression $\mathbf{E}$ that is based on another restricted pointer object P2, associated with block B2, then either the execution of B2 shall begin before the execution of $\mathbf{B}$, or the execution of $\mathbf{B} 2$ shall end prior to the assignment. If these requirements are not met, then the behavior is undefined.

5 Here an execution of $\mathbf{B}$ means that portion of the execution of the program during which storage is guaranteed to be reserved for an instance of an object that is associated with B and that has automatic storage duration. A reference to a value
means either an access to or a modification of the value. During an execution of $\mathbf{B}$, attention is confined to those references that are actually evaluated. (This excludes references that appear in unevaluated expressions, and also excludes references that are 'available', in the sense of employing visible identifiers, but do not actually appear in the text of $\mathbf{B}$.)

A translator is free to ignore any or all aliasing implications of uses of restrict.

## Examples

1. The file scope declarations
```
int * restrict a;
int * restrict b;
extern int c[];
```

assert that if an object is referenced using the value of one of $\mathbf{a}, \mathbf{b}$, or $\mathbf{c}$, then it is never referenced using the value of either of the other two.
2. The function parameter declarations in the following example

```
void f(int n, int * restrict p, int * restrict q)
{
    while (n-- > 0)
        *p++ = *q++;
}
```

assert that, during each execution of the function, if an object is referenced through one of the pointer parameters, then it is not also referenced through the other.

The benefit of the restrict qualifiers is that they enable a translator to make an effective dependence analysis of function $\mathbf{f}$ without examining any of the calls of $\mathbf{f}$ in the program. The cost is that the programmer must examine all of those calls to ensure that none give undefined behavior. For example, the second call of $\mathbf{f}$ in $g$ has undefined behavior because each of d[1] through d [49] is referenced through both p and q .

```
void g(void)
{
    extern float d[100];
    f(50, d + 50, d); // ok
    f(50, d + 1, d); // undefined behavior
}
```

3. The function parameter declarations
```
void h(int n, int * const restrict p,
    int * const q, int * const r)
{
    int i;
    for (i = 0; i < n; i++)
        p[i] = q[i] + r[i];
}
```

show how const can be used in conjunction with restrict. The const qualifiers imply, without the need to examine the body of $\mathbf{h}$, that $\mathbf{q}$ and $\mathbf{r}$ cannot become based on $\mathbf{p}$. The fact that $\mathbf{p}$ is restrict-qualified therefore implies that an object referenced through $\mathbf{p}$ is never referenced through either of $\mathbf{q}$ or $\mathbf{r}$. This is the precise assertion required to optimize the loop. Note that a call of the form $h(100, a, b, b)$ has defined behavior, which would not be true if all three of $\mathbf{p}, \mathbf{q}$, and $\mathbf{r}$ were restrict-qualified.
4. The rule limiting assignments between restricted pointers does not distinguish between a function call and an equivalent nested block. With one exception, only 'outer-to-inner', assignments between restricted pointers declared in nested blocks have defined behavior.

```
{
    int * restrict p1;
int * restrict q1;
p1 = q1; // undefined behavior
{
    int * restrict p2 = p1; // ok
    int * restrict q2 = q1; // ok
    p1 = q2; // undefined behavior
    p2 = q2; // undefined behavior
    }
}
```

The exception allows the value of a restricted pointer to be carried out of the block in which it (or, more precisely, the ordinary identifier used to designate it) is declared when that block finishes execution. For example, this permits new_vector to return a vector.

```
typedef struct { int n; float * restrict v; } vector;
vector new_vector(int n)
{
    vector t;
    t.n = n;
    t.v = malloc(n * sizeof (float));
    return t;
}
```


### 6.5.4 Function specifiers

Syntax
function-specifier:
inline

## Constraints

2 Function specifiers shall be used only in the declaration of an identifier for a function.
3 An inline definition (see below) of a function with external linkage shall not contain a definition of an object with static storage duration that can be modified, and shall not contain a reference to an identifier with internal linkage.

4 The inline function specifier shall not appear in a declaration of main.

## Semantics

5 A function declared with an inline function specifier is an inline function. Making a function an inline function suggests that calls to the function be as fast as possible by using, for example, an alternative to the usual function call mechanism known as "inline substitution". ${ }^{101}$ The extent to which such suggestions are effective is implementation-defined. ${ }^{102}$

6 Any function with internal linkage can be an inline function. For a function with external linkage, the following restrictions apply. If a function is declared with an inline function specifier, then it shall also be defined in the same translation unit.

[^43]If all of the file scope declarations for a function in a translation unit include the inline function specifier without extern, then the definition in that translation unit is an inline definition. An inline definition does not provide an external definition for the function, and does not forbid an external definition in another translation unit. An inline definition provides an alternative to an external definition, which a translator may use to implement any call to the function in the same translation unit. It is unspecified whether a call to the function uses the inline definition or the external definition.

## Examples

7 The declaration of an inline function with external linkage can result in either an external definition, or a definition available for use only within the translation unit. A file scope declaration with extern creates an external definition. The following example shows an entire translation unit.

```
inline double fahr(double t)
{
    return (9.0 * t) / 5.0 + 32.0;
}
inline double cels(double t)
{
    return (5.0 * (t - 32.0)) / 9.0;
}
/* Creates an external definition. */
extern double fahr(double);
double convert(int is_fahr, double temp)
{
    /* A translator may perform inline substitutions. */
    return is_fahr ? cels(temp) : fahr(temp);
}
```

Note that the definition of fahr is an external definition because fahr is also declared with extern, but the definition of cels is an inline definition. Because there is a call to cels, an external definition of cels in another translation unit is still required by 6.7 .

### 6.5.5 Declarators

## Syntax

```
declarator:
        pointer opt direct-declarator
direct-declarator:
    identifier
    ( declarator )
    direct-declarator [ assignment-expression opt ]
    direct-declarator [ * ]
    direct-declarator ( parameter-type-list )
    direct-declarator ( identifier-list }\mp@subsup{}{opt}{}
pointer:
    * type-qualifier-list
    * type-qualifier-list opt pointer
type-qualifier-list:
    type-qualifier
    type-qualifier-list type-qualifier
parameter-type-list:
    parameter-list
    parameter-list , ...
parameter-list:
    parameter-declaration
    parameter-list , parameter-declaration
parameter-declaration:
    declaration-specifiers declarator
    declaration-specifiers abstract-declarator opt
identifier-list:
    identifier
    identifier-list , identifier
```


## Semantics

2 Each declarator declares one identifier, and asserts that when an operand of the same form as the declarator appears in an expression, it designates a function or object with the scope, storage duration, and type indicated by the declaration specifiers.

3 A full declarator is a declarator that is not part of another declarator. The end of a full declarator is a sequence point. If the nested sequence of declarators in a full declarator contains a variable length array type, the type specified by the full 118

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declarator is said to be variably modified.
4 In the following subclauses, consider a declaration

## T D1

where $\mathbf{T}$ contains the declaration specifiers that specify a type $T$ (such as int) and D1 is a declarator that contains an identifier ident. The type specified for the identifier ident in the various forms of declarator is described inductively using this notation.
5 If, in the declaration " $T$ D1', $D 1$ has the form
identifier
then the type specified for ident is $T$.
6 If, in the declaration " $T$ D1'", $D 1$ has the form
( D )
then ident has the type specified by the declaration ' $\boldsymbol{T} \mathbf{D}$ '. Thus, a declarator in parentheses is identical to the unparenthesized declarator, but the binding of complicated declarators may be altered by parentheses.

## Implementation limits

7 The implementation shall allow the specification of types that have at least 12 pointer, array, and function declarators (in any valid combinations) modifying an arithmetic, structure, union, or incomplete type, either directly or via one or more typedefs.

Forward references: array declarators (6.5.5.2), type definitions (6.5.7).

### 6.5.5.1 Pointer declarators

## Semantics

1 If, in the declaration ' $\mathbf{T} \mathbf{D 1}$ ', $\mathbf{D} 1$ has the form

* type-qualifier-list ${ }_{\text {opt }} \mathrm{D}$
and the type specified for ident in the declaration 'T D' is "derived-declarator-typelist $T$ ', then the type specified for ident is "derived-declarator-type-list type-qualifier-list pointer to $T^{\prime \prime}$. For each type qualifier in the list, ident is a so-qualified pointer.

2 For two pointer types to be compatible, both shall be identically qualified and both shall be pointers to compatible types.

## Examples

3 The following pair of declarations demonstrates the difference between a "variable pointer to a constant value" and a "constant pointer to a variable value".

```
const int *ptr_to_constant;
int *const constant_ptr;
```

The contents of an object pointed to by ptr_to_constant shall not be modified through that pointer, but ptr_to_constant itself may be changed to point to another object. Similarly, the contents of the int pointed to by constant_ptr may be modified, but constant_ptr itself shall always point to the same location.

4 The declaration of the constant pointer constant_ptr may be clarified by including a definition for the type 'pointer to int'".

```
typedef int *int_ptr;
const int_ptr constant_ptr;
```

declares constant_ptr as an object that has type 'const-qualified pointer to int'".

### 6.5.5.2 Array declarators

## Constraints

1 The [ and ] may delimit an expression or *. If [ and ] delimit an expression (which specifies the size of an array), it shall have an integer type. If the expression is a constant expression then it shall have a value greater than zero. The element type shall not be an incomplete or function type.

2 Only ordinary identifiers (as defined in 6.1.2.3) with block scope or function prototype scope and without linkage can have a variably modified type. If an identifier is declared to be an object with static storage duration, it shall not have a variable length array type.

## Semantics

3 If, in the declaration ' $T$ D1'", D1 has the form
D[assignment-expression ${ }_{o p t}$ ]
or

```
D [*]
```

and the type specified for ident in the declaration "T $\mathbf{D}$ ' is "derived-declarator-typelist $T$ ', then the type specified for ident is "derived-declarator-type-list array of $T,{ }^{\prime} .{ }^{103}$ If the size is not present, the array type is an incomplete type. If * is used
instead of a size expression, the array type is a variable length array type of unspecified size, which can only be used in declarations with function prototype scope. If the size expression is an integer constant expression and the element type has a known constant size, the array type is not a variable length array type. Otherwise, the array type is a variable length array type. If the size expression is not a constant expression, and it is evaluated at program execution time, it shall evaluate to a value greater than zero. It is unspecified whether side effects are produced when the size expression is evaluated. The size of each instance of a variable length array type does not change during its lifetime.

For two array types to be compatible, both shall have compatible element types, and if both size specifiers are present, and are integer constant expressions, then both size specifiers shall have the same constant value. If the two array types are used in a context which requires them to be compatible, it is undefined behavior if the two size specifiers evaluate to unequal values.

## Examples

1. float fa[11], *afp[17];
declares an array of float numbers and an array of pointers to float numbers.
2. Note the distinction between the declarations
```
extern int *x;
extern int y[];
```

The first declares $\mathbf{x}$ to be a pointer to int; the second declares $\mathbf{y}$ to be an array of int of unspecified size (an incomplete type), the storage for which is defined elsewhere.
3. The following declarations demonstrate the compatibility rules for variably modified types.

[^44]```
extern int n;
extern int m;
void fcompat (void)
{
    int a[n][6][m];
    int (*p)[4][n+1];
    int c[n][n][6][m];
    int (*r)[n][n][n+1];
    p =a; // Error - not compatible because 4 != 6.
    r = c; // Compatible, but defined behavior
        // only if n == 6 and m == n+1.
}
```

4. All declarations of variably modified (VM) types must be declared at either block scope or function prototype scope. Array objects declared with the static or extern storage class specifier cannot have a variable length array (VLA) type. However, an object declared with the static storage class specifier can have a VM type (that is, a pointer to a VLA type). Finally, all identifiers declared with a VM type must be ordinary identifiers, and can not, therefore, be members of structures or unions.
```
extern int n;
int A[n]; // Error - file scope VLA.
extern int (*p2)[n]; // Error-file scope VM.
int B[100]; // OK - file scope but not VM.
void fvla(int n, int C[m][m]) // OK - VLA with prototype scope.
{
typedef int VLA[m][m] // OK - block scope typedef VLA.
struct tag {
            int (*y)[n]; // Error-y not ordinary identifier.
            int z[n]; // Error-z not ordinary identifier.
};
int D[m]; // OK - auto VLA.
static int E[m]; // Error-static block scope VLA.
extern int F[m]; // Error-F has linkage and is VLA.
int (*s)[m]; // OK -auto pointer to VLA.
```



```
// a pointer to VLA.
static int (*q)[m] = &B; // OK-q is a static block
    // pointer to VLA.
}
```

Forward references: function definitions (6.7.1), initialization (6.5.8).

### 6.5.5.3 Function declarators (including prototypes)

## Constraints

1 A function declarator shall not specify a return type that is a function type or an array type.

2 The only storage-class specifier that shall occur in a parameter declaration is register.

3 An identifier list in a function declarator that is not part of a function definition shall be empty.

4 After all rewrites, the parameters in a parameter-type-list that is part of a function definition shall not have incomplete type. ${ }^{104}$

[^45]
## Semantics

5 If, in the declaration ' $T$ D1', D1 has the form
D (parameter-type-list)
or

## D (identifier-list ${ }_{\text {opt }}$ )

and the type specified for ident in the declaration ' $\mathrm{T} \mathbf{D}$ ' is "derived-declarator-typelist $T$ ', then the type specified for ident is "derived-declarator-type-list function returning $T^{\prime}$.

6 A parameter type list specifies the types of, and may declare identifiers for, the parameters of the function. A declared parameter that is a member of a parameter type list that is not part of a function definition, may use the [*] notation in its sequence of declarator specifiers to specify a variable length array type. If the list terminates with an ellipsis (, . . .), no information about the number or types of the parameters after the comma is supplied. ${ }^{105}$ The special case of an unnamed parameter of type void as the only item in the list specifies that the function has no parameters.

7 If, in a parameter declaration, an identifier can be treated as a typedef name or as a parameter name, it shall be taken as a typedef name.

8 If the function declarator is not part of a function definition, the parameters may have incomplete type.

9 The storage-class specifier in the declaration specifiers for a parameter declaration, if present, is ignored unless the declared parameter is one of the members of the parameter type list for a function definition.

10 An identifier list declares only the identifiers of the parameters of the function. An empty list in a function declarator that is part of a function definition specifies that the function has no parameters. The empty list in a function declarator that is not part of a function definition specifies that no information about the number or types of the parameters is supplied. ${ }^{106}$

11 For two function types to be compatible, both shall specify compatible return types. ${ }^{107}$ Moreover, the parameter type lists, if both are present, shall agree in the number of parameters and in use of the ellipsis terminator; corresponding parameters

[^46]shall have compatible types. If one type has a parameter type list and the other type is specified by a function declarator that is not part of a function definition and that contains an empty identifier list, the parameter list shall not have an ellipsis terminator and the type of each parameter shall be compatible with the type that results from the application of the default argument promotions. If one type has a parameter type list and the other type is specified by a function definition that contains a (possibly empty) identifier list, both shall agree in the number of parameters, and the type of each prototype parameter shall be compatible with the type that results from the application of the default argument promotions to the type of the corresponding identifier. (In the determination of type compatibility and of a composite type, each parameter declared with function or array type is taken as having the type that results from conversion to a pointer type, as in 6.7.1, and each parameter declared with qualified type is taken as having the unqualified version of its declared type.)

## Examples

1. The declaration
```
int f(void), *fip(), (*pfi)();
```

declares a function $\mathbf{f}$ with no parameters returning an int, a function $\mathbf{f i p}$ with no parameter specification returning a pointer to an int, and a pointer pfi to a function with no parameter specification returning an int. It is especially useful to compare the last two. The binding of *fip() is *(fip()), so that the declaration suggests, and the same construction in an expression requires, the calling of a function fip, and then using indirection through the pointer result to yield an int. In the declarator (*pfi) (), the extra parentheses are necessary to indicate that indirection through a pointer to a function yields a function designator, which is then used to call the function; it returns an int.

If the declaration occurs outside of any function, the identifiers have file scope and external linkage. If the declaration occurs inside a function, the identifiers of the functions $\mathbf{f}$ and fip have block scope and either internal or external linkage (depending on what file scope declarations for these identifiers are visible), and the identifier of the pointer pfi has block scope and no linkage.
2. The declaration

```
int (*apfi[3])(int *x, int *y);
```

declares an array apfi of three pointers to functions returning int. Each of these functions has two parameters that are pointers to int. The identifiers $\mathbf{x}$ and $\mathbf{y}$ are declared for descriptive purposes only and go out of scope at the end of the declaration of apfi.
3. The declaration

```
int (*fpfi(int (*)(long), int))(int, ...);
```

declares a function fpfi that returns a pointer to a function returning an int. The function fpfi has two parameters: a pointer to a function returning an int (with one parameter of type long), and an int. The pointer returned by fpfi points to a function that has one int parameter and accepts zero or more additional arguments of any type.
4. The following prototype has a variably modified parameter.

```
void addscalar(int n, int m,
    double a[n][n*m+300], double x);
int main()
{
    double b[4][308];
    addscalar(4, 2, b, 2.17);
    return 0;
}
void addscalar(int n, int m,
    double a[n][n*m+300], double x)
{
    for (int i = 0; i < n; i++)
            for (int j = 0, k = n*m+300; j < k; j++)
                // a is a pointer to a VLA
                // with n*m+300 elements
                        a[i][j] += x;
}
```

5. The following are all compatible function prototype declarators.
```
double maximum(int n, int m, double a[n][m]);
double maximum(int n, int m, double a[*][*]);
double maximum(int n, int m, double a[ ][*]);
double maximum(int n, int m, double a[ ][m]);
```

Forward references: function definitions (6.7.1), type names (6.5.6).

### 6.5.6 Type names

## Syntax

1
type-name:
specifier-qualifier-list abstract-declarator ${ }_{\text {opt }}$
abstract-declarator:
pointer
pointer ${ }_{\text {opt }}$ direct-abstract-declarator
direct-abstract-declarator:


## Semantics

2 In several contexts, it is desired to specify a type. This is accomplished using a type name, which is syntactically a declaration for a function or an object of that type that omits the identifier. ${ }^{108}$

## Examples

3 The constructions
(a) int
(b) int *
(c) int *[3]
(d) int (*) [3]
(e) int * ()
(f) int (*) (void)
(g) int (*const []) (unsigned int, ...)
name respectively the types (a) int, (b) pointer to int, (c) array of three pointers to int, (d) pointer to an array of three ints, (e) function with no parameter specification returning a pointer to int, (f) pointer to function with no parameters returning an int, and (g) array of an unspecified number of constant pointers to functions, each with one parameter that has type unsigned int and an unspecified number of other parameters, returning an int.

[^47]
### 6.5.7 Type definitions

## Syntax

## typedef-name:

identifier

## Constraints

If a typedef name specifies a variably modified type then it shall have block scope.

## Semantics

3 In a declaration whose storage-class specifier is typedef, each declarator defines an identifier to be a typedef name that specifies the type specified for the identifier in the way described in 6.5.5. Any array size expressions associated with variable length array declarators shall be evaluated with the typedef name at the beginning of its scope upon each normal entry to the block. A typedef declaration does not introduce a new type, only a synonym for the type so specified. That is, in the following declarations:

```
typedef T type_ident;
type_ident D;
```

type_ident is defined as a typedef name with the type specified by the declaration specifiers in $\mathbf{T}$ (known as $T$ ), and the identifier in $\mathbf{D}$ has the type "derived-declarator-type-list $T$ ", where the derived-declarator-type-list is specified by the declarators of $\mathbf{D}$. A typedef name shares the same name space as other identifiers declared in ordinary declarators. If the identifier is redeclared in an inner scope or is declared as a member of a structure or union in the same or an inner scope, the type specifiers shall not be omitted in the inner declaration.

## Examples

1. After
```
typedef int MILES, KLICKSP();
typedef struct { double hi, lo; } range;
```

the constructions

```
MILES distance;
extern KLICKSP *metricp;
range x;
range z, *zp;
```

are all valid declarations. The type of distance is int, that of metricp is "pointer to function with no parameter specification returning int', and that of $\mathbf{x}$ and $\mathbf{z}$ is the specified structure; $\mathbf{z p}$ is a pointer to such a structure. The
object distance has a type compatible with any other int object.
2. After the declarations

```
typedef struct s1 { int x; } t1, *tp1;
typedef struct s2 { int x; } t2, *tp2;
```

type t 1 and the type pointed to by tp1 are compatible. Type t 1 is also compatible with type struct $\mathbf{s 1}$, but not compatible with the types struct $\mathbf{s} 2, \mathrm{t} 2$, the type pointed to by tp 2 , and int.
3. The following obscure constructions

```
typedef signed int t;
typedef int plain;
struct tag {
        unsigned t:4;
        const t:5;
        plain r:5;
};
```

declare a typedef name $t$ with type signed int, a typedef name plain with type int, and a structure with three bit-field members, one named $t$ that contains values in the range [0, 15], an unnamed const-qualified bit-field which (if it could be accessed) would contain values in at least the range $[-15,+15]$, and one named $\boldsymbol{r}$ that contains values in the range $[0,31]$ or values in at least the range $[-15,+15]$. (The choice of range is implementation-defined.) The first two bit-field declarations differ in that unsigned is a type specifier (which forces $t$ to be the name of a structure member), while const is a type qualifier (which modifies $t$ which is still visible as a typedef name). If these declarations are followed in an inner scope by

```
t f(t (t));
long t;
```

then a function $\mathbf{f}$ is declared with type "function returning signed int with one unnamed parameter with type pointer to function returning signed int with one unnamed parameter with type signed int'", and an identifier $t$ with type long.
4. On the other hand, typedef names can be used to improve code readability. All three of the following declarations of the signal function specify exactly the same type, the first without making use of any typedef names.

```
typedef void fv(int), (*pfv) (int);
void (*signal(int, void (*) (int))) (int);
fv *signal(int, fv *);
pfv signal(int, pfv);
```

5. The following is a block scope declaration of a typedef name $\mathbf{A}$ with a variable length array type.
```
void tdef(int n)
{
    typedef int A[n];
    A a;
    A *p;
    p = &a;
}
```

6. The size expression that is part of the variable length array type named by typedef name B is evaluated each time function copyt is entered. However, the size of the variable length array type does not change if the value of $\mathbf{n}$ is subsequently changed.
```
void copyt(int n)
{
    typedef int B[n]; // B is n ints, n evaluated now.
    n += 1;
    {
        B a; // a is n ints, n without += 1.
        int b[n]; // a and b}\mathrm{ are different sizes
        for (i = 1; i < n; i++)
        a[i-1] = b[i];
    }
}
```

Forward references: the signal function (7.11.1.1).

### 6.5.8 Initialization

1
Syntax
initializer:
assignment-expression
\{ initializer-list \}
\{ initializer-list , \}
initializer-list:
designation $_{\text {opt }}$ initializer
initializer-list, designation ${ }_{\text {opt }}$ initializer
designation:
designator-list $=$
designator-list:
designator
designator-list designator
designator:
[ constant-expression ]

- identifier


## Constraints

2 No initializer shall attempt to provide a value for an object not contained within the entity being initialized.

3 The type of the entity to be initialized shall be an array of unknown size or an object type that is not a variable length array type.

4 All the expressions in an initializer for an object that has static storage duration shall be constant expressions or string literals.

5 If the declaration of an identifier has block scope, and the identifier has external or internal linkage, the declaration shall have no initializer for the identifier.

If a designator has the form
[ constant-expression ]
then the current object (defined below) shall have array type and the expression shall be an integer constant expression. If the array is of unknown size, any nonnegative value is valid.

7 If a designator has the form
. identifier

Language
then the current object (defined below) shall have structure or union type and the identifier shall be a member of that type.

## Semantics

8 An initializer specifies the initial value stored in an object.
9 Except where explicitly stated otherwise, for the purposes of this subclause unnamed members of objects of structure and union type do not participate in initialization. Unnamed members of structure objects have indeterminate value even after initialization. A union object containing only unnamed members has indeterminate value even after initialization.

10 If an object that has automatic storage duration is not initialized explicitly, its value is indeterminate. If an object that has static storage duration is not initialized explicitly, then:

- if it has pointer type, it is initialized to a null pointer;
- if it has arithmetic type, it is initialized to zero;
- if it is an aggregate, every member is initialized (recursively) according to these rules;
- if it is a union, the first named member is initialized (recursively) according to these rules.

11 The initializer for a scalar shall be a single expression, optionally enclosed in braces. The initial value of the object, including unnamed members, is that of the expression; the same type constraints and conversions as for simple assignment apply, taking the type of the scalar to be the unqualified version of its declared type.

12 Each brace-enclosed initializer list has an associated current object. When no designations are present, subobjects of the current object are initialized in order according to the type of the current object: array elements in increasing subscript order, structure members in declaration order, and the first named member of a union. ${ }^{109}$ In contrast, a designation causes the following initializer to begin initialization of the subobject described by the designator. Initialization then continues forward in order, beginning with the next subobject after that described by the designator. ${ }^{110}$

[^48]13 Each designator list begins its description with the current object associated with the closest surrounding brace pair. Each item in the designator list (in order) specifies a particular member of its current object and changes the current object for the next designator (if any) to be that member. ${ }^{111}$ The current object that results at the end of the designator list is the subobject to be initialized by the following initializer.

14 The initialization shall occur in initializer list order, each initializer provided for a particular subobject overriding any previously listed initializer for the same subobject; all subobjects that are not initialized explicitly shall be initialized implicitly the same as objects that have static storage duration.

15 The initializer for a structure or union object that has automatic storage duration either shall be an initializer list as described below, or shall be a single expression that has compatible structure or union type. In the latter case, the initial value of the object is that of the expression.

16 The rest of this subclause deals with initializers for objects that have aggregate or union type.

17 An array of character type may be initialized by a character string literal, optionally enclosed in braces. Successive characters of the character string literal (including the terminating null character if there is room or if the array is of unknown size) initialize the elements of the array.

18 An array with element type compatible with wchar_t may be initialized by a wide string literal, optionally enclosed in braces. Successive codes of the wide string literal (including the terminating zero-valued code if there is room or if the array is of unknown size) initialize the elements of the array.

19 Otherwise, the initializer for an object that has aggregate type shall be a braceenclosed list of initializers for the named members of the aggregate, written in increasing subscript or member order; and the initializer for an object that has union type shall be a brace-enclosed initializer for the first named member of the union.

20 If the aggregate contains members that are aggregates or unions, or if the first member of a union is an aggregate or union, the rules apply recursively to the subaggregates or contained unions. If the initializer of a subaggregate or contained union begins with a left brace, the initializers enclosed by that brace and its matching right brace initialize the members of the subaggregate or the first member of the contained union. Otherwise, only enough initializers from the list are taken to account for the members

[^49]of the subaggregate or the first member of the contained union; any remaining initializers are left to initialize the next member of the aggregate of which the current subaggregate or contained union is a part.

21 If there are fewer initializers in a brace-enclosed list than there are members of an aggregate, or fewer characters in a string literal used to initialize an array of known size than there are elements in the array, the remainder of the aggregate shall be initialized implicitly the same as objects that have static storage duration.

22 If an array of unknown size is initialized, its size is determined by the largest indexed element with an explicit initializer. At the end of its initializer list, the array no longer has incomplete type.

23 The order in which any side effects occur among the initialization list expressions is unspecified. ${ }^{112}$

## Examples

24 1. The declaration

```
int x[] = { 1, 3, 5 };
```

defines and initializes $\mathbf{x}$ as a one-dimensional array object that has three elements, as no size was specified and there are three initializers.
2. The declaration

```
int y[4][3] = {
    { 1, 3, 5 },
    {2, 4, 6 },
    { 3, 5, 7 },
};
```

is a definition with a fully bracketed initialization: 1,3 , and 5 initialize the first row of $y$ (the array object $y[0]$ ), namely $y[0][0], y[0][1]$, and $y[0][2]$. Likewise the next two lines initialize $y[1]$ and $y[2]$. The initializer ends early, so $\mathbf{y}[3]$ is initialized with zeros. Precisely the same effect could have been achieved by

[^50]```
int y[4][3] = {
    1, 3, 5, 2, 4, 6, 3, 5, 7
};
```

The initializer for $\mathbf{y}[0]$ does not begin with a left brace, so three items from the list are used. Likewise the next three are taken successively for $\mathrm{y}[1]$ and $y[2]$.
3. The declaration

```
int z[4][3] = {
    {1 },{2},{ 3 },{4}
};
```

initializes the first column of $\mathbf{z}$ as specified and initializes the rest with zeros.
4. The declaration

```
struct { int a[3], b; } w[] = { { 1 }, 2 };
```

is a definition with an inconsistently bracketed initialization. It defines an array with two element structures: w[0].a[0] is 1 and w[1].a[0] is 2 ; all the other elements are zero.
5. The declaration

```
short q[4][3][2] = {
    { 1 },
    { 2, 3 },
    {4, 5, 6 }
};
```

contains an incompletely but consistently bracketed initialization. It defines a three-dimensional array object: $\mathrm{q}[0][0][0]$ is 1 , $\mathrm{q}[1][0][0]$ is 2 , $\mathrm{q}[1][0][1]$ is 3 , and 4,5 , and 6 initialize $q[2][0][0], q[2][0][1]$, and $\mathrm{q}[2][1][0]$, respectively; all the rest are zero. The initializer for $\mathrm{q}[0][0]$ does not begin with a left brace, so up to six items from the current list may be used. There is only one, so the values for the remaining five elements are initialized with zero. Likewise, the initializers for $q[1][0]$ and $q[2][0]$ do not begin with a left brace, so each uses up to six items, initializing their respective two-dimensional subaggregates. If there had been more than six items in any of the lists, a diagnostic message would have been issued. The same initialization result could have been achieved by:

```
short \(q[4][3][2]=\{\)
    \(1,0,0,0,0,0\),
    \(2,3,0,0,0,0\),
    4, 5, 6
\};
```

or by:

```
short q[4][3][2] = {
    {
        { 1 },
        },
        {
            {2, 3 },
        },
        {
            {4,5 },
            { 6 },
        }
};
```

in a fully bracketed form.
Note that the fully bracketed and minimally bracketed forms of initialization are, in general, less likely to cause confusion.
6. One form of initialization that completes array types involves typedef names. Given the declaration
typedef int A[]; // OK-declared with block scope
the declaration

```
A a = { 1, 2 }, b = { 3, 4, 5 };
```

is identical to

```
int a[] = { 1, 2 }, b[] = { 3, 4, 5 };
```

due to the rules for incomplete types.
7. The declaration

```
char s[] = "abc", t[3] = "abc";
```

defines 'plain'" char array objects $\mathbf{s}$ and $t$ whose elements are initialized with character string literals. This declaration is identical to

```
char s[] = { 'a', 'b', 'c', '\0' },
    t[] = { 'a', 'b', 'c' };
```

The contents of the arrays are modifiable. On the other hand, the declaration

```
char *p = "abc";
```

defines $\mathbf{p}$ with type "pointer to char' that is initialized to point to an object with type 'array of char', with length 4 whose elements are initialized with a character string literal. If an attempt is made to use $\mathbf{p}$ to modify the contents of the array, the behavior is undefined.
8. Arrays can be initialized to correspond to the elements of an enumeration by using designators:

```
enum { member_one, member_two };
const char *nm[] = {
    [member_two] = "member two",
    [member_one] = "member one",
};
```

9. Structure members can be initialized to nonzero values without depending on their order:
```
div_t answer = { .quot = 2, .rem = -1 };
```

10. Designators can be used to provide explicit initialization when unadorned initializer lists might be misunderstood:
```
struct { int a[3], b; } w[] =
    { [0].a = {1}, [1].a[0] = 2 };
```

11. Space can be "allocated" from both ends of an array by using a single designator:
```
int a[MAX] = {
    1, 3, 5, 7, 9, [MAX-5] = 8, 6, 4, 2, 0
};
```

In the above, if MAX is greater than ten, there will be some zero-valued elements in the middle; if it is less than ten, some of the values provided by the first five initializers will be overridden by the second five.
12. Any member of a union can be initialized:

$$
\text { union }\{/ * \ldots * /\} u=\{\text {.any_member }=42\} ;
$$

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Forward references: common definitions <stddef. h> (7.1.6).

### 6.6 Statements

Syntax

## statement:

labeled-statement
compound-statement
expression-statement
selection-statement
iteration-statement
jump-statement

## Semantics

2 A statement specifies an action to be performed. Except as indicated, statements are executed in sequence.

3 A full expression is an expression that is not part of another expression. Each of the following is a full expression: an initializer; the expression in an expression statement; the controlling expression of a selection statement (if or switch); the controlling expression of a while or do statement; each of the (optional) expressions of a for statement; the (optional) expression in a return statement. The end of a full expression is a sequence point.

Forward references: expression and null statements (6.6.3), selection statements (6.6.4), iteration statements (6.6.5), the return statement (6.6.6.4).

### 6.6.1 Labeled statements

Syntax
labeled-statement:
identifier : statement
case constant-expression : statement
default : statement

## Constraints

2 A case or default label shall appear only in a switch statement. Further constraints on such labels are discussed under the switch statement.

## Semantics

3 Any statement may be preceded by a prefix that declares an identifier as a label name. Labels in themselves do not alter the flow of control, which continues unimpeded across them.

Forward references: the goto statement (6.6.6.1), the switch statement (6.6.4.2).

### 6.6.2 Compound statement, or block

Syntax
compound-statement:
\{ block-item-list $\left.{ }_{o p t}\right\}$
block-item-list:
block-item
block-item-list block-item
block-item:
declaration
statement

## Semantics

2 A compound statement (also called a block) allows a set of statements to be grouped into one syntactic unit, which may have its own set of declarations and initializations (as discussed in 6.1.2.4). The initializers of objects that have automatic storage duration, and the variable length array declarators of ordinary identifiers with block scope are evaluated and the values are stored in the objects (including storing an indeterminate value in objects without an initializer) each time that the declaration is reached in the order of execution, as if it were a statement, and within each declaration in the order that declarators appear.

### 6.6.3 Expression and null statements

Syntax
expression-statement:

$$
\text { expression }_{\text {opt }} \text {; }
$$

## Semantics

2 The expression in an expression statement is evaluated as a void expression for its side effects. ${ }^{113}$

3 A null statement (consisting of just a semicolon) performs no operations.

[^51]
## Examples

1. If a function call is evaluated as an expression statement for its side effects only, the discarding of its value may be made explicit by converting the expression to a void expression by means of a cast:
```
int p(int);
/* ... */
(void)p(0);
```

2. In the program fragment
```
char *s;
/* ... */
while (*s++ != '\0')
    ;
```

a null statement is used to supply an empty loop body to the iteration statement.
3. A null statement may also be used to carry a label just before the closing \} of a compound statement.

```
while (loop1) {
        /* ... */
        while (loop2) {
            /* ... */
            if (want_out)
                    goto end_loop1;
            /* ... */
        }
        /* ... */
end_loop1: ;
}
```

Forward references: iteration statements (6.6.5).

### 6.6.4 Selection statements

Syntax
selection-statement:
if ( expression ) statement
if ( expression ) statement else statement
switch ( expression ) statement

## Semantics

2 A selection statement selects among a set of statements depending on the value of a controlling expression.

### 6.6.4.1 The if statement

## Constraints

1 The controlling expression of an if statement shall have scalar type.

## Semantics

2 In both forms, the first substatement is executed if the expression compares unequal to 0 . In the else form, the second substatement is executed if the expression compares equal to 0 . If the first substatement is reached via a label, the second substatement is not executed.

3 An else is associated with the lexically nearest preceeding if that is allowed by the grammar.

### 6.6.4.2 The switch statement

## Constraints

1 The controlling expression of a switch statement shall have integer type, and shall not cause a block to be entered by a jump from outside the block to a statement that follows a case or default label in the block (or an enclosed block) if that block contains the declaration of a variably modified object or variably modified typedef name. The expression of each case label shall be an integer constant expression. No two of the case constant expressions in the same switch statement shall have the same value after conversion. There may be at most one default label in a switch statement. (Any enclosed switch statement may have a default label or case constant expressions with values that duplicate case constant expressions in the enclosing switch statement.)

## Semantics

2 A switch statement causes control to jump to, into, or past the statement that is the switch body, depending on the value of a controlling expression, and on the presence of a default label and the values of any case labels on or in the switch body. A case or default label is accessible only within the closest enclosing switch statement.

3 The integer promotions are performed on the controlling expression. The constant expression in each case label is converted to the promoted type of the controlling expression. If a converted value matches that of the promoted controlling expression, control jumps to the statement following the matched case label. Otherwise, if there
is a default label, control jumps to the labeled statement. If no converted case constant expression matches and there is no default label, no part of the switch body is executed.

## Implementation limits

As discussed previously (5.2.4.1), the implementation may limit the number of case values in a switch statement.

## Examples

5 In the artificial program fragment

```
switch (expr)
{
    int i = 4;
    f(i);
case 0:
    i = 17;
    /* falls through into default code */
default:
    printf("%d\n", i);
}
```

the object whose identifier is $\mathbf{i}$ exists with automatic storage duration (within the block) but is never initialized, and thus if the controlling expression has a nonzero value, the call to the printf function will access an indeterminate value. Similarly, the call to the function $f$ cannot be reached.

### 6.6.5 Iteration statements

Syntax
iteration-statement:
while ( expression ) statement
do statement while ( expression ) ;
for ( expression ${ }_{\text {opt }}$; expression ${ }_{\text {opt }}$; expression ${ }_{\text {opt }}$ ) statement for ( declaration ; expression ${ }_{o p t}$; expression ${ }_{o p t}$ ) statement

## Constraints

The controlling expression of an iteration statement shall have scalar type.
3 The declaration part of a for statement shall only declare identifiers for objects having storage class auto or register.

## Semantics

4 An iteration statement causes a statement called the loop body to be executed repeatedly until the controlling expression compares equal to 0 .

### 6.6.5.1 The while statement

1 The evaluation of the controlling expression takes place before each execution of the loop body.

### 6.6.5.2 The do statement

1 The evaluation of the controlling expression takes place after each execution of the loop body.

### 6.6.5.3 The for statement

1 Except for the behavior of a continue statement in the loop body, the statement
for ( clause-1 ; expression-2 ; expression-3 ) statement
and the sequence of statements
\{
clause-1 ;
while ( expression-2 ) \{
statement
expression-3 ;
\}
\}
are equivalent (where clause-l can be an expression or a declaration). ${ }^{114}$
2 Both clause-1 and expression-3 can be omitted. If either or both are an expression, they are evaluated as a void expression. An omitted expression- 2 is replaced by a nonzero constant.

[^52]Forward references: the continue statement (6.6.6.2).

### 6.6.6 Jump statements

Syntax

## jump-statement:

goto identifier ;
continue ;
break ;
return expression $_{\text {opt }}$;

## Semantics

2 A jump statement causes an unconditional jump to another place.

### 6.6.6.1 The goto statement

## Constraints

1 The identifier in a goto statement shall name a label located somewhere in the enclosing function. A goto statement shall not cause a block to be entered by a jump from outside the block to a labeled statement in the block (or an enclosed block) if that block contains the declaration of a variably modified object or variably modified typedef name.

## Semantics

A goto statement causes an unconditional jump to the statement prefixed by the named label in the enclosing function.

## Examples

1. It is sometimes convenient to jump into the middle of a complicated set of statements. The following outline presents one possible approach to a problem based on these three assumptions:
2. The general initialization code accesses objects only visible to the current function.
3. The general initialization code is too large to warrant duplication.
4. The code to determine the next operation must be at the head of the loop. (To allow it to be reached by continue statements, for example.)
```
/* ... */
goto first_time;
for (;;) {
    // determine next operation
    /* ... */
    if (need to reinitialize) {
        // reinitialize-only code
    /* ... */
    first_time:
    // general initialization code
    /* ... */
    continue;
    }
    // handle other operations
    /* ... */
}
```

2. A goto statement is not allowed to jump past any declarations of objects with variably modified types. A jump within the block, however, is permitted.
```
goto lab3; // Error: going INTO scope of VLA.
{
    double a[n];
    a[j] = 4.4;
lab3:
    a[j] = 3.3;
    goto lab 4; // OK, going WITHIN scope of VLA.
    a[j] = 5.5;
lab4:
    a[j] = 6.6;
}
goto lab4; // Error: going INTO scope of VLA.
```

6.6.6.2 The continue statement

## Constraints

1 A continue statement shall appear only in or as a loop body.

## Semantics

2 A continue statement causes a jump to the loop-continuation portion of the smallest enclosing iteration statement; that is, to the end of the loop body. More precisely, in each of the statements
while (/* ... */) {
while (/* ... */) {
/* ... */
/* ... */
continue;
continue;
/* ... */
/* ... */
contin: ;
contin: ;
}
}
$\begin{array}{lc}\text { do }\{ & \text { for }(/ * \ldots \text { */) \{ } \\ \text { /* ... */ } & \text { /* ... */ } \\ \text { continue; } & \text { continue; } \\ \text { /* ... */ } & \text { /*... */ } \\ \text { contin: ; } & \text { contin: ; } \\ \text { \} while (/* ... */); } & \text { \} }\end{array}$
unless the continue statement shown is in an enclosed iteration statement (in which case it is interpreted within that statement), it is equivalent to goto contin; . ${ }^{115}$

### 6.6.6.3 The break statement

## Constraints

1 A break statement shall appear only in or as a switch body or loop body.

## Semantics

2 A break statement terminates execution of the smallest enclosing switch or iteration statement.

### 6.6.6.4 The return statement

## Constraints

1 A return statement with an expression shall not appear in a function whose return type is void. A return statement without an expression shall only appear in a function whose return type is void.

## Semantics

2 A return statement terminates execution of the current function and returns control to its caller. A function may have any number of return statements.

3 If a return statement with an expression is executed, the value of the expression is returned to the caller as the value of the function call expression. If the expression has a type different from the return type of the function in which it appears, the value is converted as if by assignment to an object having the return type of the function. ${ }^{116}$

4 If a return statement without an expression is executed, and the value of the function call is used by the caller, the behavior is undefined.

[^53]Working Draft, 1997-11-21, WG14/N794 J11/97-158

## Examples

In:

```
struct s { double i; } f(void);
union {
        struct {
                int f1;
                struct s f2;
            } u1;
            struct {
                struct s f3;
                int f4;
            } u2;
} g;
struct s f(void)
{
        return g.u1.f2;
}
/* ... */
g.u2.f3 = f();
```

there is no undefined behavior.

### 6.7 External definitions

## Syntax

translation-unit:
external-declaration
translation-unit external-declaration

## external-declaration:

function-definition
declaration

## Constraints

2 The storage-class specifiers auto and register shall not appear in the declaration specifiers in an external declaration.

3 There shall be no more than one external definition for each identifier declared with internal linkage in a translation unit. Moreover, if an identifier declared with internal linkage is used in an expression (other than as a part of the operand of a sizeof operator), there shall be exactly one external definition for the identifier in the translation unit.

## Semantics

4 As discussed in 5.1.1.1, the unit of program text after preprocessing is a translation unit, which consists of a sequence of external declarations. These are described as 'external" because they appear outside any function (and hence have file scope). As discussed in 6.5, a declaration that also causes storage to be reserved for an object or a function named by the identifier is a definition.

5 An external definition is an external declaration that is also a definition of a function or an object. If an identifier declared with external linkage is used in an expression (other than as part of the operand of a sizeof operator), somewhere in the entire program there shall be exactly one external definition for the identifier; otherwise, there shall be no more than one. ${ }^{17}$

[^54]
### 6.7.1 Function definitions

Syntax

## function-definition:

declaration-specifiers declarator declaration-list ${ }_{\text {opt }}$ compound-statement

## Constraints

2 The identifier declared in a function definition (which is the name of the function) shall have a function type, as specified by the declarator portion of the function definition. ${ }^{118}$

3 The return type of a function shall be void or an object type other than array type.
4 The storage-class specifier, if any, in the declaration specifiers shall be either extern or static.

5 If the declarator includes a parameter type list, the declaration of each parameter shall include an identifier (except for the special case of a parameter list consisting of a single parameter of type void, in which there shall not be an identifier). No declaration list shall follow.

6 If the declarator includes an identifier list, each declaration in the declaration list shall have at least one declarator, those declarators shall declare only identifiers from the identifier list, and every identifier in the identifier list shall be declared. An identifier declared as a typedef name shall not be redeclared as a parameter. The declarations in the declaration list shall contain no storage-class specifier other than register and no initializations.

[^55]
## Semantics

7 The declarator in a function definition specifies the name of the function being defined and the identifiers of its parameters. If the declarator includes a parameter type list, the list also specifies the types of all the parameters; such a declarator also serves as a function prototype for later calls to the same function in the same translation unit. If the declarator includes an identifier list, ${ }^{119}$ the types of the parameters shall be declared in a following declaration list.

8 If a function that accepts a variable number of arguments is defined without a parameter type list that ends with the ellipsis notation, the behavior is undefined.
9 Each parameter has automatic storage duration. Its identifier is an lvalue. ${ }^{120}$ The layout of the storage for parameters is unspecified.
10 On entry to the function all size expressions of its variably modified parameters are evaluated, and the value of each argument expression shall be converted to the type of its corresponding parameter, as if by assignment to the parameter. Array expressions and function designators as arguments are converted to pointers before the call. A declaration of a parameter as "array of type" shall be adjusted to "pointer to type," and a declaration of a parameter as "function returning type" shall be adjusted to "pointer to function returning type," as in 6.2.2.1. The resulting parameter type shall be an object type.

11 After all parameters have been assigned, the compound statement that constitutes the body of the function definition is executed.

12 If the \} that terminates a function is reached, and the value of the function call is used by the caller, the behavior is undefined.

## Examples

13 1. In the following:

```
extern int max(int a, int b)
{
    return a > b ? a : b;
}
```

[^56]extern is the storage-class specifier and int is the type specifier; $\max ($ int $a$, int $b$ ) is the function declarator; and

```
    { return a > b ? a : b; }
```

is the function body. The following similar definition uses the identifier-list form for the parameter declarations:

```
extern int max(a, b)
int a, b;
{
    return a > b ? a : b;
}
```

Here int $\mathbf{a}, \mathbf{b}$; is the declaration list for the parameters. The difference between these two definitions is that the first form acts as a prototype declaration that forces conversion of the arguments of subsequent calls to the function, whereas the second form may not.
2. To pass one function to another, one might say

```
int f(void);
/* ... */
g(f);
```

Then the definition of $\mathbf{g}$ might read

```
void g(int (*funcp) (void))
{
    /* ... */ (*funcp) () /* or funcp() ... */
}
```

or, equivalently,

```
void g(int func(void))
{
    /* ... */ func() /* or (*func)() ... */
}
```


### 6.7.2 External object definitions

## Semantics

1 If the declaration of an identifier for an object has file scope and an initializer, the declaration is an external definition for the identifier.

2 A declaration of an identifier for an object that has file scope without an initializer, and without a storage-class specifier or with the storage-class specifier static, constitutes a tentative definition. If a translation unit contains one or more tentative definitions for an identifier, and the translation unit contains no external definition for that identifier, then the behavior is exactly as if the translation unit contains a file scope declaration of that identifier, with the composite type as of the end of the translation unit, with an initializer equal to 0 .

3 If the declaration of an identifier for an object is a tentative definition and has internal linkage, the declared type shall not be an incomplete type.

## Examples

```
int i1 = 1; // definition, external linkage
static int i2 = 2; // definition, internal linkage
extern int i3 = 3; // definition, external linkage
int i4; // tentative definition, external linkage
static int i5; // tentative definition, internal linkage
int i1; // valid tentative definition, refers to previous
int i2; // 6.1.2.2 renders undefined, linkage disagreement
int i3; // valid tentative definition, refers to previous
int i4; // valid tentative definition, refers to previous
int i5; // 6.1.2.2 renders undefined, linkage disagreement
extern int i1; // refers to previous, whose linkage is external
extern int i2; // refers to previous, whose linkage is internal
extern int i3; // refers to previous, whose linkage is external
extern int i4; // refers to previous, whose linkage is external
extern int i5; // refers to previous, whose linkage is internal
```

2. If at the end of the translation unit containing

## int i[];

the array $\mathbf{i}$ still has incomplete type, the array is assumed to have one element. This element is initialized to zero on program startup.

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### 6.8 Preprocessing directives

## Syntax

preprocessing-file:
group $_{\text {opt }}$
group:
group-part
group group-part
group-part:
pp-tokens ${ }_{\text {opt }}$ new-line
if-section
control-line
if-section:
if-group elif-groups ${ }_{\text {opt }}$ else-group $_{\text {opt }}$ endif-line
if-group:
\# if constant-expression new-line group ${ }_{\text {opt }}$
\# ifdef identifier new-line group ${ }_{\text {opt }}$
\# ifndef identifier new-line group opt
elif-groups:
elif-group
elif-groups elif-group
elif-group:
\# elif constant-expression new-line group ${ }_{\text {opt }}$
else-group:
\# else new-line group $_{\text {opt }}$
endif-line:
\# endif new-line
control-line:
 lparen:
the left-parenthesis character without preceding white-space
replacement-list:
pp-tokens ${ }_{\text {opt }}$
pp-tokens:
preprocessing-token
pp-tokens preprocessing-token
new-line:
the new-line character

## Description

2 A preprocessing directive consists of a sequence of preprocessing tokens that begins with a \# preprocessing token that (at the start of translation phase 4) is either the first character in the source file (optionally after white space containing no new-line characters) or that follows white space containing at least one new-line character, and is ended by the next new-line character. ${ }^{121}$ A new-line character ends the preprocessing directive even if it occurs within what would otherwise be an invocation of a function-like macro.

[^57]
## Constraints

3 The only white-space characters that shall appear between preprocessing tokens within a preprocessing directive (from just after the introducing \# preprocessing token through just before the terminating new-line character) are space and horizontal-tab (including spaces that have replaced comments or possibly other white-space characters in translation phase 3).

4 In the definition of an object-like macro, if the first character of a replacement list is not a character required by subclause 5.2 .1 , then there shall be white-space separation between the identifier and the replacement list. ${ }^{122}$

## Semantics

5 The implementation can process and skip sections of source files conditionally, include other source files, and replace macros. These capabilities are called preprocessing, because conceptually they occur before translation of the resulting translation unit.

6 The preprocessing tokens within a preprocessing directive are not subject to macro expansion unless otherwise stated.

## Examples

7 In:

```
#define EMPTY
EMPTY # include <file.h>
```

the sequence of preprocessing tokens on the second line is not a preprocessing directive, because it does not begin with a \# at the start of translation phase 4, even though it will do so after the macro EMPTY has been replaced.

[^58]
### 6.8.1 Conditional inclusion

## Constraints

1 The expression that controls conditional inclusion shall be an integer constant expression except that: it shall not contain a cast; identifiers (including those lexically identical to keywords) are interpreted as described below; ${ }^{123}$ and it may contain unary operator expressions of the form

```
    defined identifier
or
    defined ( identifier )
```

which evaluate to 1 if the identifier is currently defined as a macro name (that is, if it is predefined or if it has been the subject of a \#define preprocessing directive without an intervening \#undef directive with the same subject identifier), 0 if it is not.

## Semantics

2 Preprocessing directives of the forms

$$
\begin{aligned}
& \text { \# if constant-expression new-line group } \\
& \text { \# elift } \\
& \text { constant-expression new-line group opt }
\end{aligned}
$$

check whether the controlling constant expression evaluates to nonzero.
3 Prior to evaluation, macro invocations in the list of preprocessing tokens that will become the controlling constant expression are replaced (except for those macro names modified by the defined unary operator), just as in normal text. If the token defined is generated as a result of this replacement process or use of the defined unary operator does not match one of the two specified forms prior to macro replacement, the behavior is undefined. After all replacements due to macro expansion and the defined unary operator have been performed, all remaining identifiers are replaced with the pp-number 0 , and then each preprocessing token is converted into a token. The resulting tokens compose the controlling constant expression which is evaluated according to the rules of 6.4 , except that all signed integer types and all unsigned integer types act as if they have the same representation as, respectively, the types intmax_t and uintmax_t defined in the header <inttypes.h>. This includes interpreting character constants, which may involve converting escape sequences into execution character set members. Whether the numeric value for these

[^59]character constants matches the value obtained when an identical character constant occurs in an expression (other than within a \#if or \#elif directive) is implementation-defined. ${ }^{124}$ Also, whether a single-character character constant may have a negative value is implementation-defined.

4 Preprocessing directives of the forms

$$
\begin{aligned}
& \text { \# ifdef identifier new-line } \text { group }_{\text {opt }} \\
& \text { \# ifndef identifier new-line group opt }
\end{aligned}
$$

check whether the identifier is or is not currently defined as a macro name. Their conditions are equivalent to \#if defined identifier and \#if !defined identifier respectively.

5 Each directive's condition is checked in order. If it evaluates to false (zero), the group that it controls is skipped: directives are processed only through the name that determines the directive in order to keep track of the level of nested conditionals; the rest of the directives' preprocessing tokens are ignored, as are the other preprocessing tokens in the group. Only the first group whose control condition evaluates to true (nonzero) is processed. If none of the conditions evaluates to true, and there is a \#else directive, the group controlled by the \#else is processed; lacking a \#else directive, all the groups until the \#endif are skipped. ${ }^{125}$

Forward references: macro replacement (6.8.3), source file inclusion (6.8.2), largest integer types (7.4.1.5).
124. Thus, the constant expression in the following \#if directive and if statement is not guaranteed to evaluate to the same value in these two contexts.

```
#if 'z' - 'a' == 25
if ('z' - 'a' == 25)
```

125. As indicated by the syntax, a preprocessing token shall not follow a \#else or \#endif directive before the terminating new-line character. However, comments may appear anywhere in a source file, including within a preprocessing directive.

### 6.8.2 Source file inclusion

## Constraints

1 A \#include directive shall identify a header or source file that can be processed by the implementation.

## Semantics

2 A preprocessing directive of the form
\# include <h-char-sequence> new-line
searches a sequence of implementation-defined places for a header identified uniquely by the specified sequence between the < and > delimiters, and causes the replacement of that directive by the entire contents of the header. How the places are specified or the header identified is implementation-defined.

3 A preprocessing directive of the form
\# include "q-char-sequence" new-line
causes the replacement of that directive by the entire contents of the source file identified by the specified sequence between the " delimiters. The named source file is searched for in an implementation-defined manner. If this search is not supported, or if the search fails, the directive is reprocessed as if it read

## \# include <h-char-sequence> new-line

with the identical contained sequence (including > characters, if any) from the original directive.

A preprocessing directive of the form
\# include pp-tokens new-line
(that does not match one of the two previous forms) is permitted. The preprocessing tokens after include in the directive are processed just as in normal text. (Each identifier currently defined as a macro name is replaced by its replacement list of preprocessing tokens.) The directive resulting after all replacements shall match one of the two previous forms. ${ }^{126}$ The method by which a sequence of preprocessing tokens between a < and a > preprocessing token pair or a pair of " characters is combined into a single header name preprocessing token is implementation-defined.

[^60]5 The implementation shall provide unique mappings for sequences consisting of one or more letters or digits (as defined in 5.2.1) followed by a period (.) and a single letter. The first character shall be a letter. The implementation may ignore the distinctions of alphabetical case and restrict the mapping to eight significant characters before the period.

6 A \#include preprocessing directive may appear in a source file that has been read because of a \#include directive in another file, up to an implementation-defined nesting limit (see 5.2.4.1).

## Examples

1. The most common uses of \#include preprocessing directives are as in the following:
```
#include <stdio.h>
#include "myprog.h"
```

2. This illustrates macro-replaced \#include directives:
```
#if VERSION == 1
    #define INCFILE "vers1.h"
#elif VERSION == 2
    #define INCFILE "vers2.h" // and so on
#else
    #define INCFILE "versN.h"
#endif
#include INCFILE
```

Forward references: macro replacement (6.8.3).

### 6.8.3 Macro replacement

## Constraints

1 Two replacement lists are identical if and only if the preprocessing tokens in both have the same number, ordering, spelling, and white-space separation, where all white-space separations are considered identical.

2 An identifier currently defined as a macro without use of lparen (an object-like macro) shall not be redefined by another \#define preprocessing directive unless the second definition is an object-like macro definition and the two replacement lists are identical.

An identifier currently defined as a macro using lparen (a function-like macro) shall not be redefined by another \#define preprocessing directive unless the second definition is a function-like macro definition that has the same number and spelling of parameters, and the two replacement lists are identical.

4 If the identifier-list in the macro definition does not end with an ellipsis, the number of arguments, including those arguments consisting of no preprocessing tokens, in an invocation of a function-like macro shall agree with the number of parameters in the macro definition. Otherwise, there shall be more arguments in the invocation than there are parameters in the macro definition (excluding the . . .). There shall exist a ) preprocessing token that terminates the invocation.

5 The identifier __VA_ARGS__ shall only occur in the replacement-list of a \#define preprocessing directive using the ellipsis notation in the arguments.

6 A parameter identifier in a function-like macro shall be uniquely declared within its scope.

## Semantics

7 The identifier immediately following the define is called the macro name. There is one name space for macro names. Any white-space characters preceding or following the replacement list of preprocessing tokens are not considered part of the replacement list for either form of macro.

8 If a \# preprocessing token, followed by an identifier, occurs lexically at the point at which a preprocessing directive could begin, the identifier is not subject to macro replacement.
9 A preprocessing directive of the form
\# define identifier replacement-list new-line
defines an object-like macro that causes each subsequent instance of the macro name ${ }^{127}$ to be replaced by the replacement list of preprocessing tokens that constitute the remainder of the directive. The replacement list is then rescanned for more macro names as specified below.

A preprocessing directive of the form
\# define identifier lparen identifier-list ${ }_{\text {opt }}$ ) replacement-list new-line
\# define identifier lparen . . . ) replacement-list new-line
\# define identifier lparen identifier-list , ... ) replacement-list new-line
defines a function-like macro with arguments, similar syntactically to a function call. The parameters are specified by the optional list of identifiers, whose scope extends
127. Since, by macro-replacement time, all character constants and string literals are preprocessing tokens, not sequences possibly containing identifier-like subsequences (see 5.1.1.2, translation phases), they are never scanned for macro names or parameters.
from their declaration in the identifier list until the new-line character that terminates the \#define preprocessing directive. Each subsequent instance of the function-like macro name followed by a ( as the next preprocessing token introduces the sequence of preprocessing tokens that is replaced by the replacement list in the definition (an invocation of the macro). The replaced sequence of preprocessing tokens is terminated by the matching ) preprocessing token, skipping intervening matched pairs of left and right parenthesis preprocessing tokens. Within the sequence of preprocessing tokens making up an invocation of a function-like macro, new-line is considered a normal white-space character.

11 The sequence of preprocessing tokens bounded by the outside-most matching parentheses forms the list of arguments for the function-like macro. The individual arguments within the list are separated by comma preprocessing tokens, but comma preprocessing tokens between matching inner parentheses do not separate arguments. If there are sequences of preprocessing tokens within the list of arguments that would otherwise act as preprocessing directives, the behavior is undefined.

12 If there is a ... in the identifier-list in the macro definition, then the trailing arguments, including any separating comma preprocessing tokens, are merged to form a single item: the variable arguments. The number of arguments so combined is such that, following merger, the number of arguments is one more than the number of parameters in the macro definition (excluding the . . .).

### 6.8.3.1 Argument substitution

1 After the arguments for the invocation of a function-like macro have been identified, argument substitution takes place. A parameter in the replacement list, unless preceded by a \# or \#\# preprocessing token or followed by a \#\# preprocessing token (see below), is replaced by the corresponding argument after all macros contained therein have been expanded. Before being substituted, each argument's preprocessing tokens are completely macro replaced as if they formed the rest of the preprocessing file; no other preprocessing tokens are available.

2 An identifier__VA_ARGS__ that occurs in the replacement list shall be treated as if
it were a parameter, and the variable arguments shall form the preprocessing tokens used to replace it.

### 6.8.3.2 The \# operator

## Constraints

1 Each \# preprocessing token in the replacement list for a function-like macro shall be followed by a parameter as the next preprocessing token in the replacement list.

## Semantics

2 If, in the replacement list, a parameter is immediately preceded by a \# preprocessing token, both are replaced by a single character string literal preprocessing token that contains the spelling of the preprocessing token sequence for the corresponding argument. Each occurrence of white space between the argument's preprocessing tokens becomes a single space character in the character string literal. White space before the first preprocessing token and after the last preprocessing token comprising the argument is deleted. Otherwise, the original spelling of each preprocessing token in the argument is retained in the character string literal, except for special handling for producing the spelling of string literals and character constants: a $\backslash$ character is inserted before each " and \ character of a character constant or string literal (including the delimiting " characters). If the replacement that results is not a valid character string literal, the behavior is undefined. The character string literal corresponding to an empty argument is "". The order of evaluation of \# and \#\# operators is unspecified.

### 6.8.3.3 The \#\# operator

## Constraints

1 A \#\# preprocessing token shall not occur at the beginning or at the end of a replacement list for either form of macro definition.

## Semantics

2 If, in the replacement list, a parameter is immediately preceded or followed by a \#\# preprocessing token, the parameter is replaced by the corresponding argument's preprocessing token sequence; however, if an argument consists of no preprocessing tokens, the parameter is replaced by a placemarker preprocessing token instead.

3 For both object-like and function-like macro invocations, before the replacement list is reexamined for more macro names to replace, each instance of a \#\# preprocessing token in the replacement list (not from an argument) is deleted and the preceding preprocessing token is concatenated with the following preprocessing token (placemarker preprocessing tokens are handled specially: concatenation of two placemarkers results in a single placemarker preprocessing token; concatenation of a placemarker with a non-placemarker preprocessing token results in the non-placemarker preprocessing token). If the result is not a valid preprocessing token, the behavior is undefined. The resulting token is available for further macro replacement. The order of evaluation of \#\# operators is unspecified.

## Examples

```
#define hash_hash # ## #
#define mkstr(a) # a
#define in_between(a) mkstr(a)
#define join(c, d) in_between(c hash_hash d)
char p[] = join(x, y); // equivalent to
    // char p[] = "x ## y";
```

The expansion produces, at various stages:

```
join(x, y)
in_between(x hash_hash y)
in_between(x ## y)
mkstr(x ## y)
"x ## y"
```

In other words, expanding hash_hash produces a new token, consisting of two adjacent sharp signs, but this new token is not the catenation operator.

### 6.8.3.4 Rescanning and further replacement

1 After all parameters in the replacement list have been substituted and \# and \#\# processing has taken place, all placemarker preprocessing tokens are removed, then the resulting preprocessing token sequence is rescanned with all subsequent preprocessing tokens of the source file for more macro names to replace.

2 If the name of the macro being replaced is found during this scan of the replacement list (not including the rest of the source file's preprocessing tokens), it is not replaced. Further, if any nested replacements encounter the name of the macro being replaced, it is not replaced. These nonreplaced macro name preprocessing tokens are no longer available for further replacement even if they are later (re)examined in contexts in which that macro name preprocessing token would otherwise have been replaced.

3 The resulting completely macro-replaced preprocessing token sequence is not processed as a preprocessing directive even if it resembles one, but all pragma unary operator expressions within it are then processed as specified in 6.8 .9 below.

### 6.8.3.5 Scope of macro definitions

A macro definition lasts (independent of block structure) until a corresponding \#undef directive is encountered or (if none is encountered) until the end of translation phase 4.

2 A preprocessing directive of the form
\# undef identifier new-line
causes the specified identifier no longer to be defined as a macro name. It is ignored if the specified identifier is not currently defined as a macro name.

## Examples

3 1. The simplest use of this facility is to define a "manifest constant," as in

```
#define TABSIZE 100
int table[TABSIZE];
```

2. The following defines a function-like macro whose value is the maximum of its arguments. It has the advantages of working for any compatible types of the arguments and of generating in-line code without the overhead of function calling. It has the disadvantages of evaluating one or the other of its arguments a second time (including side effects) and generating more code than a function if invoked several times. It also cannot have its address taken, as it has none.
```
#define max(a, b) ((a) > (b) ? (a) : (b))
```

The parentheses ensure that the arguments and the resulting expression are bound properly.
3. To illustrate the rules for redefinition and reexamination, the sequence

```
#define x 3
#define f(a) f(x * (a))
#undef x
#define x 2
#define g f
#define z z[0]
#define h g(~
#define m(a) a(w)
#define w 0,1
#define t(a) a
#define p() int
#define q(x) x
#define r(x,y) x ## y
#define str(x) # x
f(y+1) + f(f(z)) % t(t(g)(0) + t)(1);
g(x+(3,4)-w) | h 5) & m
    (f)^m(m);
p() i[q()] = {q(1),r(2,3),r(4,),r(,5),r(,) };
char c[2][6] = { str(hello), str() };
```

results in

```
f(2 * (y+1)) + f(2 * (f(2 * (z[0])))) % f(2 * (0)) + t(1);
f(2 * (2+(3,4)-0,1)) | f(2 * (~ 5)) & f(2 * (0,1))^m(0,1);
int i[] = { 1, 23, 4, 5, };
char c[2][6] = { "hello", "" };
```

4. To illustrate the rules for creating character string literals and concatenating tokens, the sequence
```
#define str(s) # s
#define xstr(s) str(s)
#define debug(s, t) printf("x" # s "= %d, x" # t "= %s", \
x ## s, x ## t)
#define INCFILE(n) vers ## n // from previous #include example
#define glue(a, b) a ## b
#define xglue(a, b) glue(a, b)
#define HIGHLOW "hello"
#define LOW LOW ", world"
debug(1, 2);
fputs(str(strncmp("abc\0d", "abc", '\4') // this goes away
    == 0) str(: @\n), s);
#include xstr(INCFILE(2).h)
glue(HIGH, LOW);
xglue(HIGH, LOW)
```

results in

```
printf("x" "1" "= %d, x" "2" "= %s", x1, x2);
fputs(
    "strncmp(\"abc\\0d\", \"abc\", '\\4') == 0" ": @\n",
    s);
#include "vers2.h" (after macro replacement, before file access)
"hello";
"hello" ", world"
```

or, after concatenation of the character string literals,

```
printf("x1= %d, x2= %s", x1, x2);
fputs(
    "strncmp(\"abc\\0d\", \"abc\", '\\4') == 0: @\n",
    s);
#include "vers2.h" (after macro replacement, before file access)
"hello";
"hello, world"
```

Space around the \# and \#\# tokens in the macro definition is optional.
5. To illustrate the rules for
placemarker \#\# placemarker
the sequence

```
#define t(x,y,z) x ## y ## z
int j[] = { t(1,2,3), t(,4,5), t(6,,7), t(8,9,),
    t(10,,), t(,11,), t(,,12), t(,,) };
```

results in

```
int j[] = { 123, 45, 67, 89,
    10, 11, 12, };
```

6. To demonstrate the redefinition rules, the following sequence is valid.
```
#define OBJ_LIKE (1-1)
#define OBJ_LIKE /* white space */ (1-1) /* other */
#define FUNC_LIKE(a) ( a )
#define FUNC_LIKE( a )( /* note the white space */ \
    a /* other stuff on this line
                */ )
```

But the following redefinitions are invalid:

```
#define OBJ_LIKE (0) /* different token sequence */
#define OBJ_LIKE (1 - 1) /* different white space */
#define FUNC_LIKE (b) ( a ) /* different parameter usage */
#define FUNC_LIKE (b) ( b ) /* different parameter spelling */
```

7. Finally, to show the variable argument list macro facilities:
```
#define debug(...) fprintf(stderr, __VA_ARGS__)
#define showlist(...) puts(#__VA_ARGS__)
#define report(test, ...) ((test)?puts(#test):\
    printf(__VA_ARGS__))
debug("Flag");
debug("X = %d\n", x);
showlist(The first, second, and third items.);
report(x>y, "x is %d but y is %d", x, y);
```

results in

```
fprintf(stderr, "Flag" );
fprintf(stderr, "X = %d\n", x );
puts( "The first, second, and third items." );
((x>y) ?puts("x>y") :
    printf("x is %d but y is %d", x, y));
```


### 6.8.4 Line control

## Constraints

1 The string literal of a \#line directive, if present, shall be a character string literal.

## Semantics

2 The line number of the current source line is one greater than the number of new-line characters read or introduced in translation phase 1 (5.1.1.2) while processing the source file to the current token.

3 A preprocessing directive of the form
\# line digit-sequence new-line
causes the implementation to behave as if the following sequence of source lines begins with a source line that has a line number as specified by the digit sequence (interpreted as a decimal integer). The digit sequence shall not specify zero, nor a number greater than 2147483647.

4 A preprocessing directive of the form

```
# line digit-sequence "s-char-sequence opt" new-line
```

sets the line number similarly and changes the presumed name of the source file to be the contents of the character string literal.

5 A preprocessing directive of the form
\# line pp-tokens new-line
(that does not match one of the two previous forms) is permitted. The preprocessing tokens after line on the directive are processed just as in normal text (each identifier currently defined as a macro name is replaced by its replacement list of preprocessing tokens). The directive resulting after all replacements shall match one of the two previous forms and is then processed as appropriate.

### 6.8.5 Error directive

## Semantics

1 A preprocessing directive of the form

```
# error pp-tokens opt new-line
```

causes the implementation to produce a diagnostic message that includes the specified sequence of preprocessing tokens.

### 6.8.6 Pragma directive

## Semantics

1 A preprocessing directive of the form
\# pragma pp-tokens ${ }_{\text {opt }}$ new-line
where the preprocessing token STDC does not immediately follow the pragma on the directive causes the implementation to behave in a manner which it shall document. The behavior might cause translation to fail or the resulting program to behave in a non-conforming manner. Any such pragma that is not recognized by the implementation is ignored.

2 If the preprocessing token STDC does immediately follow the pragma on the directive, then no macro replacements are performed on the directive, and the directive shall have one of the following forms whose meaning is described elsewhere:

```
#pragma STDC FP_CONTRACT on-off-switch
#pragma STDC FENV_ACCESS on-off-switch
#pragma STDC CX_LIMITED_RANGE on-off-switch
```

on-off-switch: one of
ON OFF DEFAULT
Forward references: the FP_CONTRACT pragma (7.7.2), the FENV_ACCESS pragma (7.6.1), the CX_LIMITED_RANGE pragma (7.8.1).

### 6.8.7 Null directive

## Semantics

1 A preprocessing directive of the form
\# new-line
has no effect.

### 6.8.8 Predefined macro names

1 The following macro names shall be defined by the implementation:
__LINE__ The line number of the current source line (a decimal constant).
__FILE__ The presumed name of the source file (a character string literal).
__DATE__ The date of translation of the source file (a character string literal of the form "Mmm dd yyyy", where the names of the months are the same as those generated by the asctime function, and the first character of dd is a space character if the value is less than 10). If the date of
translation is not available, an implementation-defined valid date shall be supplied.
$\qquad$ The time of translation of the source file (a character string literal of the form "hh:mm:ss" as in the time generated by the asctime function). If the time of translation is not available, an implementation-defined valid time shall be supplied.
__STDC__ The decimal constant 1 , intended to indicate a conforming implementation.
__STDC_VERSION__ The decimal constant 199901L. ${ }^{128}$
2 The following macro names are conditionally defined by the implementation: 1912.nr:c 0u+000m'unu

STDC_IEC_559__ The decimal constant 1, intended to indicate conformance to the specifications in Annex F (IEC 559 floating-point arithmetic).
__STDC_IEC_559_COMPLEX__ The decimal constant 1, intended to indicate adherence to the specifications in informative Annex G (IEC 559 compatible complex arithmetic).

3 The values of the predefined macros (except for __LINE__ and __FILE__) remain constant throughout the translation unit.

4 None of these macro names, nor the identifier defined, shall be the subject of a \#define or a \#undef preprocessing directive. All predefined macro names shall begin with a leading underscore followed by an uppercase letter or a second underscore.

Forward references: the asctime function (7.16.3.1).

[^61]
### 6.8.9 Pragma operator

## Semantics

1 A unary operator expression of the form:
_Pragma ( string-literal )
is processed as follows. The string-literal is destringized by deleting the L prefix, if present, deleting the leading and trailing double-quotes, replacing each escape sequence \" by a double-quote, and replacing each escape sequence $\backslash \backslash$ by a single backslash. The resulting sequence of characters is processed through translation phase 3 to produce preprocessing tokens that are executed as if they were the pp-tokens in a pragma directive. The original four preprocessing tokens in the unary operator expression are removed.

## Examples

2 A directive of the form:
\#pragma list on "..\listing.dir"
can also be expressed as:
_Pragma ( "listing on \".. <br>listing.dir\"" )
The latter form is processed in the same way whether it appears literally as shown, or results from macro replacement, as in:

```
#define LISTING(x) PRAGMA(listing on #x)
#define PRAGMA(x) _Pragma(#x)
LISTING ( ..\listing.dir )
```


### 6.9 Future language directions

### 6.9.1 Character escape sequences

1 Lowercase letters as escape sequences are reserved for future standardization. Other characters may be used in extensions.

### 6.9.2 Storage-class specifiers

1 The placement of a storage-class specifier other than at the beginning of the declaration specifiers in a declaration is an obsolescent feature.

### 6.9.3 Function declarators

1 The use of function declarators with empty parentheses (not prototype-format parameter type declarators) is an obsolescent feature.

### 6.9.4 Function definitions

1 The use of function definitions with separate parameter identifier and declaration lists (not prototype-format parameter type and identifier declarators) is an obsolescent feature.

### 6.9.5 Pragma directives

1 Pragmas whose first pp-token is STDC are reserved for future standardization.

## 7. Library

### 7.1 Introduction

### 7.1.1 Definitions of terms

1 A string is a contiguous sequence of characters terminated by and including the first null character. A 'pointer to" a string is a pointer to its initial (lowest addressed) character. The 'length' of a string is the number of characters preceding the null character and its "value" is the sequence of the values of the contained characters, in order.

2 A letter is a printing character in the execution character set corresponding to any of the 52 required lowercase and uppercase letters in the source character set, listed in 5.2.1.

3 The decimal-point character is the character used by functions that convert floatingpoint numbers to or from character sequences to denote the beginning of the fractional part of such character sequences. ${ }^{129}$ It is represented in the text and examples by a period, but may be changed by the setlocale function.

4 A wide character is a code value (a binary encoded integer) of an object of type wchar_t that corresponds to a member of the extended character set. ${ }^{130}$

5 A null wide character is a wide character with code value zero.
6 A wide string is a contiguous sequence of wide characters terminated by and including the first null wide character. A pointer to $a$ wide string is a pointer to its initial (lowest addressed) wide character. The length of a wide string is the number of wide characters preceding the null wide character and the value of a wide string is the sequence of code values of the contained wide characters, in order.

7 A shift sequence is a contiguous sequence of bytes within a multibyte string that (potentially) causes a change in shift state. (See subclause 5.2.1.2.) A shift sequence shall not have a corresponding wide character; it is instead taken to be an adjunct to

[^62]130. An equivalent definition can be found in subclause 6.1.3.4.
an adjacent multibyte character. ${ }^{131}$
Forward references: character handling (7.3), the setlocale function (7.5.1.1).

### 7.1.2 Standard headers

1 Each library function is declared, with a type that includes a prototype, in a header, ${ }^{132}$ whose contents are made available by the \#include preprocessing directive. The header declares a set of related functions, plus any necessary types and additional macros needed to facilitate their use. Declarations of types described in this clause shall not include type qualifiers, unless explicity stated otherwise.

2 The standard headers are

```
<assert.h>
<errno.h>
<inttypes.h>
<locale.h>
<signal.h>
<stddef.h>
<string.h>
<wchar.h>
```

```
<complex.h>
<wctype.h>
```

<ctype.h>
<fenv.h> <float.h>
<iso646.h> <limits.h>
<math.h> <setjmp.h>
<stdarg.h> <stdbool.h>
<stdio.h> <stdlib.h>
<tgmath.h> <time.h>

If a file with the same name as one of the above < and > delimited sequences, not provided as part of the implementation, is placed in any of the standard places for a source file to be included, the behavior is undefined.

4 Standard headers may be included in any order; each may be included more than once in a given scope, with no effect different from being included only once, except that the effect of including <assert. h> depends on the definition of NDEBUG. If used, a header shall be included outside of any external declaration or definition, and it shall first be included before the first reference to any of the functions or objects it declares, or to any of the types or macros it defines. However, if an identifier is declared or defined in more than one header, the second and subsequent associated headers may be included after the initial reference to the identifier. The program shall not have any macros with names lexically identical to keywords currently defined prior to the

[^63]inclusion.
5 Any definition of an object-like macro described in this clause shall expand to code that is fully protected by parentheses where necessary, so that it groups in an arbitrary expression as if it were a single identifier.

6 Any declaration of a library function shall have external linkage.
Forward references: diagnostics (7.2).

### 7.1.3 Reserved identifiers

1 Each header declares or defines all identifiers listed in its associated subclause, and optionally declares or defines identifiers listed in its associated future library directions subclause and identifiers which are always reserved either for any use or for use as file scope identifiers.

- All identifiers that begin with an underscore and either an uppercase letter or another underscore are always reserved for any use.
- All identifiers that begin with an underscore are always reserved for use as macros and as identifiers with file scope in both the ordinary and tag name spaces.
- Each macro name in any of the following subclauses (including the future library directions) is reserved for use as specified if any of its associated headers is included; unless explicitly stated otherwise (see 7.1.8).
- All identifiers with external linkage in any of the following subclauses (including the future library directions) are always reserved for use as identifiers with external linkage. ${ }^{133}$
- Each identifier with file scope listed in any of the following subclauses (including the future library directions) is reserved for use as macro and as an identifier with file scope in the same name space if any of its associated headers is included.

No other identifiers are reserved. If the program declares or defines an identifier that is reserved in that context (other than as allowed by 7.1.8), the behavior is undefined. ${ }^{134}$

[^64]3 If the program removes (with \#undef) any macro definition of an identifier in the first group listed above, the behavior is undefined.

### 7.1.4 Errors <errno.h>

1 The header <errno. h> defines several macros, all relating to the reporting of error conditions.

2 The macros are
EDOM
EILSEQ
ERANGE
which expand to integer constant expressions with type int, distinct positive values, and which are suitable for use in \#if preprocessing directives; and
errno
which expands to a modifiable lvalue ${ }^{135}$ that has type int, the value of which is set to a positive error number by several library functions. It is unspecified whether errno is a macro or an identifier declared with external linkage. If a macro definition is suppressed in order to access an actual object, or a program defines an identifier with the name errno, the behavior is undefined.

3 The value of errno is zero at program startup, but is never set to zero by any library function. ${ }^{136}$ The value of errno may be set to nonzero by a library function call whether or not there is an error, provided the use of errno is not documented in the description of the function in this International Standard.

4 Additional macro definitions, beginning with $\mathbf{E}$ and a digit or $\mathbf{E}$ and an uppercase letter, ${ }^{137}$ may also be specified by the implementation.

[^65]
### 7.1.5 Limits <float. $\mathrm{h}>$ and <limits. h >

1 The headers <float.h> and <limits.h> define several macros that expand to various limits and parameters.

2 The macros, their meanings, and the constraints (or restrictions) on their values are listed in 5.2.4.2.

### 7.1.6 Common definitions <stddef.h>

1 The following types and macros are defined in the standard header <stddef.h>. Some are also defined in other headers, as noted in their respective subclauses.

2 The types are
ptrdiff_t
which is the signed integer type of the result of subtracting two pointers;
size_t
which is the unsigned integer type of the result of the sizeof operator; and

```
wchar_t
```

which is an integer type whose range of values can represent distinct codes for all members of the largest extended character set specified among the supported locales; the null character shall have the code value zero and each member of the basic character set defined in 5.2 .1 shall have a code value equal to its value when used as the lone character in an integer character constant.

3 The macros are

## NULL

which expands to an implementation-defined null pointer constant; and
offsetof (type, member-designator)
which expands to an integer constant expression that has type size_t, the value of which is the offset in bytes, to the structure member (designated by memberdesignator), from the beginning of its structure (designated by type). The memberdesignator shall be such that given

```
static type t;
```

then the expression $\&(t$.member-designator) evaluates to an address constant. (If the specified member is a bit-field, the behavior is undefined.)

Forward references: localization (7.5).

### 7.1.7 Boolean type and values <stdbool . h >

1 The header <stdbool. h > defines one type and three macros.
2 The type is
bool
which is an integer type that promotes to int or unsigned int, and that is suitable to be used as the type of a bit-field. A bit-field of any width and type bool shall be capable for representing the value $1 .{ }^{138}$

3 The macros are
true
which expands to the decimal constant 1 , false
which expands to the decimal constant 0 , and
__bool_true_false_are_defined
which expands to the decimal constant 1 . The macros are suitable for use in \#if preprocessing directives.

[^66]
### 7.1.8 Use of library functions

1 Each of the following statements applies unless explicitly stated otherwise in the detailed descriptions that follow. If an argument to a function has an invalid value (such as a value outside the domain of the function, or a pointer outside the address space of the program, or a null pointer) or a type (after promotion) not expected by a function with variable number of arguments, the behavior is undefined. If a function argument is described as being an array, the pointer actually passed to the function shall have a value such that all address computations and accesses to objects (that would be valid if the pointer did point to the first element of such an array) are in fact valid. Any function declared in a header may be additionally implemented as a function-like macro defined in the header, so if a library function is declared explicitly when its header is included, one of the techniques shown below can be used to ensure the declaration is not affected by such a macro. Any macro definition of a function can be suppressed locally by enclosing the name of the function in parentheses, because the name is then not followed by the left parenthesis that indicates expansion of a macro function name. For the same syntactic reason, it is permitted to take the address of a library function even if it is also defined as a macro. ${ }^{139}$ The use of \#undef to remove any macro definition will also ensure that an actual function is referred to. Any invocation of a library function that is implemented as a macro shall expand to code that evaluates each of its arguments exactly once, fully protected by parentheses where necessary, so it is generally safe to use arbitrary expressions as arguments. ${ }^{140}$ Likewise, those function-like macros described in the following subclauses may be invoked in an expression anywhere a function with a compatible return type could be called. ${ }^{141}$ All object-like macros listed as expanding to integer
139. This means that an implementation must provide an actual function for each library function, even if it also provides a macro for that function.
140. Such macros might not contain the sequence points that the corresponding function calls do.
141. Because external identifiers and some macro names beginning with an underscore are reserved, implementations may provide special semantics for such names. For example, the identifier _BUILTIN_abs could be used to indicate generation of in-line code for the abs function. Thus, the appropriate header could specify
\#define abs(x) _BUILTIN_abs(x)
for a compiler whose code generator will accept it.
In this manner, a user desiring to guarantee that a given library function such as abs will be a genuine function may write

## \#undef abs

whether the implementation's header provides a macro implementation of abs or a built-in implementation. The prototype for the function, which precedes and is hidden by any macro definition, is thereby revealed also.
constant expressions shall additionally be suitable for use in \#if preprocessing directives.

2 Provided that a library function can be declared without reference to any type defined in a header, it is also permissible to declare the function and use it without including its associated header.

3 There is a sequence point immediately before a library function return.
4 The functions in the standard library are not guaranteed to be reentrant and may modify objects with static storage duration. ${ }^{142}$

## Examples

5 The function atoi may be used in any of several ways:

- by use of its associated header (possibly generating a macro expansion)

```
#include <stdlib.h>
const char *str;
/* ... */
i = atoi(str);
```

- by use of its associated header (assuredly generating a true function reference)

```
#include <stdlib.h>
#undef atoi
const char *str;
/* ... */
i = atoi(str);
```

```
#include <stdlib.h>
```

\#include <stdlib.h>
const char *str;
const char *str;
/* ... */
/* ... */
i = (atoi) (str);

```
i = (atoi) (str);
```

or

- by explicit declaration

142. Thus, a signal handler cannot, in general, call standard library functions.

Working Draft, 1997-11-21, WG14/N794 J11/97-158

```
extern int atoi (const char *);
const char *str;
/* ... */
i = atoi(str);
```


### 7.2 Diagnostics <assert. h >

1 The header <assert. h > defines the assert macro and refers to another macro, NDEBUG
which is not defined by <assert.h>. If NDEBUG is defined as a macro name at the point in the source file where <assert.h> is included, the assert macro is defined simply as

```
#define assert(ignore) ((void)0)
```

2 The assert macro shall be implemented as a macro, not as an actual function. If the macro definition is suppressed in order to access an actual function, the behavior is undefined.

### 7.2.1 Program diagnostics

### 7.2.1.1 The assert macro

## Synopsis

```
#include <assert.h>
void assert(int expression);
```


## Description

2 The assert macro puts diagnostic tests into programs. When it is executed, if expression is false (that is, compares equal to 0), the assert macro writes information about the particular call that failed (including the text of the argument, the name of the source file, and the source line number - the latter are respectively the values of the preprocessing macros __FILE__ and __LINE__ and the identifier __func__) on the standard error file in an implementation-defined format. ${ }^{143}$ It then calls the abort function.

## Returns

3 The assert macro returns no value.

[^67]Working Draft, 1997-11-21, WG14/N794 J11/97-158

Forward references: the abort function (7.14.4.1).

### 7.3 Character handling <ctype.h>

1 The header <ctype.h> declares several functions useful for testing and mapping characters. ${ }^{144}$ In all cases the argument is an int, the value of which shall be representable as an unsigned char or shall equal the value of the macro EOF. If the argument has any other value, the behavior is undefined.

2 The behavior of these functions is affected by the current locale. Those functions that have locale-specific aspects only when not in the "C" locale are noted below.

3 The term printing character refers to a member of a locale-specific set of characters, each of which occupies one printing position on a display device; the term control character refers to a member of a locale-specific set of characters that are not printing characters. ${ }^{145}$

Forward references: EOF (7.13.1), localization (7.5).

### 7.3.1 Character testing functions

1 The functions in this subclause return nonzero (true) if and only if the value of the argument $\mathbf{c}$ conforms to that in the description of the function.

### 7.3.1.1 The isalnum function

Synopsis

```
#include <ctype.h>
int isalnum(int c);
```


## Description

2 The isalnum function tests for any character for which isalpha or isdigit is true.

[^68]
### 7.3.1.2 The isalpha function

Synopsis

```
#include <ctype.h>
int isalpha(int c);
```


## Description

2 The isalpha function tests for any character for which isupper or islower is true, or any character that is one of a locale-specific set of characters for which none of iscntrl, isdigit, ispunct, or isspace is true. ${ }^{146}$ In the "C" locale, isalpha returns true only for the characters for which isupper or islower is true.

### 7.3.1.3 The isblank function

Synopsis

```
#include <ctype.h>
int isblank(int c);
```


## Description

2 The isblank function tests for any character for that is a standard blank character or is one of a locale-specific set of characters, for which isalnum is false. The standard blank characters are the following: space (' '), and horizontal tab (' $\backslash \mathbf{t}^{\prime}$ ). In the "C" locale, isblank returns true only for the standard blank characters.

### 7.3.1.4 The iscntrl function

## Synopsis

```
#include <ctype.h>
int iscntrl(int c);
```


## Description

The iscntrl function tests for any control character.

[^69]
### 7.3.1.5 The isdigit function

Synopsis

```
#include <ctype.h>
int isdigit(int c);
```


## Description

2 The isdigit function tests for any decimal-digit character (as defined in 5.2.1).

### 7.3.1.6 The isgraph function

Synopsis

```
#include <ctype.h>
int isgraph(int c);
```


## Description

2 The isgraph function tests for any printing character except space (' ' ) .

### 7.3.1.7 The islower function

Synopsis

```
#include <ctype.h>
int islower(int c);
```


## Description

2 The islower function tests for any character that is a lowercase letter or is one of a locale-specific set of characters for which none of iscntrl, isdigit, ispunct, or isspace is true. In the "C" locale, islower returns true only for the characters defined as lowercase letters (as defined in 5.2.1).

### 7.3.1.8 The isprint function

Synopsis

```
#include <ctype.h>
    int isprint(int c);
```


## Description

The isprint function tests for any printing character including space (' ').

### 7.3.1.9 The ispunct function

Synopsis

```
#include <ctype.h>
int ispunct(int c);
```


## Description

2 The ispunct function tests for any printing character that is one of a locale-specific set of characters for which neither isspace nor isalnum is true.

### 7.3.1.10 The isspace function

## Synopsis

```
#include <ctype.h>
int isspace(int c);
```


## Description

2 The isspace function tests for any character that is a standard white-space character or is one of a locale-specific set of characters for which isalnum is false. The standard white-space characters are the following: space (' '), form feed ('\f'), new-line ( ${ }^{\prime} \backslash \mathbf{n}^{\prime}$ ), carriage return ( ${ }^{\prime} \backslash \mathbf{r}^{\prime}$ ), horizontal tab ( ${ }^{\prime} \backslash \mathrm{t}^{\prime}$ ), and vertical tab (' $\left.\backslash \mathbf{v}^{\prime}\right)$. In the "C" locale, isspace returns true only for the standard white-space characters.

### 7.3.1.11 The isupper function

Synopsis

```
#include <ctype.h>
int isupper(int c);
```


## Description

2 The isupper function tests for any character that is an uppercase letter or is one of a locale-specific set of characters for which none of iscntrl, isdigit, ispunct, or isspace is true. In the " $\mathbf{C}$ " locale, isupper returns true only for the characters defined as uppercase letters (as defined in 5.2.1).

### 7.3.1.12 The isxdigit function

Synopsis

```
#include <ctype.h>
int isxdigit(int c);
```


## Description

2 The isxdigit function tests for any hexadecimal-digit character (as defined in 6.1.3.1).

### 7.3.2 Character case mapping functions

### 7.3.2.1 The tolower function

Synopsis
1 \#include <ctype.h>
int tolower (int c);

## Description

2 The tolower function converts an uppercase letter to a corresponding lowercase letter.

## Returns

3 If the argument is a character for which isupper is true and there are one or more corresponding characters, as specified by the current locale, for which islower is true, the tolower function returns one of the corresponding characters (always the same one for any given locale); otherwise, the argument is returned unchanged.

### 7.3.2.2 The toupper function

Synopsis

```
#include <ctype.h>
int toupper(int c);
```


## Description

2 The toupper function converts a lowercase letter to a corresponding uppercase letter.

## Returns

3 If the argument is a character for which islower is true and there are one or more corresponding characters, as specified by the current locale, for which isupper is true, the toupper function returns one of the corresponding characters (always the same one for any given locale); otherwise, the argument is returned unchanged.

### 7.4 Integer types <inttypes.h>

1 The header <inttypes.h> defines sets of typedef names for integer types having specified widths, and defines corresponding sets of macros. It also defines macros that specify limits of integer types corresponding to typedef names defined in other standard headers, and declares four functions for converting numeric character strings to greatest-width integers.

2 Typedef names are defined in the following categories:
— integer types having certain exact widths;
— integer types having at least certain specified widths;

- fastest integer types having at least certain specified widths;
— integer types wide enough to hold pointers to objects;
- integer types having greatest width.
(Some of these typedef names may denote the same type.)
3 Corresponding macros specify limits of the defined types, construct suitable character constants, and provide conversion specifiers for use with the formatted input/output functions.

4 For each typedef name described herein that can be defined as a type existing in the implementation, ${ }^{147}$ <inttypes.h> shall define that typedef name, and it shall define the associated macros. Conversely, for each typedef name described herein that cannot be defined as a type existing in the implementation, <inttypes.h> shall not define that typedef name, nor shall it define the associated macros.

### 7.4.1 Typedef names for integer types

1 When typedef names differing only in the absence or presence of the initial $\mathbf{u}$ are defined, they shall denote corresponding signed and unsigned types as described in subclause 6.1.2.5.

[^70]
### 7.4.1.1 Exact-width integer types

1 Each of the following typedef names designates an integer type that has exactly the specified width. These typedef names have the general form of int $n \_t$ or uint $n \_t$ where $n$ is the required width. For example, uint8_t denotes an unsigned integer type that has a width of exactly 8 bits.

2 The following designate exact-width signed integer types:
int8_t int16_t int32_t int64_t

3 The following designate exact-width unsigned integer types:

```
uint8_t uint16_t uint32_t uint64_t
```

(Any of these types might not exist.)

### 7.4.1.2 Minimum-width integer types

1 Each of the following typedef names designates an integer type that has at least the specified width, such that no integer type of lesser size has at least the specified width. These typedef names have the general form of int_least $n \_t$ or uint_least $n \_t$ where $n$ is the minimum required width. For example, int_least32_t denotes a signed integer type that has a width of at least 32 bits.

2 The following designate minimum-width signed integer types:

```
int_least8_t int_least16_t
int_least32_t int_least64_t
```

3 The following designate minimum-width unsigned integer types:

```
uint_least8_t uint_least16_t
uint_least32_t uint_least64_t
```

(These types must exist.)

### 7.4.1.3 Fastest minimum-width integer types

1 Each of the following typedef names designates an integer type that is usually fastest ${ }^{148}$ to operate with among all integer types that have at least the specified width. These typedef names have the general form of int_fast $n \_t$ or uint_fast $n \_t$ where $n$ is the minimum required width. For example, int_fast16_t denotes the

[^71]fastest signed integer type that has a width of at least 16 bits.
2 The following designate fastest minimum-width signed integer types:

```
int_fast8_t int_fast16_t
int_fast32_t int_fast64_t
```

3 The following designate fastest minimum-width unsigned integer types:

```
uint_fast8_t uint_fast16_t
uint_fast32_t uint_fast64_t
```

(These types must exist.)

### 7.4.1.4 Integer types capable of holding object pointers

1 The following typedef name designates a signed integer type with the property that any valid pointer to void can be converted to this type, then converted back to pointer to void, and the result will compare equal to the original pointer:
intptr_t

2 The following typedef name designates an unsigned integer type with the property that any valid pointer to void can be converted to this type, then converted back to pointer to void, and the result will compare equal to the original pointer:

```
uintptr_t
```

(Either or both of these types might not exist.)

### 7.4.1.5 Greatest-width integer types

1 The following typedef name designates a signed integer type capable of representing any value of any signed integer type:

## intmax_t

2 The following typedef name designates an unsigned integer type capable of representing any value of any unsigned integer type:
uintmax_t
(These types must exist.)

### 7.4.2 Limits of specified-width integer types

1 The following object-like macros ${ }^{149}$ specify the minimum and maximum limits of integer types corresponding to the typedef names defined in <inttypes.h>. Each macro name corresponds to a similar typedef name in subclause 7.4.1.

2 Each instance of any defined macro shall be replaced by a constant expression suitable for use in \#if preprocessing directives, and this expression shall have the same type as would an expression that is an object of the corresponding type converted according to the integer promotions. Its implementation-defined value shall be equal to or greater in magnitude (absolute value) than the corresponding value given below, with the same sign.

### 7.4.2.1 Limits of exact-width integer types

- minimum values of exact-width signed integer types

INT8_MIN -127
INT16_MIN -32767
INT32_MIN -2147483647
INT64_MIN -9223372036854775807
(The value must be either that given or exactly 1 less.)

- maximum values of exact-width signed integer types

INT8_MAX
+127
INT16_MAX +32767
INT32_MAX +2147483647
INT64_MAX +9223372036854775807
(The value must be exactly that given.)

- maximum values of exact-width unsigned integer types

UINT8_MAX
255
UINT16_MAX
65535
UINT32_MAX
4294967295
UINT64_MAX
18446744073709551615
(The value must be exactly that given.)

[^72]
### 7.4.2.2 Limits of minimum-width integer types

- minimum values of minimum-width signed integer types

| INT_LEAST8_MIN | -127 |
| :--- | ---: |
| INT_LEAST16_MIN | -32767 |
| INT_LEAST32_MIN | $\mathbf{- 2 1 4 7 4 8 3 6 4 7}$ |
| INT_LEAST64_MIN | -9223372036854775807 |

- maximum values of minimum-width signed integer types

| INT_LEAST8_MAX | $\mathbf{+ 1 2 7}$ |
| :--- | ---: |
| INT_LEAST16_MAX | $\mathbf{+ 3 2 7 6 7}$ |
| INT_LEAST32_MAX | $\mathbf{+ 2 1 4 7 4 8 3 6 4 7}$ |
| INT_LEAST64_MAX | +9223372036854775807 |

- maximum values of minimum-width unsigned integer types

UINT_LEAST8_MAX 255
UINT_LEAST16_MAX 65535
UINT_LEAST32_MAX 4294967295
UINT_LEAST64_MAX 18446744073709551615
7.4.2.3 Limits of fastest minimum-width integer types

- minimum values of fastest minimum-width signed integer types

| INT_FAST8_MIN | -127 |
| :--- | ---: |
| INT_FAST16_MIN | -32767 |
| INT_FAST32_MIN | -2147483647 |
| INT_FAST64_MIN | -9223372036854775807 |

- maximum values of fastest minimum-width signed integer types

| INT_FAST8_MAX | +127 |
| :--- | ---: |
| INT_FAST16_MAX | +32767 |
| INT_FAST32_MAX | +2147483647 |
| INT_FAST64_MAX | +9223372036854775807 |

- maximum values of fastest minimum-width unsigned integer types

UINT_FAST8_MAX
255
UINT_FAST16_MAX
65535
UINT_FAST32_MAX
4294967295
UINT_FAST64_MAX
18446744073709551615

### 7.4.2.4 Limits of integer types capable of holding object pointers

- minimum value of pointer-holding signed integer type INTPTR_MIN -32767
- maximum value of pointer-holding signed integer type
INTPTR_MAX +32767
- maximum value of pointer-holding unsigned integer type

UINTPTR_MAX 65535

### 7.4.2.5 Limits of greatest-width integer types

- minimum value of greatest-width signed integer type

INTMAX_MIN -9223372036854775807

- maximum value of greatest-width signed integer type

INTMAX_MAX +9223372036854775807

- maximum value of greatest-width unsigned integer type

UINTMAX_MAX
18446744073709551615

### 7.4.3 Macros for integer constants

1 The following function-like macros ${ }^{150}$ expand to integer constants suitable for initializing objects that have integer types corresponding to typedef names defined in <inttypes.h>. Each macro name corresponds to a similar typedef name in subclause 7.4.1.2 or 7.4.1.5.

2 The argument in any instance of these macros shall be a decimal, octal, or hexadecimal constant (as defined in subclause 6.1.3.2) with a value that does not exceed the limits for the corresponding type.

### 7.4.3.1 Macros for minimum-width integer constants

1 Each of the following macros expands to an integer constant having the value specified by its argument and a type with at least the specified width. These macro names have the general form of INT $n \_\mathbf{C}$ or UINT $n \_\mathbf{C}$ where $n$ is the minimum required width. For example, UINT64_C ( $0 \times 123$ ) might expand to the integer constant $0 \times 123 \mathrm{LLL}$.

2 The following expand to integer constants that have signed integer types:

[^73]```
INT8_C(value) INT16_C(value)
INT32_C(value) INT64_C(value)
```

3 The following expand to integer constants that have unsigned integer types:

```
UINT8_C(value) UINT16_C(value)
UINT32_C(value) UINT64_C(value)
```


### 7.4.3.2 Macros for greatest-width integer constants

1 The following macro expands to an integer constant having the value specified by its argument and the type intmax_t:

## INTMAX_C (value)

2 The following macro expands to an integer constant having the value specified by its argument and the type uintmax_t:

## UINTMAX_C (value)

### 7.4.4 Macros for format specifiers

1 Each of the following object-like macros ${ }^{151}$ expands to a string literal containing a conversion specifier, possibly modified by a prefix such as $\mathbf{h h}, \mathbf{h}, \mathbf{l}$, or ll, suitable for use within the format argument of a formatted input/output function when converting the corresponding integer type. These macro names have the general form of PRI (character string literals for the fprintf family) or SCN (character string literals for the fscanf family, ${ }^{152}$ followed by the conversion specifier, followed by a name corresponding to a similar typedef name in subclause 7.4.1. For example, PRIdFAST32 can be used in a format string to print the value of an integer of type int_fast32_t.

2 The fprintf macros for signed integers are:

[^74]| PRId8 | PRId16 | PRId32 | PRId64 |
| :--- | :--- | :--- | :--- |
| PRIdLEAST8 | PRIdLEAST16 | PRIdLEAST32 | PRIdLEAST64 |
| PRIdFAST8 | PRIdFAST16 | PRIdFAST32 | PRIdFAST64 |
| PRIdMAX | PRIdPTR |  |  |
| PRIi8 |  |  | PRIi64 |
| PRIiLEAST8 | PRIi16 | PRIi32 | PRIiLEAST16 |
| PRIiFAST8 | PRIiFAST16 | PRIiLEAST32 | PRAST32 |
| PRIiMAX | PRIiPTR |  |  |

3 The fprintf macros for unsigned integers are:

| PRIO8 | PRIO16 | PRIO32 | PRIO64 |
| :---: | :---: | :---: | :---: |
| PRIOLEAST8 | PRIOLEAST16 | PRIOLEAST32 | PRIOLEAST64 |
| PRIOFAST8 | PRIoFAST16 | PRIOFAST32 | PRIOFAST64 |
| PRIOMAX | PRIOPTR |  |  |
| PRIU8 | PRIu16 | PRIu32 | PRIU64 |
| PRIuLEAST8 | PRIuLEAST16 | PRIULEAST32 | PRIULEAST64 |
| PRIuFAST8 | PRIUFAST16 | PRIUFAST32 | PRIUFAST64 |
| PRIUMAX | PRIUPTR |  |  |
| PRIx 8 | PRIx16 | PRIx32 | PRIx64 |
| PRIxLEAST8 | PRIxLEAST16 | PRIxLEAST32 | PRIxLEAST64 |
| PRIxFAST8 | PRIxFAST16 | PRIxFAST32 | PRIxFAST64 |
| PRIxMAX | PRIxPTR |  |  |
| PRIX8 | PRIX16 | PRIX32 | PRIX64 |
| PRIXLEAST8 | PRIXLEAST16 | PRIXLEAST32 | PRIXLEAST64 |
| PRIXFAST8 | PRIXFAST16 | PRIXFAST32 | PRIXFAST64 |
| PRIXMAX | PRIXPTR |  |  |

4 The fscanf macros for signed integers are:

| SCNd8 | SCNd16 | SCNd32 | SCNd64 |
| :--- | :--- | :--- | :--- |
| SCNdLEAST8 | SCNdLEAST16 | SCNdLEAST32 | SCNDLEAST64 |
| SCNdFAST8 | SCNdFAST16 | SCNdFAST32 | SCNDFAST64 |
| SCNdMAX | SCNdPTR |  |  |
| SCNi8 |  |  | SCNi64 |
| SCNiLEAST8 | SCNi16 | SCNi32 | SCNiLEAST16 |
| SCNiFAST8 | SCNiFAST16 | SCNiLEAST32 | SCNiFAST32 |
| SCNiMAX | SCNiPTR |  | SCNiFAST64 |

The fscanf macros for unsigned integers are:

| SCNo8 | SCNo16 | SCNo32 | SCN064 |
| :---: | :---: | :---: | :---: |
| SCNOLEAST8 | SCNoLEAST16 | SCNOLEAST32 | SCNoLEAST64 |
| SCNoFAST8 | SCNoFAST16 | SCNoFAST32 | SCNoFAST64 |
| SCNoMAX | SCNoPTR |  |  |
| SCNu8 | SCNu16 | SCNu32 | SCNu64 |
| SCNuLEAST8 | SCNuLEAST16 | SCNuLEAST32 | SCNuLEAST64 |
| SCNuFAST8 | SCNuFAST16 | SCNuFAST32 | SCNuFAST64 |
| SCNuMAX | SCNuPTR |  |  |
| SCNx8 | SCNx16 | SCNx 32 | SCNx 64 |
| SCNxLEAST8 | SCNxLEAST16 | SCNxLEAST32 | SCNxLEAST64 |
| SCNxFAST8 | SCNxFAST16 | SCNxFAST32 | SCNxFAST64 |
| SCNxMAX | SCNxPTR |  |  |

6 Because the default argument promotions do not affect pointer parameters, there might not exist suitable fscanf format specifiers for some of the typedef names defined in this header. Consequently, as a special exception to the requirement that the implementation shall define all macros associated with each typedef name defined in this header, in such a case the problematic fscanf macros may be left undefined.

Examples

```
#include <inttypes.h>
#include <wchar.h>
int main(void)
{
    uintmax_t i = UINTMAX_MAX; // this type always exists
    wprintf(L"The largest integer value is %020"
            PRIxMAX "\n", i);
    return 0;
}
```


### 7.4.5 Limits of other integer types

1 The following object-like macros ${ }^{151}$ specify the minimum and maximum limits of integer types corresponding to typedef names defined in other standard headers.

2 Each instance of these macros shall be replaced by a constant expression suitable for use in \#if preprocessing directives, and this expression shall have the same type as would an expression that is an object of the corresponding type converted according to the integer promotions. Its implementation-defined value shall be equal to or greater in magnitude (absolute value) than the corresponding value given below, with the same sign.

- limits of ptrdiff_t

PTRDIFF_MIN
PTRDIFF_MAX

- limits of sig_atomic_t

SIG_ATOMIC_MIN
SIG_ATOMIC_MAX

- limit of size_t

SIZE_MAX

- limits of wchar_t

WCHAR_MIN
WCHAR_MAX

- limits of wint_t

WINT_MIN
WINT_MAX
-65535
+65535
see below
see below

65535
see below
see below
see below
see below

If sig_atomic_t is defined as a signed integer type, the value of SIG_ATOMIC_MIN shall be no greater than -127 and the value of SIG_ATOMIC_MAX shall be no less than 127; otherwise, sig_atomic_t is defined as an unsigned integer type, and the value of SIG_ATOMIC_MIN shall be 0 and the value of SIG_ATOMIC_MAX shall be no less than 255.

4 If wchar_t is defined as a signed integer type, the value of WCHAR_MIN shall be no greater than - 127 and the value of WCHAR_MAX shall be no less than 127; otherwise, wchar_t is defined as an unsigned integer type, and the value of WCHAR_MIN shall be 0 and the value of WCHAR_MAX shall be no less than 255.

If wint_t is defined as a signed integer type, the value of WINT_MIN shall be no greater than -32767 and the value of WINT_MAX shall be no less than 32767; otherwise, wint_t is defined as an unsigned integer type, and the value of WINT_MIN shall be 0 and the value of WINT_MAX shall be no less than 65535.

### 7.4.6 Conversion functions for greatest-width integer types

### 7.4.6.1 The strtoimax function

Synopsis

```
#include <inttypes.h>
intmax_t strtoimax(const char * restrict nptr,
    char ** restrict endptr, int base);
```


## Description

2 The strtoimax function is equivalent to strtol, except that the initial portion of the string is converted to intmax_t representation.

## Returns

3 The strtoimax function returns the converted value, if any. If no conversion could be performed zero is returned. If the correct value is outside the range of representable values, INTMAX_MAX or INTMAX_MIN is returned (according to the sign of the value), and the value of the macro ERANGE is stored in errno.

### 7.4.6.2 The strtoumax function

Synopsis

```
#include <inttypes.h>
uintmax_t strtoumax(const char * restrict nptr,
    char ** restrict endptr, int base);
```


## Description

2 The strtoumax function is equivalent to strtoul, except that the initial portion of the string is converted to uintmax_t representation.

## Returns

3 The strtoumax function returns the converted value, if any. If no conversion could be performed zero is returned. If the correct value is outside the range of representable values, UINTMAX_MAX is returned, and the value of the macro ERANGE is stored in errno.

### 7.4.6.3 The westoimax function

Synopsis

```
#include <stddef.h> // for wchar_t
#include <inttypes.h>
intmax_t wcstoimax(const wchar_t * restrict nptr,
    wchar_t ** restrict endptr, int base);
```


## Description

2 The wcstoimax function is equivalent to wcstol, except that the initial portion of the wide string is converted to intmax_t representation.

## Returns

3 The wcstoimax function returns the converted value, if any. If no conversion could be performed zero is returned. If the correct value is outside the range of representable values, INTMAX_MAX or INTMAX_MIN is returned (according to the sign of the value), and the value of the macro ERANGE is stored in errno.

### 7.4.6.4 The westoumax function

Synopsis

```
#include <stddef.h> // for wchar_t
#include <inttypes.h>
uintmax_t wcstoumax(const wchar_t * restrict nptr,
    wchar_t ** restrict endptr, int base);
```


## Description

2 The wcstoumax function is equivalent to wcstoul, except that the initial portion of the wide string is converted to uintmax_t representation.

## Returns

3 The wcstoumax function returns the converted value, if any. If no conversion could be performed zero is returned. If the correct value is outside the range of representable values, UINTMAX_MAX is returned, and the value of the macro ERANGE is stored in errno.

### 7.5 Localization <locale.h>

1 The header <locale.h> declares two functions, one type, and defines several macros.

2 The type is

```
struct lconv
```

which contains members related to the formatting of numeric values. The structure shall contain at least the following members, in any order. The semantics of the members and their normal ranges is explained in 7.5.2.1. In the "C" locale, the members shall have the values specified in the comments.

```
char *decimal_point; // "."
char *thousands_sep; // ""
char *grouping; // ""
char *int_curr_symbol; // ""
char *currency_symbol; // ""
char *mon_decimal_point; // ""
char *mon_thousands_sep; // ""
char *mon_grouping; // ""
char *positive_sign; // ""
char *negative_sign; // ""
char int_frac_digits; // CHAR_MAX
char frac_digits; // CHAR_MAX
char p_cs_precedes; // CHAR_MAX
char p_sep_by_space; // CHAR_MAX
char n_cs_precedes; // CHAR_MAX
char n_sep_by_space; // CHAR_MAX
char p_sign_posn; // CHAR_MAX
char n_sign_posn; // CHAR_MAX
```

3 The macros defined are NULL (described in 7.1.6); and ${ }^{153}$

[^75]```
IC_ALI
LC_COLIATE
LC_CTYPE
LC_MONETARY
LC_NUMERIC
LC_TIME
```

which expand to integer constant expressions with distinct values, suitable for use as the first argument to the setlocale function. Additional macro definitions, beginning with the characters LC_ and an uppercase letter, ${ }^{154}$ may also be specified by the implementation.

### 7.5.1 Locale control

### 7.5.1.1 The setlocale function

## Synopsis

```
#include <locale.h>
char *setlocale(int category, const char *locale);
```


## Description

2 The setlocale function selects the appropriate portion of the program's locale as specified by the category and locale arguments. The setlocale function may be used to change or query the program's entire current locale or portions thereof. The value LC_ALL for category names the program's entire locale; the other values for category name only a portion of the program's locale. LC_COLLATE affects the behavior of the strcoll and strxfrm functions. LC_CTYPE affects the behavior of the character handling functions ${ }^{155}$ and the multibyte functions. LC_MONETARY affects the monetary formatting information returned by the localeconv function. LC_NUMERIC affects the decimal-point character for the formatted input/output functions and the string conversion functions, as well as the nonmonetary formatting information returned by the localeconv function. LC_TIME affects the behavior of the strftime function.

3 A value of " C " for locale specifies the minimal environment for C translation; a value of "" for locale specifies the locale-specific native environment. Other implementation-defined strings may be passed as the second argument to setlocale.

[^76]4 At program startup, the equivalent of

```
setlocale(LC_ALL, "C");
```

is executed.
5 The implementation shall behave as if no library function calls the setlocale function.

## Returns

6 If a pointer to a string is given for locale and the selection can be honored, the setlocale function returns a pointer to the string associated with the specified category for the new locale. If the selection cannot be honored, the setlocale function returns a null pointer and the program's locale is not changed.

7 A null pointer for locale causes the setlocale function to return a pointer to the string associated with the category for the program's current locale; the program's locale is not changed. ${ }^{156}$

8 The pointer to string returned by the setlocale function is such that a subsequent call with that string value and its associated category will restore that part of the program's locale. The string pointed to shall not be modified by the program, but may be overwritten by a subsequent call to the setlocale function.

Forward references: formatted input/output functions (7.13.6), the multibyte character functions (7.14.7), the multibyte string functions (7.14.8), string conversion functions (7.14.1), the strcoll function (7.15.4.3), the strftime function (7.16.3.6), the strxfrm function (7.15.4.5).

### 7.5.2 Numeric formatting convention inquiry

### 7.5.2. The localeconv function

## Synopsis

```
#include <locale.h>
struct lconv *localeconv(void);
```

[^77]
## Description

2 The localeconv function sets the components of an object with type struct lconv with values appropriate for the formatting of numeric quantities (monetary and otherwise) according to the rules of the current locale.

3 The members of the structure with type char * are pointers to strings, any of which (except decimal_point) can point to "", to indicate that the value is not available in the current locale or is of zero length. Apart from grouping and mon_grouping, the strings shall start and end in the initial shift state. The members with type char are nonnegative numbers, any of which can be CHAR_MAX to indicate that the value is not available in the current locale. The members include the following:

## char *decimal_point

The decimal-point character used to format nonmonetary quantities.

## char *thousands_sep

The character used to separate groups of digits before the decimal-point character in formatted nonmonetary quantities.

## char *grouping

A string whose elements indicate the size of each group of digits in formatted nonmonetary quantities.
char *int_curr_symbol
The international currency symbol applicable to the current locale. The first three characters contain the alphabetic international currency symbol in accordance with those specified in ISO 4217:1987. The fourth character (immediately preceding the null character) is the character used to separate the international currency symbol from the monetary quantity.

## char *currency_symbol

The local currency symbol applicable to the current locale.
char *mon_decimal_point
The decimal-point used to format monetary quantities.
char *mon_thousands_sep
The separator for groups of digits before the decimal-point in formatted monetary quantities.
char *mon_grouping
A string whose elements indicate the size of each group of digits in formatted monetary quantities.

## char *positive_sign

The string used to indicate a nonnegative-valued formatted monetary quantity.

## char *negative_sign

The string used to indicate a negative-valued formatted monetary quantity.

## char int_frac_digits

The number of fractional digits (those after the decimal-point) to be displayed in an internationally formatted monetary quantity.

## char frac_digits

The number of fractional digits (those after the decimal-point) to be displayed in a formatted monetary quantity.

## char p_cs_precedes

Set to 1 or 0 if the currency_symbol respectively precedes or succeeds the value for a nonnegative formatted monetary quantity.

## char p_sep_by_space

Set to 1 or 0 if the currency_symbol respectively is or is not separated by a space from the value for a nonnegative formatted monetary quantity.

## char n_cs_precedes

Set to 1 or 0 if the currency_symbol respectively precedes or succeeds the value for a negative formatted monetary quantity.

## char n_sep_by_space

Set to 1 or 0 if the currency_symbol respectively is or is not separated by a space from the value for a negative formatted monetary quantity.

## char p_sign_posn

Set to a value indicating the positioning of the positive_sign for a nonnegative formatted monetary quantity.

## char n_sign_posn

Set to a value indicating the positioning of the negative_sign for a negative formatted monetary quantity.

4 The elements of grouping and mon_grouping are interpreted according to the following:

CHAR_MAX No further grouping is to be performed.
$0 \quad$ The previous element is to be repeatedly used for the remainder of the digits.
other The integer value is the number of digits that comprise the current group. The next element is examined to determine the size of the next group of digits before the current group.

5 The value of p_sign_posn and n_sign_posn is interpreted according to the following:

0 Parentheses surround the quantity and currency_symbol.
1 The sign string precedes the quantity and currency_symbol.
2 The sign string succeeds the quantity and currency_symbol.
3 The sign string immediately precedes the currency_symbol.
4 The sign string immediately succeeds the currency_symbol.
6 The implementation shall behave as if no library function calls the localeconv function.

## Returns

7 The localeconv function returns a pointer to the filled-in object. The structure pointed to by the return value shall not be modified by the program, but may be overwritten by a subsequent call to the localeconv function. In addition, calls to the setlocale function with categories LC_ALL, LC_MONETARY, or LC_NUMERIC may overwrite the contents of the structure.

## Examples

8 The following table illustrates the rules which may well be used by four countries to format monetary quantities.

| Country | Positive format | Negative format | International format |
| :--- | :--- | :--- | :--- |
|  |  |  |  |
| Italy | L.1.234 | -L.1.234 | ITL.1.234 |
| Netherlands | F 1.234,56 | F -1.234,56 | NLG 1.234,56 |
| Norway | kr1.234,56 | kr1.234,56- | NOK 1.234,56 |
| Switzerland | SFrs.1,234.56 | SFrs.1,234.56C | CHF 1,234.56 |

9 For these four countries, the respective values for the monetary members of the structure returned by localeconv are:

|  | Italy | Netherlands | Norway | Switzerland |
| :---: | :---: | :---: | :---: | :---: |
| int_curr_symbol | "ITL. ${ }^{\text {P }}$ | "NLG " | "NOK " | "CHF " |
| currency_symbol | "L." | "F" | "kr" | "SFrs." |
| mon_decimal_point | " | ", " | ", " | "." |
| mon_thousands_sep | ". " | "." | "." | ", " |
| mon_grouping | "\3" | "\3" | "\3" | "\3" |
| positive_sign | " " | " ${ }^{\prime \prime}$ | " ${ }^{\prime \prime}$ | " " |
| negative_sign | "-" | "-" | "-" | "C" |
| int_frac_digits | 0 | 2 | 2 | 2 |
| frac_digits | 0 | 2 | 2 | 2 |
| p_cs_precedes | 1 | 1 | 1 | 1 |
| p_sep_by_space | 0 | 1 | 0 | 0 |
| n_cs_precedes | 1 | 1 | 1 | 1 |
| n_sep_by_space | 0 | 1 | 0 | 0 |
| p_sign_posn | 1 | 1 | 1 | 1 |
| n_sign_posn | 1 | 4 | 2 | 2 |

### 7.6 Floating-point environment <fenv.h>

1 The header <fenv.h> declares two types and several macros and functions to provide access to the floating-point environment. The floating-point environment refers collectively to any floating-point status flags and control modes supported by the implementation. ${ }^{157}$ A floating-point status flag is a system variable whose value is set as a side effect of the arithmetic to provide auxiliary information. A floating-point control mode is a system variable whose value may be set by the user to affect the subsequent behavior of the arithmetic.

2 Certain programming conventions support the intended model of use for the floatingpoint environment: ${ }^{158}$

- a function call must not alter its caller's modes, clear its caller's flags, nor depend on the state of its caller's flags unless the function is so documented;

157. This header is designed to support the exception status flags and directed-rounding control modes required by IEC 559, and other similar floating-point state information. Also it is designed to facilitate code portability among all systems.
158. With these conventions, a programmer can safely assume default modes (or be unaware of them). The responsibilities associated with accessing the floating-point environment fall on the programmer or program that does so explicitly.

- a function call is assumed to require default modes, unless its documentation promises otherwise or unless the function is known not to use floating-point;
- a function call is assumed to have the potential for raising floating-point exceptions, unless its documentation promises otherwise, or unless the function is known not to use floating-point.

3 The type
fenv_t
represents the entire floating-point environment.
4 The type
fexcept_t
represents the floating-point exception flags collectively, including any status the implementation associates with the flags.

5 Each of the macros
FE_DIVBYZERO
FE_INEXACT
FE_INVALID
FE_OVERFLOW
FE_UNDERFLOW
is defined if and only if the implementation supports the exception by means of the functions in 7.6.2. The defined macros expand to integer constant expressions with values such that bitwise ORs of all combinations of the macros result in distinct values.

6 The macro

## FE_ALL_EXCEPT

is simply the bitwise OR of all exception macros defined by the implementation.
7 Each of the macros
FE_DOWNWARD
FE_TONEAREST
FE_TOWARDZERO
FE_UPWARD
is defined if and only if the implementation supports getting and setting the represented rounding direction by means of the fegetround and fesetround functions. The defined macros expand to integer constant expressions whose values are distinct nonnegative values. ${ }^{159}$
Library

The macro

```
FE_DFL_ENV
```

represents the default floating-point environment - the one installed at program startup - and has type pointer to const-qualified fenv_t. It can be used as an argument to <fenv.h> functions that manage the floating-point environment.

9 Additional macro definitions, beginning with $\mathbf{F E}$ _ and having type pointer to constqualified fenv_t, may also be specified by the implementation.

### 7.6.1 The FENV_ACCESS pragma

## Synopsis

```
#include <fenv.h>
    #pragma STDC FENV_ACCESS on-off-switch
```


## Description

2 The FENV_ACCESS pragma provides a means to inform the implementation when a program might access the floating-point environment to test flags or run under nondefault modes. ${ }^{160}$ The pragma can occur either outside external declarations or preceding all explicit declarations and statements inside a compound statement. When outside external declarations, the pragma takes effect from its occurrence until another FENV_ACCESS pragma is encountered, or until the end of the translation unit. When inside a compound statement, the pragma takes effect from its occurrence until another FENV_ACCESS pragma is encountered (within a nested compound statement), or until the end of the compound statement; at the end of a compound statement the state for the pragma is restored to its condition just before the compound statement. If this pragma is used in any other context, the behavior is undefined. If part of a program tests flags or runs under non-default mode settings, but was translated with the state for the FENV_ACCESS pragma off, then the behavior of that program is undefined. The default state (on or off) for the pragma is implementation-defined.

[^78]
## Examples

```
#include <fenv.h>
void f(double x)
{
    #pragma STDC FENV_ACCESS ON
    void g(double);
    void h(double);
    /* ... */
    g(x + 1);
    h(x + 1);
    /* ... */
}
```

If the function $\mathbf{g}$ might depend on status flags set as a side effect of the first $\mathbf{x}+\mathbf{1}$, or if the second $\mathbf{x}+1$ might depend on control modes set as a side effect of the call to function $\mathbf{g}$, then the program must contain an appropriately placed invocation of \#pragma STDC FENV_ACCESS ON. ${ }^{161}$

### 7.6.2 Exceptions

1 The following functions provide access to the exception flags. ${ }^{162}$ The int input argument for the functions represents a subset of floating-point exceptions, and can be constructed by bitwise ORs of the exception macros, for example FE_OVERFLOW FE_INEXACT. For other argument values the behavior of these functions is undefined.

### 7.6.2.1 The feclearexcept function

Synopsis

```
#include <fenv.h>
void feclearexcept(int excepts);
```

[^79]
## Description

2 The feclearexcept function clears the supported exceptions represented by its argument.

### 7.6.2.2 The fegetexceptflag function

Synopsis

```
#include <fenv.h>
void fegetexceptflag(fexcept_t *flagp,
    int excepts);
```


## Description

2 The fegetexceptflag function stores an implementation-defined representation of the exception flags indicated by the argument excepts in the object pointed to by the argument flagp.

### 7.6.2.3 The feraiseexcept function

Synopsis

```
#include <fenv.h>
void feraiseexcept(int excepts);
```


## Description

2 The feraiseexcept function raises the supported exceptions represented by its argument. ${ }^{163}$ The order in which these exceptions are raised is unspecified, except as stated in F.7.6. Whether the feraiseexcept function additionally raises the inexact exception whenever it raises the overflow or underflow exception is implementation-defined.

### 7.6.2.4 The fesetexceptflag function

Synopsis

```
#include <fenv.h>
void fesetexceptflag(const fexcept_t *flagp,
        int excepts);
```

[^80]
## Description

2 The fesetexceptflag function sets the complete status for those exception flags indicated by the argument excepts, according to the representation in the object pointed to by flagp. The value of $\boldsymbol{*} \mathbf{f l a g p}$ must have been set by a previous call to fegetexceptflag whose second argument represented at least those exceptions represented by the argument excepts; if not, the effect on the indicated exception flags is undefined. This function does not raise exceptions, but only sets the state of the flags.

### 7.6.2.5 The fetestexcept function

## Synopsis

```
#include <fenv.h>
int fetestexcept(int excepts);
```


## Description

2 The fetestexcept function determines which of a specified subset of the exception flags are currently set. The excepts argument specifies the exception flags to be queried. ${ }^{164}$

## Returns

3 The fetestexcept function returns the value of the bitwise OR of the exception macros corresponding to the currently set exceptions included in excepts.

## Examples

4 Call $\mathbf{f}$ if invalid is set, then $\mathbf{g}$ if overflow is set:

[^81]```
#include <fenv.h>
/* ... */
{
    #pragma STDC FENV_ACCESS ON
    int set_excepts;
    // maybe raise exceptions
    set_excepts =
                fetestexcept (FE_INVALID | FE_OVERFLOW);
    if (set_excepts & FE_INVALID) f();
    if (set_excepts & FE_OVERFLOW) g();
    /* ... */
}
```


### 7.6.3 Rounding

1 The fegetround and fesetround functions provide control of rounding direction modes.

### 7.6.3.1 The fegetround function

Synopsis
1 \#include <fenv.h>
int fegetround (void);
Description
2 The fegetround function gets the current rounding direction.

## Returns

3 The fegetround function returns the value of the rounding direction macro representing the current rounding direction.

### 7.6.3.2 The fesetround function

Synopsis
1 \#include <fenv.h>
int fesetround(int round);
Description
2 The fesetround function establishes the rounding direction represented by its argument round. If the argument does not match a rounding direction macro, the rounding direction is not changed.

## Returns

3 The fesetround function returns a nonzero value if and only if the argument matches a rounding direction macro (that is, if and only if the requested rounding direction can be established).

## Examples

4 Save, set, and restore the rounding direction. Report an error and abort if setting the rounding direction fails.

```
#include <fenv.h>
#include <assert.h>
/* ... */
{
    #pragma STDC FENV_ACCESS ON
    int save_round;
    int setround_ok;
    save_round = fegetround();
    setround_ok = fesetround(FE_UPWARD);
    assert(setround_ok);
    /* ... */
    fesetround(save_round);
    /* ... */
}
```


### 7.6.4 Environment

1 The functions in this section manage the floating-point environment - status flags and control modes - as one entity.

### 7.6.4.1 The fegetenv function

Synopsis

```
#include <fenv.h>
void fegetenv(fenv_t *envp);
```


## Description

2 The fegetenv function stores the current floating-point environment in the object pointed to by envp.

### 7.6.4.2 The feholdexcept function

Synopsis

```
#include <fenv.h>
int feholdexcept(fenv_t *envp);
```


## Description

2 The feholdexcept function saves the current environment in the object pointed to by envp, clears the exception flags, and installs a non-stop (continue on exceptions) mode, if available, for all exceptions. ${ }^{165}$

## Returns

3 The feholdexcept function returns nonzero if and only if non-stop exception handling was successfully installed.

### 7.6.4.3 The fesetenv function

## Synopsis

```
#include <fenv.h>
void fesetenv(const fenv_t *envp);
```


## Description

2 The fesetenv function establishes the floating-point environment represented by the object pointed to by envp. The argument envp must point to an object set by a call to fegetenv or feholdexcept, or equal the macro FE_DFL_ENV or an implementation-defined environment macro. Note that fesetenv merely installs the state of the exception flags represented through its argument, and does not raise these exceptions.

[^82]
### 7.6.4.4 The feupdateenv function

Synopsis

```
#include <fenv.h>
void feupdateenv(const fenv_t *envp);
```


## Description

2 The feupdateenv function saves the currently raised exceptions in its automatic storage, installs the environment represented through envp, and then raises the saved exceptions. The argument envp must point to an object set by a call to feholdexcept or fegetenv, or equal the macro FE_DFL_ENV or an implementation-defined environment macro.

## Examples

3 Hide spurious underflow exceptions:

```
#include <fenv.h>
double f(double x)
{
    #pragma STDC FENV_ACCESS ON
    double result;
    fenv_t save_env;
    feholdexcept (&save_env);
    // compute result
    if (/* test spurious underflow */)
            feclearexcept (FE_UNDERFLOW);
    feupdateenv(&save_env);
    return result;
}
```


### 7.7 Mathematics <math.h>

1 The header <math.h> declares two types and several mathematical functions and defines several macros. Most synopses specify a function which takes one or more double arguments and returns a double value; for each such function, there are functions with the same name but with $\mathbf{f}$ and $\mathbf{l}$ suffixes which are corresponding functions with float and long double arguments and return values. ${ }^{166}$ Some synopses specify a function which takes a double argument and returns an integertype value; for each such function, there are functions with the same name but with $\mathbf{f}$ and 1 suffixes which are corresponding functions with float and long double arguments. Integer arithmetic functions and conversion functions are discussed later.

2 The types

```
float_t
double_t
```

are floating types at least as wide as float and double, respectively, and such that double_t is at least as wide as float_t. If FLT_EVAL_METHOD equals 0, float_t and double_t are float and double, respectively; if FLT_EVAL_METHOD equals 1, they are both double; if FLT_EVAL_METHOD equals 2, they are both long double; and for other values of FLT_EVAL_METHOD, they are otherwise implementation-defined. ${ }^{167}$

The macro

## HUGE_VAL

expands to a positive double constant expression, not necessarily representable as a float. The macros

```
HUGE_VALF
HUGE_VALL
```

are respectively float and long double analogs of HUGE_VAL. ${ }^{168}$
166. Particularly on systems with wide expression evaluation, a <math.h> function might pass arguments and return values in wider format than the synopsis prototype indicates.
167. The types float_t and double_t are intended to be the implementation's most efficient types at least as wide as float and double, respectively. For FLT_EVAL_METHOD equal 0, 1, or 2, the type float_t is the narrowest type used by the implementation to evaluate floating expressions.
168. HUGE_VAL, HUGE_VALF, and HUGE_VALL can be positive infinities in an implementation that supports infinities.

## INFINITY

expands to a constant expression of type float representing an implementationdefined positive or unsigned infinity, if available, else to a positive constant of type float that overflows at translation time.

5 The macro
NAN
is defined if and only if the implementation supports quiet NaNs for the float type. It expands to a constant expression of type float representing an implementationdefined quiet NaN .

The macros

```
FP_INFINITE
FP_NAN
FP_NORMAL
FP_SUBNORMAL
FP_ZERO
```

are for number classification. They represent the mutually exclusive kinds of floatingpoint values. They expand to integer constant expressions with distinct values.

7 The macro

```
FP_FAST_FMA
```

is optionally defined. If defined, it indicates the fma function generally executes about as fast as a multiply and an add of double operands. ${ }^{169}$ The macros

```
FP_FAST_FMAF
FP_FAST_FMAL
```

are, respectively, float and long double analogs of FP_FAST_FMA.

[^83]```
FP_ILOGB0
FP_ILOGBNAN
```

expand to integer constant expressions whose values are returned by ilogb( $\mathbf{x}$ ) if $\mathbf{x}$ is zero or NaN, respectively. The value of FP_ILOGBO shall be either INT_MIN or -INT_MAX. The value of FP_ILOGBNAN shall be either INT_MAX or INT_MIN.

9 The macro

## DECIMAL_DIG

expands to an integer constant expression whose value is implementation-defined. It represents a number of decimal digits supported by conversion between decimal and all internal floating-point formats. ${ }^{170}$

## Recommended practice

10 Conversion from (at least) double to decimal with DECIMAL_DIG digits and back is the identity function. ${ }^{171}$

### 7.7.1 Treatment of error conditions

1 The behavior of each of the functions in <math. h> is specified for all representable values of its input arguments, except where stated otherwise.

2 For all functions, a domain error occurs if an input argument is outside the domain over which the mathematical function is defined. The description of each function lists any required domain errors; an implementation may define additional domain errors, provided that such errors are consistent with the mathematical definition of the function. ${ }^{172}$ On a domain error, the function returns an implementation-defined value; whether the integer expression errno acquires the value EDOM is implementationdefined.

3 Similarly, a range error occurs if the mathematical result of the function cannot be represented in an object of the specified type, due to extreme magnitude. A floating

[^84]result overflows if the magnitude of the mathematical result is finite but so large that the mathematical result cannot be represented, without extraordinary roundoff error, in an object of the specified type. If a floating result overflows and default rounding is in effect, or if the mathematical result is an exact infinity (for example $\log (0.0)$ ), then the function returns the value of the macro HUGE_VAL, HUGE_VALF, or HUGE_VALL appropriate to the specified result type, with the same sign as the correct value of the function; whether errno acquires the value ERANGE when a range error occurs is implementation-defined. The result underflows if the magnitude of the mathematical result is so small that the mathematical result cannot be represented, without extraordinary roundoff error, in an object of the specified type. ${ }^{173}$ If the result underflows, the function returns a value whose magnitude is no greater than the smallest normalized positive number in the specified type and is otherwise implementation-defined; whether errno acquires the value ERANGE is implementation-defined.

### 7.7.2 The FP_CONTRACT pragma

## Synopsis

```
#include <math.h>
#pragma STDC FP_CONTRACT on-off-switch
```


## Description

The FP_CONTRACT pragma can be used to allow (if the state is on) or disallow (if the state is off) the implementation to contract expressions (6.3). Each pragma can occur either outside external declarations or preceding all explicit declarations and statements inside a compound statement. When outside external declarations, the pragma takes effect from its occurrence until another FP_CONTRACT pragma is encountered, or until the end of the translation unit. When inside a compound statement, the pragma takes effect from its occurrence until another FP_CONTRACT pragma is encountered (within a nested compound statement), or until the end of the compound statement; at the end of a compound statement the state for the pragma is restored to its condition just before the compound statement. If this pragma is used in any other context, the behavior is undefined. The default state (on or off) for the pragma is implementation-defined.

[^85]
### 7.7.3 Classification macros

1 In the synopses in this subclause, real-floating indicates that the argument must be an expression of real floating type. The result is undefined if an argument is not of real floating type.

### 7.7.3.1 The fpclassify macro

Synopsis

```
#include <math.h>
int fpclassify(real-floating x);
```


## Description

2 The fpclassify macro classifies its argument value as NaN , infinite, normal, subnormal, or zero. First, an argument represented in a format wider than its semantic type is converted to its semantic type. Then classification is based on the type of the argument. ${ }^{174}$

## Returns

3 The fpclassify macro returns the value of the number classification macro appropriate to the value of its argument.

## Examples

4 The fpclassify macro might be implemented in terms of ordinary functions as

```
    #define fpclassify(x) \
            ((sizeof (x) == sizeof (float)) ? \
                fpclassifyf(x) \
            : (sizeof (x) == sizeof (double)) ? \
                fpclassifyd(x) \
    : __fpclassifyl(x))
```

[^86]
### 7.7.3.2 The signbit macro

Synopsis

```
#include <math.h>
int signbit(real-floating x);
```


## Description

2 The signbit macro determines whether the sign of its argument value is negative. ${ }^{175}$

## Returns

3 The signbit macro returns a nonzero value if and only if the sign of its argument value is negative.

### 7.7.3.3 The isfinite macro

Synopsis

```
#include <math.h>
int isfinite(real-floating x);
```


## Description

2 The isfinite macro determines whether its argument has a finite value (zero, subnormal, or normal, and not infinite or NaN). First, an argument represented in a format wider than its semantic type is converted to its semantic type. Then determination is based on the type of the argument.

## Returns

3 The isfinite macro returns a nonzero value if and only if its argument has a finite value.

### 7.7.3.4 The isnormal macro

Synopsis

```
#include <math.h>
int isnormal(real-floating x);
```

175. The signbit macro reports the sign of all values, including infinities, zeros, and NaNs.

## Description

2 The isnormal macro determines whether its argument value is normal (neither zero, subnormal, infinite, nor NaN). First, an argument represented in a format wider than its semantic type is converted to its semantic type. Then determination is based on the type of the argument.

## Returns

3 The isnormal macro returns a nonzero value if and only if its argument has a normal value.

### 7.7.3.5 The isnan macro

## Synopsis

```
#include <math.h>
int isnan(real-floating x);
```


## Description

2 The isnan macro determines whether its argument value is a NaN. First, an argument represented in a format wider than its semantic type is converted to its semantic type. Then determination is based on the type of the argument. ${ }^{176}$

## Returns

3 The isnan macro returns a nonzero value if and only if its argument has a NaN value.

### 7.7.3.6 The isinf macro

## Synopsis

```
#include <math.h>
int isinf(real-floating x);
```


## Description

2 The isinf macro determines whether its argument value is an infinity (positive or negative). First, an argument represented in a format wider than its semantic type is converted to its semantic type. Then determination is based on the type of the argument.

[^87]
## Returns

3 The isinf macro returns a nonzero value if and only if its argument has an infinite value.

### 7.7.4 Trigonometric functions

### 7.7.4.1 The acos function

Synopsis

```
#include <math.h>
double acos(double x);
```


## Description

2 The acos function computes the principal value of the arc cosine of $\mathbf{x}$. A domain error occurs for arguments not in the range $[-1,+1]$.

## Returns

3 The acos function returns the arc cosine in the range $[0, \pi]$ radians.

### 7.7.4.2 The asin function

Synopsis

```
#include <math.h>
double asin(double x);
```


## Description

2 The asin function computes the principal value of the arc sine of $\mathbf{x}$. A domain error occurs for arguments not in the range $[-1,+1]$.

## Returns

3 The asin function returns the arc sine in the range $[-\pi / 2,+\pi / 2]$ radians.

### 7.7.4.3 The atan function

Synopsis

```
#include <math.h>
double atan(double x);
```


## Description

2 The atan function computes the principal value of the arc tangent of $\mathbf{x}$.

## Returns

3 The atan function returns the arc tangent in the range $[-\pi / 2,+\pi / 2]$ radians.

### 7.7.4. 4 The atan2 function

Synopsis
1 \#include <math.h>
double atan2 (double $y$, double $x$ );

## Description

2 The atan2 function computes the principal value of the arc tangent of $\mathbf{y} / \mathbf{x}$, using the signs of both arguments to determine the quadrant of the return value. A domain error may occur if both arguments are zero.

## Returns

3 The atan2 function returns the arc tangent of $\mathbf{y} / \mathbf{x}$, in the range $[-\pi,+\pi]$ radians.

### 7.7.4.5 The cos function

## Synopsis

```
#include <math.h>
double cos(double x);
```


## Description

2 The cos function computes the cosine of $\mathbf{x}$ (measured in radians).

## Returns

3 The cos function returns the cosine value.

### 7.7.4.6 The $\sin$ function

Synopsis

```
#include <math.h>
double sin(double x);
```


## Description

2 The $\boldsymbol{\operatorname { s i n }}$ function computes the sine of $\mathbf{x}$ (measured in radians).

## Returns

3 The $\boldsymbol{\operatorname { s i n }}$ function returns the sine value.

### 7.7.4.7 The $\tan$ function

Synopsis

```
#include <math.h>
double tan(double x);
```

Description
2 The tan function returns the tangent of $\mathbf{x}$ (measured in radians).

## Returns

3 The $\tan$ function returns the tangent value.

### 7.7.5 Hyperbolic functions

### 7.7.5.1 The cosh function

Synopsis

```
#include <math.h>
double cosh(double x);
```


## Description

2 The cosh function computes the hyperbolic cosine of $\mathbf{x}$. A range error occurs if the magnitude of $\mathbf{x}$ is too large.

## Returns

3 The cosh function returns the hyperbolic cosine value.

### 7.7.5.2 The sinh function

Synopsis

```
#include <math.h>
double sinh(double x);
```


## Description

2 The sinh function computes the hyperbolic sine of $\mathbf{x}$. A range error occurs if the magnitude of $\mathbf{x}$ is too large.

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## Returns

3 The sinh function returns the hyperbolic sine value.

### 7.7.5.3 The tanh function

Synopsis
1 \#include <math.h>

```
double tanh(double x);
```

Description
2 The tanh function computes the hyperbolic tangent of $\mathbf{x}$.

## Returns

3 The tanh function returns the hyperbolic tangent value.

### 7.7.5.4 The acosh function

Synopsis
1

```
#include <math.h>
double acosh(double x);
```


## Description

2 The acosh function computes the (nonnegative) arc hyperbolic cosine of $\mathbf{x}$. A domain error occurs for arguments less than 1.

## Returns

3 The acosh function returns the arc hyperbolic cosine in the range [ $0,+\infty$ ].

### 7.7.5.5 The asinh function

Synopsis

```
#include <math.h>
double asinh(double x);
```


## Description

The asinh function computes the arc hyperbolic sine of $\mathbf{x}$.

## Returns

3 The asinh function returns the arc hyperbolic sine value.

### 7.7.5.6 The atanh function

Synopsis

```
#include <math.h>
double atanh(double x);
```

Description
2 The atanh function computes the arc hyperbolic tangent of $\mathbf{x}$. A domain error occurs for arguments not in the range $[-1,+1]$.

## Returns

3 The atanh function returns the arc hyperbolic tangent value.

### 7.7.6 Exponential and logarithmic functions

### 7.7.6. The exp function

Synopsis
1 \#include <math.h>
double exp(double x);
Description
2 The exp function computes the exponential function of $\mathbf{x}$. A range error occurs if the magnitude of $\mathbf{x}$ is too large.

## Returns

3 The exp function returns the exponential value.

### 7.7.6.2 The frexp function

Synopsis

```
#include <math.h>
double frexp(double value, int *exp);
```

Description
2 The frexp function breaks a floating-point number into a normalized fraction and an integral power of 2. It stores the integer in the int object pointed to by exp.

## Returns

3 The frexp function returns the value $\mathbf{x}$, such that $\mathbf{x}$ is a double with magnitude in the interval $\left[1 / 2,1\right.$ ) or zero, and value equals $\mathbf{x} \times 2^{* e x p}$. If value is zero, both parts of the result are zero.

### 7.7.6.3 The ldexp function

Synopsis

2 The ldexp function multiplies a floating-point number by an integral power of 2. A range error may occur.

## Returns

3 The ldexp function returns the value of $\mathbf{x} \times 2^{\exp }$.
7.7.6.4 The $\log$ function

Synopsis

```
#include <math.h>
double log(double x);
```


## Description

2 The $\log$ function computes the natural logarithm of $\mathbf{x}$. A domain error occurs if the argument is negative. A range error may occur if the argument is zero.

## Returns

3 The log function returns the natural logarithm value.
7.7.6.5 The $\log 10$ function

Synopsis

```
#include <math.h>
    double log10(double x);
```


## Description

2 The log10 function computes the base-ten logarithm of $\mathbf{x}$. A domain error occurs if the argument is negative. A range error may occur if the argument is zero.

## Returns

3 The log10 function returns the base-ten logarithm value.

### 7.7.6.6 The modf function

Synopsis

2 The modf function breaks the argument value into integral and fractional parts, each of which has the same sign as the argument. It stores the integral part as a double in the object pointed to by iptr.

## Returns

3 The modf function returns the value of the signed fractional part of value.

### 7.7.6.7 The exp2 function

Synopsis
1 \#include <math.h>
double exp2 (double x);

## Description

2 The exp2 function computes the base-2 exponential of $\mathbf{x}$ : $2^{\mathbf{x}}$. A range error occurs if the magnitude of $\mathbf{x}$ is too large.

## Returns

3 The exp2 function returns the base-2 exponential value.

### 7.7.6.8 The expm1 function

Synopsis

```
#include <math.h>
double expm1 (double x);
```


## Description

2 The expm1 function computes the base-e exponential of the argument, minus 1: $e^{\mathbf{x}}-1 .{ }^{177}$ A range error occurs if $\mathbf{x}$ is too large.
177. For small magnitude $\mathbf{x}, \operatorname{expm} 1(\mathbf{x})$ is expected to be more accurate than $\exp (\mathbf{x})-1$.

## Returns

3 The expm1 function returns the value of $e^{\mathbf{x}}-1$.

### 7.7.6.9 The $\log 1 p$ function

Synopsis

```
#include <math.h>
double log1p(double x);
```


## Description

2 The $\log 1 \mathrm{p}$ function computes the base-e logarithm of 1 plus the argument. ${ }^{178} \mathrm{~A}$ domain error occurs if the argument is less than -1 . A range error may occur if the argument equals -1 .

## Returns

3 The $\log 1 \mathrm{p}$ function returns the value of the base-e logarithm of 1 plus the argument.

### 7.7.6.10 The $\log 2$ function

## Synopsis

```
#include <math.h>
double log2(double x);
```


## Description

2 The $\log 2$ function computes the base- 2 logarithm of $\mathbf{x}$. A domain error occurs if the argument is less than zero. A range error may occur if the argument is zero.

## Returns

3 The $\log 2$ function returns the base- 2 logarithm value.

### 7.7.6.11 The logb function

Synopsis

```
#include <math.h>
double logb(double x);
```

[^88]
## Description

2 The logb function extracts the exponent of $\mathbf{x}$, as a signed integer value in the format of $\mathbf{x}$. If $\mathbf{x}$ is subnormal it is treated as though it were normalized; thus for positive finite $\mathbf{x}$,

```
1\leqx }\times\mathrm{ FLT_RADIX }\mp@subsup{}{}{-logb(x)}< FLT_RADIX
```

A range error may occur if the argument is zero.

## Returns

3 The logb function returns the signed exponent of its argument.

### 7.7.6.12 The scalbn function

Synopsis

```
#include <math.h>
double scalbn(double x, int n);
```


## Description

2 The scalbn function computes $\mathbf{x} \times$ FLT_RADIX ${ }^{\text {n }}$ efficiently, not normally by computing FLT_RADIX ${ }^{\mathbf{n}}$ explicitly. A range error may occur.

## Returns

3 The scalbn function returns the value of $\mathbf{x} \times$ FLT_RADIX ${ }^{n}$.
7.7.6.13 The scalbln function

Synopsis

```
#include <math.h>
double scalbln(double x, long int n);
```


## Description

2 The scalbln function is equivalent to the scalbn function, except that the integer argument is long int.

### 7.7.6.14 The ilogb function

Synopsis

```
#include <math.h>
int ilogb(double x);
```


## Description

2 The ilogb function extracts the exponent of $\mathbf{x}$ as a signed int value. It is equivalent to (int) logb(x), for finite nonzero $\mathbf{x}$; it computes the value FP_ILOGBO if $\mathbf{x}$ is zero; it computes the value INT_MAX if $\mathbf{x}$ is infinite; and it computes the value $\mathbf{F P}$ _ILOGBNAN if $\mathbf{x}$ is a NaN. A range error may occur if $\mathbf{x}$ is 0 .

## Returns

3 The ilogb function returns the exponent of $\mathbf{x}$ as a signed int value.

### 7.7.7 Power and absolute value functions

### 7.7.7.1 The fabs function

Synopsis

```
#include <math.h>
double fabs(double x);
```


## Description

2 The fabs function computes the absolute value of a floating-point number $\mathbf{x}$.

## Returns

3 The fabs function returns the absolute value of $\mathbf{x}$.

### 7.7.7.2 The pow function

Synopsis

```
#include <math.h>
double pow(double x, double y);
```


## Description

2 The pow function computes $\mathbf{x}$ raised to the power $\mathbf{y}$. A domain error occurs if $\mathbf{x}$ is negative and $\mathbf{y}$ is finite and not an integer value. A domain error occurs if the result cannot be represented when $\mathbf{x}$ is zero and $\mathbf{y}$ is less than or equal to zero. A range error may occur.

## Returns

3 The pow function returns the value of $\mathbf{x}$ raised to the power $\mathbf{y}$.

### 7.7.7.3 The sqrt function

Synopsis

```
#include <math.h>
double sqrt(double x);
```


## Description

2 The sqrt function computes the nonnegative square root of $\mathbf{x}$. A domain error occurs if the argument is less than zero.

## Returns

3 The sqrt function returns the value of the square root.
7.7.7.4 The cbrt function

Synopsis

```
#include <math.h>
double cbrt(double x);
```


## Description

2 The cbrt function computes the real cube root of $\mathbf{x}$.

## Returns

3 The cbrt function returns the value of the cube root.

### 7.7.7.5 The hypot function

Synopsis

```
#include <math.h>
double hypot(double x, double y);
```


## Description

2 The hypot function computes the square root of the sum of the squares of $\mathbf{x}$ and $\mathbf{y}$, without undue overflow or underflow. A range error may occur.

## 3 <br> Returns

4 The hypot function returns the value of the square root of the sum of the squares.

### 7.7.8 Error and gamma functions

1 7.7.8.1 The erf function
Synopsis

```
#include <math.h>
double erf(double x);
```

Description
2 The erf function computes the error function of $\mathbf{x}$.

## Returns

3 The erf function returns the error function value.

### 7.7.8.2 The erfc function

Synopsis

```
#include <math.h>
double erfc(double x);
```


## Description

2 The erfc function computes the complementary error function of $\mathbf{x}$. A range error occurs if $\mathbf{x}$ is too large.

## Returns

3 The erfc function returns the complementary error function value.

### 7.7.8.3 The gamma function

Synopsis

```
#include <math.h>
double gamma(double x);
```


## Description

2 The gamma function computes the gamma function of $\mathbf{x}: \Gamma(\mathbf{x})$. A domain error occurs if $\mathbf{x}$ is a negative integer or zero. A range error may occur if the magnitude of $\mathbf{x}$ is too large.

## Returns

3 The gamma function returns the gamma function value.

### 7.7.8.4 The lgamma function

Synopsis

```
#include <math.h>
double lgamma(double x);
```


## Description

2 The lgamma function computes the natural logarithm of the absolute value of gamma of $\mathbf{x}: \log _{e}(|\Gamma(\mathbf{x})|)$. A range error occurs if $\mathbf{x}$ is too large or if $\mathbf{x}$ is a negative integer or zero.

## Returns

3 The lgamma function returns the value of the natural logarithm of the absolute value of gamma of $\mathbf{x}$.

### 7.7.9 Nearest integer functions

### 7.7.9.1 The ceil function

Synopsis

```
#include <math.h>
double ceil(double x);
```


## Description

2 The ceil function computes the smallest integer value not less than $\mathbf{x}$.

## Returns

3 The ceil function returns the smallest integer value not less than $\mathbf{x}$, expressed as a double.
7.7.9.2 The floor function

Synopsis

```
#include <math.h>
double floor(double x);
```

Description
The floor function computes the largest integer value not greater than $\mathbf{x}$.

## Returns

3 The floor function returns the largest integer value not greater than $\mathbf{x}$, expressed as a double.

### 7.7.9.3 The nearbyint function

Synopsis
1 \#include <math.h>
double nearbyint (double x);

## Description

2 The nearbyint function differs from the rint function (7.7.9.4) only in that the nearbyint function does not raise the inexact exception. (See F.9.6.3-F.9.6.4.)

## Returns

3 The nearbyint function returns the rounded integer value.

### 7.7.9.4 The rint function

## Synopsis

```
#include <math.h>
    double rint(double x);
```


## Description

2 The rint function rounds its argument to an integer value in floating-point format, using the current rounding direction.

## Returns

3 The rint function returns the rounded integer value.

### 7.7.9.5 The lrint function

Synopsis

```
#include <math.h>
    long int lrint(double x);
```


## Description

2 The lrint function rounds its argument to the nearest integer value, rounding according to the current rounding direction. If the rounded value is outside the range of long int, the numeric result is unspecified. A range error may occur if the magnitude of $\mathbf{x}$ is too large.

## Returns

3 The lrint function returns the rounded integer value, using the current rounding direction.

### 7.7.9.6 The llrint function

## Synopsis

1 \#include <math.h>
long long llrint (double x);

## Description

2 The llrint function is equivalent to the lrint function, except that the returned value has type long long.

### 7.7.9.7 The round function

Synopsis

```
#include <math.h>
double round(double x);
```


## Description

2 The round function rounds its argument to the nearest integer value in floating-point format, rounding halfway cases away from zero, regardless of the current rounding direction.

## Returns

3 The round function returns the rounded integer value.

### 7.7.9.8 The lround function

Synopsis

```
#include <math.h>
long int lround(double x);
```


## Description

2 The lround function rounds its argument to the nearest integer value, rounding halfway cases away from zero, regardless of the current rounding direction. If the rounded value is outside the range of long int, the numeric result is unspecified. A range error may occur if the magnitude of $\mathbf{x}$ is too large.

## Returns

3 The lround function returns the rounded integer value.

### 7.7.9.9 The llround function

Synopsis

```
#include <math.h>
long long llround(double x);
```


## Description

2 The llround function is equivalent to the lround function, except that the returned value has type long long.
7.7.9.10 The trunc function

Synopsis

```
#include <math.h>
double trunc(double x);
```


## Description

2 The trunc function rounds its argument to the integer value, in floating format, nearest to but no larger in magnitude than the argument.

## Returns

3 The trunc function returns the truncated integer value.

### 7.7.10 Remainder functions

### 7.7.10.1 The fmod function

Synopsis

```
#include <math.h>
double fmod(double x, double y);
```


## Description

2 The fmod function computes the floating-point remainder of $\mathbf{x} / \mathbf{y}$.

## Returns

3 The fmod function returns the value $\mathbf{x}-n \times \mathbf{y}$, for some integer $n$ such that, if $\mathbf{y}$ is nonzero, the result has the same sign as $\mathbf{x}$ and magnitude less than the magnitude of $\mathbf{y}$. If $\mathbf{y}$ is zero, whether a domain error occurs or the fmod function returns zero is
implementation-defined.

### 7.7.10.2 The remainder function

Synopsis

```
#include <math.h>
    double remainder(double x, double y);
```


## Description

2 The remainder function computes the remainder $\mathbf{x} R E M \mathbf{y}$ required by IEC 559. ${ }^{179}$

## Returns

3 The remainder function returns the value of $\mathbf{x} R E M \mathbf{y}$.

### 7.7.10.3 The remquo function

Synopsis

```
#include <math.h>
double remquo(double x, double y, int *quo);
```


## Description

2 The remquo function computes the same remainder as the remainder function. In the object pointed to by quo it stores a value whose sign is the sign of $\mathbf{x} / \mathbf{y}$ and whose magnitude is congruent $\bmod 2^{n}$ to the magnitude of the integral quotient of $\mathbf{x} / \mathbf{y}$, where $n$ is an implementation-defined integer at least 3 .

## Returns

3 The remquo function returns the value of $\mathbf{x} R E M \mathbf{y}$.

## 179. "When <br> $$
y \neq 0,
$$

the remainder $r=x$ REM $y$ is defined regardless of the rounding mode by the mathematical relation $r=x-y \times n$, where n is the integer nearest the exact value of $x / y$; whenever $|n-x / y|=1 / 2$, then $n$ is even. Thus, the remainder is always exact. If $r=0$, its sign shall be that of $x$." This definition is applicable for all implementations.

### 7.7.11 Manipulation functions

### 7.7.11.1 The copysign function

Synopsis

```
#include <math.h>
    double copysign(double x, double y);
```


## Description

2 The copysign function produces a value with the magnitude of $\mathbf{x}$ and the sign of $\mathbf{y}$. It produces a NaN (with the sign of $\mathbf{y}$ ) if $\mathbf{x}$ is a NaN . On implementations that represent a signed zero but do not treat negative zero consistently in arithmetic operations, the copysign function regards the sign of zero as positive.

## Returns

3 The copysign function returns a value with the magnitude of $\mathbf{x}$ and the sign of $\mathbf{y}$.

### 7.7.11.2 The nan function

## Synopsis

```
#include <math.h>
double nan(const char *tagp);
```


## Description

2 If the implementation supports quiet NaNs for the double type, then the call nan("n-char-sequence") is equivalent to strtod("NAN(n-char-sequence)", (char**) NULL); the call nan("") is equivalent to strtod("NAN()", (char**) NULL). If tagp does not point to an n-char-sequence string then the result NaN's content is unspecified. If the implementation does not support quiet NaNs for the double type, a call to the nan function is unspecified.

## Returns

3 The nan function returns a quiet NaN , if available, with content indicated through tagp.

### 7.7.11.3 The nextafter function

Synopsis

```
#include <math.h>
double nextafter(double x, double y);
```


## Description

2 The nextafter function determines the next representable value, in the type of the function, after $\mathbf{x}$ in the direction of $\mathbf{y}$, where $\mathbf{x}$ and $\mathbf{y}$ are first converted to the type of the function. ${ }^{180}$ The nextafter function returns $\mathbf{y}$ if $\mathbf{x}$ equals $\mathbf{y}$.

## Returns

3 The nextafter function returns the next representable value in the specified format after $\mathbf{x}$ in the direction of $\mathbf{y}$.

### 7.7.11.4 The nextafterx function

Synopsis

```
#include <math.h>
    double nextafterx(double x, long double y);
```


## Description

2 The nextafterx function is equivalent to the nextafter function except that the second parameter has type long double. ${ }^{181}$
7.7.12 Maximum, minimum, and positive difference functions

### 7.7.12. 1 The fdim function

Synopsis

```
#include <math.h>
double fdim(double x, double y);
```

Description
2 The fdim function determines the positive difference between its arguments:

```
x - y if x > y
+0 if x \leq y
```

A range error may occur.

[^89]
## Returns

3 The fdim function returns the positive difference value.

### 7.7.12.2 The fmax function

Synopsis

```
#include <math.h>
double fmax(double x, double y);
```


## Description

2 The fmax function determines the maximum numeric value of its arguments. ${ }^{182}$

## Returns

3 The fmax function returns the maximum numeric value of its arguments.

### 7.7.12.3 The fmin function

Synopsis

```
#include <math.h>
double fmin(double x, double y);
```


## Description

2 The fmin function determines the minimum numeric value of its arguments. ${ }^{183}$

## Returns

3 The fmin function returns the minimum numeric value of its arguments.
182. NaN arguments are treated as missing data: if one argument is a NaN and the other numeric, then fmax chooses the numeric value. See F.9.9.2.
183. fmin is analogous to fmax in its treatment of NaNs.

### 7.7.13 Floating multiply-add

### 7.7.13.1 The fma function

Synopsis

```
#include <math.h>
double fma(double x, double y, double z);
```


## Description

2 The fma function computes the sum $\mathbf{z}$ plus the product $\mathbf{x}$ times $\mathbf{y}$, rounded as one ternary operation: it computes the sum $\mathbf{z}$ plus the product $\mathbf{x}$ times $\mathbf{y}$ (as if) to infinite precision and rounds once to the result format, according to the rounding mode characterized by the value of FLT_ROUNDS.

## Returns

3 The fma function returns the sum $\mathbf{z}$ plus the product $\mathbf{x}$ times $\mathbf{y}$, rounded as one ternary operation.

### 7.7.14 Comparison macros

1 The relational and equality operators support the usual mathematical relationships between numeric values. For any ordered pair of numeric values exactly one of the relationships - less, greater, and equal - is true. Relational operators may raise the invalid exception when argument values are NaNs. For a NaN and a numeric value, or for two NaNs, just the unordered relationship is true. ${ }^{184}$ The following subclauses provide macros that are quiet (non exception raising) versions of the relational operators, and other comparison macros that facilitate writing efficient code that accounts for NaNs without suffering the invalid exception. In the synopses in this subclause, real-floating indicates that the argument must be an expression of real floating type.

[^90]
### 7.7.14.1 The isgreater macro

Synopsis

## Description

```
\#include <math.h>
int isgreater (real-floating \(\mathbf{x}\), real-floating \(\mathbf{y}\) );
```

The isgreater macro determines whether its first argument is greater than it second argument. The value of isgreater $(\mathbf{x}, \mathbf{y})$ is always equal to $(\mathbf{x})>(\mathbf{y})$; however, unlike $(\mathbf{x})>(\mathbf{y})$, isgreater $(\mathbf{x}, \mathbf{y})$ does not raise the invalid exception when $\mathbf{x}$ and $\mathbf{y}$ are unordered.

## Returns

3 The isgreater macro returns the value of $(\mathbf{x})>(\mathrm{y})$.

### 7.7.14.2 The isgreaterequal macro

Synopsis

```
#include <math.h>
int isgreaterequal(real-floating x, real-floating y);
```


## Description

2 The isgreaterequal macro determines whether its first argument is greater than or equal to its second argument. The value of isgreaterequal $(\mathbf{x}, \mathbf{y})$ is always equal to ( $\mathbf{x}$ ) $>=(\mathbf{y})$; however, unlike ( $\mathbf{x}$ ) >= ( $\mathbf{y}$ ), isgreaterequal ( $\mathbf{x}, \mathrm{y}$ ) does not raise the invalid exception when $\mathbf{x}$ and $\mathbf{y}$ are unordered.

## Returns

3 The isgreaterequal macro returns the value of $(x)>=(y)$.

### 7.7.14.3 The isless macro

Synopsis

```
#include <math.h>
    int isless(real-floating x, real-floating y);
```


## Description

2 The isless macro determines whether its first argument is less than its second argument. The value of isless $(\mathbf{x}, \mathbf{y})$ is always equal to $(\mathbf{x})<(\mathbf{y})$; however, unlike $(\mathbf{x})<(\mathbf{y})$, isless $(\mathbf{x}, \mathbf{y})$ does not raise the invalid exception when $\mathbf{x}$ and $\mathbf{y}$ are unordered.

## Returns

3 The isless macro returns the value of (x) < (y).

### 7.7.14.4 The islessequal macro

Synopsis

```
#include <math.h>
int islessequal(real-floating \mathbf{x, real-floating y);}
```

Description
2 The islessequal macro determines whether its first argument is less than or equal to its second argument. The value of islessequal ( $\mathbf{x}, \mathbf{y}$ ) is always equal to $(x)<=(y)$; however, unlike $(x)<=(y)$, islessequal $(x, y)$ does not raise the invalid exception when $\mathbf{x}$ and $\mathbf{y}$ are unordered.

## Returns

3 The islessequal macro returns the value of (x) <= (y).
7.7.14.5 The islessgreater macro

Synopsis

```
#include <math.h>
int islessgreater(real-floating x, real-floating y);
```


## Description

2 The islessgreater macro determines whether its first argument is less than or greater than its second argument. The islessgreater $(\mathbf{x}, \mathbf{y})$ macro is similar to $(x)<(y)| |(x)>(y)$; however, islessgreater $(x, y)$ does not raise the invalid exception when $\mathbf{x}$ and $\mathbf{y}$ are unordered (nor does it evaluate $\mathbf{x}$ and $\mathbf{y}$ twice).

## Returns

3 The islessgreater macro returns the value of $(x)<(y)| |(x)>(y)$.
7.7.14.6 The isunordered macro

Synopsis

```
#include <math.h>
int isunordered(real-floating \mathbf{x,}}\mathrm{ real-floating y);
```

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## Description

2 The isunordered macro determines whether its arguments are unordered.

## Returns

The isunordered macro returns 1 if its arguments are unordered and 0 otherwise.

### 7.8 Complex arithmetic <complex.h>

1 The header <complex.h> defines macros and declares functions that support complex arithmetic. Each synopsis specifies a function with one or two double complex parameters and returning a double complex or double value; for each such function, there are similar functions with the same name but with $\mathbf{f}$ and $\mathbf{l}$ suffixes. The $\mathbf{f}$ suffix indicates that $\mathbf{f l o a t}$ (instead of double) is the corresponding real type for the parameters and result. Similarly the 1 suffix indicates that long double is the corresponding real type for the parameters and result.

2 The macro
_Complex_I
expands to a constant expression of type const float complex, with the value of the imaginary unit. ${ }^{185}$ The macro
_Imaginary_I
is defined if and only if the implementation supports imaginary types; ${ }^{186}$ it expands to a constant expression of type const float imaginary, with the value of the imaginary unit.

The macro

## I

is defined to be _Complex_I or, if defined, _Imaginary_I. Notwithstanding the provisions of subclause 7.1.3, it is permitted to undefine the macro $I$.

### 7.8.1 The CX_LIMITED_RANGE pragma

## Synopsis

```
#include <complex.h>
#pragma STDC CX_LIMITED_RANGE on-off-switch
```


## Description

2 The usual mathematical formula for multiplication of two complex numbers and the one for division by a complex number are problematic because of their treatment of infinities and because of undue overflow and underflow. The CX_LIMITED_RANGE pragma can be used to inform the implementation that (where the state is on) the usual

[^91]mathematical formulas for multiplication and division are acceptable. ${ }^{187}$ The pragma can occur either outside external declarations or preceding all explicit declarations and statements inside a compound statement. When outside external declarations, the pragma takes effect from its occurrence until another CX_LIMITED_RANGE pragma is encountered, or until the end of the translation unit. When inside a compound statement, the pragma takes effect from its occurrence until another CX_LIMITED_RANGE pragma is encountered (within a nested compound statement), or until the end of the compound statement; at the end of a compound statement the state for the pragma is restored to its condition just before the compound statement. If this pragma is used in any other context, the behavior is undefined. The default state for the pragma is off.

### 7.8.2 Complex functions

1 Values are interpreted as radians, not degrees. An implementation may set errno but is not required to.

### 7.8.2.1 Branch cuts

1 Some of the functions below have branch cuts, across which the function is discontinuous. For implementations with a signed zero (including all IEC 559 implementations) that follow the specification of Annex $G$, the sign of zero distinguishes one side of a cut from another so the function is continuous (except for format limitations) as the cut is approached from either side. For example, for the square root function, which has a branch cut along the negative real axis, the top of the cut, with imaginary part +0 , maps to the positive imaginary axis, and the bottom of the cut, with imaginary part -0 , maps to the negative imaginary axis.

Implementations that do not support a signed zero (see Annex F) cannot distinguish the sides of branch cuts. These implementations must map a cut so the function is continuous as the cut is approached coming around the finite endpoint of the cut in a counter clockwise direction. (Branch cuts for the functions specified here have just one finite endpoint.) For example, for the square root function, coming counter clockwise around the finite endpoint of the cut along the negative real axis approaches the cut from above, so the cut maps to the positive imaginary axis.
187. The purpose of the pragma is to allow the implementation to use the formulas

$$
\begin{gathered}
\left(\mathrm{x}+\mathrm{y}^{*} i\right) *\left(\mathrm{u}+\mathrm{v}^{*} i\right)=\left(\mathrm{x} * \mathrm{u}-\mathrm{y}^{*} \mathrm{v}\right)+\left(\mathrm{y} * \mathrm{u}+\mathrm{x}^{*} \mathrm{v}\right)^{*} i \\
\left(\mathrm{x}+\mathrm{y}^{* i}\right) /\left(\mathrm{u}+\mathrm{v}^{*} i\right)=\left(\mathrm{x}^{*} \mathrm{u}+\mathrm{y}^{*} \mathrm{v}\right) /\left(\mathrm{u}^{*} \mathrm{u}+\mathrm{v}^{*} \mathrm{v}\right)+ \\
\left(\left(\mathrm{y}^{*} \mathrm{u}-\mathrm{x}^{*} \mathrm{v}\right) /\left(\mathrm{u}^{*} \mathrm{u}+\mathrm{v}^{*} \mathrm{v}\right)\right)^{*} i
\end{gathered}
$$

where the programmer can determine they are safe.

### 7.8.2.2 The cacos function

Synopsis

```
#include <complex.h>
double complex cacos(double complex z);
```


## Description

2 The cacos function computes the complex arc cosine of $\mathbf{z}$, with branch cuts outside the interval $[-1,1]$ along the real axis.

## Returns

3 The cacos function returns the complex arc cosine value, in the range of a strip mathematically unbounded along the imaginary axis and in the interval $[0, \pi]$ along the real axis.

### 7.8.2.3 The casin function

Synopsis

```
#include <complex.h>
double complex casin(double complex z);
```


## Description

2 The casin function computes the complex arc sine of $\mathbf{z}$, with branch cuts outside the interval $[-1,1]$ along the real axis.

## Returns

3 The casin function returns the complex arc sine value, in the range of a strip mathematically unbounded along the imaginary axis and in the interval $[-\pi / 2, \pi / 2]$ along the real axis.

### 7.8.2. The catan function

Synopsis

```
#include <complex.h>
double complex catan(double complex z);
```


## Description

2 The catan function computes the complex arc tangent of $\mathbf{z}$, with branch cuts outside the interval $[-i, i]$ along the imaginary axis.

## Returns

3 The catan function returns the complex arc tangent value, in the range of a strip mathematically unbounded along the imaginary axis and in the interval $[-\pi / 2, \pi / 2]$ along the real axis.

### 7.8.2.5 The ccos function

Synopsis

```
#include <complex.h>
double complex ccos(double complex z);
```


## Description

2 The ccos function computes the complex cosine of $\mathbf{z}$.

## Returns

3 The ccos function returns the complex cosine value.

### 7.8.2.6 The csin function

Synopsis

```
#include <complex.h>
    double complex csin(double complex z);
```

Description
2 The csin function computes the complex sine of $\mathbf{z}$.

## Returns

3 The csin function returns the complex sine value.

### 7.8.2.7 The ctan function

Synopsis

```
#include <complex.h>
double complex ctan(double complex z);
```


## Description

2 The ctan function computes the complex tangent of $\mathbf{z}$.

## Returns

3 The ctan function returns the complex tangent value.

### 7.8.2.8 The cacosh function

Synopsis

```
#include <complex.h>
double complex cacosh(double complex z);
```


## Description

2 The cacosh function computes the complex arc hyperbolic cosine of $\mathbf{z}$, with a branch cut at values less than 1 along the real axis.

## Returns

3 The cacosh function returns the complex arc hyperbolic cosine value, in the range of a half-strip of non-negative values along the real axis and in the interval $[-i \pi, i \pi]$ along the imaginary axis.

### 7.8.2.9 The casinh function

Synopsis

```
#include <complex.h>
double complex casinh(double complex z);
```


## Description

2 The casinh function computes the complex arc hyperbolic sine of $\mathbf{z}$, with branch cuts outside the interval $[-i, i]$ along the imaginary axis.

## Returns

3 The casinh function returns the complex arc hyperbolic sine value, in the range of a strip mathematically unbounded along the real axis and in the interval $[-i \pi / 2, i \pi / 2]$ along the imaginary axis.

### 7.8.2.10 The catanh function

Synopsis

```
#include <complex.h>
double complex catanh(double complex z);
```


## Description

2 The catanh function computes the complex arc hyperbolic tangent of $\mathbf{z}$, with branch cuts outside the interval $[-1,1]$ along the real axis.

## Returns

3 The catanh function returns the complex arc hyperbolic tangent value, in the range of a strip mathematically unbounded along the real axis and in the interval $[-i \pi / 2, i \pi / 2]$ along the imaginary axis.

### 7.8.2.11 The ccosh function

Synopsis
1 \#include <complex.h>
double complex ccosh (double complex z);
Description
2 The ccosh function computes the complex hyperbolic cosine of $\mathbf{z}$.

## Returns

3 The ccosh function returns the complex hyperbolic cosine value.

### 7.8.2.12 The csinh function

Synopsis

```
#include <complex.h>
double complex csinh(double complex z);
```


## Description

2 The csinh function computes the complex hyperbolic sine of $\mathbf{z}$.

## Returns

3 The csinh function returns the complex hyperbolic sine value.

### 7.8.2.13 The ctanh function

Synopsis

```
#include <complex.h>
double complex ctanh(double complex z);
```


## Description

2 The ctanh function computes the complex hyperbolic tangent of $\mathbf{z}$.

## Returns

3 The ctanh function returns the complex hyperbolic tangent value.

### 7.8.2.14 The cexp function

Synopsis

```
#include <complex.h>
double complex cexp(double complex z);
```


## Description

2 The cexp function computes the complex base-e exponential of $\mathbf{z}$.

## Returns

3 The cexp function returns the complex base-e exponential value.

### 7.8.2.15 The clog function

Synopsis

```
#include <complex.h>
double complex clog(double complex z);
```


## Description

2 The clog function computes the complex natural (base-e) logarithm of $\mathbf{z}$, with a branch cut along the negative real axis.

## Returns

3 The clog function returns the complex natural logarithm value, in the range of a strip mathematically unbounded along the real axis and in the interval $[-i \pi, i \pi]$ along the imaginary axis.

### 7.8.2.16 The csqrt function

Synopsis

```
#include <complex.h>
double complex csqrt(double complex z);
```


## Description

2 The csqrt function computes the complex square root of $\mathbf{z}$, with a branch cut along the negative real axis.

## Returns

3 The csqrt function returns the complex square root value, in the range of the right half-plane (including the imaginary axis).

### 7.8.2.17 The cabs function

Synopsis

```
#include <complex.h>
double cabs(double complex z);
```


## Description

2 The cabs function computes the complex absolute value (also called norm, modulus, or magnitude) of $\mathbf{z}$.

## Returns

3 The cabs function returns the complex absolute value.

### 7.8.2.18 The cpow function

Synopsis

```
#include <complex.h>
double complex cpow(double complex x, double complex y);
```


## Description

2 The cpow function computes the complex power function $\mathbf{x}^{\mathbf{y}}$, with a branch cut for the first parameter along the negative real axis.

## Returns

3 The cpow function returns the complex power function value.

### 7.8.2.19 The carg function

Synopsis

```
#include <complex.h>
double carg(double complex z);
```


## Description

2 The carg function computes the argument (also called phase angle) of $\mathbf{z}$, with a branch cut along the negative real axis.

## Returns

3 The carg function returns the value of the argument in the range $[-\pi, \pi]$.

### 7.8.2.20 The conj function

Synopsis

```
#include <complex.h>
double complex conj(double complex z);
```


## Description

2 The conj function computes the complex conjugate of $\mathbf{z}$, by reversing the sign of its imaginary part.

## Returns

3 The conj function returns the complex conjugate value.

### 7.8.2.21 The cimag function

Synopsis

```
#include <complex.h>
double cimag(double complex z);
```


## Description

2 The cimag function computes the imaginary part of $\mathbf{z}$. ${ }^{188}$

## Returns

3 The cimag function returns the imaginary part value (as a real).

[^92]
### 7.8.2.22 The cproj function

Synopsis

```
#include <complex.h>
double complex cproj(double complex z);
```


## Description

2 The cproj function computes a projection of $\mathbf{z}$ onto the Riemann sphere: $\mathbf{z}$ projects to $\mathbf{z}$ except that all complex infinities (even those with one infinite part and one NaN part) project to positive infinity on the real axis. If $\mathbf{z}$ has an infinite part, then $\operatorname{cproj}(z)$ is equivalent to

INFINITY + I * copysign(0.0, cimag(z))

## Returns

3 The cproj function returns the value of the projection onto the Riemann sphere.

### 7.8.2.23 The creal function

Synopsis
\#include <complex.h>
double creal (double complex $z$ );

## Description

2 The creal function computes the real part of $\mathbf{z}$.

## Returns

3 The creal function returns the real part value.

### 7.9 Type-generic math <tgmath.h>

1 The header <tgmath.h> includes the headers <math.h> and <complex.h> and defines several type-generic macros.

### 7.9.1 Type-generic macros

1 Of the <math. $\mathrm{h}>$ and <complex.h> functions without an $\mathbf{f}$ (float) or $\mathbf{l}$ (long double) suffix, several have one or more parameters whose corresponding real type is double. For each such function, except modf, there is a corresponding typegeneric macro. ${ }^{189}$ The parameters whose corresponding real type is double in the function synopsis are generic parameters. Use of the macro invokes a function whose corresponding real type and type-domain are determined by the arguments for the generic parameters. ${ }^{190}$
2 Use of the macro invokes a function whose generic parameters have the corresponding real type determined as follows:

- First, if any argument for generic parameters has type long double, the type determined is long double.
- Otherwise, if any argument for generic parameters has type double or is of integer type, the type determined is double.
- Otherwise, the type determined is float.

3 For each unsuffixed function in <math. $\mathrm{h}>$ for which there is a function in <complex.h> with the same name except for a c prefix, the corresponding typegeneric macro (for both functions) has the same name as the function in <math.h>. The corresponding type-generic macro for fabs and cabs is fabs.

[^93]| <math.h> function | <complex.h> function | type-generic macro |
| :---: | :---: | :---: |
| acos | cacos | acos |
| asin | casin | asin |
| atan | catan | atan |
| acosh | cacosh | acosh |
| asinh | casinh | asinh |
| atanh | catanh | atanh |
| cos | ccos | cos |
| sin | csin | sin |
| tan | ctan | $\tan$ |
| cosh | ccosh | cosh |
| sinh | csinh | sinh |
| tanh | ctanh | tanh |
| exp | cexp | exp |
| log | clog | log |
| pow | cpow | pow |
| sqrt | csqrt | sqrt |
| fabs | cabs | fabs |

4 If at least one argument for a generic parameter is complex, then use of the macro invokes a complex function; otherwise, use of the macro invokes a real function.

5 For each unsuffixed function in <math.h> without a c-prefixed counterpart in <complex.h>, the corresponding type-generic macro has the same name as the function. These type-generic macros are:

| atan2 | cbrt | ceil |
| :--- | :--- | :--- |
| copysign | erf | erfc |
| exp2 | expm1 | fdim |
| floor | fma | fmax |
| fmin | fmod | frexp |
| gamma | hypot | ilogb |
| ldexp | lgamma | llrint |
| llround | log10 | log1p |
| log2 | logb | lrint |
| lround | nearbyint | nextafter |
| nextafterx | remainder | remquo |
| rint | round | scalbn |
| scalbln | trunc |  |

6 If all arguments for generic parameters are real, then use of the macro invokes a real function; otherwise, use of the macro results in undefined behavior.

7 For each unsuffixed function in <complex.h> that is not a c-prefixed counterpart to a function in <math. h$\rangle$, the corresponding type-generic macro has the same name as the function. These type-generic macros are:

```
carg cimag conj
cproj creal
```

Use of the macro with any real or complex argument invokes a complex function.

## Examples

9 With the declarations

```
#include <tgmath.h>
int n;
float f;
double d;
long double ld;
float complex fc;
double complex dc;
long double complex ldc;
```

functions invoked by use of type-generic macros are shown in the following table:
macro use
$\exp (n)$
$\operatorname{acosh}(f)$
$\sin (d)$
atan(ld)
$\log (f c)$
sqrt (dc)
pow (ldc, f)
remainder ( $\mathrm{n}, \mathrm{n}$ )
nextafter (d, f)
nextafterx (f, ld)
copysign ( $n$, ld)
ceil(fc)
rint (dc)
fmax(ldc, ld)
carg ( n )
cproj(f)
creal (d)
cimag (ld)
cabs (fc)
carg (dc)
cproj(ldc)
invokes
$\exp (n)$, the function
$\operatorname{acoshf}(f)$
$\sin (d)$, the function
atanl (ld)
clogf(fc)
csqrt (dc)
cpowl(ldc, f)
remainder ( $n, n$ ), the function
nextafter ( $\mathrm{d}, \mathrm{f}$ ), the function
nextafterxf ( $f$, ld)
copysignl (n, ld)
undefined behavior
undefined behavior
undefined behavior
$\operatorname{carg}(n)$, the function
cprojf(f)
creal (d), the function
cimagl(ld)
cabsf(fc)
carg (dc), the function
cprojl(ldc)

### 7.10 Nonlocal jumps <setjmp.h>

1 The header <setjmp. h > defines the macro setjmp, and declares one function and one type, for bypassing the normal function call and return discipline. ${ }^{191}$

2 The type declared is

```
jmp_buf
```

which is an array type suitable for holding the information needed to restore a calling environment.

3 It is unspecified whether setjmp is a macro or an identifier declared with external linkage. If a macro definition is suppressed in order to access an actual function, or a program defines an external identifier with the name setjmp, the behavior is undefined.

### 7.10.1 Save calling environment

### 7.10.1.1 The set jmp macro

Synopsis

```
#include <setjmp.h>
int setjmp(jmp_buf env);
```


## Description

2 The setjmp macro saves its calling environment in its jmp_buf argument for later use by the longjmp function.

## Returns

3 If the return is from a direct invocation, the setjmp macro returns the value zero. If the return is from a call to the longjmp function, the setjmp macro returns a nonzero value.

## Environmental restriction

4 An invocation of the setjmp macro shall appear only in one of the following contexts:

- the entire controlling expression of a selection or iteration statement;

[^94]- one operand of a relational or equality operator with the other operand an integer constant expression, with the resulting expression being the entire controlling expression of a selection or iteration statement;
- the operand of a unary ! operator with the resulting expression being the entire controlling expression of a selection or iteration statement; or
- the entire expression of an expression statement (possibly cast to void).

If the invocation appears in any other context, the behavior is undefined.

### 7.10.2 Restore calling environment

### 7.10.2.1 The long jmp function

## Synopsis

```
#include <setjmp.h>
void longjmp(jmp_buf env, int val);
```


## Description

2 The longjmp function restores the environment saved by the most recent invocation of the setjmp macro in the same invocation of the program, with the corresponding jmp_buf argument. If there has been no such invocation, or if the function containing the invocation of the setjmp macro has terminated execution ${ }^{192}$ in the interim, the behavior is undefined.

3 All accessible objects have values as of the time longjmp was called, except that the values of objects of automatic storage duration that are local to the function containing the invocation of the corresponding setjmp macro that do not have volatile-qualified type and have been changed between the setjmp invocation and longjmp call are indeterminate.

## Returns

4 After longjmp is completed, program execution continues as if the corresponding invocation of the setjmp macro had just returned the value specified by val. The longjmp function cannot cause the set jmp macro to return the value 0 ; if val is 0 , the set jmp macro returns the value 1 .

[^95]
## Examples

5 The longjmp function that returns control back to the point of the setjmp invocation might cause memory associated with a variable length array object to be squandered.

```
#include <setjmp.h>
jmp_buf buf;
void g(int n);
void h(int n);
int n = 6;
void f(void)
{
    int x[n]; // OK, f is not terminated.
    setjmp(buf);
    g(n);
}
void g(int n)
{
    int a[n]; // a may remain allocated.
    h(n);
}
void h(int n)
{
    int b[n]; // b may remain allocated.
    longjmp(buf,2); // might cause memory loss.
}
```


### 7.11 Signal handling <signal.h>

1 The header <signal.h> declares a type and two functions and defines several macros, for handling various signals (conditions that may be reported during program execution).

2 The type defined is
sig_atomic_t
which is the (possibly volatile-qualified) integer type of an object that can be accessed as an atomic entity, even in the presence of asynchronous interrupts.

3 The macros defined are

```
SIG_DFL
SIG_ERR
SIG_IGN
```

which expand to positive integer constant expressions with type int and distinct values that have type compatible with the second argument to and the return value of the signal function, and whose value compares unequal to the address of any declarable function; and the following, which expand to positive integer constant expressions with distinct values that are the signal numbers, each corresponding to the specified condition:

SIGABRT abnormal termination, such as is initiated by the abort function
SIGFPE an erroneous arithmetic operation, such as zero divide or an operation resulting in overflow

SIGILL detection of an invalid function image, such as an illegal instruction
SIGINT receipt of an interactive attention signal
SIGSEGV an invalid access to storage
SIGTERM a termination request sent to the program
4 An implementation need not generate any of these signals, except as a result of explicit calls to the raise function. Additional signals and pointers to undeclarable functions, with macro definitions beginning, respectively, with the letters SIG and an uppercase letter or with SIG_ and an uppercase letter, ${ }^{193}$ may also be specified by the
193. See 'future library directions" (7.20.6). The names of the signal numbers reflect the following terms (respectively): abort, floating-point exception, illegal instruction, interrupt, segmentation violation, and termination.
implementation. The complete set of signals, their semantics, and their default handling is implementation-defined; all signal numbers shall be positive.

### 7.11.1 Specify signal handling

### 7.11.1.1 The signal function

Synopsis

```
#include <signal.h>
void (*signal(int sig, void (*func) (int))) (int);
```


## Description

2 The signal function chooses one of three ways in which receipt of the signal number sig is to be subsequently handled. If the value of func is SIG_DFL, default handling for that signal will occur. If the value of func is SIG_IGN, the signal will be ignored. Otherwise, func shall point to a function to be called when that signal occurs. An invocation of such a function because of a signal, or (recursively) of any further functions called by that invocation (other than functions in the standard library), is called a signal handler.

3 When a signal occurs and func points to a function, it is implementation-defined whether the equivalent of signal(sig, SIG_DFL); is executed or the implementation prevents some implementation-defined set of signals (at least including sig) from occuring until the current signal handling has completed; in the case of SIGILL, the implementation may alternatively define that no action is taken. Then the equivalent of (*func) (sig) ; is executed. If and when the function returns, if the value of sig is SIGFPE, SIGILL, SIGSEGV, or any other implementationdefined value corresponding to a computational exception, the behavior is undefined; otherwise the program will resume execution at the point it was interrupted.

4 If the signal occurs as the result of calling the abort or raise function, the signal handler shall not call the raise function.

5 If the signal occurs other than as the result of calling the abort or raise function, the behavior is undefined if the signal handler refers to any object with static storage duration other than by assigning a value to an object declared as volatile $\boldsymbol{s i g}$ _atomic_t, or the signal handler calls any function in the standard library other than the abort function or the signal function with the first argument equal to the signal number corresponding to the signal that caused the invocation of the handler. Furthermore, if such a call to the signal function results in a SIG_ERR return, the value of errno is indeterminate. ${ }^{194}$

6 At program startup, the equivalent of
signal(sig, SIG_IGN);
may be executed for some signals selected in an implementation-defined manner; the equivalent of

```
signal(sig, SIG_DFL);
```

is executed for all other signals defined by the implementation.
7 The implementation shall behave as if no library function calls the signal function.

## Returns

8 If the request can be honored, the signal function returns the value of func for the most recent successful call to signal for the specified signal sig. Otherwise, a value of SIG_ERR is returned and a positive value is stored in errno.

Forward references: the abort function (7.14.4.1), the exit function (7.14.4.3).

### 7.11.2 Send signal

### 7.11.2.1 The raise function

Synopsis

```
#include <signal.h>
int raise(int sig);
```


## Description

2 The raise function carries out the actions described in subclause 7.11.1.1 for the signal sig. If a signal handler is called, the raise function shall not return until after the signal handler does.

## Returns

3 The raise function returns zero if successful, nonzero if unsuccessful.
194. If any signal is generated by an asynchronous signal handler, the behavior is undefined.

### 7.12 Variable arguments <stdarg. $\mathrm{h}>$

1 The header <stdarg. h > declares a type and defines four macros, for advancing through a list of arguments whose number and types are not known to the called function when it is translated.

2 A function may be called with a variable number of arguments of varying types. As described in 6.7.1, its parameter list contains one or more parameters. The rightmost parameter plays a special role in the access mechanism, and will be designated parmN in this description.

3 The type declared is

```
va_list
```

which is an object type suitable for holding information needed by the macros va_start, va_arg, va_end, and va_copy. If access to the varying arguments is desired, the called function shall declare an object (referred to as ap in this subclause) having type va_list. The object ap may be passed as an argument to another function; if that function invokes the va_arg macro with parameter ap, the value of ap in the calling function is indeterminate and shall be passed to the va_end macro prior to any further reference to ap. ${ }^{195}$

### 7.12.1 Variable argument list access macros

1 The va_start, va_arg, and va_copy macros described in this subclause shall be implemented as macros, not functions. It is unspecified whether va_end is a macro or an identifier declared with external linkage. If a macro definition is suppressed in order to access an actual function, or a program defines an external identifier with the name va_end, the behavior is undefined. Each invocation of the va_start or va_copy macros shall be matched by a corresponding invocation of the va_end macro in the function accepting a varying number of arguments.

### 7.12.1.1 The va_start macro

## Synopsis

```
#include <stdarg.h>
    void va_start(va_list ap, parmN);
```

[^96]
## Description

2 The va_start macro shall be invoked before any access to the unnamed arguments.
3 The va_start macro initializes ap for subsequent use by va_arg and va_end. va_start shall not be invoked again for the same ap without an intervening invocation of va_end for the same ap.

4 The parameter parmN is the identifier of the rightmost parameter in the variable parameter list in the function definition (the one just before the , ...). If the parameter parm $N$ is declared with the register storage class, with a function or array type, or with a type that is not compatible with the type that results after application of the default argument promotions, the behavior is undefined.

## Returns

5 The va_start macro returns no value.

### 7.12.1.2 The va_arg macro

## Synopsis

```
#include <stdarg.h>
type va_arg(va_list ap, type);
```


## Description

2 The va_arg macro expands to an expression that has the type and value of the next argument in the call. The parameter ap shall be the same as the va_list ap initialized by va_start. Each invocation of va_arg modifies ap so that the values of successive arguments are returned in turn. The parameter type is a type name specified such that the type of a pointer to an object that has the specified type can be obtained simply by postfixing a * to type. If there is no actual next argument, or if type is not compatible with the type of the actual next argument (as promoted according to the default argument promotions), the behavior is undefined.

## Returns

3 The first invocation of the va_arg macro after that of the va_start macro returns the value of the argument after that specified by parmN. Successive invocations return the values of the remaining arguments in succession.

### 7.12.1.3 The va_copy macro

Synopsis

```
#include <stdarg.h>
void va_copy(va_list dest, va_list src);
```


## Description

2 The va_copy macro makes the va_list dest be a copy of the va_list src, as if the va_start macro had been applied to it followed by the same sequence of uses of the va_arg macro as had previously been used to reach the present state of src.

## Returns

3 The va_copy macro returns no value.

### 7.12.1.4 The va_end macro

Synopsis

```
#include <stdarg.h>
void va_end(va_list ap);
```


## Description

2 The va_end macro facilitates a normal return from the function whose variable argument list was referred to by the expansion of va_start that initialized the va_list ap. The va_end macro may modify ap so that it is no longer usable (without an intervening invocation of va_start). If there is no corresponding invocation of the va_start macro, or if the va_end macro is not invoked before the return, the behavior is undefined.

## Returns

3 The va_end macro returns no value.

## Examples

4 The function $\mathbf{f 1}$ gathers into an array a list of arguments that are pointers to strings (but not more than MAXARGS arguments), then passes the array as a single argument to function $\mathbf{f 2}$. The number of pointers is specified by the first argument to $\mathbf{f 1}$.

```
#include <stdarg.h>
#define MAXARGS 31
void f1(int n_ptrs, ...)
{
    va_list ap;
    char *array[MAXARGS];
    int ptr_no = 0;
    if (n_ptrs > MAXARGS)
            n_ptrs = MAXARGS;
    va_start(ap, n_ptrs);
    while (ptr_no < n_ptrs)
            array[ptr_no++] = va_arg(ap, char *);
        va_end(ap);
    f2(n_ptrs, array);
}
```

Each call to $\mathbf{f 1}$ shall have visible the definition of the function or a declaration such as

```
void f1(int, ...);
``` The function \(\mathbf{f 3}\) is similar, but saves the status of the variable argument list after the indicated number of arguments; after \(\mathbf{f 2}\) has been called once with the whole list, the trailing part of the list is gathered again and passed to function \(£ 4\).
```

\#include <stdarg.h>
\#define MAXARGS 31
void f3(int n_ptrs, int f4_after, ...)
{
va_list ap, ap_save;
char *array[MAXARGS];
int ptr_no = 0;
if (n_ptrs > MAXARGS)
n_ptrs = MAXARGS;
va_start(ap, n_ptrs);
while (ptr_no < n_ptrs) {
array[ptr_no++] = va_arg(ap, char *);
if (ptr_no == f4_after)
va_copy(ap_save, ap);
}
va_end(ap);
f2(n_ptrs, array);
// Now process the saved copy.
n_ptrs -= f4_after;
ptr_no = 0;
while (ptr_no < n_ptrs)
array[ptr_no++] = va_arg(ap_save, char *);
va_end(ap_save);
f4(n_ptrs, array);
}

```

\subsection*{7.13 Input/output <stdio. h>}

\subsection*{7.13.1 Introduction}

1 The header <stdio. h> declares three types, several macros, and many functions for performing input and output.

2 The types declared are size_t (described in 7.1.6);
FILE
which is an object type capable of recording all the information needed to control a stream, including its file position indicator, a pointer to its associated buffer (if any), an error indicator that records whether a read/write error has occurred, and an end-offile indicator that records whether the end of the file has been reached; and
fpos_t
which is an object type other than an array type capable of recording all the information needed to specify uniquely every position within a file.

3 The macros are NULL (described in 7.1.6);
_IOFBF
_IOLBF
_IONBF
which expand to integer constant expressions with distinct values, suitable for use as the third argument to the setvbuf function;

BUFSIZ
which expands to an integer constant expression, which is the size of the buffer used by the setbuf function;

EOF
which expands to an integer constant expression, with type int and a negative value, that is returned by several functions to indicate end-of-file, that is, no more input from a stream;

\section*{FOPEN_MAX}
which expands to an integer constant expression that is the minimum number of files that the implementation guarantees can be open simultaneously;

FILENAME_MAX
which expands to an integer constant expression that is the size needed for an array of char large enough to hold the longest file name string that the implementation
guarantees can be opened; \({ }^{196}\)

\section*{L_tmpnam}
which expands to an integer constant expression that is the size needed for an array of char large enough to hold a temporary file name string generated by the tmpnam function;
```

SEEK_CUR
SEEK_END
SEEK_SET

```
which expand to integer constant expressions with distinct values, suitable for use as the third argument to the fseek function;

TMP_MAX
which expands to an integer constant expression that is the minimum number of unique file names that shall be generated by the tmpnam function;
```

stderr
stdin
stdout

```
which are expressions of type "pointer to FILE" that point to the FILE objects associated, respectively, with the standard error, input, and output streams.

4 The header <wchar. h> declares a number of functions useful for wide-character input and output. The wide-character input/output functions described in that subclause provide operations analogous to most of those described here, except that the fundamental units internal to the program are wide characters. The external representation (in the file) is a sequence of "generalized" multibyte characters, as described further in subclause 7.13.3.

5 The input/output functions are given the following collective terms:
- The wide-character input functions - those functions described in these subclauses that perform input into wide characters and wide strings: fgetwc, fgetws, getwc, getwchar, fwscanf, wscanf, vfwscanf, and vwscanf.
196. If the implementation imposes no practical limit on the length of file name strings, the value of FILENAME_MAX should instead be the recommended size of an array intended to hold a file name string. Of course, file name string contents are subject to other system-specific constraints; therefore all possible strings of length FILENAME_MAX cannot be expected to be opened successfully.
- The wide-character output functions - those functions described in these subclauses that perform output from wide characters and wide strings: fputwc, fputws, putwc, putwchar, fwprintf, wprintf, vfwprintf, and vwprintf.
- The wide-character input/output functions - the union of the ungetwc function, the wide-character input functions, and the wide-character output functions.
- The byte input/output functions - those functions described in these subclauses that perform input/output: fgetc, fgets, fprintf, fputc, fputs, fread, fscanf, fwrite, getc, getchar, gets, printf, putc, putchar, puts, scanf, ungetc, vfprintf, vfscanf, vprintf, and vscanf.

Forward references: files (7.13.3), the fseek function (7.13.9.2), streams (7.13.2), the tmpnam function (7.13.4.4), <wchar.h> (7.19).

\subsection*{7.13.2 Streams}

1 Input and output, whether to or from physical devices such as terminals and tape drives, or whether to or from files supported on structured storage devices, are mapped into logical data streams, whose properties are more uniform than their various inputs and outputs. Two forms of mapping are supported, for text streams and for binary streams. \({ }^{197}\)

2 A text stream is an ordered sequence of characters composed into lines, each line consisting of zero or more characters plus a terminating new-line character. Whether the last line requires a terminating new-line character is implementation-defined. Characters may have to be added, altered, or deleted on input and output to conform to differing conventions for representing text in the host environment. Thus, there need not be a one-to-one correspondence between the characters in a stream and those in the external representation. Data read in from a text stream will necessarily compare equal to the data that were earlier written out to that stream only if: the data consist only of printable characters and the control characters horizontal tab and new-line; no new-line character is immediately preceded by space characters; and the last character is a new-line character. Whether space characters that are written out immediately before a new-line character appear when read in is implementation-defined.

3 A binary stream is an ordered sequence of characters that can transparently record internal data. Data read in from a binary stream shall compare equal to the data that
197. An implementation need not distinguish between text streams and binary streams. In such an implementation, there need be no new-line characters in a text stream nor any limit to the length of a line.
were earlier written out to that stream, under the same implementation. Such a stream may, however, have an implementation-defined number of null characters appended to the end of the stream.

4 Each stream has an orientation. After a stream is associated with an external file, but before any operations are performed on it, the stream is without orientation. Once a wide-character input/output function has been applied to a stream without orientation, the stream becomes wide-oriented. Similarly, once a byte input/output function has been applied to a stream without orientation, the stream becomes byte-oriented. Only a call to the freopen function or the fwide function can otherwise alter the orientation of a stream. (A successful call to freopen removes any orientation.) \({ }^{198}\)

5 Byte input/output functions shall not be applied to a wide-oriented stream; and widecharacter input/output functions shall not be applied to a byte-oriented stream. The remaining stream operations do not affect and are not affected by a stream's orientation, except for the following additional restrictions:
- Binary wide-oriented streams have the file-positioning restrictions ascribed to both text and binary streams.
- For wide-oriented streams, after a successful call to a file-positioning function that leaves the file position indicator prior to the end-of-file, a wide-character output function can overwrite a partial multibyte character; any file contents beyond the byte(s) written are henceforth indeterminate.

Each wide-oriented stream has an associated mbstate_t object that stores the current parse state of the stream. A successful call to fgetpos stores a representation of the value of this mbstate_t object as part of the value of the fpos_t object. A later successful call to fsetpos using the same stored fpos_t value restores the value of the associated mbstate_t object as well as the position within the controlled stream.

\section*{Environmental limits}

7 An implementation shall support text files with lines containing at least 254 characters, including the terminating new-line character. The value of the macro BUFSIZ shall be at least 256 .
198. The three predefined streams stdin, stdout, and stderr are unoriented at program startup.

Forward references: freopen (7.13.5.4), fwide (7.19.3.10), mbstate_t (7.18.1), fgetpos (7.13.9.1), and fsetpos (7.13.9.3).

\subsection*{7.13.3 Files}

1 A stream is associated with an external file (which may be a physical device) by opening a file, which may involve creating a new file. Creating an existing file causes its former contents to be discarded, if necessary. If a file can support positioning requests (such as a disk file, as opposed to a terminal), then a file position indicator associated with the stream is positioned at the start (character number zero) of the file, unless the file is opened with append mode in which case it is implementation-defined whether the file position indicator is initially positioned at the beginning or the end of the file. The file position indicator is maintained by subsequent reads, writes, and positioning requests, to facilitate an orderly progression through the file.

2 Binary files are not truncated, except as defined in 7.13.5.3. Whether a write on a text stream causes the associated file to be truncated beyond that point is implementationdefined.

3 When a stream is unbuffered, characters are intended to appear from the source or at the destination as soon as possible. Otherwise characters may be accumulated and transmitted to or from the host environment as a block. When a stream is fully buffered, characters are intended to be transmitted to or from the host environment as a block when a buffer is filled. When a stream is line buffered, characters are intended to be transmitted to or from the host environment as a block when a new-line character is encountered. Furthermore, characters are intended to be transmitted as a block to the host environment when a buffer is filled, when input is requested on an unbuffered stream, or when input is requested on a line buffered stream that requires the transmission of characters from the host environment. Support for these characteristics is implementation-defined, and may be affected via the setbuf and setvbuf functions.

4 A file may be disassociated from a controlling stream by closing the file. Output streams are flushed (any unwritten buffer contents are transmitted to the host environment) before the stream is disassociated from the file. The value of a pointer to a FILE object is indeterminate after the associated file is closed (including the standard text streams). Whether a file of zero length (on which no characters have been written by an output stream) actually exists is implementation-defined.

The file may be subsequently reopened, by the same or another program execution, and its contents reclaimed or modified (if it can be repositioned at its start). If the main function returns to its original caller, or if the exit function is called, all open files are closed (hence all output streams are flushed) before program termination. Other paths to program termination, such as calling the abort function, need not close all files properly.

6 The address of the FILE object used to control a stream may be significant; a copy of a FILE object may not necessarily serve in place of the original.

7 At program startup, three text streams are predefined and need not be opened explicitly - standard input (for reading conventional input), standard output (for writing conventional output), and standard error (for writing diagnostic output). When opened, the standard error stream is not fully buffered; the standard input and standard output streams are fully buffered if and only if the stream can be determined not to refer to an interactive device.

8 Functions that open additional (nontemporary) files require a file name, which is a string. The rules for composing valid file names are implementation-defined. Whether the same file can be simultaneously open multiple times is also implementationdefined.

9 Although both text and binary wide-oriented streams are conceptually sequences of wide characters, the external file associated with a wide-oriented stream is a sequence of multibyte characters, generalized as follows:
- Multibyte encodings within files may contain embedded null bytes (unlike multibyte encodings valid for use internal to the program).
- A file need not begin nor end in the initial shift state. \({ }^{199}\)

10 Moreover, the encodings used for multibyte characters may differ among files. Both the nature and choice of such encodings are implementation-defined.

11 The wide-character input functions read multibyte characters from the stream and convert them to wide characters as if they were read by successive calls to the getwc function. Each conversion occurs as if by a call to the mbrtowc function, with the conversion state described by the stream's own mbstate_t object. The byte input functions read characters from the stream as if by successive calls to the fgetc function.

12 The wide-character output functions convert wide characters to multibyte characters and write them to the stream as if they were written by successive calls to the fputwc function. Each conversion occurs as if by a call to the wertomb function, with the conversion state described by the stream's own mbstate_t object. The byte output functions write characters to the stream as if by successive calls to the
199. Setting the file position indicator to end-of-file, as with fseek (file, 0, SEEK_END), has undefined behavior for a binary stream (because of possible trailing null characters) or for any stream with state-dependent encoding that does not assuredly end in the initial shift state.

\section*{fputc function.}

13 An encoding error occurs if the character sequence presented to the underlying mbrtowc function does not form a valid (generalized) multibyte character, or if the code value passed to the underlying wcrtomb does not correspond to a valid (generalized) multibyte character. The wide-character input/output functions and the byte input/output functions store the value of the macro EILSEQ in errno if and only if an encoding error occurs.

\section*{Environmental limits}

14 The value of FOPEN_MAX shall be at least eight, including the three standard text streams.

Forward references: the exit function (7.14.4.3), the fgetc function (7.13.7.1), the fopen function (7.13.5.3), the fputc function (7.13.7.3), the setbuf function (7.13.5.5), the setvbuf function (7.13.5.6), the fgetwc function (7.19.3.1), the fputwc function (7.19.3.3), conversion state (7.19.7), the mbrtowc function (7.19.7.3.2), the wartomb function (7.19.7.3.3).

\subsection*{7.13.4 Operations on files}

\subsection*{7.13.4. 1 The remove function}

Synopsis
1 \#include <stdio.h>

Description
2 The remove function causes the file whose name is the string pointed to by filename to be no longer accessible by that name. A subsequent attempt to open that file using that name will fail, unless it is created anew. If the file is open, the behavior of the remove function is implementation-defined.

\section*{Returns}

3 The remove function returns zero if the operation succeeds, nonzero if it fails.

\subsection*{7.13.4.2 The rename function}

Synopsis
```

\#include <stdio.h>
int rename(const char *old, const char *new);

```

\section*{Description}

2 The rename function causes the file whose name is the string pointed to by old to be henceforth known by the name given by the string pointed to by new. The file named old is no longer accessible by that name. If a file named by the string pointed to by new exists prior to the call to the rename function, the behavior is implementation-defined.

\section*{Returns}

3 The rename function returns zero if the operation succeeds, nonzero if it fails, \({ }^{200}\) in which case if the file existed previously it is still known by its original name.

\subsection*{7.13.4.3 The tmpfile function}

\section*{Synopsis}
```

\#include <stdio.h>
FILE *tmpfile(void);

```

\section*{Description}

2 The tmpfile function creates a temporary binary file that will automatically be removed when it is closed or at program termination. If the program terminates abnormally, whether an open temporary file is removed is implementation-defined. The file is opened for update with "wb+" mode.

\section*{Returns}

3 The tmpfile function returns a pointer to the stream of the file that it created. If the file cannot be created, the tmpfile function returns a null pointer.

Forward references: the fopen function (7.13.5.3).

\subsection*{7.13.4.4 The tmpnam function}

Synopsis
```

\#include <stdio.h>
char *tmpnam(char *s);

```

\footnotetext{
200. Among the reasons the implementation may cause the rename function to fail are that the file is open or that it is necessary to copy its contents to effectuate its renaming.
}

\section*{Description}

2 The tmpnam function generates a string that is a valid file name and that is not the same as the name of an existing file. \({ }^{201}\)

3 The tmpnam function generates a different string each time it is called, up to TMP_MAX times. If it is called more than TMP_MAX times, the behavior is implementation-defined.

4 The implementation shall behave as if no library function calls the tmpnam function.

\section*{Returns}

5 If the argument is a null pointer, the tmpnam function leaves its result in an internal static object and returns a pointer to that object. Subsequent calls to the tmpnam function may modify the same object. If the argument is not a null pointer, it is assumed to point to an array of at least L_tmpnam chars; the tmpnam function writes its result in that array and returns the argument as its value.

\section*{Environmental limits}

6 The value of the macro TMP_MAX shall be at least 25.

\subsection*{7.13.5 File access functions}

\subsection*{7.13.5.1 The fclose function}

Synopsis
```

\#include <stdio.h>
int fclose(FILE *stream);

```

\section*{Description}

2 The fclose function causes the stream pointed to by stream to be flushed and the associated file to be closed. Any unwritten buffered data for the stream are delivered to the host environment to be written to the file; any unread buffered data are discarded. The stream is disassociated from the file. If the associated buffer was automatically allocated, it is deallocated.

\footnotetext{
201. Files created using strings generated by the tmpnam function are temporary only in the sense that their names should not collide with those generated by conventional naming rules for the implementation. It is still necessary to use the remove function to remove such files when their use is ended, and before program termination.
}

\section*{Returns}

3 The fclose function returns zero if the stream was successfully closed, or EOF if any errors were detected.

\subsection*{7.13.5.2 The fflush function}

Synopsis
1 \#include <stdio.h>
int fflush(FILE *stream);

\section*{Description}

2 If stream points to an output stream or an update stream in which the most recent operation was not input, the \(\mathbf{f f l u s h}\) function causes any unwritten data for that stream to be delivered to the host environment to be written to the file; otherwise, the behavior is undefined.

3 If stream is a null pointer, the fflush function performs this flushing action on all streams for which the behavior is defined above.

\section*{Returns}

4 The \(\boldsymbol{f f l u s h}\) function sets the error indicator for the stream and returns EOF if a write error occurs, otherwise it returns zero.

Forward references: the fopen function (7.13.5.3), the ungetc function (7.13.7.11).

\subsection*{7.13.5.3 The fopen function}

Synopsis
```

\#include <stdio.h>
FILE *fopen(const char * restrict filename,
const char * restrict mode);

```

\section*{Description}

2 The fopen function opens the file whose name is the string pointed to by filename, and associates a stream with it.

3 The argument mode points to a string. If the string is one of the following, the file is open in the indicated mode. Otherwise, the behavior is undefined. \({ }^{202}\)

\footnotetext{
202. If the string begins with one of the above sequences, the implementation might choose to ignore the remaining characters, or it might use them to select different kinds of a file (some of which might not conform to the properties in 7.13.2).
}
\begin{tabular}{ll}
\(\mathbf{r}\) & open text file for reading \\
\(\mathbf{w}\) & truncate to zero length or create text file for writing \\
\(\mathbf{a}\) & append; open or create text file for writing at end-of-file \\
\(\mathbf{r b}\) & \begin{tabular}{l} 
open binary file for reading
\end{tabular} \\
\(\mathbf{w b}\) & truncate to zero length or create binary file for writing \\
\(\mathbf{a b}\) & append; open or create binary file for writing at end-of-file \\
\(\mathbf{r}+\) & open text file for update (reading and writing) \\
\(\mathbf{w +}\) & truncate to zero length or create text file for update \\
\(\mathbf{a +}\) & append; open or create text file for update, writing at end-of-file \\
\(\mathbf{r + b}\) & or \(\mathbf{r b +}\) open binary file for update (reading and writing) \\
\(\mathbf{w + b}\) & or \\
\(\mathbf{w b +}+\) truncate to zero length or create binary file for update
\end{tabular}

4 Opening a file with read mode (' \(\mathbf{r}^{\prime}\) as the first character in the mode argument) fails if the file does not exist or cannot be read.

5 Opening a file with append mode (' \(\mathbf{a}^{\prime}\) as the first character in the mode argument) causes all subsequent writes to the file to be forced to the then current end-of-file, regardless of intervening calls to the fseek function. In some implementations, opening a binary file with append mode ( \({ }^{\prime} \mathbf{b}^{\prime}\) as the second or third character in the above list of mode argument values) may initially position the file position indicator for the stream beyond the last data written, because of null character padding.

6 When a file is opened with update mode \(\left({ }^{\prime}+{ }^{\prime}\right.\) as the second or third character in the above list of mode argument values), both input and output may be performed on the associated stream. However, output shall not be directly followed by input without an intervening call to the \(\mathbf{f f l u s h}\) function or to a file positioning function (fseek, fsetpos, or rewind), and input shall not be directly followed by output without an intervening call to a file positioning function, unless the input operation encounters end-of-file. Opening (or creating) a text file with update mode may instead open (or create) a binary stream in some implementations.

7 When opened, a stream is fully buffered if and only if it can be determined not to refer to an interactive device. The error and end-of-file indicators for the stream are cleared.

\section*{Returns}

8 The fopen function returns a pointer to the object controlling the stream. If the open operation fails, fopen returns a null pointer.

Forward references: file positioning functions (7.13.9).

\subsection*{7.13.5.4 The freopen function}

Synopsis
1 \#include <stdio.h>
FILE *freopen (const char * restrict filename, const char * restrict mode, FILE * restrict stream);

\section*{Description}

2 The freopen function opens the file whose name is the string pointed to by filename and associates the stream pointed to by stream with it. The mode argument is used just as in the fopen function. \({ }^{203}\)

3 The freopen function first attempts to close any file that is associated with the specified stream. Failure to close the file successfully is ignored. The error and end-of-file indicators for the stream are cleared.

\section*{Returns}

4 The freopen function returns a null pointer if the open operation fails. Otherwise, freopen returns the value of stream.

\subsection*{7.13.5.5 The setbuf function}

Synopsis
```

\#include <stdio.h>
void setbuf(FILE * restrict stream,
char * restrict buf);

```

Description
2 Except that it returns no value, the setbuf function is equivalent to the setvbuf function invoked with the values _IOFBF for mode and BUFSIZ for size, or (if buf is a null pointer), with the value _IONBF for mode.

\footnotetext{
203. The primary use of the freopen function is to change the file associated with a standard text stream (stderr, stdin, or stdout), as those identifiers need not be modifiable lvalues to which the value returned by the fopen function may be assigned.
}

\section*{Returns}

3 The setbuf function returns no value.
Forward references: the setvbuf function (7.13.5.6).

\subsection*{7.13.5.6 The setvbuf function}

\section*{Synopsis}
```

\#include <stdio.h>
int setvbuf(FILE * restrict stream,
char * restrict buf,
int mode, size_t size);

```

\section*{Description}

2 The setvbuf function may be used only after the stream pointed to by stream has been associated with an open file and before any other operation (other than an unsuccessful call to setvbuf) is performed on the stream. The argument mode determines how stream will be buffered, as follows: _IOFBF causes input/output to be fully buffered; _IOLBF causes input/output to be line buffered; _IONBF causes input/output to be unbuffered. If buf is not a null pointer, the array it points to may be used instead of a buffer allocated by the setvbuf function \({ }^{204}\) and the argument size specifies the size of the array; otherwise, size may determine the size of a buffer allocated by the setvbuf function. The contents of the array at any time are indeterminate.

\section*{Returns}

3 The setvbuf function returns zero on success, or nonzero if an invalid value is given for mode or if the request cannot be honored.

\footnotetext{
204. The buffer must have a lifetime at least as great as the open stream, so the stream should be closed before a buffer that has automatic storage duration is deallocated upon block exit.
}

\subsection*{7.13.6 Formatted input/output functions}

1 The formatted input/output functions \({ }^{205}\) shall behave as if there is a sequence point after the actions associated with each specifier.

\subsection*{7.13.6.1 The fprintf function}

Synopsis
```

\#include <stdio.h>
int fprintf(FILE * restrict stream,
const char * restrict format, ...);

```

\section*{Description}

2 The fprintf function writes output to the stream pointed to by stream, under control of the string pointed to by format that specifies how subsequent arguments are converted for output. If there are insufficient arguments for the format, the behavior is undefined. If the format is exhausted while arguments remain, the excess arguments are evaluated (as always) but are otherwise ignored. The fprintf function returns when the end of the format string is encountered.

3 The format shall be a multibyte character sequence, beginning and ending in its initial shift state. The format is composed of zero or more directives: ordinary multibyte characters (not \%), which are copied unchanged to the output stream; and conversion specifications, each of which results in fetching zero or more subsequent arguments. Each conversion specification is introduced by the character \%. After the \%, the following appear in sequence:
- Zero or more flags (in any order) that modify the meaning of the conversion specification.
- An optional minimum field width. If the converted value has fewer characters than the field width, it will be padded with spaces (by default) on the left (or right, if the left adjustment flag, described later, has been given) to the field width. The field width takes the form of an asterisk * (described later) or a decimal integer. \({ }^{206}\)
- An optional precision that gives the minimum number of digits to appear for the \(\mathbf{d}\), \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), and \(\mathbf{x}\) conversions, the number of digits to appear after the decimalpoint character for \(\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}\), and \(\mathbf{F}\) conversions, the maximum number of significant digits for the \(\boldsymbol{g}\) and \(\mathbf{G}\) conversions, or the maximum number of

\footnotetext{
205. The printf functions perform writes to memory for the \%n specifier.
206. Note that 0 is taken as a flag, not as the beginning of a field width.
}
characters to be written from a string in s conversions. The precision takes the form of a period (.) followed either by an asterisk * (described later) or by an optional decimal integer; if only the period is specified, the precision is taken as zero. If a precision appears with any other conversion specifier, the behavior is undefined.
- An optional \(\mathbf{h h}\) specifying that a following \(\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), or \(\mathbf{x}\) conversion specifier applies to a signed char or unsigned char argument (the argument will have been promoted according to the integer promotions, and its value shall be converted to signed char or unsigned char before printing); an optional \(\mathbf{h}\) specifying that a following \(\mathbf{d}\), \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), or \(\mathbf{x}\) conversion specifier applies to a short int or unsigned short int argument (the argument will have been promoted according to the integer promotions, and its value shall be converted to short int or unsigned short int before printing); an optional \(h\) specifying that a following \(\mathbf{n}\) conversion specifier applies to a pointer to a short int argument; an optional \(\mathbf{l}\) (ell) specifying that a following \(\mathbf{d}\), \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), or \(\mathbf{x}\) conversion specifier applies to a long int or unsigned long int argument; an optional \(\mathbf{l l}\) (ell-ell) specifying that a following \(\mathbf{d}\), \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), or \(\mathbf{x}\) conversion specifier applies to a long long int or unsigned long long int argument; an optional \(\mathbf{l}\) specifying that a following \(\mathbf{n}\) conversion specifier applies to a pointer to a long int argument; an optional 11 specifying that a following n conversion specifier applies to a pointer to a long long int argument; an optional l specifying that a following \(\mathbf{c}\) conversion specifier applies to a wint_t argument; an optional \(\mathbf{l}\) specifying that a following \(\mathbf{s}\) conversion specifier applies to a pointer to a wchar_t argument; an optional \(\mathbf{l}\) which has no effect on a following a, A, \(\mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}\), or \(\mathbf{G}\) conversion specifier; or an optional \(\mathbf{L}\) specifying that a following \(\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}\), or \(\mathbf{G}\) conversion specifier applies to a long double argument. If an \(\mathbf{h h}, \mathbf{h}, \mathbf{l}, \mathbf{l l}\), or \(\mathbf{L}\) appears with any other conversion specifier, the behavior is undefined.
- A character that specifies the type of conversion to be applied.

As noted above, a field width, or precision, or both, may be indicated by an asterisk. In this case, an int argument supplies the field width or precision. The arguments specifying field width, or precision, or both, shall appear (in that order) before the argument (if any) to be converted. A negative field width argument is taken as a flag followed by a positive field width. A negative precision argument is taken as if the precision were omitted.

The flag characters and their meanings are
- The result of the conversion will be left-justified within the field. (It will be right-justified if this flag is not specified.)
\(+\quad\) The result of a signed conversion will always begin with a plus or minus sign. (It will begin with a sign only when a negative value is converted if this flag is not specified. \({ }^{207}\)
space If the first character of a signed conversion is not a sign, or if a signed conversion results in no characters, a space will be prefixed to the result. If the space and + flags both appear, the space flag will be ignored.
\# The result is to be converted to an "alternate form." For o conversion, it increases the precision, if and only if necessary, to force the first digit of the result to be a zero (if the value and precision are both 0 , a single 0 is printed). For \(\mathbf{x}\) ( or \(\mathbf{x}\) ) conversion, a nonzero result will have \(\mathbf{0 x}\) (or \(\mathbf{0 x}\) ) prefixed to it. For \(\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}\), and \(\mathbf{G}\) conversions, the result will always contain a decimal-point character, even if no digits follow it. (Normally, a decimal-point character appears in the result of these conversions only if a digit follows it.) For \(\mathbf{g}\) and \(\mathbf{G}\) conversions, trailing zeros are not removed from the result. For other conversions, the behavior is undefined.
\(\mathbf{0}\) For \(\mathbf{d}\), \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}, \mathbf{x}, \mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}\), and \(\mathbf{G}\) conversions, leading zeros (following any indication of sign or base) are used to pad to the field width; no space padding is performed. If the \(\mathbf{0}\) and - flags both appear, the \(\mathbf{0}\) flag will be ignored. For \(\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), and \(\mathbf{x}\) conversions, if a precision is specified, the \(\mathbf{0}\) flag will be ignored. For other conversions, the behavior is undefined.

The conversion specifiers and their meanings are
d,i The int argument is converted to signed decimal in the style [-]dddd. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it will be expanded with leading zeros. The default precision is 1 . The result of converting a zero value with a precision of zero is no characters.
\(\mathbf{0}, \mathbf{u}, \mathbf{x}, \mathbf{x}\) The unsigned int argument is converted to unsigned octal (0), unsigned decimal ( \(\mathbf{u}\) ), or unsigned hexadecimal notation ( \(\mathbf{x}\) or \(\mathbf{x}\) ) in the style \(d d d d\); the letters abcdef are used for \(\mathbf{x}\) conversion and the letters ABCDEF for \(\mathbf{x}\) conversion. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it will be expanded with leading zeros. The default precision is 1 . The result of converting a zero value with a precision of zero is no
207. The results of all floating conversions of a negative zero, and of negative values that round to zero, include a minus sign.
characters.
\(\mathbf{f , F} \quad\) A double argument representing a floating-point number is converted to decimal notation in the style [-]ddd.ddd, where the number of digits after the decimal-point character is equal to the precision specification. If the precision is missing, it is taken as 6 ; if the precision is zero and the \# flag is not specified, no decimal-point character appears. If a decimal-point character appears, at least one digit appears before it. The value is rounded to the appropriate number of digits.

A double argument representing an infinity is converted to one of the styles [-]inf or [-]infinity - which style is implementation-defined. A double argument representing a NaN is converted in one of the styles [-]nan or [-]nan (n-char-sequence) - which style, and the meaning of any \(n\)-char-sequence, is implementation-defined. The \(\mathbf{F}\) conversion specifier produces INF, INFINITY, or NAN instead of inf, infinity, or nan, respectively. \({ }^{208}\)
\(\mathbf{e , E} \quad\) A double argument representing a floating-point number is converted in the style [-]d. \(d d d \mathbf{e} \pm d d\), where there is one digit before the decimal-point character (which is nonzero if the argument is nonzero) and the number of digits after it is equal to the precision; if the precision is missing, it is taken as 6; if the precision is zero and the \# flag is not specified, no decimal-point character appears. The value is rounded to the appropriate number of digits. The \(\mathbf{E}\) conversion specifier will produce a number with \(\mathbf{E}\) instead of \(\mathbf{e}\) introducing the exponent. The exponent always contains at least two digits, and only as many more digits as necessary to represent the exponent. If the value is zero, the exponent is zero.

A double argument representing an infinity or a NaN is converted in the style of an \(\mathbf{f}\) or \(\mathbf{F}\) conversion specifier.
g,G A double argument representing a floating-point number is converted in style \(\mathbf{f}\) or \(\mathbf{e}\) (or in style \(\mathbf{F}\) or \(\mathbf{E}\) in the case of a \(\mathbf{G}\) conversion specifier), with the precision specifying the number of significant digits. If the precision is zero, it is taken as 1 . The style used depends on the value converted; style e (or \(\mathbf{E}\) ) will be used only if the exponent resulting from such a conversion is less than -4 or greater than or equal to the precision. Trailing zeros are removed from the fractional portion of the result; a

\footnotetext{
208. When applied to infinite and NaN values, the \(\boldsymbol{-}, \boldsymbol{+}\), and space flag characters have their usual meaning; the \# and 0 flag characters have no effect.
}
decimal-point character appears only if it is followed by a digit. A double argument representing an infinity or NaN is converted in the style of an \(\mathbf{f}\) or \(\mathbf{F}\) conversion specifier.
a,A A double argument representing a floating-point number is converted in the style \([-] 0 \mathbf{x} h . h h h h \mathbf{p} \pm d\). The number of hexadecimal digits \(h\) after the decimal-point character is equal to the precision; if the precision is missing and FLT_RADIX is a power of 2 , then the precision is sufficient for an exact representation of the value; if the precision is missing and FLT_RADIX is not a power of 2, then the precision is sufficient to distinguish \({ }^{209}\) values of type double, except that trailing zeros may be omitted. The hexadecimal digit to the left of the decimal-point character is nonzero for normalized floating-point numbers and is otherwise unspecified; \({ }^{210}\) if the precision is zero and the \# flag is not specified, no decimal-point character appears. The letters abcdef are used for a conversion and the letters \(\operatorname{ABCDEF}\) for \(\mathbf{A}\) conversion. The a conversion specifier will produce a number with \(\mathbf{x}\) and \(\mathbf{p}\) and the \(\mathbf{A}\) conversion specifier will produce a number with \(\mathbf{X}\) and \(\mathbf{P}\). The exponent always contains at least one digit, and only as many more digits as necessary to represent the decimal exponent of 2 . If the value is zero, the exponent is zero. A double argument representing an infinity or a NaN is converted in the style of an \(\mathbf{f}\) or \(\mathbf{F}\) conversion specifier.

C
s

If no 1 qualifier is present, the int argument is converted to an unsigned char, and the resulting character is written. Otherwise, the wint_t argument is converted as if by an ls conversion specification with no precision and an argument that points to the initial element of a two-element array of wchar_t, the first element containing the wint_t argument to the lc conversion specification and the second a null wide character.

If no 1 qualifier is present, the argument shall be a pointer to the initial element of an array of character type. \({ }^{211}\) Characters from the array are
209. The precision \(p\) is sufficient to distinguish values of the source type if
\[
16^{p-1}>b^{n}
\]
where \(b\) is \(\operatorname{FLT}\) _RADIX and \(n\) is the number of base- \(b\) digits in the significand of the source type. A smaller \(p\) might suffice depending on the implementation's scheme for determining the digit to the left of the decimal-point character.
210. Binary implementations can choose the hexadecimal digit to the left of the decimal-point character so that subsequent digits align to nibble (4-bit) boundaries.
211. No special provisions are made for multibyte characters.
written up to (but not including) a terminating null character. If the precision is specified, no more than that many characters are written. If the precision is not specified or is greater than the size of the array, the array shall contain a null character.

If an 1 qualifier is present, the argument shall be a pointer to the initial element of an array of wchar_t type. Wide characters from the array are converted to multibyte characters (each as if by a call to the wcrtomb function, with the conversion state described by an mbstate_t object initialized to zero before the first wide character is converted) up to and including a terminating null wide character. The resulting multibyte characters are written up to (but not including) the terminating null character (byte). If no precision is specified, the array shall contain a null wide character. If a precision is specified, no more than that many characters (bytes) are written (including shift sequences, if any), and the array shall contain a null wide character if, to equal the multibyte character sequence length given by the precision, the function would need to access a wide character one past the end of the array. In no case is a partial multibyte character written. \({ }^{212}\)
p The argument shall be a pointer to void. The value of the pointer is converted to a sequence of printable characters, in an implementationdefined manner.
\(\mathbf{n} \quad\) The argument shall be a pointer to signed integer into which is written the number of characters written to the output stream so far by this call to fprintf. No argument is converted, but one is consumed. If the conversion specification with this conversion specifier is not one of \(\% \mathrm{n}\), \(\% \ln , \% \ln , \% \mathrm{hn}\), or \(\% \mathrm{hhn}\), the behavior is undefined.
\%
A \% is written. No argument is converted. The complete conversion specification shall be \(\% \%\).
7 If a conversion specification is invalid, the behavior is undefined. \({ }^{213}\)
8 If any argument is, or points to, a union or an aggregate (except for an array of char type using \%s conversion, an array of wchar_t type using \%ls conversion, or a pointer using \(\%\) p conversion), the behavior is undefined.

\footnotetext{
212. Redundant shift sequences may result if multibyte characters have a state-dependent encoding.
213. See "future library directions" (7.20.5).
}

9 In no case does a nonexistent or small field width cause truncation of a field; if the result of a conversion is wider than the field width, the field is expanded to contain the conversion result.

10 For a and A conversions, if FLT_RADIX is a power of 2, the value is correctly rounded to a hexadecimal floating number with the given precision.

\section*{Recommended practice}

11 If FLT_RADIX is not a power of 2, the result is one of the two adjacent numbers in hexadecimal floating style with the given precision, with the extra stipulation that the error have a correct sign for the current rounding direction.

12 For \(\mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}\), and \(\mathbf{G}\) conversions, if the number of significant decimal digits is at most DECIMAL_DIG, then the result is correctly rounded. \({ }^{214}\) If the number of significant decimal digits is more than DECIMAL_DIG but the source value is exactly representable with DECIMAL_DIG digits, then the result is an exact representation with trailing zeros. Otherwise, the source value is bounded by two adjacent decimal strings \(L<U\), both having DECIMAL_DIG significant digits; the value of the resultant decimal string \(D\) satisfies \(L \leq D \leq U\), with the extra stipulation that the error have a correct sign for the current rounding direction.

\section*{Returns}

13 The fprintf function returns the number of characters transmitted, or a negative value if an output error occurred.

\section*{Environmental limit}

14 The minimum value for the maximum number of characters produced by any single conversion shall be 4095.

\section*{Examples}

15 To print a date and time in the form 'Sunday, July 3, 10:02', followed by \(\pi\) to five decimal places:

\footnotetext{
214. For binary-to-decimal conversion, the result format's values are the numbers representable with the given format specifier. The number of significant digits is determined by the format specifier, and in the case of fixed-point conversion by the source value as well.
}
```

\#include <math.h>
\#include <stdio.h>
/* ... */
char *weekday, *month; // pointers to strings
int day, hour, min;
fprintf(stdout, "%s, %s %d, %.2d:%.2d\n",
weekday, month, day, hour, min);
fprintf(stdout, "pi = %.5f\n", 4 * atan(1.0));

```

16 In this example, multibyte characters do not have a state-dependent encoding, and the multibyte members of the extended character set each consist of two bytes, the first of which is denoted here by a \(\$ 0\) and the second by an uppercase letter.

17 Given the following wide string with length seven,
```

static wchar_t wstr[] = L"$OX$OYabc$OZ$OW";

```
the seven calls
```

fprintf(stdout, "|1234567890123|\n");
fprintf(stdout, "|%13ls|\n", wstr);
fprintf(stdout, "|%-13.9ls|\n", wstr);
fprintf(stdout, "|%13.101s|\n", wstr);
fprintf(stdout, "|%13.11s|\n", wstr);
fprintf(stdout, "|%13.15ls|\n", \&wstr[2]);
fprintf(stdout, "|%13lc|\n", wstr[5]);

```

Forward references: conversion state (7.19.7), the wartomb function (7.19.7.3.3).

\subsection*{7.13.6.2 The fscanf function}

Synopsis
```

\#include <stdio.h>
int fscanf(FILE * restrict stream,
const char * restrict format, ...);

```

\section*{Description}

2 The fscanf function reads input from the stream pointed to by stream, under control of the string pointed to by format that specifies the admissible input sequences and how they are to be converted for assignment, using subsequent arguments as pointers to the objects to receive the converted input. If there are insufficient arguments for the format, the behavior is undefined. If the format is exhausted while arguments remain, the excess arguments are evaluated (as always) but are otherwise ignored.

3 The format shall be a multibyte character sequence, beginning and ending in its initial shift state. The format is composed of zero or more directives: one or more whitespace characters; an ordinary multibyte character (neither \% nor a white-space character); or a conversion specification. Each conversion specification is introduced by the character \%. After the \%, the following appear in sequence:
- An optional assignment-suppressing character *.
- An optional nonzero decimal integer that specifies the maximum field width.
- An optional hh, h, lell), or \(1 \mathbf{l}\) (ell-ell), or \(\mathbf{L}\) indicating the size of the receiving object. The conversion specifiers \(\mathbf{d}\), \(\mathbf{i}\), and \(\mathbf{n}\) shall be preceded by \(\mathbf{h h}\) if the corresponding argument is a pointer to signed char rather than a pointer to int, by \(\mathbf{h}\) if it is a pointer to short int rather than a pointer to int, by \(\mathbf{l}\) if it is a pointer to long int, or by \(l l\) if it is a pointer to long long int. Similarly, the conversion specifiers \(\mathbf{o}, \mathbf{u}\), and \(\mathbf{x}\) shall be preceded by \(\mathbf{h h}\) if the corresponding argument is a pointer to unsigned char rather than a pointer to unsigned int, by \(h\) if it is a pointer to unsigned short int rather than a pointer to unsigned int, by 1 if it is a pointer to unsigned long int, or by 11 if it is a pointer to unsigned long long int. The conversion specifiers \(\mathbf{a}, \mathbf{e}, \mathbf{f}\), and \(\mathbf{g}\) shall be preceded by \(\mathbf{l}\) if the corresponding argument is a pointer to double rather than a pointer to float, or by \(L\) if it is a pointer to long double. Finally, the conversion specifiers \(\mathbf{c}, \mathbf{s}\), and [ shall be preceded by \(\mathbf{l}\) if the corresponding argument is a pointer to wchar_t rather than a pointer to a character type. If an \(\mathbf{h h}, \mathbf{h}, \mathbf{l}, \mathbf{l l}\), or \(\mathbf{L}\) appears with any other conversion specifier, the behavior is undefined.
- A character that specifies the type of conversion to be applied. The valid conversion specifiers are described below.

The fscanf function executes each directive of the format in turn. If a directive fails, as detailed below, the fscanf function returns. Failures are described as input failures (if an encoding error occurs or due to the unavailability of input characters), or matching failures (due to inappropriate input).

5 A directive composed of white-space character(s) is executed by reading input up to the first non-white-space character (which remains unread), or until no more characters can be read.

6 A directive that is an ordinary multibyte character is executed by reading the next characters of the stream. If one of the characters differs from one comprising the directive, the directive fails, and the differing and subsequent characters remain unread.

7 A directive that is a conversion specification defines a set of matching input sequences, as described below for each specifier. A conversion specification is executed in the following steps:

8 Input white-space characters (as specified by the isspace function) are skipped, unless the specification includes a [, c, or \(\mathbf{n}\) specifier. \({ }^{215}\)

9 An input item is read from the stream, unless the specification includes an \(\mathbf{n}\) specifier. An input item is defined as the longest sequence of input characters which does not exceed any specified field width and which is, or is a prefix of, a matching input sequence. The first character, if any, after the input item remains unread. If the length of the input item is zero, the execution of the directive fails; this condition is a matching failure unless end-of-file, an encoding error, or a read error prevented input from the stream, in which case it is an input failure.

10 Except in the case of a \% specifier, the input item (or, in the case of a \%n directive, the count of input characters) is converted to a type appropriate to the conversion specifier. If the input item is not a matching sequence, the execution of the directive fails: this condition is a matching failure. Unless assignment suppression was indicated by a *, the result of the conversion is placed in the object pointed to by the first argument following the format argument that has not already received a conversion result. If this object does not have an appropriate type, or if the result of the conversion cannot be represented in the object, the behavior is undefined.

11 The conversion specifiers and their meanings are:
d Matches an optionally signed decimal integer, whose format is the same as expected for the subject sequence of the strtol function with the value 10 for the base argument. The corresponding argument shall be a pointer to signed integer.

\footnotetext{
215. These white-space characters are not counted against a specified field width.
}
\(\mathbf{a} \mathbf{e}, \mathbf{f}, \mathbf{g}\) Matches an optionally signed floating-point constant, whose format is the same as expected for the subject string of the strtod function. The corresponding argument shall be a pointer to floating.
s Matches a sequence of non-white-space characters. \({ }^{216}\) If no 1 qualifier is present, the corresponding argument shall be a pointer to the initial element of a character array large enough to accept the sequence and a terminating null character, which will be added automatically.

If an 1 qualifier is present, the input shall be a sequence of multibyte characters that begins in the initial shift state. Each multibyte character is converted to a wide character as if by a call to the mbrtowc function, with the conversion state described by an mbstate_t object initialized to zero before the first multibyte character is converted. The corresponding argument shall be a pointer to the initial element of an array of wchar_t large enough to accept the sequence and the terminating null wide character, which will be added automatically.

\footnotetext{
216. No special provisions are made for multibyte characters in the matching rules used by any of the conversion specifiers s, [, c - the extent of the input field is still determined on a byte-by-byte basis. The resulting field must nevertheless be a sequence of multibyte characters that begins in the initial shift state.
}
[ Matches a nonempty sequence of characters from a set of expected characters (the scanset). If no 1 qualifier is present, the corresponding argument shall be a pointer to the initial element of a character array large enough to accept the sequence and a terminating null character, which will be added automatically.

If an 1 qualifier is present, the input shall be a sequence of multibyte characters that begins in the initial shift state. Each multibyte character is converted to a wide character as if by a call to the mbrtowc function, with the conversion state described by an mbstate_t object initialized to zero before the first multibyte character is converted. The corresponding argument shall be a pointer to the initial element of an array of wchar_t large enough to accept the sequence and the terminating null wide character, which will be added automatically. The conversion specifier includes all subsequent characters in the format string, up to and including the matching right bracket (]). The characters between the brackets (the scanlist) comprise the scanset, unless the character after the left bracket is a circumflex ( \(\wedge\) ), in which case the scanset contains all characters that do not appear in the scanlist between the circumflex and the right bracket. If the conversion specifier begins with [] or [^], the right bracket character is in the scanlist and the next right bracket character is the matching right bracket that ends the specification; otherwise the first right bracket character is the one that ends the specification. If a character is in the scanlist and is not the first, nor the second where the first character is a \(\wedge\), nor the last character, the behavior is implementation-defined.
c Matches a sequence of characters of exactly the number specified by the field width ( 1 if no field width is present in the directive). If no 1 qualifier is present, the corresponding argument shall be a pointer to the initial element of a character array large enough to accept the sequence. No null character is added.

If an 1 qualifier is present, the input shall be a sequence of multibyte characters that begins in the initial shift state. Each multibyte character in the sequence is converted to a wide character as if by a call to the mbrtowc function, with the conversion state described by an mbstate_t object initialized to zero before the first multibyte character is converted. The corresponding argument shall be a pointer to the initial element of an array of wchar_t large enough to accept the resulting sequence of wide characters. No null wide character is added.
p Matches an implementation-defined set of sequences, which should be the same as the set of sequences that may be produced by the \(\% \mathrm{p}\) conversion
of the fprintf function. The corresponding argument shall be a pointer to a pointer to void. The interpretation of the input item is implementation-defined. If the input item is a value converted earlier during the same program execution, the pointer that results shall compare equal to that value; otherwise the behavior of the \(\% \mathrm{p}\) conversion is undefined.
n

응

No input is consumed. The corresponding argument shall be a pointer to signed integer into which is to be written the number of characters read from the input stream so far by this call to the fscanf function. Execution of a \(\%\) directive does not increment the assignment count returned at the completion of execution of the fscanf function. No argument is converted, but one is consumed. If the conversion specification with this conversion specifier is not one of \(\% \mathbf{n}, \% \ln , \% \ln\), \(\% \mathrm{hn}\), or \%hhn, the behavior is undefined.

Matches a single \%; no conversion or assignment occurs. The complete conversion specification shall be \(\% \%\).

12 If a conversion specification is invalid, the behavior is undefined. \({ }^{217}\)
13 The conversion specifiers \(\mathbf{A}, \mathbf{E}, \mathbf{G}\), and \(\mathbf{x}\) are also valid and behave the same as, respectively, a, e, g, and \(\mathbf{x}\).

14 If end-of-file is encountered during input, conversion is terminated. If end-of-file occurs before any characters matching the current directive have been read (other than leading white space, where permitted), execution of the current directive terminates with an input failure; otherwise, unless execution of the current directive is terminated with a matching failure, execution of the following directive (if any) is terminated with an input failure.

15 Trailing white space (including new-line characters) is left unread unless matched by a directive. The success of literal matches and suppressed assignments is not directly determinable other than via the \%n directive.

16 If conversion terminates on a conflicting input character, the offending input character is left unread in the input stream. \({ }^{218}\)

\footnotetext{
217. See "future library directions" (7.20.5).
218. fscanf pushes back at most one input character onto the input stream. Therefore, some sequences that are acceptable to strtod, strtol, etc., are unacceptable to fscanf.
}

\section*{Returns}

17 The fscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the fscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\section*{Examples}
1. The call:
```

\#include <stdio.h>
/* ... */
int n, i; float x; char name[50];
n = fscanf(stdin, "%d%f%s", \&i, \&x, name);

```
with the input line:

\section*{25 54.32E-1 thompson}
will assign to \(\mathbf{n}\) the value 3 , to \(\mathbf{i}\) the value 25 , to \(\mathbf{x}\) the value 5.432 , and name will contain thompson \(\backslash 0\).
2. The call:
```

\#include <stdio.h>
/* ... */
int i; float x; char name[50];
fscanf(stdin, "%2d%f%*d %[0123456789]", \&i, \&x, name);

```
with input:
567890123 56a72
will assign to \(\mathbf{i}\) the value 56 and to \(\mathbf{x}\) the value 789.0 , will skip 0123, and name will contain \(56 \backslash 0\). The next character read from the input stream will be a.
3. To accept repeatedly from stdin a quantity, a unit of measure, and an item name:
```

\#include <stdio.h>
/* ... */
int count; float quant; char units[21], item[21];
while (!feof(stdin) \&\& !ferror(stdin)) {
count = fscanf(stdin, "%f%20s of %20s",
\&quant, units, item);
fscanf(stdin,"%*[^\n]");
}

```

If the stdin stream contains the following lines:
```

2 quarts of oil
-12.8degrees Celsius
lots of luck
10.0LBS of
dirt
100ergs of energy

```
the execution of the above example will be analogous to the following assignments:
```

quant = 2; strcpy(units, "quarts"); strcpy(item, "oil");
count = 3;
quant = -12.8; strcpy(units, "degrees");
count = 2; // "C" fails to match "०"
count = 0; // "l" fails to match "%£"
quant = 10.0; strcpy(units, "LBS"); strcpy(item, "dirt");
count = 3;
count = 0; // "100e" fails to match "%f"
count = EOF;

```
4. In:
```

\#include <stdio.h>
/* ... */
int d1, d2, n1, n2, i;
i = sscanf("123", "%d%n%n%d", \&d1, \&n1, \&n2, \&d2);

```
the value 123 is assigned to d 1 and the value 3 to n 1 . Because \(\% \mathrm{n}\) can never get an input failure the value of 3 is also assigned to \(\mathbf{n} 2\). The value of \(\mathbf{d 2}\) is not affected. The value 1 is assigned to \(\mathbf{i}\).

In these examples, multibyte characters do have a state-dependent encoding, and multibyte members of the extended character set consist of two bytes, the first of which is denoted here by a \(\$ 0\) and the second by an uppercase letter, but are only recognized as such when in the alternate shift state. The shift sequences are denoted by Library
\(\uparrow\) and \(\downarrow\), in which the first causes entry into the alternate shift state.
1. After the call:
```

\#include <stdio.h>
/* ... */
char str[50];
fscanf(stdin, "a%s", str);

```
with the input line:
\(a \uparrow \$ 0 X \$ 0 \mathrm{Y} \downarrow \mathrm{bc}\)
str will contain \(\uparrow \mathbf{\$ 0 X} \mathbf{\$ 0 Y} \downarrow \mathbf{0}\) assuming that none of the bytes of the shift sequences (or of the multibyte characters, in the more general case) appears to be a single-byte white-space character.
2. In contrast, after the call:
```

\#include <stdio.h>
\#include <stddef.h>
/* ... */
wchar_t wstr[50];
fscanf(stdin, "a%ls", wstr);

```
with the same input line, wstr will contain the two wide characters that correspond to \(\mathbf{\$ 0 X}\) and \(\mathbf{\$ 0 Y}\) and a terminating null wide character.
3. However, the call:
```

\#include <stdio.h>
\#include <stddef.h>
/* ... */
wchar_t wstr[50];
fscanf(stdin, "a`\$0x\downarrow%ls", wstr);

```
with the same input line will return zero due to a matching failure against the \(\downarrow\) sequence in the format string.
4. Assuming that the first byte of the multibyte character \(\mathbf{\$ 0} \mathbf{X}\) is the same as the first byte of the multibyte character \(\mathbf{\$ 0 Y}\), after the call:
```

\#include <stdio.h>
\#include <stddef.h>
/* ... */
wchar_t wstr[50];
fscanf(stdin, "a`\$0Y\downarrow%ls", wstr);

```
with the same input line, zero will again be returned, but stdin will be left with a partially consumed multibyte character.

Forward references: the strtod function (7.14.1.4), the strtol function (7.14.1.5), the strtoul function (7.14.1.6), conversion state (7.19.7), the wcrtomb function (7.19.7.3.3).

\subsection*{7.13.6.3 The printf function}

Synopsis
```

\#include <stdio.h>
int printf(const char * restrict format, ...);

```

\section*{Description}

2 The printf function is equivalent to fprintf with the argument stdout interposed before the arguments to printf.

\section*{Returns}

3 The printf function returns the number of characters transmitted, or a negative value if an output error occurred.

\subsection*{7.13.6.4 The scanf function}

\section*{Synopsis}
```

\#include <stdio.h>
int scanf(const char * restrict format, ...);

```

\section*{Description}

2 The scanf function is equivalent to fscanf with the argument stdin interposed before the arguments to scanf.

\section*{Returns}

3 The scanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the scanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\subsection*{7.13.6.5 The sprintf function}

Synopsis
```

\#include <stdio.h>
int sprintf(char * restrict s,
const char * restrict format, ...);

```

\section*{Description}

2 The sprintf function is equivalent to fprintf, except that the argument s specifies an array into which the generated output is to be written, rather than to a stream. A null character is written at the end of the characters written; it is not counted as part of the returned sum. If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

3 The sprintf function returns the number of characters written in the array, not counting the terminating null character.

\subsection*{7.13.6.6 The snprintf function}

\section*{Synopsis}
```

\#include <stdio.h>
int snprintf(char * restrict s, size_t n,
const char * restrict format, ...);

```

\section*{Description}

2 The snprintf function is equivalent to fprintf, except that the argument s specifies an array into which the generated output is to be written, rather than to a stream. If \(\mathbf{n}\) is zero, nothing is written, and \(\mathbf{s}\) may be a null pointer. Otherwise, output characters beyond the \(\mathbf{n}-1\) st are discarded rather than being written to the array, and a null character is written at the end of the characters actually written into the array. If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

3 The snprintf function returns the number of characters that would have been written had \(\mathbf{n}\) been sufficiently large, not counting the terminating null character. Thus, the null-terminated output has been completely written if and only if the returned value is less than \(\mathbf{n}\).

\subsection*{7.13.6.7 The sscanf function}

Synopsis
```

\#include <stdio.h>
int sscanf(const char * restrict s,
const char * restrict format, ...);

```

\section*{Description}

2 The sscanf function is equivalent to \(\mathbf{f s c a n f}\), except that the argument \(\mathbf{s}\) specifies a string from which the input is to be obtained, rather than from a stream. Reaching the end of the string is equivalent to encountering end-of-file for the fscanf function. If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

3 The sscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the sscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\subsection*{7.13.6.8 The vfprintf function}

Synopsis
```

\#include <stdarg.h>
\#include <stdio.h>
int vfprintf(FILE * restrict stream,
const char * restrict format,
va_list arg);

```

\section*{Description}

2 The vfprintf function is equivalent to fprintf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vfprintf function does not invoke the va_end macro. \({ }^{219}\)

\footnotetext{
219. As the functions vfprintf, vfscanf, vprintf, vscanf, vsnprintf, vsprintf, and vsscanf invoke the va_arg macro, the value of arg after the return is indeterminate.
}

\section*{Returns}

3 The vfprintf function returns the number of characters transmitted, or a negative value if an output error occurred.

Examples
4 The following shows the use of the vfprintf function in a general error-reporting routine.
```

\#include <stdarg.h>
\#include <stdio.h>
void error(char *function_name, char *format, ...)
{
va_list args;
va_start(args, format);
// print out name of function causing error
fprintf(stderr, "ERROR in %s: ", function_name);
// print out remainder of message
vfprintf(stderr, format, args);
va_end(args);
}

```

\subsection*{7.13.6.9 The vprintf function}

Synopsis
```

\#include <stdarg.h>
\#include <stdio.h>
int vprintf(const char * restrict format,
va_list arg);

```

Description
2 The vprintf function is equivalent to printf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vprintf function does not invoke the va_end macro. \({ }^{219}\)

\section*{Returns}

3 The vprintf function returns the number of characters transmitted, or a negative value if an output error occurred.

\subsection*{7.13.6.10 The vsprintf function}

Synopsis
```

\#include <stdarg.h>
\#include <stdio.h>
int vsprintf(char * restrict s,
const char * restrict format,
va_list arg);

```

\section*{Description}

2 The vsprintf function is equivalent to sprintf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vsprintf function does not invoke the va_end macro. \({ }^{219}\) If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

3 The vsprintf function returns the number of characters written in the array, not counting the terminating null character.

\subsection*{7.13.6.11 The vsnprintf function}

\section*{Synopsis}
```

\#include <stdarg.h>
\#include <stdio.h>
int vsprintf(char * restrict s, size_t n,
const char * restrict format,
va_list arg);

```

\section*{Description}

2 The vsnprintf function is equivalent to snprintf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vsnprintf function does not invoke the va_end macro. \({ }^{219}\) If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

3 The vsnprintf function returns the number of characters that would have been written had \(\mathbf{n}\) been sufficiently large, not counting the terminating null character. Thus, the null-terminated output has been completely written if and only if the returned value is less than \(n\).

\subsection*{7.13.6.12 The vfscanf function}

Synopsis
```

\#include <stdarg.h>
\#include <stdio.h>
int vfscanf(FILE * restrict stream,
const char * restrict format,
va_list arg);

```

\section*{Description}

2 The vfscanf function is equivalent to fscanf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vfscanf does not invoke the va_end macro. \({ }^{219}\)

\section*{Returns}

3 The vfscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the vfscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\subsection*{7.13.6.13 The vscanf function}

Synopsis
```

\#include <stdarg.h>
\#include <stdio.h>
int vscanf(const char * restrict format,
va_list arg);

```

\section*{Description}

2 The vscanf function is equivalent to scanf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vscanf does not invoke the va_end macro. \({ }^{219}\)

\section*{Returns}

3 The vscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the vscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\subsection*{7.13.6.14 The vsscanf function}

Synopsis
```

\#include <stdarg.h>
\#include <stdio.h>
int vsscanf(const char * restrict s,
const char * restrict format,
va_list arg);

```

\section*{Description}

2 The vsscanf function is equivalent to sscanf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vsscanf does not invoke the va_end macro. \({ }^{219}\)

\section*{Returns}

3 The vsscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the vscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\subsection*{7.13.7 Character input/output functions}

\subsection*{7.13.7.1 The fgetc function}

Synopsis
```

\#include <stdio.h>
int fgetc(FILE *stream);

```

\section*{Description}

2 If a next character is present from the input stream pointed to by stream, the fgetc function obtains that character as an unsigned char converted to an int and advances the associated file position indicator for the stream (if defined).

\section*{Returns}

3 The fgetc function returns the next character from the input stream pointed to by stream. If the stream is at end-of-file, the end-of-file indicator for the stream is set and fgetc returns EOF. If a read error occurs, the error indicator for the stream is
set and fgetc returns EOF. \({ }^{220}\)

\subsection*{7.13.7.2 The fgets function}

Synopsis
```

\#include <stdio.h>
char *fgets(char * restrict s, int n,
FILE * restrict stream);

```

\section*{Description}

2 The fgets function reads at most one less than the number of characters specified by \(\mathbf{n}\) from the stream pointed to by stream into the array pointed to by \(\mathbf{s}\). No additional characters are read after a new-line character (which is retained) or after end-of-file. A null character is written immediately after the last character read into the array.

\section*{Returns}

3 The fgets function returns \(\mathbf{s}\) if successful. If end-of-file is encountered and no characters have been read into the array, the contents of the array remain unchanged and a null pointer is returned. If a read error occurs during the operation, the array contents are indeterminate and a null pointer is returned.

\subsection*{7.13.7.3 The fputc function}

Synopsis
```

\#include <stdio.h>
int fputc(int c, FILE *stream);

```

\section*{Description}

2 The fputc function writes the character specified by \(\mathbf{c}\) (converted to an unsigned char) to the output stream pointed to by stream, at the position indicated by the associated file position indicator for the stream (if defined), and advances the indicator appropriately. If the file cannot support positioning requests, or if the stream was opened with append mode, the character is appended to the output stream.

\footnotetext{
220. An end-of-file and a read error can be distinguished by use of the feof and ferror functions.
}

\section*{Returns}

3 The fputc function returns the character written. If a write error occurs, the error indicator for the stream is set and fputc returns EOF.

\subsection*{7.13.7.4 The fputs function}

Synopsis
1 \#include <stdio.h>
int fputs (const char * restrict s, FILE * restrict stream);

\section*{Description}

2 The fputs function writes the string pointed to by \(\mathbf{s}\) to the stream pointed to by stream. The terminating null character is not written.

\section*{Returns}

3 The fputs function returns EOF if a write error occurs; otherwise it returns a nonnegative value.

\subsection*{7.13.7.5 The getc function}

Synopsis
```

\#include <stdio.h>
int getc(FILE *stream);

```

\section*{Description}

2 The getc function is equivalent to fgetc, except that if it is implemented as a macro, it may evaluate stream more than once, so the argument should never be an expression with side effects.

\section*{Returns}

3 The getc function returns the next character from the input stream pointed to by stream. If the stream is at end-of-file, the end-of-file indicator for the stream is set and getc returns EOF. If a read error occurs, the error indicator for the stream is set and getc returns EOF.

\subsection*{7.13.7.6 The getchar function}

Synopsis

\section*{Description}

2 The getchar function is equivalent to getc with the argument stdin.

\section*{Returns}

3 The getchar function returns the next character from the input stream pointed to by stdin. If the stream is at end-of-file, the end-of-file indicator for the stream is set and getchar returns EOF. If a read error occurs, the error indicator for the stream is set and getchar returns EOF.

\subsection*{7.13.7.7 The gets function}

\section*{Synopsis}
```

\#include <stdio.h>
char *gets(char *s);

```

\section*{Description}

2 The gets function reads characters from the input stream pointed to by stdin, into the array pointed to by \(\mathbf{s}\), until end-of-file is encountered or a new-line character is read. Any new-line character is discarded, and a null character is written immediately after the last character read into the array.

\section*{Returns}

3 The gets function returns \(\mathbf{s}\) if successful. If end-of-file is encountered and no characters have been read into the array, the contents of the array remain unchanged and a null pointer is returned. If a read error occurs during the operation, the array contents are indeterminate and a null pointer is returned.

\subsection*{7.13.7. 8 The putc function}

\section*{Synopsis}
```

\#include <stdio.h>
int putc(int c, FILE *stream);

```

\section*{Description}

2 The putc function is equivalent to fputc, except that if it is implemented as a macro, it may evaluate stream more than once, so the argument should never be an expression with side effects.

\section*{Returns}

3 The putc function returns the character written. If a write error occurs, the error indicator for the stream is set and putc returns EOF.

\subsection*{7.13.7.9 The putchar function}

Synopsis
1 \#include <stdio.h>
int putchar(int c);

\section*{Description}

2 The putchar function is equivalent to putc with the second argument stdout.

\section*{Returns}

3 The putchar function returns the character written. If a write error occurs, the error indicator for the stream is set and putchar returns EOF.

\subsection*{7.13.7.10 The puts function}

Synopsis
```

\#include <stdio.h>
int puts(const char *s);

```

\section*{Description}

2 The puts function writes the string pointed to by \(\mathbf{s}\) to the stream pointed to by stdout, and appends a new-line character to the output. The terminating null character is not written.

\section*{Returns}

3 The puts function returns EOF if a write error occurs; otherwise it returns a nonnegative value.

\subsection*{7.13.7.11 The ungetc function}

Synopsis
```

\#include <stdio.h>
int ungetc(int c, FILE *stream);

```

\section*{Description}

2 The ungetc function pushes the character specified by \(\mathbf{c}\) (converted to an unsigned char) back onto the input stream pointed to by stream. The pushedback characters will be returned by subsequent reads on that stream in the reverse order of their pushing. A successful intervening call (with the stream pointed to by stream) to a file positioning function (fseek, fsetpos, or rewind) discards any pushed-back characters for the stream. The external storage corresponding to the stream is unchanged.

3 One character of pushback is guaranteed. If the ungetc function is called too many times on the same stream without an intervening read or file positioning operation on that stream, the operation may fail.

4 If the value of \(\mathbf{c}\) equals that of the macro EOF, the operation fails and the input stream is unchanged.

5 A successful call to the ungetc function clears the end-of-file indicator for the stream. The value of the file position indicator for the stream after reading or discarding all pushed-back characters shall be the same as it was before the characters were pushed back. For a text stream, the value of its file position indicator after a successful call to the ungetc function is unspecified until all pushed-back characters are read or discarded. For a binary stream, its file position indicator is decremented by each successful call to the ungetc function; if its value was zero before a call, it is indeterminate after the call.

\section*{Returns}

6 The ungetc function returns the character pushed back after conversion, or EOF if the operation fails.

Forward references: file positioning functions (7.13.9).

\subsection*{7.13.8 Direct input/output functions}

\subsection*{7.13.8. The fread function}

Synopsis
```

\#include <stdio.h>
size_t fread(void * restrict ptr,
size_t size, size_t nmemb,
FILE * restrict stream);

```

\section*{Description}

2 The fread function reads, into the array pointed to by ptr, up to nmemb elements whose size is specified by size, from the stream pointed to by stream. The file position indicator for the stream (if defined) is advanced by the number of characters successfully read. If an error occurs, the resulting value of the file position indicator for the stream is indeterminate. If a partial element is read, its value is indeterminate.

\section*{Returns}

3 The fread function returns the number of elements successfully read, which may be less than nmemb if a read error or end-of-file is encountered. If size or nmemb is zero, fread returns zero and the contents of the array and the state of the stream remain unchanged.

\subsection*{7.13.8.2 The fwrite function}

\section*{Synopsis}
```

\#include <stdio.h>
size_t fwrite(const void * restrict ptr,
size_t size, size_t nmemb,
FILE * restrict stream);

```

\section*{Description}

2 The fwrite function writes, from the array pointed to by ptr, up to nmemb elements whose size is specified by size, to the stream pointed to by stream. The file position indicator for the stream (if defined) is advanced by the number of characters successfully written. If an error occurs, the resulting value of the file position indicator for the stream is indeterminate.

\section*{Returns}

3 The fwrite function returns the number of elements successfully written, which will be less than nmemb only if a write error is encountered.

\subsection*{7.13.9 File positioning functions}

\subsection*{7.13.9.1 The fgetpos function}

Synopsis
```

\#include <stdio.h>
int fgetpos(FILE * restrict stream,
fpos_t * restrict pos);

```

\section*{Description}

2 The fgetpos function stores the current value of the file position indicator for the stream pointed to by stream in the object pointed to by pos. The value stored contains unspecified information usable by the fsetpos function for repositioning the stream to its position at the time of the call to the fgetpos function.

\section*{Returns}

3 If successful, the fgetpos function returns zero; on failure, the fgetpos function returns nonzero and stores an implementation-defined positive value in errno.

Forward references: the fsetpos function (7.13.9.3).

\subsection*{7.13.9.2 The fseek function}

Synopsis
```

\#include <stdio.h>
int fseek(FILE *stream, long int offset, int whence);

```

Description
2 The fseek function sets the file position indicator for the stream pointed to by stream. If a read or write error occurs, the error indicator for the stream is set and fseek fails.

3 For a binary stream, the new position, measured in characters from the beginning of the file, is obtained by adding offset to the position specified by whence. The specified position is the beginning of the file if whence is SEEK_SET, the current value of the file position indicator if SEEK_CUR, or end-of-file if SEEK_END. A binary stream need not meaningfully support fseek calls with a whence value of SEEK_END.

4 For a text stream, either offset shall be zero, or offset shall be a value returned by an earlier successful call to the ftell function on a stream associated with the same file and whence shall be SEEK_SET.

5 After determining the new position, a successful call to the fseek function undoes any effects of the ungetc function on the stream, clears the end-of-file indicator for the stream, and then establishes the new position.

\section*{Returns}

6 The fseek function returns nonzero only for a request that cannot be satisfied.
Forward references: the ftell function (7.13.9.4).

\subsection*{7.13.9.3 The fsetpos function}

Synopsis
1 \#include <stdio.h>
int fsetpos(FILE *stream, const fpos_t *pos);

\section*{Description}

2 The fsetpos function sets the file position indicator for the stream pointed to by stream according to the value of the object pointed to by pos, which shall be a value obtained from an earlier successful call to the fgetpos function on a stream associated with the same file. If a read or write error occurs, the error indicator for the stream is set and fsetpos fails.

3 A successful call to the fsetpos function clears the end-of-file indicator for the stream and undoes any effects of the ungetc function on the same stream. After an fsetpos call, the next operation on an update stream may be either input or output.

\section*{Returns}

4 If successful, the fsetpos function returns zero; on failure, the fsetpos function returns nonzero and stores an implementation-defined positive value in errno.

\subsection*{7.13.9.4 The ftell function}

\section*{Synopsis}
```

\#include <stdio.h>
long int ftell(FILE *stream);

```

\section*{Description}

2 The ftell function obtains the current value of the file position indicator for the stream pointed to by stream. For a binary stream, the value is the number of characters from the beginning of the file. For a text stream, its file position indicator contains unspecified information, usable by the fseek function for returning the file position indicator for the stream to its position at the time of the ftell call; the difference between two such return values is not necessarily a meaningful measure of the number of characters written or read.

\section*{Returns}

3 If successful, the ftell function returns the current value of the file position indicator for the stream. On failure, the ftell function returns -1 L and stores an implementation-defined positive value in errno.

\subsection*{7.13.9.5 The rewind function}

Synopsis
1 \#include <stdio.h>
void rewind (FILE *stream);

\section*{Description}

2 The rewind function sets the file position indicator for the stream pointed to by stream to the beginning of the file. It is equivalent to
(void) fseek (stream, OL, SEEK_SET)
except that the error indicator for the stream is also cleared.

\section*{Returns}

3 The rewind function returns no value.

\subsection*{7.13.10 Error-handling functions}

\subsection*{7.13.10.1 The clearerr function}

Synopsis
```

\#include <stdio.h>
void clearerr(FILE *stream);

```

\section*{Description}

2 The clearerr function clears the end-of-file and error indicators for the stream pointed to by stream.

\section*{Returns}

3 The clearerr function returns no value.

\subsection*{7.13.10.2 The feof function}

Synopsis
```

\#include <stdio.h>
int feof(FILE *stream);

```

\section*{Description}

2 The feof function tests the end-of-file indicator for the stream pointed to by stream.

\section*{Returns}

3 The feof function returns nonzero if and only if the end-of-file indicator is set for stream.

\subsection*{7.13.10.3 The ferror function}

Synopsis
```

\#include <stdio.h>
int ferror(FILE *stream);

```

\section*{Description}

2 The ferror function tests the error indicator for the stream pointed to by stream.

\section*{Returns}

3 The ferror function returns nonzero if and only if the error indicator is set for stream.

\subsection*{7.13.10.4 The perror function}

Synopsis
```

\#include <stdio.h>
void perror(const char *s);

```

\section*{Description}

2 The perror function maps the error number in the integer expression errno to an error message. It writes a sequence of characters to the standard error stream thus: first (if \(\mathbf{s}\) is not a null pointer and the character pointed to by \(\mathbf{s}\) is not the null character), the string pointed to by \(\mathbf{s}\) followed by a colon (:) and a space; then an appropriate error message string followed by a new-line character. The contents of the error message strings are the same as those returned by the strerror function with argument errno.

\section*{Returns}

3 The perror function returns no value.
Forward references: the strerror function (7.15.6.2).

\subsection*{7.14 General utilities <stdlib.h>}

1 The header <stdlib.h> declares five types and several functions of general utility, and defines several macros. \({ }^{221}\)

2 The types declared are size_t and wchar_t (both described in 7.1.6),
div_t
which is a structure type that is the type of the value returned by the div function,
ldiv_t
which is a structure type that is the type of the value returned by the ldiv function, and
```

lldiv_t

```
which is a structure type that is the type of the value returned by the lldiv function.
The macros defined are NULL (described in 7.1.6);
EXIT_FAILURE
and
EXIT_SUCCESS
which expand to integer expressions that may be used as the argument to the exit function to return unsuccessful or successful termination status, respectively, to the host environment;

RAND_MAX
which expands to an integer constant expression, the value of which is the maximum value returned by the rand function; and

\section*{MB_CUR_MAX}
which expands to a positive integer expression whose value is the maximum number of bytes in a multibyte character for the extended character set specified by the current locale (category LC_CTYPE), and whose value is never greater than MB_LEN_MAX.
221. See "future library directions" (7.20.6).

\subsection*{7.14.1 String conversion functions}

1 The functions atof, atoi, atol, and atoll need not affect the value of the integer expression errno on an error. If the value of the result cannot be represented, the behavior is undefined.

\subsection*{7.14.1.1 The atof function}

Synopsis
```

\#include <stdlib.h>
double atof(const char *nptr);

```

\section*{Description}

2 The atof function converts the initial portion of the string pointed to by nptr to double representation. Except for the behavior on error, it is equivalent to
strtod (nptr, (char **) NULL)

\section*{Returns}

3 The atof function returns the converted value.
Forward references: the strtod function (7.14.1.5).

\subsection*{7.14.1.2 The atoi function}

Synopsis
```

\#include <stdlib.h>
int atoi(const char *nptr);

```

\section*{Description}

2 The atoi function converts the initial portion of the string pointed to by nptr to int representation. Except for the behavior on error, it is equivalent to
(int) strtol (nptr, (char **) NULL, 10)

\section*{Returns}

The atoi function returns the converted value.
Forward references: the strtol function (7.14.1.8).

\subsection*{7.14.1.3 The atol function}

Synopsis
```

\#include <stdlib.h>
long int atol(const char *nptr);

```

\section*{Description}

2 The atol function converts the initial portion of the string pointed to by nptr to long int representation. Except for the behavior on error, it is equivalent to
```

strtol(nptr, (char **)NULL, 10)

```

\section*{Returns}

3 The atol function returns the converted value.
Forward references: the strtol function (7.14.1.8).

\subsection*{7.14.1.4 The atoll function}

Synopsis
```

\#include <stdlib.h>
long long int atoll(const char *nptr);

```

\section*{Description}

2 The atoll function is equivalent to the atol function, except that it converts the initial portion of the string pointed to by nptr to long long int representation. Except for the behavior on error, it is equivalent to
```

strtoll(nptr, (char **)NULL, 10)

```

\section*{Returns}

3 The atoll function returns the converted value.
Forward references: the strtoll function (7.14.1.9).

\subsection*{7.14.1.5 The strtod function}

Synopsis
```

\#include <stdlib.h>
double strtod(const char * restrict nptr,
char ** restrict endptr);

```

\section*{Description}

2 The strtod function converts the initial portion of the string pointed to by nptr to double representation. First, it decomposes the input string into three parts: an initial, possibly empty, sequence of white-space characters (as specified by the isspace function), a subject sequence resembling a floating-point constant; and a final string of one or more unrecognized characters, including the terminating null character of the input string. Then, it attempts to convert the subject sequence to a floating-point number, and returns the result.

3 The expected form of the subject sequence is an optional plus or minus sign, then one of the following:
- a nonempty sequence of decimal digits optionally containing a decimal-point character, then an optional exponent part as defined in 6.1.3.1;
- a \(\mathbf{0 x}\) or \(\mathbf{0 x}\), then a nonempty sequence of hexadecimal digits optionally containing a decimal-point character, then an optional binary-exponent part as defined in 6.1.3.1, where either the decimal-point character or the binary-exponent part is present;
- one of INF or INFINITY, ignoring case
- one of NAN or NAN ( \(n\)-char-sequence \({ }_{o p t}\) ), ignoring case in the NAN part, where:
\[
\begin{aligned}
& \text { n-char-sequence: } \\
& \quad \text { digit } \\
& \text { nondigit } \\
& \text { n-char-sequence digit } \\
& \text { n-char-sequence nondigit }
\end{aligned}
\]
but no floating suffix. The subject sequence is defined as the longest initial subsequence of the input string, starting with the first non-white-space character, that is of the expected form. The subject sequence contains no characters if the input string is not of the expected form.

If the subject sequence has the expected form for a floating-point number, the sequence of characters starting with the first digit or the decimal-point character (whichever occurs first) is interpreted as a floating constant according to the rules of 6.1.3.1, except that the decimal-point character is used in place of a period, and that if neither an exponent part, a binary-exponent part, nor a decimal-point character appears, a decimal point is assumed to follow the last digit in the string. A character sequence INF or INFINITY is interpreted as an infinity, if representable in the double type, else like a floating constant that is too large for the range of double. A character sequence NAN or NAN (n-char-sequence \({ }_{\text {opt }}\) ), is interpreted as a quiet NaN, if supported in the double type, else like a subject sequence part that does not have the
expected form; the meaning of the \(n\)-char sequences is implementation-defined. \({ }^{222}\) If the subject sequence begins with a minus sign, the value resulting from the conversion is negated. \({ }^{223} \mathrm{~A}\) pointer to the final string is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

5 In other than the "C" locale, additional locale-specific subject sequence forms may be accepted.

6 If the subject sequence is empty or does not have the expected form, no conversion is performed; the value of nptr is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

7 If the subject sequence has the hexadecimal form and FLT_RADIX is a power of 2, then the value resulting from the conversion is correctly rounded.

\section*{Recommended practice}

8 If the subject sequence has the hexadecimal form and FLT_RADIX is not a power of 2 , then the result is one of the two numbers in the appropriate internal format that are adjacent to the hexadecimal floating source value, with the extra stipulation that the error have a correct sign for the current rounding direction.

9 If the subject sequence has the decimal form and at most DECIMAL_DIG (defined in <math. \(\mathrm{h}>\) ) significant digits, then the value resulting from the conversion is correctly rounded. If the subject sequence \(D\) has the decimal form and more than DECIMAL_DIG significant digits, consider the two bounding, adjacent decimal strings \(L\) and \(U\), both having DECIMAL_DIG significant digits, such that the values of \(L, D\), and \(U\) satisfy \(L \leq D \leq U\). The result of conversion is one of the (equal or adjacent) values that would be obtained by correctly rounding \(L\) and \(U\) according to the current rounding direction, with the extra stipulation that the error with respect to \(D\) has a correct sign for the current rounding direction. \({ }^{224}\)

\footnotetext{
222. An implementation may use the \(n\)-char-sequence to determine extra information to be represented in the NaN's significand.
223. The strtod function honors the sign of zero if the arithmetic supports signed zeros.
224. DECIMAL_DIG, defined in <math.h>, is recommended to be sufficiently large that \(L\) and \(U\) will usually round to the same internal floating value, but if not will round to adjacent values.
}

\section*{Returns}

10 The strtod function returns the converted value, if any. If no conversion could be performed, zero is returned. If the correct value is outside the range of representable values, plus or minus HUGE_VAL is returned (according to the sign of the value), and the value of the macro ERANGE is stored in errno. If the result underflows (7.7.1), the function returns a value whose magnitude is no greater than the smallest normalized positive number in the result type; whether errno acquires the value ERANGE is implementation-defined.

\subsection*{7.14.1.6 The strtof function}

Synopsis
1 \#include <stdlib.h>
float strtof(const char * restrict nptr, char ** restrict endptr);

\section*{Description}

2 The strtof function is similar to the strtod function, expect the returned value has type float and plus or minus HUGE_VALF is returned for values outside the range.

\subsection*{7.14.1.7 The strtold function}

Synopsis
```

\#include <stdlib.h>
long double strtold(const char * restrict nptr,
char ** restrict endptr);

```

\section*{Description}

2 The strtold function is similar to the strtod function, expect the returned value has type long double and plus or minus HUGE_VALL is returned for values outside the range.

\subsection*{7.14.1.8 The strtol function}

Synopsis
```

\#include <stdlib.h>
long int strtol(const char * restrict nptr,
char ** restrict endptr, int base);

```

\section*{Description}

2 The strtol function converts the initial portion of the string pointed to by nptr to long int representation. First, it decomposes the input string into three parts: an initial, possibly empty, sequence of white-space characters (as specified by the isspace function), a subject sequence resembling an integer represented in some radix determined by the value of base, and a final string of one or more unrecognized characters, including the terminating null character of the input string. Then, it attempts to convert the subject sequence to an integer, and returns the result.

3 If the value of base is zero, the expected form of the subject sequence is that of an integer constant as described in 6.1.3.2, optionally preceded by a plus or minus sign, but not including an integer suffix. If the value of base is between 2 and 36, the expected form of the subject sequence is a sequence of letters and digits representing an integer with the radix specified by base, optionally preceded by a plus or minus sign, but not including an integer suffix. The letters from \(\mathbf{a}\) (or \(\mathbf{A}\) ) through \(\mathbf{z}\) (or \(\mathbf{z}\) ) are ascribed the values 10 to 35 ; only letters whose ascribed values are less than that of base are permitted. If the value of base is 16 , the characters \(0 x\) or \(0 x\) may optionally precede the sequence of letters and digits, following the sign if present.

4 The subject sequence is defined as the longest initial subsequence of the input string, starting with the first non-white-space character, that is of the expected form. The subject sequence contains no characters if the input string is empty or consists entirely of white space, or if the first non-white-space character is other than a sign or a permissible letter or digit.

5 If the subject sequence has the expected form and the value of base is zero, the sequence of characters starting with the first digit is interpreted as an integer constant according to the rules of 6.1.3.2. If the subject sequence has the expected form and the value of base is between 2 and 36 , it is used as the base for conversion, ascribing to each letter its value as given above. If the subject sequence begins with a minus sign, the value resulting from the conversion is negated. A pointer to the final string is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

6 In other than the "C" locale, additional locale-specific subject sequence forms may be accepted.

7 If the subject sequence is empty or does not have the expected form, no conversion is performed; the value of nptr is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

\section*{Returns}

8 The strtol function returns the converted value, if any. If no conversion could be performed, zero is returned. If the correct value is outside the range of representable values, LONG_MAX or LONG_MIN is returned (according to the sign of the value), and the value of the macro ERANGE is stored in errno.

\subsection*{7.14.1.9 The strtoll function}

Synopsis
```

\#include <stdlib.h>
long long int strtoll(const char * restrict nptr,
char ** restrict endptr, int base);

```

\section*{Description}

2 The strtoll is equivalent to the strtol function, except that it converts the initial portion of the string pointed to by nptr to long long int representation.

\section*{Returns}

3 The strtoll function returns the converted value, if any. If no conversion could be performed, zero is returned. If the correct value is outside the range of representable values, LLONG_MAX or LLONG_MIN is returned (according to the sign of the value), and the value of the macro ERANGE is stored in errno.

\subsection*{7.14.1.10 The strtoul function}

\section*{Synopsis}
```

\#include <stdlib.h>
unsigned long int strtoul(
const char * restrict nptr,
char ** restrict endptr,
int base);

```

\section*{Description}

2 The strtoul function converts the initial portion of the string pointed to by nptr to unsigned long int representation. First, it decomposes the input string into three parts: an initial, possibly empty, sequence of white-space characters (as specified by the isspace function), a subject sequence resembling an unsigned integer represented in some radix determined by the value of base, and a final string of one or more unrecognized characters, including the terminating null character of the input string. Then, it attempts to convert the subject sequence to an unsigned integer, and returns the result.

3 If the value of base is zero, the expected form of the subject sequence is that of an integer constant as described in 6.1.3.2, optionally preceded by a plus or minus sign, but not including an integer suffix. If the value of base is between 2 and 36, the expected form of the subject sequence is a sequence of letters and digits representing an integer with the radix specified by base, optionally preceded by a plus or minus sign, but not including an integer suffix. The letters from \(\mathbf{a}\) (or \(\mathbf{A}\) ) through \(\mathbf{z}\) (or \(\mathbf{z}\) ) are ascribed the values 10 to 35 ; only letters whose ascribed values are less than that of base are permitted. If the value of base is 16 , the characters \(0 x\) or \(0 x\) may optionally precede the sequence of letters and digits, following the sign if present.

4 The subject sequence is defined as the longest initial subsequence of the input string, starting with the first non-white-space character, that is of the expected form. The subject sequence contains no characters if the input string is empty or consists entirely of white space, or if the first non-white-space character is other than a sign or a permissible letter or digit.

5 If the subject sequence has the expected form and the value of base is zero, the sequence of characters starting with the first digit is interpreted as an integer constant according to the rules of 6.1.3.2. If the subject sequence has the expected form and the value of base is between 2 and 36, it is used as the base for conversion, ascribing to each letter its value as given above. If the subject sequence begins with a minus sign, the value resulting from the conversion is negated. A pointer to the final string is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

6 In other than the "C" locale, additional locale-specific subject sequence forms may be accepted.

7 If the subject sequence is empty or does not have the expected form, no conversion is performed; the value of nptr is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

\section*{Returns}

8 The strtoul function returns the converted value, if any. If no conversion could be performed, zero is returned. If the correct value is outside the range of representable values, ULONG_MAX is returned, and the value of the macro ERANGE is stored in errno.

\subsection*{7.14.1.11 The strtoull function}

Synopsis
```

\#include <stdlib.h>
unsigned long long int strtoull(
const char * restrict nptr,
char ** restrict endptr,
int base);

```

\section*{Description}

2 The strtoull is equivalent to the strtoul function, except that it converts the initial portion of the string pointed to by nptr to unsigned long long int representation.

\section*{Returns}

3 The strtoull function returns the converted value, if any. If no conversion could be performed, zero is returned. If the correct value is outside the range of representable values, ULLONG_MAX is returned, and the value of the macro ERANGE is stored in errno.

\subsection*{7.14.2 Pseudo-random sequence generation functions}

\subsection*{7.14.2.1 The rand function}

Synopsis
```

\#include <stdlib.h>
int rand(void);

```

\section*{Description}

2 The rand function computes a sequence of pseudo-random integers in the range 0 to RAND_MAX.

3 The implementation shall behave as if no library function calls the rand function.

\section*{Returns}

4 The rand function returns a pseudo-random integer.
Environmental limit
5 The value of the RAND_MAX macro shall be at least 32767.

\subsection*{7.14.2.2 The srand function}

Synopsis
```

\#include <stdlib.h>
void srand(unsigned int seed);

```

\section*{Description}

2 The srand function uses the argument as a seed for a new sequence of pseudorandom numbers to be returned by subsequent calls to rand. If srand is then called with the same seed value, the sequence of pseudo-random numbers shall be repeated. If rand is called before any calls to srand have been made, the same sequence shall be generated as when srand is first called with a seed value of 1 .

3 The implementation shall behave as if no library function calls the srand function.

\section*{Returns}

The srand function returns no value.

\section*{Examples}

5 The following functions define a portable implementation of rand and srand.
```

static unsigned long int next = 1;
int rand(void) // RAND_MAX assumed to be 32767
{
next = next * 1103515245 + 12345;
}
void srand(unsigned int seed)
{
next = seed;
}

```
    return (unsigned int) (next/65536) \% 32768;

\subsection*{7.14.3 Memory management functions}

1 The order and contiguity of storage allocated by successive calls to the calloc, malloc, and realloc functions is unspecified. The pointer returned if the allocation succeeds is suitably aligned so that it may be assigned to a pointer to any type of object and then used to access such an object or an array of such objects in the space allocated (until the space is explicitly freed or reallocated). Each such allocation shall yield a pointer to an object disjoint from any other object. The pointer returned points to the start (lowest byte address) of the allocated space. If the space cannot be allocated, a null pointer is returned. If the size of the space requested is zero, the behavior is implementation-defined: either a null pointer is returned, or the behavior is as if the size were some nonzero value, except that the returned pointer shall not be used to access an object. The value of a pointer that refers to freed space is indeterminate.

\subsection*{7.14.3.1 The calloc function}

\section*{Synopsis}

1 \#include <stdlib.h> void *calloc (size_t nmemb, size_t size);

\section*{Description}

2 The calloc function allocates space for an array of nmemb objects, each of whose size is size. The space is initialized to all bits zero. \({ }^{225}\)

\section*{Returns}

3 The calloc function returns either a null pointer or a pointer to the allocated space.

\subsection*{7.14.3.2 The free function}

\section*{Synopsis}
```

\#include <stdlib.h>
void free(void *ptr);

```

\footnotetext{
225. Note that this need not be the same as the representation of floating-point zero or a null pointer constant.
}

\section*{Description}

2 The free function causes the space pointed to by ptr to be deallocated, that is, made available for further allocation. If ptr is a null pointer, no action occurs. Otherwise, if the argument does not match a pointer earlier returned by the calloc, malloc, or realloc function, or if the space has been deallocated by a call to free or realloc, the behavior is undefined.

\section*{Returns}

3 The free function returns no value.

\subsection*{7.14.3.3 The malloc function}

\section*{Synopsis}
```

\#include <stdlib.h>
void *malloc(size_t size);

```

\section*{Description}

2 The malloc function allocates space for an object whose size is specified by size and whose value is indeterminate.

\section*{Returns}

3 The malloc function returns either a null pointer or a pointer to the allocated space.

\subsection*{7.14.3.4 The realloc function}

\section*{Synopsis}
```

\#include <stdlib.h>
void *realloc(void *ptr, size_t size);

```

\section*{Description}

2 The realloc function changes the size of the object pointed to by ptr to the size specified by size. The contents of the object shall be unchanged up to the lesser of the new and old sizes. If the new size is larger, the value of the newly allocated portion of the object is indeterminate. If ptr is a null pointer, the realloc function behaves like the malloc function for the specified size. Otherwise, if ptr does not match a pointer earlier returned by the calloc, malloc, or realloc function, or if the space has been deallocated by a call to the free or realloc function, the behavior is undefined. If the space cannot be allocated, the object pointed to by ptr is unchanged. If size is zero and ptr is not a null pointer, the object it points to is freed.

\section*{Returns}

3 The realloc function returns either a null pointer or a pointer to the possibly moved allocated space. If a pointer is returned which does not compare equal to ptr, then the object has moved and ptr is a pointer that refers to freed space.

\subsection*{7.14.4 Communication with the environment}

\subsection*{7.14.4.1 The abort function}

Synopsis
1 \#include <stdlib.h> void abort (void);

\section*{Description}

2 The abort function causes abnormal program termination to occur, unless the signal SIGABRT is being caught and the signal handler does not return. Whether open output streams are flushed or open streams closed or temporary files removed is implementation-defined. An implementation-defined form of the status unsuccessful termination is returned to the host environment by means of the function call raise (SIGABRT).

\section*{Returns}

3 The abort function cannot return to its caller.

\subsection*{7.14.4.2 The atexit function}

Synopsis
```

\#include <stdlib.h>
int atexit(void (*func) (void));

```

\section*{Description}

2 The atexit function registers the function pointed to by func, to be called without arguments at normal program termination.

\section*{Implementation limits}

3 The implementation shall support the registration of at least 32 functions.

\section*{Returns}

4 The atexit function returns zero if the registration succeeds, nonzero if it fails.

Forward references: the exit function (7.14.4.3).

\subsection*{7.14.4.3 The exit function}

Synopsis
1 \#include <stdlib.h>
void exit(int status);

\section*{Description}

2 The exit function causes normal program termination to occur. If more than one call to the exit function is executed by a program, the behavior is undefined.

3 First, all functions registered by the atexit function are called, in the reverse order of their registration. \({ }^{226}\)

4 Next, all open streams with unwritten buffered data are flushed, all open streams are closed, and all files created by the tmpfile function are removed.

5 Finally, control is returned to the host environment. If the value of status is zero or EXIT_SUCCESS, an implementation-defined form of the status successful termination is returned. If the value of status is EXIT_FAILURE, an implementation-defined form of the status unsuccessful termination is returned. Otherwise the status returned is implementation-defined.

\section*{Returns}

6 The exit function cannot return to its caller.

\subsection*{7.14.4.4 The getenv function}

\section*{Synopsis}
```

\#include <stdlib.h>
char *getenv(const char *name);

```

\section*{Description}

2 The getenv function searches an environment list, provided by the host environment, for a string that matches the string pointed to by name. The set of environment names and the method for altering the environment list are implementation-defined.

3 The implementation shall behave as if no library function calls the getenv function.
226. Each function is called as many times as it was registered.

\section*{Returns}

4 The getenv function returns a pointer to a string associated with the matched list member. The string pointed to shall not be modified by the program, but may be overwritten by a subsequent call to the getenv function. If the specified name cannot be found, a null pointer is returned.

\subsection*{7.14.4.5 The system function}

Synopsis
```

\#include <stdlib.h>
int system(const char *string);

```

\section*{Description}

2 If string is a null pointer, the system function determines whether the host environment has a command processor. If string is not a null pointer, the system function passes the string pointed to by string to that command processor to be executed in a manner which the implementation shall document; this might then cause the program calling system to behave in a non-conforming manner or to terminate.

\section*{Returns}

3 If the argument is a null pointer, the system function returns nonzero only if a command processor is available. If the argument is not a null pointer, and the system function does return, it returns an implementation-defined value.

\subsection*{7.14.5 Searching and sorting utilities}

1 These utilities make use of a comparison function.
2 The implementation shall ensure that the second argument of the comparison function (when called from bsearch), or both arguments (when called from qsort), shall be pointers to elements of the array. \({ }^{227}\) The first argument when called from bsearch shall equal key.

3 The comparison function shall not alter the contents of the array. The implementation may reorder elements of the array between calls to the comparison function, but shall not alter the contents of any individual element.
227. That is, if the value passed is \(\mathbf{p}\), then the following expressions are always non-zero:
```

((char *p)p - (char *)base) % size == 0
(char *)p >= (char *)base

```

4 When the same object (consisting of size bytes, irrespective of its current position in the array) is passed more than once to the comparison function, the results shall be consistent with one another. That is, for qsort they shall define a total ordering on the array, and for bsearch the same object shall always compare the same way with the key.

5 A sequence point occurs immediately before and immediately after each call to the comparison function, and also between any call to the comparison function and any movement of the objects passed as arguments to that call.

\subsection*{7.14.5.1 The bsearch function}

\section*{Synopsis}
```

\#include <stdlib.h>
void *bsearch(const void *key, const void *base,
size_t nmemb, size_t size,
int (*compar) (const void *, const void *));

```

\section*{Description}

The bsearch function searches an array of nmemb objects, the initial element of which is pointed to by base, for an element that matches the object pointed to by key. The size of each element of the array is specified by size.

3 The comparison function pointed to by compar is called with two arguments that point to the key object and to an array element, in that order. The function shall return an integer less than, equal to, or greater than zero if the key object is considered, respectively, to be less than, to match, or to be greater than the array element. The array shall consist of: all the elements that compare less than, all the elements that compare equal to, and all the elements that compare greater than the key object, in that order. \({ }^{228}\)

\section*{Returns}

4 The bsearch function returns a pointer to a matching element of the array, or a null pointer if no match is found. If two elements compare as equal, which element is matched is unspecified.

\footnotetext{
228. In practice, the entire array is sorted according to the comparison function.
}

\subsection*{7.14.5.2 The qsort function}

Synopsis
```

\#include <stdlib.h>
void qsort (void *base, size_t nmemb, size_t size,
int (*compar) (const void *, const void *));

```

\section*{Description}

2 The qsort function sorts an array of nmemb objects, the initial element of which is pointed to by base. The size of each object is specified by size.

3 The contents of the array are sorted into ascending order according to a comparison function pointed to by compar, which is called with two arguments that point to the objects being compared. The function shall return an integer less than, equal to, or greater than zero if the first argument is considered to be respectively less than, equal to, or greater than the second.

4 If two elements compare as equal, their order in the sorted array is unspecified.

\section*{Returns}

5 The qsort function returns no value.

\subsection*{7.14.6 Integer arithmetic functions}

\subsection*{7.14.6.1 The abs function}

Synopsis
```

\#include <stdlib.h>
int abs(int j);

```

\section*{Description}

2 The abs function computes the absolute value of an integer \(\mathbf{j}\). If the result cannot be represented, the behavior is undefined. \({ }^{229}\)

\section*{Returns}

The abs function returns the absolute value.
229. The absolute value of the most negative number cannot be represented in two's complement.

\subsection*{7.14.6.2 The div function}

Synopsis
```

\#include <stdlib.h>
div_t div(int numer, int denom);

```

\section*{Description}

2 The div function computes the quotient and remainder of the division of the numerator numer by the denominator denom. The returned quotient is the algebraic quotient with any fractional part discarded. If the result cannot be represented, the behavior is undefined; otherwise, quot * denom + rem shall equal numer.

\section*{Returns}

3 The div function returns a structure of type div_t, comprising both the quotient and the remainder. The structure shall contain the following members, in either order:
```

int quot; // quotient
int rem; // remainder

```

\subsection*{7.14.6.3 The labs function}

Synopsis
```

\#include <stdlib.h>
long int labs(long int j);

```

\section*{Description}

2 The labs function is equivalent to the abs function, except that the argument and the returned value each have type long int.

\subsection*{7.14.6.4 The llabs function}

Synopsis
```

\#include <stdlib.h>
long long int llabs(long long int j);

```

\section*{Description}

2 The llabs function is equivalent to the abs function, except that the argument and the returned value each have type long long int.

\subsection*{7.14.6.5 The ldiv function}

\section*{Synopsis}
```

\#include <stdlib.h>
ldiv_t ldiv(long int numer, long int denom);

```

\section*{Description}

2 The ldiv function is equivalent to the div function, except that the arguments and the members of the returned structure (which has type ldiv_t) all have type long int.

\subsection*{7.14.6.6 The lldiv function}

Synopsis
```

\#include <stdlib.h>
lldiv_t lldiv(long long int numer,
long long int denom);

```

\section*{Description}

2 The lldiv function is equivalent to the div function, except that the arguments and the members of the returned structure (which has type lldiv_t) all have type long long int.

\subsection*{7.14.7 Multibyte character functions}

1 The behavior of the multibyte character functions is affected by the LC_CTYPE category of the current locale. For a state-dependent encoding, each function is placed into its initial state by a call for which its character pointer argument, \(\mathbf{s}\), is a null pointer. Subsequent calls with \(\mathbf{s}\) as other than a null pointer cause the internal state of the function to be altered as necessary. A call with \(\mathbf{s}\) as a null pointer causes these functions to return a nonzero value if encodings have state dependency, and zero otherwise. \({ }^{230}\) Changing the LC_CTYPE category causes the shift state of these functions to be indeterminate.

\footnotetext{
230. If the locale employs special bytes to change the shift state, these bytes do not produce separate wide character codes, but are grouped with an adjacent multibyte character.
}

\subsection*{7.14.7.1 The mblen function}

Synopsis
```

\#include <stdlib.h>
int mblen(const char *s, size_t n);

```

\section*{Description}

2 If \(\boldsymbol{s}\) is not a null pointer, the mblen function determines the number of bytes contained in the multibyte character pointed to by s. Except that the shift state of the mbtowc function is not affected, it is equivalent to
```

mbtowc((wchar_t *)0, s, n);

```

3 The implementation shall behave as if no library function calls the mblen function.

\section*{Returns}

4 If \(\boldsymbol{s}\) is a null pointer, the mblen function returns a nonzero or zero value, if multibyte character encodings, respectively, do or do not have state-dependent encodings. If \(\boldsymbol{s}\) is not a null pointer, the mblen function either returns 0 (if \(\mathbf{s}\) points to the null character), or returns the number of bytes that are contained in the multibyte character (if the next \(\mathbf{n}\) or fewer bytes form a valid multibyte character), or returns - 1 (if they do not form a valid multibyte character).

Forward references: the mbtowc function (7.14.7.2).

\subsection*{7.14.7.2 The mbtowc function}

Synopsis
```

\#include <stdlib.h>
int mbtowc(wchar_t * restrict pwc,
const char * restrict s,
size_t n);

```

\section*{Description}

2 If \(\mathbf{s}\) is not a null pointer, the mbtowc function determines the number of bytes that are contained in the multibyte character pointed to by s. It then determines the code for the value of type wchar_t that corresponds to that multibyte character. (The value of the code corresponding to the null character is zero.) If the multibyte character is valid and pwc is not a null pointer, the mbtowc function stores the code in the object pointed to by pwc. At most \(\mathbf{n}\) bytes of the array pointed to by \(\mathbf{s}\) will be examined.

3 The implementation shall behave as if no library function calls the mbtowc function.

\section*{Returns}

If \(\mathbf{s}\) is a null pointer, the mbtowc function returns a nonzero or zero value, if multibyte character encodings, respectively, do or do not have state-dependent encodings. If \(\boldsymbol{s}\) is not a null pointer, the mbtowc function either returns 0 (if \(\mathbf{s}\) points to the null character), or returns the number of bytes that are contained in the converted multibyte character (if the next \(\mathbf{n}\) or fewer bytes form a valid multibyte character), or returns - 1 (if they do not form a valid multibyte character).

4 In no case will the value returned be greater than \(\mathbf{n}\) or the value of the MB_CUR_MAX macro.

\subsection*{7.14.7.3 The wctomb function}

Synopsis
```

\#include <stdlib.h>
int wctomb(char *s, wchar_t wchar);

```

\section*{Description}

2 The wctomb function determines the number of bytes needed to represent the multibyte character corresponding to the code whose value is wchar (including any change in shift state). It stores the multibyte character representation in the array object pointed to by \(\mathbf{s}\) (if \(\mathbf{s}\) is not a null pointer). At most MB_CUR_MAX characters are stored. If the value of wchar is zero, the wctomb function is left in the initial shift state.

3 The implementation shall behave as if no library function calls the wctomb function.

\section*{Returns}

4 If \(\mathbf{s}\) is a null pointer, the wctomb function returns a nonzero or zero value, if multibyte character encodings, respectively, do or do not have state-dependent encodings. If \(\mathbf{s}\) is not a null pointer, the wctomb function returns -1 if the value of wchar does not correspond to a valid multibyte character, or returns the number of bytes that are contained in the multibyte character corresponding to the value of wchar.

5 In no case will the value returned be greater than the value of the MB_CUR_MAX macro.

\subsection*{7.14.8 Multibyte string functions}

1 The behavior of the multibyte string functions is affected by the LC_CTYPE category of the current locale.

\subsection*{7.14.8. 1 The mbstowcs function}

Synopsis
```

\#include <stdlib.h>
size_t mbstowcs(wchar_t * restrict pwcs,
const char * restrict s,
size_t n);

```

\section*{Description}

2 The mbstowcs function converts a sequence of multibyte characters that begins in the initial shift state from the array pointed to by \(\mathbf{s}\) into a sequence of corresponding codes and stores not more than \(\mathbf{n}\) codes into the array pointed to by pwcs. No multibyte characters that follow a null character (which is converted into a code with value zero) will be examined or converted. Each multibyte character is converted as if by a call to the mbtowc function, except that the shift state of the mbtowc function is not affected.

3 No more than \(\mathbf{n}\) elements will be modified in the array pointed to by pwcs. If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

4 If an invalid multibyte character is encountered, the mbstowcs function returns (size_t)-1. Otherwise, the mbstowcs function returns the number of array elements modified, not including a terminating zero code, if any. \({ }^{231}\)

\subsection*{7.14.8.2 The wcstombs function}

Synopsis
```

\#include <stdlib.h>
size_t wcstombs(char * restrict s,
const wchar_t * restrict pwcs,
size_t n);

```
231. The array will not be null- or zero-terminated if the value returned is \(\mathbf{n}\).

\section*{Description}

2 The wcstombs function converts a sequence of codes that correspond to multibyte characters from the array pointed to by pwcs into a sequence of multibyte characters that begins in the initial shift state and stores these multibyte characters into the array pointed to by \(\mathbf{s}\), stopping if a multibyte character would exceed the limit of \(\mathbf{n}\) total bytes or if a null character is stored. Each code is converted as if by a call to the wctomb function, except that the shift state of the wctomb function is not affected.

3 No more than \(\mathbf{n}\) bytes will be modified in the array pointed to by \(\mathbf{s}\). If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

4 If a code is encountered that does not correspond to a valid multibyte character, the wcstombs function returns (size_t)-1. Otherwise, the wcstombs function returns the number of bytes modified, not including a terminating null character, if any. \({ }^{231}\)

\subsection*{7.15 String handling <string.h>}

\subsection*{7.15.1 String function conventions}

1 The header <string.h> declares one type and several functions, and defines one macro useful for manipulating arrays of character type and other objects treated as arrays of character type. \({ }^{232}\) The type is size_t and the macro is NULL (both described in 7.1.6). Various methods are used for determining the lengths of the arrays, but in all cases a char * or void * argument points to the initial (lowest addressed) character of the array. If an array is accessed beyond the end of an object, the behavior is undefined.

2 Where an argument declared as size_t \(\boldsymbol{n}\) specifies the length of the array for a function, \(\mathbf{n}\) can have the value zero on a call to that function. Unless explicitly stated otherwise in the description of a particular function in this subclause, pointer arguments on such a call must still have valid values, as described in subclause 7.1.8. On such a call, a function that locates a character finds no occurrence, a function that compares two character sequences returns zero, and a function that copies characters copies zero characters.

\subsection*{7.15.2 Copying functions}

\subsection*{7.15.2.1 The memcpy function}

\section*{Synopsis}
```

\#include <string.h>
void *memcpy(void * restrict s1,
const void * restrict s2,
size_t n);

```

\section*{Description}

2 The memcpy function copies \(\mathbf{n}\) characters from the object pointed to by \(\mathbf{s} 2\) into the object pointed to by s1. If copying takes place between objects that overlap, the behavior is undefined.

\footnotetext{
232. See "future library directions" (7.20.7).
}

\section*{Returns}

3 The memcpy function returns the value of \(\boldsymbol{s} 1\).

\subsection*{7.15.2.2 The memmove function}

Synopsis
```

\#include <string.h>
void *memmove(void *s1, const void *s2, size_t n);

```

\section*{Description}

2 The memmove function copies \(\mathbf{n}\) characters from the object pointed to by \(\boldsymbol{s} 2\) into the object pointed to by s1. Copying takes place as if the \(\mathbf{n}\) characters from the object pointed to by s2 are first copied into a temporary array of \(\mathbf{n}\) characters that does not overlap the objects pointed to by s1 and s2, and then the \(\mathbf{n}\) characters from the temporary array are copied into the object pointed to by \(\mathbf{s 1}\).

\section*{Returns}

3 The memmove function returns the value of \(\mathbf{s} 1\).

\subsection*{7.15.2.3 The strcpy function}

Synopsis
```

\#include <string.h>
char *strcpy(char * restrict s1,
const char * restrict s2);

```

\section*{Description}

2 The strcpy function copies the string pointed to by \(\boldsymbol{s} 2\) (including the terminating null character) into the array pointed to by \(\mathbf{s 1}\). If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

3 The strcpy function returns the value of \(\boldsymbol{s} 1\).

\subsection*{7.15.2.4 The strncpy function}

Synopsis
```

\#include <string.h>
char *strncpy(char * restrict s1,
const char * restrict s2,
size_t n);

```

\section*{Description}

2 The strncpy function copies not more than \(\mathbf{n}\) characters (characters that follow a null character are not copied) from the array pointed to by \(\boldsymbol{s} 2\) to the array pointed to by s1. \({ }^{233}\) If copying takes place between objects that overlap, the behavior is undefined.

3 If the array pointed to by s2 is a string that is shorter than \(\mathbf{n}\) characters, null characters are appended to the copy in the array pointed to by \(\mathbf{s} \mathbf{1}\), until \(\mathbf{n}\) characters in all have been written.

\section*{Returns}

4 The strncpy function returns the value of \(\boldsymbol{s} \mathbf{1}\).

\subsection*{7.15.3 Concatenation functions}

\subsection*{7.15.3.1 The strcat function}

Synopsis
```

\#include <string.h>
char *strcat(char * restrict s1,
const char * restrict s2);

```

\section*{Description}

2 The strcat function appends a copy of the string pointed to by s2 (including the terminating null character) to the end of the string pointed to by s1. The initial character of \(\boldsymbol{s} \mathbf{2}\) overwrites the null character at the end of \(\boldsymbol{s} \mathbf{1}\). If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

The strcat function returns the value of \(\boldsymbol{s} 1\).

\footnotetext{
233. Thus, if there is no null character in the first \(\mathbf{n}\) characters of the array pointed to by \(\boldsymbol{s} \mathbf{2}\), the result will not be null-terminated.
}

\subsection*{7.15.3.2 The strncat function}

Synopsis
```

\#include <string.h>
char *strncat(char * restrict s1,
const char * restrict s2,
size_t n);

```

\section*{Description}

2 The strncat function appends not more than \(\mathbf{n}\) characters (a null character and characters that follow it are not appended) from the array pointed to by \(\boldsymbol{s} 2\) to the end of the string pointed to by \(\mathbf{s 1}\). The initial character of \(\mathbf{s 2}\) overwrites the null character at the end of s1. A terminating null character is always appended to the result. \({ }^{234}\) If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

3 The strncat function returns the value of \(\mathbf{s} 1\).
Forward references: the strlen function (7.15.6.3).

\subsection*{7.15.4 Comparison functions}

1 The sign of a nonzero value returned by the comparison functions memcmp, strcmp, and strncmp is determined by the sign of the difference between the values of the first pair of characters (both interpreted as unsigned char) that differ in the objects being compared.

\subsection*{7.15.4.1 The memcmp function}

Synopsis
```

\#include <string.h>
int memcmp(const void *s1, const void *s2, size_t n);

```

\footnotetext{
234. Thus, the maximum number of characters that can end up in the array pointed to by \(\mathbf{s} 1\) is strlen (s1) \(\mathrm{n}+1\).
}

\section*{Description}

2 The memcmp function compares the first \(\mathbf{n}\) characters of the object pointed to by \(\mathbf{s} \mathbf{1}\) to the first \(\mathbf{n}\) characters of the object pointed to by \(\mathbf{s} \mathbf{2} .{ }^{235}\)

\section*{Returns}

3 The memcmp function returns an integer greater than, equal to, or less than zero, accordingly as the object pointed to by \(\mathbf{s 1}\) is greater than, equal to, or less than the object pointed to by \(\mathbf{s} 2\).

\subsection*{7.15.4.2 The stramp function}

\section*{Synopsis}

1 \#include <string.h>
int strcmp (const char *s1, const char *s2);

\section*{Description}

2 The strcmp function compares the string pointed to by \(\mathbf{s} 1\) to the string pointed to by s2.

\section*{Returns}

3 The stremp function returns an integer greater than, equal to, or less than zero, accordingly as the string pointed to by \(\mathbf{s} 1\) is greater than, equal to, or less than the string pointed to by \(\mathbf{s} 2\).

\subsection*{7.15.4.3 The strcoll function}

Synopsis
```

\#include <string.h>
int strcoll(const char *s1, const char *s2);

```

\section*{Description}

The strcoll function compares the string pointed to by s1 to the string pointed to by s2, both interpreted as appropriate to the LC_COLLATE category of the current locale.
235. The contents of "holes" used as padding for purposes of alignment within structure objects are indeterminate. Strings shorter than their allocated space and unions may also cause problems in comparison.

\section*{Returns}

2 The strcoll function returns an integer greater than, equal to, or less than zero, accordingly as the string pointed to by \(\mathbf{s} 1\) is greater than, equal to, or less than the string pointed to by s2 when both are interpreted as appropriate to the current locale.

\subsection*{7.15.4.4 The strncmp function}

\section*{Synopsis}
```

\#include <string.h>
int strncmp(const char *s1, const char *s2, size_t n);

```

\section*{Description}

2 The strncmp function compares not more than \(\mathbf{n}\) characters (characters that follow a null character are not compared) from the array pointed to by \(\mathbf{s} 1\) to the array pointed to by \(\mathbf{s} 2\).

\section*{Returns}

3 The strncmp function returns an integer greater than, equal to, or less than zero, accordingly as the possibly null-terminated array pointed to by \(\mathbf{s 1}\) is greater than, equal to, or less than the possibly null-terminated array pointed to by \(\mathbf{s 2}\).

\subsection*{7.15.4.5 The strxfrm function}

Synopsis
```

\#include <string.h>
size_t strxfrm(char * restrict s1,
const char * restrict s2,
size_t n);

```

\section*{Description}

2 The strxfrm function transforms the string pointed to by s2 and places the resulting string into the array pointed to by s1. The transformation is such that if the strcmp function is applied to two transformed strings, it returns a value greater than, equal to, or less than zero, corresponding to the result of the strcoll function applied to the same two original strings. No more than \(\mathbf{n}\) characters are placed into the resulting array pointed to by \(\mathbf{s} \mathbf{1}\), including the terminating null character. If \(\mathbf{n}\) is zero, \(\mathbf{s} \mathbf{1}\) is permitted to be a null pointer. If copying takes place between objects that overlap, the behavior is undefined.

\section*{Returns}

3 The strxfrm function returns the length of the transformed string (not including the terminating null character). If the value returned is \(\mathbf{n}\) or more, the contents of the array pointed to by s1 are indeterminate.

\section*{Examples}

4 The value of the following expression is the size of the array needed to hold the transformation of the string pointed to by \(\mathbf{s}\).
```

1 + strxfrm(NULL, s, 0)

```

\subsection*{7.15.5 Search functions}

\subsection*{7.15.5.1 The memchr function}

\section*{Synopsis}
```

\#include <string.h>
void *memchr(const void *s, int c, size_t n);

```

\section*{Description}

2 The memchr function locates the first occurrence of \(\mathbf{c}\) (converted to an unsigned char) in the initial \(\mathbf{n}\) characters (each interpreted as unsigned char) of the object pointed to by \(\mathbf{s}\).

\section*{Returns}

3 The memchr function returns a pointer to the located character, or a null pointer if the character does not occur in the object.

\subsection*{7.15.5.2 The strchr function}

Synopsis
```

\#include <string.h>
char *strchr(const char *s, int c);

```

\section*{Description}

2 The strchr function locates the first occurrence of \(\mathbf{c}\) (converted to a char) in the string pointed to by \(\mathbf{s}\). The terminating null character is considered to be part of the string.

\section*{Returns}

3 The strchr function returns a pointer to the located character, or a null pointer if the character does not occur in the string.

\subsection*{7.15.5.3 The strespn function}

Synopsis
1 \#include <string.h> size_t strcspn(const char *s1, const char *s2);

\section*{Description}

2 The strcspn function computes the length of the maximum initial segment of the string pointed to by \(\mathbf{s} 1\) which consists entirely of characters not from the string pointed to by s2.

\section*{Returns}

3 The strcspn function returns the length of the segment.

\subsection*{7.15.5.4 The strpbrk function}

Synopsis
```

\#include <string.h>
char *strpbrk(const char *s1, const char *s2);

```

\section*{Description}

2 The strpbrk function locates the first occurrence in the string pointed to by \(\mathbf{s} \mathbf{1}\) of any character from the string pointed to by \(\mathbf{s} 2\).

\section*{Returns}

3 The strpbrk function returns a pointer to the character, or a null pointer if no character from s2 occurs in s1.

\subsection*{7.15.5.5 The strrchr function}

Synopsis
```

\#include <string.h>
char *strrchr(const char *s, int c);

```

\section*{Description}

2 The strrchr function locates the last occurrence of (converted to a char) in the string pointed to by s. The terminating null character is considered to be part of the string.

\section*{Returns}

3 The strrchr function returns a pointer to the character, or a null pointer if \(\mathbf{c}\) does not occur in the string.

\subsection*{7.15.5.6 The strspn function}

Synopsis
1 \#include <string.h>
```

size_t strspn(const char *s1, const char *s2);

```

\section*{Description}

2 The strspn function computes the length of the maximum initial segment of the string pointed to by \(\mathbf{s} \mathbf{1}\) which consists entirely of characters from the string pointed to by \(\mathbf{s} 2\).

\section*{Returns}

3 The strspn function returns the length of the segment.

\subsection*{7.15.5.7 The strstr function}

Synopsis
```

\#include <string.h>
char *strstr(const char *s1, const char *s2);

```

\section*{Description}

2 The strstr function locates the first occurrence in the string pointed to by s1 of the sequence of characters (excluding the terminating null character) in the string pointed to by \(\mathbf{s} 2\)

\section*{Returns}

3 The strstr function returns a pointer to the located string, or a null pointer if the string is not found. If s2 points to a string with zero length, the function returns s1.

\subsection*{7.15.5.8 The strtok function}

\section*{Synopsis}
```

\#include <string.h>
char *strtok(char * restrict s1,
const char * restrict s2);

```

\section*{Description}

2 A sequence of calls to the strtok function breaks the string pointed to by \(\mathbf{s} 1\) into a sequence of tokens, each of which is delimited by a character from the string pointed to by s2. The first call in the sequence has s1 as its first argument, and is followed by calls with a null pointer as their first argument. The separator string pointed to by s2 may be different from call to call.

3 The first call in the sequence searches the string pointed to by \(\mathbf{s} 1\) for the first character that is not contained in the current separator string pointed to by s2. If no such character is found, then there are no tokens in the string pointed to by s1 and the strtok function returns a null pointer. If such a character is found, it is the start of the first token.

4 The strtok function then searches from there for a character that is contained in the current separator string. If no such character is found, the current token extends to the end of the string pointed to by \(\mathbf{s 1}\), and subsequent searches for a token will return a null pointer. If such a character is found, it is overwritten by a null character, which terminates the current token. The strtok function saves a pointer to the following character, from which the next search for a token will start.

5 Each subsequent call, with a null pointer as the value of the first argument, starts searching from the saved pointer and behaves as described above.

6 The implementation shall behave as if no library function calls the strtok function.

\section*{Returns}

7 The strtok function returns a pointer to the first character of a token, or a null pointer if there is no token.

\section*{Examples}
```

\#include <string.h>
static char str[] = "?a???b, , ,\#c";
char *t;
t = strtok(str, "?"); // t points to the token "a"
t = strtok(NULL, ","); // t points to the token "??\mathbf{b"}
t = strtok (NULL, "\#,"); // t points to the token "c"
t = strtok(NULL, "?"); // t is a null pointer

```

\subsection*{7.15.6 Miscellaneous functions}

\subsection*{7.15.6.1 The memset function}

Synopsis
```

\#include <string.h>
void *memset(void *s, int c, size_t n);

```

\section*{Description}

2 The memset function copies the value of \(\mathbf{c}\) (converted to an unsigned char) into each of the first \(\mathbf{n}\) characters of the object pointed to by \(\mathbf{s}\).

\section*{Returns}

3 The memset function returns the value of \(\boldsymbol{s}\).

\subsection*{7.15.6.2 The strerror function}

Synopsis
```

\#include <string.h>
char *strerror(int errnum);

```

\section*{Description}

2 The strerror function maps the number in errnum to a message string. Typically, the values for errnum come from errno, but strerror shall map any value of type int to a message.

3 The implementation shall behave as if no library function calls the strerror function.

\section*{Returns}

4 The strerror function returns a pointer to the string, the contents of which are locale-specific. The array pointed to shall not be modified by the program, but may be overwritten by a subsequent call to the strerror function.

Working Draft, 1997-11-21, WG14/N794 J11/97-158

\subsection*{7.15.6.3 The strlen function}

Synopsis
```

\#include <string.h>
size_t strlen(const char *s);

```

\section*{Description}

2 The strlen function computes the length of the string pointed to by s.

\section*{Returns}

3 The strlen function returns the number of characters that precede the terminating null character.

\subsection*{7.16 Date and time <time. h>}

\subsection*{7.16.1 Components of time}

1 The header <time. h > defines four macros, and declares four types and several functions for manipulating time. Many functions deal with a calendar time that represents the current date (according to the Gregorian calendar) and time. Some functions deal with local time, which is the calendar time expressed for some specific time zone, and with Daylight Saving Time, which is a temporary change in the algorithm for determining local time. The local time zone and Daylight Saving Time are implementation-defined.

2 The macros defined are NULL (described in 7.1.6);
CLOCKS_PER_SEC
which expands to a constant expression with the type clock_t described below, and which is the number per second of the value returned by the clock function;
```

_NO_LEAP_SECONDS

```
(described in 7.16.2.4); and
_LOCALTIME // must be outside the range [-14400, +14400]
(described in 7.16.2.4).
3 The types declared are size_t (described in 7.1.6);
```

clock_t

```
and
```

time_t

```
which are arithmetic types capable of representing times;
struct tm
which holds the components of a calendar time, called the broken-down time; and
```

struct tmx

```
which is an extended version of struct tm.
The tm structure shall contain at least the following members, in any order. The
semantics of the members and their normal ranges are expressed in the comments. \({ }^{236}\)
```

int tm_sec; // seconds after the minute - [0, 60]
int tm_min; // minutes after the hour - [0, 59]
int tm_hour; // hours since midnight - [0, 23]
int tm_mday; // day of the month - [1,31]
int tm_mon; // months since January - [0, 11]
int tm_year; // years since 1900
int tm_wday; // days since Sunday - [0, 6]
int tm_yday; // days since January 1 - [0, 365]
int tm_isdst; // Daylight Saving Time flag

```

The value of tm_isdst is positive if Daylight Saving Time is in effect, zero if Daylight Saving Time is not in effect, and negative if the information is not available.

The tmx structure shall contain all the members of struct tm in a manner such that all these members are part of a common initial subsequence. In addition, it contains the members:
```

int tm_version; // version number
int tm_zone; // time zone offset in minutes
// from UTC [-1439, +1439]
int tm_leapsecs; // number of leap seconds applied
void *tm_ext; // extension block
size_t tm_extlen; // size of the extension block

```

The meaning of tm_isdst is also different: it is the positive number of minutes of offset if Daylight Saving Time is in effect, zero if Daylight Saving Time is not in effect, and -1 if the information is not available. A positive value for tm_zone indicates a time that is ahead of Coordinated Universal Time (UTC). The implementation or a future version of this International Standard may include further members in a separate object. If so, the tm_ext member shall point to this object and the tm_extlen object shall be its size. Otherwise, the tm_ext member shall be a null pointer and the value of the tm_extlen object is unspecified.

\footnotetext{
236. The range \([0,60]\) for \(t m \_s e c\) allows for a positive leap second.
}

\subsection*{7.16.2 Time manipulation functions}

\subsection*{7.16.2.1 The clock function}

Synopsis
```

\#include <time.h>
clock_t clock(void);

```

\section*{Description}

2 The clock function determines the processor time used.

\section*{Returns}

3 The clock function returns the implementation's best approximation to the processor time used by the program since the beginning of an implementation-defined era related only to the program invocation. To determine the time in seconds, the value returned by the clock function should be divided by the value of the macro CLOCKS_PER_SEC. If the processor time used is not available or its value cannot be represented, the function returns the value (clock_t)-1. \({ }^{237}\)

\subsection*{7.16.2.2 The difftime function}

Synopsis
```

\#include <time.h>
double difftime(time_t time1, time_t time0);

```

\section*{Description}

2 The difftime function computes the difference between two calendar times: time1 - time0.

\section*{Returns}

3 The difftime function returns the difference expressed in seconds as a double.

\footnotetext{
237. In order to measure the time spent in a program, the clock function should be called at the start of the program and its return value subtracted from the value returned by subsequent calls.
}

\subsection*{7.16.2.3 The mktime function}

\section*{Synopsis}
```

\#include <time.h>
time_t mktime(struct tm *timeptr);

```

\section*{Description}

2 The mktime function converts the broken-down time, expressed as local time, in the structure pointed to by timeptr into a calendar time value with the same encoding as that of the values returned by the time function. The original values of the tm_wday and tm_yday components of the structure are ignored, and the original values of the other components are not restricted to the ranges indicated above. \({ }^{238}\) On successful completion, the values of the tm_wday and tm_yday components of the structure are set appropriately, and the other components are set to represent the specified calendar time, but with their values forced to the ranges indicated above; the


3 The normalization process shall be as described in subclause 7.16.2.6.
4 If the call is successful, a second call to the mktime function with the resulting struct tm value shall always leave it unchanged and return the same value as the first call. Furthermore, if the normalized time is exactly representable as a time_t value, then the normalized broken-down time and the broken-down time generated by converting the result of the mktime function by a call to localtime shall be identical.

\section*{Returns}

5 The mktime function returns the specified calendar time encoded as a value of type time_t. If the calendar time cannot be represented, the function returns the value (time_t)-1. The normalization process shall be as described in subclause 7.16.2.6.

6 If the call is successful, a second call to the mktime function with the resulting struct tm value shall always leave it unchanged and return the same value as the first call. Furthermore, if the normalized time is exactly representable as a time_t value, then the normalized broken-down time and the broken-down time generated by converting the result of the mktime function by a call to localtime shall be identical.

\footnotetext{
238. Thus, a positive or zero value for \(\mathbf{t m}\) _isdst causes the mktime function to presume initially that Daylight Saving Time, respectively, is or is not in effect for the specified time. A negative value causes it to attempt to determine whether Daylight Saving Time is in effect for the specified time.
}

\section*{Examples}

What day of the week is July 4, 2001?
```

\#include <stdio.h>
\#include <time.h>
static const char *const wday[] = {
"Sunday", "Monday", "Tuesday", "Wednesday",
"Thursday", "Friday", "Saturday", "-unknown-"
};
struct tm time_str;
/* ... */
time_str.tm_year = 2001 - 1900;
time_str.tm_mon = 7-1;
time_str.tm_mday = 4;
time_str.tm_hour = 0;
time_str.tm_min = 0;
time_str.tm_sec = 1;
time_str.tm_isdst = -1;
if (mktime(\&time_str) == (time_t)-1)
time_str.tm_wday = 7;
printf("%s\n", wday[time_str.tm_wday]);

```

\subsection*{7.16.2.4 The mkxtime function}

Synopsis
```

\#include <time.h>
time_t mkxtime(struct tmx *timeptr);

```

Description
2 The mkxtime function has the same behavior and result as the mktime function except that it takes into account of the additional members.

3 If the value of the tm_version member is not 1 , the behavior is undefined. If the implementation cannot determine the relationship between local time and UTC, it shall set the tm_zone member of the pointed-to structure to _LOCALTIME. Otherwise, if the tm_zone member was _LOCALTIME, it shall be set to the offset of local time from UTC, include the effects of the value of the tm_isdst member; otherwise, the original value of the tm_isdst member does not affect the result.

4 If the tm_leapsecs member is equal to _NO_LEAP_SECONDS, then the implementation shall determine the number of leap seconds that apply and set the member accordingly (or use 0 if it cannot determine it); otherwise, it shall use the number of leap seconds given. The tm_leapsecs member shall then be set to the Library
number of leap seconds actually applied to produce the value represented by the structure, or to _NO_LEAP_SECONDS if it was not possible to determine it.

5 If the call is successful, a second call to the mkxtime function with the resulting struct tmx value shall always leave it unchanged and return the same value as the first call. Furthermore, if the normalized time is exactly representable as a time_t value, then the normalized broken-down time and the broken-down time generated by converting the result of the mkxtime function by a call to zonetime (with zone set to the value of the tm_zone member) shall be identical.

\subsection*{7.16.2.5 The time function}

\section*{Synopsis}

1 \#include <time.h>
```

time_t time(time_t *timer);

```

\section*{Description}

2 The time function determines the current calendar time. The encoding of the value is unspecified.

\section*{Returns}

3 The time function returns the implementation's best approximation to the current calendar time. The value (time_t)-1 is returned if the calendar time is not available. If timer is not a null pointer, the return value is also assigned to the object it points to.

\subsection*{7.16.2.6 Normalization of broken-down times}

1 A broken-down time is normalized by the mkxtime function in the following manner. A broken-down time is normalized by the mktime function in the same manner, but as if the struct \(t \boldsymbol{m}\) structure had been replaced by a struct \(\operatorname{tmx}\) structure containing the same values except:
tm_version is 1
tm_zone is _LOCALTIME
tm_leapsecs is _NO_LEAP_SECONDS
tm_isdst is \(-1,0\), or an implementation-defined positive value according to whether the original member is less than, equal to, or greater than zero

2 If any of the following members is outside the indicated range (where L is LONG_MAX/8), the behavior is undefined:
\begin{tabular}{ll} 
tm_year & {\([-\mathrm{L} / 366,+\mathrm{L} / 366]\)} \\
tm_mon & {\([-\mathrm{L} / 31,+\mathrm{L} / 31]\)} \\
tm_mday & {\([-\mathrm{L},+\mathrm{L}]\)} \\
tm_hour & {\([-\mathrm{L} / 3600,+\mathrm{L} / 3600]\)} \\
tm_min & {\([-\mathrm{L} / 60,+\mathrm{L} / 60]\)} \\
tm_sec & {\([-\mathrm{L},+\mathrm{L}]\)} \\
tm_leapsecs & {\([-\mathrm{L},+\mathrm{L}]\) or _NO_LEAP_SECONDS } \\
tm_zone & {\([-\mathrm{L} / 60,+\mathrm{L} / 60]\)} \\
tm_isdst & {\([-\mathrm{L} / 60,+\mathrm{L} / 60]\) or _LOCALTTME }
\end{tabular}

The tm_version member shall be 1 .
3 Values S and D shall be determined as follows:
```

\#define QUOT (a,b) ((a)>0 ? (a)/(b) : -(((b)-(a)-1)/(b)))
\#define REM(a,b) ((a)-(b) *QUOT (a,b))
SS = tm_hour*3600 + tm_min*60 +tm_sec +
(tm_leapsecs == _NO_LEAP_SECONDS ? X1 :
tm_leapsecs) -
(tm_zone == _LOCALTIME ? X2 : tm_zone) * 60;
// X1 is the appropriate number of leap seconds, determined by
// the implementation, or 0 if it cannot be determined.
// x2 is the appropriate offset from local time to UTC,
// determined by the implementation, or
// (tm_isdst >= 0 ? tm_isdst : 0)
M = REM(tm_mon, 12);
Y = tm_year + 1900 + QUOT(tm_mon, 12);
Z = Y - (M < 2 ? 1 : 0);
D = Y*365 + (Z/400)*97 + (Z%400)/4 +
M[(int []) {0,31,59,90,120,151,181,212,243,273,
304,335}] +
tm_mday + QUOT(SS, 86400);
S = REM(SS, 86400);

```

4 The normalized broken-down time shall produce the same values of \(\mathbf{S}\) and \(\mathbf{D}\) (though
possibly different values of \(\mathbf{M}, \mathbf{Y}\), and \(\mathbf{Z}\) ) as the original broken-down time. \({ }^{239}\)

\subsection*{7.16.3 Time conversion functions}

1 Except for the strftime and strfxtime functions, these functions each return a pointer to one of two types of static objects: a broken-down time structure or an array of char. Execution of any of the functions that return a pointer to one of these object types may overwrite the information in any object of the same type pointed to by the value returned from any previous call to any of them. The implementation shall behave as if no other library functions call these functions.

\subsection*{7.16.3.1 The asctime function}

\section*{Synopsis}
```

\#include <time.h>
char *asctime(const struct tm *timeptr);

```

\section*{Description}

2 The asctime function converts the broken-down time in the structure pointed to by timeptr into a string in the form

Sun Sep 16 01:03:52 1973\n\0
using the equivalent of the following algorithm.
```

char *asctime(const struct tm *timeptr)
{
static const char wday_name[7][3] = {
"Sun", "Mon", "Tue", "Wed", "Thu", "Fri", "Sat"
};
static const char mon_name[12][3] = {
"Jan", "Feb", "Mar", "Apr", "May", "Jun",
"Jul", "Aug", "Sep", "Oct", "Nov", "Dec"
};
static char result[26];

```
239. The effect of the above rules is to consistently use the Gregorian calendar, regardless of which calendar was in use in which year. In particular, the years 1100 and -300 are not leap years, while the years 1200 and -400 are (these 4 years correspond to \(t m \_y e a r ~ v a l u e s ~ o f ~-800, ~-2200, ~-700, ~\) and -2300 respectively, and the last of these is 401 B.C.E.). In the normalized broken-down time, tm_wday is equal to QUOT ( \(\mathrm{D}-2,7\) ).
```

    sprintf(result, "%.3s %.3s%3d %.2d:%.2d:%.2d %d\n",
    wday_name[timeptr->tm_wday],
    mon_name[timeptr->tm_mon],
    timeptr->tm_mday, timeptr->tm_hour,
    timeptr->tm_min, timeptr->tm_sec,
    1900 + timeptr->tm_year);
    return result;
}

```

Returns
3 The asctime function returns a pointer to the string.

\subsection*{7.16.3.2 The ctime function}

Synopsis
1 \#include <time.h> char *ctime (const time_t *timer);

\section*{Description}

2 The ctime function converts the calendar time pointed to by timer to local time in the form of a string. It is equivalent to
```

asctime(localtime(timer))

```

\section*{Returns}

3 The ctime function returns the pointer returned by the asctime function with that broken-down time as argument.

Forward references: the localtime function (7.16.3.4).

\subsection*{7.16.3.3 The gmtime function}

Synopsis
```

\#include <time.h>
struct tm *gmtime(const time_t *timer);

```

\section*{Description}

2 The gmtime function converts the calendar time pointed to by timer into a brokendown time, expressed as UTC.

\section*{Returns}

3 The gmtime function returns a pointer to that object, or a null pointer if UTC is not available.

\subsection*{7.16.3.4 The localtime function}

Synopsis
1 \#include <time.h>
struct tm *localtime (const time_t *timer);

\section*{Description}

2 The localtime function converts the calendar time pointed to by timer into a broken-down time, expressed as local time.

\section*{Returns}

3 The localtime function returns a pointer to that object.

\subsection*{7.16.3.5 The zonetime function}

Synopsis
1 \#include <time.h> struct tmx *zonetime (const time_t *timer, int zone);

\section*{Description}

2 The zonetime function converts the calendar time pointed to by timer into a broken-down time as represented in the specified time zone. The tm_version member is set to 1 . If the implementation cannot determine the relationship between local time and UTC, it shall set the tm_zone member to _LOCALTIME; otherwise, it shall set the tm_zone member to the value of zone unless the latter is _LOCALTIME, in which case it shall set it to the offset of local time from UTC. The value shall include the effect of Daylight Saving Time, if in effect. The tm_leapsecs member shall be set to the number of leap seconds (the UTC-UT1 offset) applied in the result \({ }^{240}\) if it can be determined, or to the value _NO_LEAP_SECONDS if it cannot (and so none were applied).

\footnotetext{
240. If the tm_sec member is set to 60 , that leap second shall not be included in the value of tm_leapsecs.
}

\section*{Returns}
7.16.3.6 The strftime function

Synopsis
```

\#include <time.h>
size_t strftime(char * restrict s,
size_t maxsize,
const char * restrict format,
const struct tm * restrict timeptr);

```

\section*{Description}

2 The strftime function places characters into the array pointed to by \(\mathbf{s}\) as controlled by the string pointed to by format. The format shall be a multibyte character sequence, beginning and ending in its initial shift state. The format string consists of zero or more conversion specifiers and ordinary multibyte characters. A conversion specifier consists of a \% character followed by a character that determines the behavior of the conversion specifier. All ordinary multibyte characters (including the terminating null character) are copied unchanged into the array. If copying takes place between objects that overlap, the behavior is undefined. No more than maxsize characters are placed into the array. Each conversion specifier is replaced by appropriate characters as described in the following list. The appropriate characters are determined by the LC_TIME category of the current locale and by the values contained in the structure pointed to by timeptr.
\%a is replaced by the locale's abbreviated weekday name. [tm_wday]
\%A is replaced by the locale's full weekday name. [tm_wday]
\(\% \mathbf{b}\) is replaced by the locale's abbreviated month name. [tm_mon]
\(\% \mathbf{B}\) is replaced by the locale's full month name. [tm_mon]
\% \(\mathbf{C}\) is replaced by the locale's appropriate date and time representation. [all specified in 7.16.1]
\%d is replaced by the day of the month as a decimal number (01-31). [tm_mday]
\(\% \mathbf{F}\) is equivalent to \(" \% \mathbf{Y}-\frac{\%}{m}-\frac{\%}{\circ}\) " (the ISO 8601 date format). [tm_year, tm_mon, tm_mday]
\(\% \mathrm{~g}\) is replaced by the last 2 digits of the week-based year (see below) as a decimal number (00-99). [tm_year, tm_wday, tm_yday]
\% \(\mathbf{G}\) is replaced by the week-based year (see below) as a decimal number (e.g., 1997). [tm_year, tm_wday, tm_yday]
\(\% \mathrm{H}\) is replaced by the hour (24-hour clock) as a decimal number (00-23). [tm_hour]
\(\% I\) is replaced by the hour (12-hour clock) as a decimal number (01-12). [tm_hour]
\% \(\mathbf{j}\) is replaced by the day of the year as a decimal number (001-366). [tm_yday]
\%m is replaced by the month as a decimal number (01-12). [tm_mon]
\(\% \mathbf{M}\) is replaced by the minute as a decimal number (00-59). [tm_min]
\%p is replaced by the locale's equivalent of the AM/PM designations associated with a 12 -hour clock. [tm_hour]
\(\% \mathbf{S}\) is replaced by the second as a decimal number (00-60). [tm_sec]
\(\% \mathbf{T}\) is equivalent to "\%H:\%M:\%S" (the ISO 8601 time format). [tm_hour, tm_min, tm_sec]
\%u is replaced by the weekday as a decimal number (1-7), where Monday is 1 (the ISO 8601 weekday number). [tm_wday]
\%U is replaced by the week number of the year (the first Sunday as the first day of week 1) as a decimal number (00-53). [tm_year, tm_wday, tm_yday]
\(\% \mathbf{V}\) is replaced by the ISO 8601 week number (see below) as a decimal number (01-53). [tm_year, tm_wday, tm_yday]
\(\% \mathbf{w}\) is replaced by the weekday as a decimal number ( \(0-6\) ), where Sunday is 0 . [tm_wday]
\(\% \mathbf{W}\) is replaced by the week number of the year (the first Monday as the first day of week 1) as a decimal number (00-53). [tm_year, tm_wday, tm_yday]
\(\% \mathbf{x}\) is replaced by the locale's appropriate date representation. [all specified in 7.16.1]
\(\% \mathbf{x}\) is replaced by the locale's appropriate time representation. [all specified in 7.16.1]
\(\% \mathbf{y}\) is replaced by the last 2 digits of the year as a decimal number (00-99). [tm_year]
\(\% \mathbf{Y}\) is replaced by the whole year as a decimal number (e.g., 1997). [tm_year]
\(\% \mathbf{z}\) is replaced by the offset from UTC in the form "-0430" (meaning 4 hours 30 minutes behind UTC, west of Greenwich). This is the ISO 8601 format. [tm_isdst]
\(\% \mathbf{Z}\) is replaced by the time zone name or abbreviation, or by no characters if no time zone is determinable. [tm_isdst]
\(\% \%\) is replaced by \%.
\(\% \mathbf{g}\), \(\% \mathbf{G}\), and \(\% \mathbf{V}\) give values according to the ISO 8601 week-based year. In this system, weeks begin on a Monday and week 1 of the year is the week that includes both January 4th and the first Thursday of the year. If the first Monday of January is the 2 nd , 3 rd , or 4 th, the preceding days are part of the last week of the preceding year; thus Saturday 2nd January 1999 has \(\% \mathbf{G}==1998\) and \(\% \mathbf{V}==53\). If December 29th, 30th, or 31 st is a Monday, it and any following days are part of week 1 of the following year. Thus, Tuesday 30th December 1997 has \(\% \mathbf{G}==1998\) and \(\% \mathbf{V}==1\).

4 If a conversion specifier is not one of the above, the behavior is undefined.
5 The characters placed in the array by each conversion specifier depend on one or more members of the structure pointed to by timeptr, as specified in brackets in the description. If this value is outside the normal range, the characters stored are unspecified.

6 In the "C" locale the replacement strings for the following specifiers are:
```

%a the first three characters of %A.
%A one of "Sunday", "Monday", ..., "Saturday".
%b the first three characters of %B.
%B}\mathrm{ one of "January", "February", ..., "December".
%\mathbf{C equivalent to "%A %B %d %T %Y".}
%P one of "am" or "pm".
%\mathbf{x}}\mathrm{ equivalent to "%A %B %d %Y".
%X equivalent to %T
%Z implementation-defined.

```

\section*{Returns}

7 If the total number of resulting characters including the terminating null character is not more than maxsize, the strftime function returns the number of characters placed into the array pointed to by \(\mathbf{s}\) not including the terminating null character. Otherwise, zero is returned and the contents of the array are indeterminate.

\subsection*{7.16.3.7 The strfxtime function}

Synopsis
```

\#include <time.h>
size_t strfxtime(char * restrict s,
size_t maxsize,
const char * restrict format,
const struct tmx * restrict timeptr);

```

\section*{Description}

2 The behavior and result of the strfxtime is identical to that of the strftime function, except that the timeptr parameter has a different type, and the \(\% \mathbf{z}\) and \(\% \mathbf{Z}\) conversion specifiers depend on both the tm_zone and tm_isdst members.

Working Draft, 1997-11-21, WG14/N794 J11/97-158

\subsection*{7.17 Alternative spellings <iso646.h>}

1 The header <iso646.h> defines the following eleven macros (on the left) that expand to the corresponding tokens (on the right):
\begin{tabular}{ll} 
and & \(\& \&\) \\
and_eq & \(\&=\) \\
bitand & \(\&\) \\
bitor & \(\mid\) \\
compl & \(\sim\) \\
not & \(!\) \\
not_eq & \(!=\) \\
or & \(|\mid\) \\
or_eq & \(\mid=\) \\
xor & \(\wedge\) \\
xor_eq & \(\wedge=\)
\end{tabular}

\subsection*{7.18 Wide-character classification and mapping utilities <wctype.h>}

\subsection*{7.18.1 Introduction}

1 The header <wctype.h> declares three data types, one macro, and many functions. \({ }^{241}\)

2 The types declared are
wint_t
which is an integer type unchanged by default argument promotions that can hold any value corresponding to members of the extended character set, as well as at least one value that does not correspond to any member of the extended character set. (See WEOF below. \({ }^{242}\)
```

wctrans_t

```
which is a scalar type that can hold values which represent locale-specific character mappings, and
wctype_t
which is a scalar type that can hold values which represent locale-specific character classifications.

3 The macro defined is
WEOF
which expands to a constant expression of type wint_t whose value does not correspond to any member of the extended character set. \({ }^{243}\) It is accepted (and returned) by several functions in this subclause to indicate end-of-file, that is, no more input from a stream. It is also used as a wide-character value that does not correspond to any member of the extended character set.

4 The functions declared are grouped as follows:
- Functions that provide wide-character classification;
- Extensible functions that provide wide-character classification;
241. See "future library directions'" (7.20).
242. wchar_t and wint_t can be the same integer type.
243. The value of the macro WEOF may differ from that of EOF and need not be negative.
- Functions that provide wide-character case mapping;
- Extensible functions that provide wide-character mapping.

5 For all functions described in this subclause that accept an argument of type wint_t, the value shall be representable as a wchar_t or shall equal the value of the macro WEOF. If this argument has any other value, the behavior is undefined.

6 The behavior of these functions is affected by the LC_CTYPE category of the current locale.

\subsection*{7.18.2 Wide-character classification utilities}

1 The header <wctype.h> declares several functions useful for classifying wide characters.

2 The term printing wide character refers to a member of a locale-specific set of wide characters, each of which occupies at least one printing position on a display device. The term control wide character refers to a member of a locale-specific set of wide characters that are not printing wide characters.

\subsection*{7.18.2.1 Wide-character classification functions}

1 The functions in this subclause return nonzero (true) if and only if the value of the argument wc conforms to that in the description of the function.

2 Except for the iswgraph and iswpunct functions with respect to printing, whitespace, wide characters other than \(\mathbf{L}^{\prime} \quad\), each of the following eleven functions returns true for each wide character that corresponds (as if by a call to the wctob function) to a character (byte) for which the respectively matching character testing function from subclause 7.3.1 returns true. \({ }^{244}\)

3 Forward References: the wctob function (7.19.7.1.2).

\subsection*{7.18.2.1.1 The iswalnum function}

\footnotetext{
244. For example, if the expression isalpha(wctob(wc)) evaluates to true, then the call iswalpha(wc) must also return true. But, if the expression isgraph (wctob(wc)) evaluates to true (which cannot occur for wc \(==\mathrm{L}^{\prime}\) ' of course), then either iswgraph (wc) or iswprint (wc) \&\& iswspace (wc) must be true, but not both.
}

Synopsis
```

\#include <wctype.h>
int iswalnum(wint_t wc);

```

\section*{Description}

2 Th iswalnum function tests for any wide character for which iswalpha or iswdigit is true.

\subsection*{7.18.2.1.2 The iswalpha function}

Synopsis
```

\#include <wctype.h>
int iswalpha(wint_t wc);

```

\section*{Description}

2 The iswalpha function tests for any wide character for which iswupper or iswlower is true, or any wide character that is one of a locale-specific set of wide characters for which none of iswentrl, iswdigit, iswpunct, or iswspace is true. \({ }^{245}\)

\subsection*{7.18.2.1.3 The iswblank function}

Synopsis
```

\#include <ctype.h>
int iswblank(int c);

```

Description
2 The iswblank function tests for any wide character for that is a standard blank wide character or is one of a locale-specific set of wide characters, for which iswalnum is false. The standard blank wide characters are the following: space ( \(\mathbf{L}^{\prime}\) '), and horizontal tab ( \(\left.\mathbf{L}^{\prime} \backslash \mathrm{t}^{\prime}\right)\). In the "C" locale, iswblank returns true only for the standard blank characters.

\footnotetext{
245. The functions iswlower and iswupper test true or false separately for each of these additional wide characters; all four combinations are possible.
}

\subsection*{7.18.2.1.4 The iswantrl function}

Synopsis

2 The iswentrl function tests for any control wide character.

\subsection*{7.18.2.1.5 The iswdigit function}

Synopsis
```

\#include <wctype.h>
int iswdigit(wint_t wc);

```

\section*{Description}

2 The iswdigit function tests for any wide character that corresponds to a decimaldigit character (as defined in subclause 5.2.1).

\subsection*{7.18.2.1.6 The iswgraph function}

Synopsis
```

\#include <wctype.h>
int iswgraph(wint_t wc);

```

\section*{Description}

2 The iswgraph function tests for any wide character for which iswprint is true and iswspace is false. \({ }^{246}\)

\subsection*{7.18.2.1.7 The iswlower function}

Synopsis
```

\#include <wctype.h>
int iswlower(wint_t wc);

```

\footnotetext{
246. Note that the behavior of the iswgraph and iswpunct functions may differ from their matching functions in subclause 7.3.1 with respect to printing, white-space, basic execution characters other than' \({ }^{\prime}\).
}

\section*{Description}

2 The iswlower function tests for any wide character that corresponds to a lowercase letter or is one of a locale-specific set of wide characters for which none of iswentrl, iswdigit, iswpunct, or iswspace is true.
7.18.2.1.8 The iswprint function

Synopsis
```

\#include <wctype.h>
int iswprint(wint_t wc);

```

\section*{Description}

2 The iswprint function tests for any printing wide character.

\subsection*{7.18.2.1.9 The iswpunct function}

Synopsis
```

\#include <wctype.h>
int iswpunct(wint_t wc);

```

Description
2 The iswpunct function tests for any printing wide character that is one of a localespecific set of wide characters for which neither iswspace nor iswalnum is true. \({ }^{246}\)

\subsection*{7.18.2.1.10 The iswspace function}

Synopsis
```

\#include <wctype.h>
int iswspace(wint_t wc);

```

\section*{Description}

2 The iswspace function tests for any wide character that corresponds to a localespecific set of wide characters for which none of iswalnum, iswgraph, or iswpunct is true.

\subsection*{7.18.2.1.11 The iswupper function}

Synopsis
```

\#include <wctype.h>
int iswupper(wint_t wc);

```

\section*{Description}

2 The iswupper function tests for any wide character that corresponds to an uppercase letter or is one of a locale-specific set of wide characters for which none of iswentrl, iswdigit, iswpunct, or iswspace is true.

\subsection*{7.18.2.1.12 The iswxdigit function}

Synopsis
```

\#include <wctype.h>
int iswxdigit(wint_t wc);

```

\section*{Description}

2 The iswxdigit function tests for any wide character that corresponds to a hexadecimal-digit character (as defined in subclause 6.1.3.2).

\subsection*{7.18.2.2 Extensible wide-character classification functions}

1 The functions wctype and iswctype provide extensible wide-character classification as well as testing equivalent to that performed by the functions described in the previous subclause (4.5.2.1).

\subsection*{7.18.2.2.1 The wctype function}

\section*{Synopsis}
```

\#include <wctype.h>
wctype_t wctype(const char *property);

```

\section*{Description}

2 The wctype function constructs a value with type wctype_t that describes a class of wide characters identified by the string argument property.

3 The eleven strings listed in the description of the iswctype function shall be valid in all locales as property arguments to the wctype function.

\section*{Returns}

4 If property identifies a valid class of wide characters according to the LC_CTYPE category of the current locale, the wctype function returns a nonzero value that is valid as the second argument to the iswctype function; otherwise, it returns zero.

\subsection*{7.18.2.2.2 The iswctype function}

Synopsis

2 The iswctype function determines whether the wide character wc has the property described by desc. The current setting of the LC_CTYPE category shall be the same as during the call to wctype that returned the value desc.

3 Each of the following eleven expressions has a truth-value equivalent to the call to the wide-character testing function (4.5.2.1) in the comment that follows the expression:
\begin{tabular}{ll} 
iswctype(wc, wctype("alnum")) & // iswalnum(wc) \\
iswctype(wc, wctype("alpha")) & // iswalpha(wc) \\
iswctype(wc, wctype("cntrl")) & // iswcntrl(wc) \\
iswctype(wc, wctype("digit")) & // iswdigit(wc) \\
iswctype(wc, wctype("graph")) & // iswgraph(wc) \\
iswctype(wc, wctype("lower")) & // iswlower(wc) \\
iswctype(wc, wctype("print")) & // iswprint(wc) \\
iswctype(wc, wctype("punct")) & // iswpunct(wc) \\
iswctype(wc, wctype("space")) & // iswspace(wc) \\
iswctype(wc, wctype("upper")) & // iswupper(wc) \\
iswctype(wc, wctype("xdigit")) & // iswxdigit(wc)
\end{tabular}

\section*{Returns}

4 The iswctype function returns nonzero (true) if and only if the value of the wide character wc has the property described by desc.

\subsection*{7.18.3 Wide-character mapping utilities}

1 The header <wctype.h> declares several functions useful for mapping wide characters.

\subsection*{7.18.3.1 Wide-character case-mapping functions}

\subsection*{7.18.3.1.1 The towlower function}

Synopsis
```

\#include <wctype.h>
wint_t towlower(wint_t wc);

```

\section*{Description}

2 The towlower function converts an uppercase letter to a corresponding lowercase letter.

\section*{Returns}

3 If the argument is a wide character for which iswupper is true and there are one or more corresponding wide characters, as specified by the current locale, for which iswlower is true, the towlower function returns one of the corresponding wide characters (always the same one for any given locale); otherwise, the argument is returned unchanged.

\subsection*{7.18.3.1.2 The towupper function}

\section*{Synopsis}
```

\#include <wctype.h>
wint_t towupper(wint_t wc);

```

\section*{Description}

2 The towupper function converts a lowercase letter to a corresponding uppercase letter.

\section*{Returns}

3 If the argument is a wide character for which iswlower is true and there are one or more corresponding characters, as specified by the current locale, for which iswupper is true, the towupper function returns one of the corresponding characters (always the same one for any given locale); otherwise, the argument is returned unchanged.

\subsection*{7.18.3.2 Extensible wide-character mapping functions}

1 The functions wctrans and towctrans provide extensible wide-character mapping as well as case mapping equivalent to that performed by the functions described in the previous subclause (7.18.3.1).

\subsection*{7.18.3.2.1 The wctrans function}

\section*{Synopsis}
```

\#include <wctype.h>
wctrans_t wctrans(const char *property);

```

\section*{Description}

2 The wctrans function constructs a value with type wctrans_t that describes a mapping between wide characters identified by the string argument property.

3 The two strings listed in the description of the towctrans function shall be valid in all locales as property arguments to the wctrans function.

\section*{Returns}

4 If property identifies a valid mapping of wide characters according to the LC_CTYPE category of the current locale, the wctrans function returns a nonzero value that is valid as the second argument to the towctrans function; otherwise, it returns zero.

\subsection*{7.18.3.2.2 The towctrans function}

\section*{Synopsis}
```

\#include <wctype.h>
wint_t towctrans(wint_t wc, wctrans_t desc);

```

\section*{Description}

2 The towctrans function maps the wide character wc using the mapping described by desc. The current setting of the LC_CTYPE category shall be the same as during the call to wctrans that returned the value desc.

3 Each of the following two expressions behaves the same as the call to the widecharacter case-mapping function (7.18.3.1) in the comment that follows the expression:
```

towctrans(wc, wctrans("tolower")) /* towlower(wc) */
towctrans(wc, wctrans("toupper")) /* towupper(wc) */

```

\section*{Returns}

4 The towctrans function returns the mapped value of wc using the mapping described by desc.

\subsection*{7.19 Extended multibyte and wide-character utilities <wchar. h>}

\subsection*{7.19.1 Introduction}

1 The header <wchar. h> declares four data types, one tag, four macros, and many functions. \({ }^{247}\)

2 The types declared are wchar_t and size_t (both described in subclause 7.1.6),
mbstate_t
3 which is an object type other than an array type that can hold the conversion state information necessary to convert between sequences of multibyte characters and wide characters, and
wint_t
4 described in subclause 7.18.1.
5 The tag \(t \mathrm{~m}\) is declared as naming an incomplete structure type, the contents of which are described in subclause 7.15.1.

6 The macros defined are NULL (described in subclause 7.1.6),

\section*{WCHAR_MAX}

7 which is the maximum value representable by an object of type wchar_t, \({ }^{248}\)
WCHAR_MIN
8 which is the minimum value representable by an object of type wchar_t, and
WEOF
9 described in subclause 7.18.1.
10 The functions declared are grouped as follows:
- Functions that perform input and output of wide characters, or multibyte characters, or both;
- Functions that provide wide-string numeric conversion;
- Functions that perform general wide-string manipulation;
247. See "future library directions" (7.20).
248. The values WCHAR_MAX and WCHAR_MIN do not necessarily correspond to members of the extended character set.
- A function for wide-string date and time conversion; and
- Functions that provide extended capabilities for conversion between multibyte and wide-character sequences.

11 Unless explicitly stated otherwise, if the execution of a function described in this subclause causes copying to take place between objects that overlap, the behavior is undefined.

\subsection*{7.19.2 Formatted wide-character input/output functions}

1 The formatted wide-character input/output functions \({ }^{249}\) shall behave as if there is a sequence point after the actions associated with each specifier.

\subsection*{7.19.2.1 The fwprintf function}

Synopsis
```

\#include <stdio.h>
\#include <wchar.h>
int fwprintf(FILE * restrict stream,
const wchar_t * restrict format, ...);

```

\section*{Description}

2 The fwprintf function writes output to the stream pointed to by stream, under control of the wide string pointed to by format that specifies how subsequent arguments are converted for output. If there are insufficient arguments for the format, the behavior is undefined. If the format is exhausted while arguments remain, the excess arguments are evaluated (as always) but are otherwise ignored. The fwprintf function returns when the end of the format string is encountered. The fwprintf function returns when the end of the format string is encountered.

3 The format is composed of zero or more directives: ordinary wide characters (not \%), and conversion specifications. The processing of conversion specifications is as if they were replaced in the format string by wide-character strings that are each the result of fetching zero or more subsequent arguments and converting them, if applicable, according to the corresponding conversion specifier. The expanded wide-character format string is then written to the output stream.

4 Each conversion specification is introduced by the wide character \%. After the \%, the following appear in sequence:

\footnotetext{
249. The fwprintf functions perform writes to memory for the \%n specifier.
}
- Zero or more flags (in any order) that modify the meaning of the conversion specification.
- An optional minimum field width. If the converted value has fewer wide characters than the field width, it is padded with spaces (by default) on the left (or right, if the left adjustment flag, described later, has been given) to the field width. The field width takes the form of an asterisk * (described later) or a decimal integer. \({ }^{250}\)
- An optional precision that gives the minimum number of digits to appear for the d, \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), and \(\mathbf{x}\) conversions, the number of digits to appear after the decimalpoint character for \(\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}\), and \(\mathbf{F}\) conversions, the maximum number of significant digits for the \(\mathbf{g}\) and \(\mathbf{G}\) conversions, or the maximum number of wide characters to be written from a string in s conversions. The precision takes the form of a period (.) followed either by an asterisk * (described later) or by an optional decimal integer; if only the period is specified, the precision is taken as zero. If a precision appears with any other conversion specifier, the behavior is undefined.
- An optional hh specifying that a following \(\mathbf{d}\), \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), or \(\mathbf{x}\) conversion specifier applies to a signed char or unsigned char argument (the argument will have been promoted according to the integer promotions, and its value shall be converted to signed char or unsigned char before printing); an optional \(\mathbf{h}\) specifying that a following \(\mathbf{d}\), \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), or \(\mathbf{X}\) conversion specifier applies to a short int or unsigned short int argument (the argument will have been promoted according to the integer promotions, and its value shall be converted to short int or unsigned short int before printing); an optional \(h\) specifying that a following \(\mathbf{n}\) conversion specifier applies to a pointer to a short int argument; an optional \(\mathbf{l}\) (ell) specifying that a following \(\mathbf{d}\), \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), or \(\mathbf{x}\) conversion specifier applies to a long int or unsigned long int argument; an optional \(1 \mathbf{l}\) (ell-ell) specifying that a following \(\mathbf{d}\), \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), or \(\mathbf{x}\) conversion specifier applies to a long long int or unsigned long long int argument; an optional 1 specifying that a following \(\mathbf{n}\) conversion specifier applies to a pointer to a long int argument; an optional 11 specifying that a following \(\mathbf{n}\) conversion specifier applies to a pointer to a long long int argument; an optional 1 specifying that a following \(\mathbf{c}\) conversion specifier applies to a wint_t argument; an optional \(\mathbf{l}\) specifying that a following \(\mathbf{s}\) conversion specifier applies to a pointer to a wchar_t argument; an optional \(\mathbf{l}\) which has no effect on a following \(\mathbf{a}, \mathbf{A}\), \(\mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}\), or \(\mathbf{G}\) conversion specifier; or an optional \(\mathbf{L}\) specifying that a

\footnotetext{
250. Note that 0 is taken as a flag, not as the beginning of a field width.
}
following \(\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}\), or \(\mathbf{G}\) conversion specifier applies to a long double argument. If an \(\mathbf{h h}, \mathbf{h}, \mathbf{l}, \mathbf{l l}\), or \(\mathbf{L}\) appears with any other conversion specifier, the behavior is undefined.
- A wide character that specifies the type of conversion to be applied.

As noted above, a field width, or precision, or both, may be indicated by an asterisk. In this case, an int argument supplies the field width or precision. The arguments specifying field width, or precision, or both, shall appear (in that order) before the argument (if any) to be converted. A negative field width argument is taken as a - flag followed by a positive field width. A negative precision argument is taken as if the precision were omitted.

The flag wide characters and their meanings are
- \(\quad\) The result of the conversion is left-justified within the field. (It is rightjustified if this flag is not specified.)
\(+\quad\) The result of a signed conversion always begins with a plus or minus sign. (It begins with a sign only when a negative value is converted if this flag is not specified.) \({ }^{251}\)
space If the first wide character of a signed conversion is not a sign, or if a signed conversion results in no wide characters, a space is prefixed to the result. If the space and + flags both appear, the space flag is ignored.
\# The result is to be converted to an 'alternate form'. For o conversion, it increases the precision, if and only if necessary, to force the first digit of the result to be a zero (if the value and precision are both 0 , a single 0 is printed). For \(\mathbf{x}\) ( or \(\mathbf{x}\) ) conversion, a nonzero result has \(\mathbf{0 x}\) (or \(\mathbf{0 x}\) ) prefixed to it. For \(\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}\), and \(\mathbf{G}\) conversions, the result always contains a decimal-point wide character, even if no digits follow it. (Normally, a decimal-point wide character appears in the result of these conversions only if a digit follows it.) For \(\mathbf{g}\) and \(\mathbf{G}\) conversions, trailing zeros are not removed from the result. For other conversions, the behavior is undefined.
\(\mathbf{0} \quad\) For \(\mathbf{d}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}, \mathbf{x}, \mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}\), and \(\mathbf{G}\) conversions, leading zeros (following any indication of sign or base) are used to pad to the field width; no space padding is performed. If the 0 and - flags both appear,
251. The results of all floating conversions of a negative zero, and of negative values that round to zero, include a minus sign.
the \(\mathbf{0}\) flag is ignored. For \(\mathbf{d}\), \(\mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}\), and \(\mathbf{x}\) conversions, if a precision is specified, the \(\mathbf{0}\) flag is ignored. For other conversions, the behavior is undefined.

7 The conversion specifiers and their meanings are
\(\mathbf{d} \mathbf{i} \quad\) The int argument is converted to signed decimal in the style [-]dddd. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it is expanded with leading zeros. The default precision is 1 . The result of converting a zero value with a precision of zero is no wide characters.
\(\mathbf{0}, \mathbf{u}, \mathbf{x}, \mathbf{x}\) The unsigned int argument is converted to unsigned octal (0), unsigned decimal ( \(\mathbf{u}\) ), or unsigned hexadecimal notation ( \(\mathbf{x}\) or \(\mathbf{x}\) ) in the style \(d d d d\); the letters abcdef are used for \(\mathbf{x}\) conversion and the letters ABCDEF for \(\mathbf{x}\) conversion. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it is expanded with leading zeros. The default precision is 1. The result of converting a zero value with a precision of zero is no wide characters.
\(\mathbf{f , F} \quad\) A double argument representing a floating-point number is converted to decimal notation in the style [-]ddd.ddd, where the number of digits after the decimal-point wide character is equal to the precision specification. If the precision is missing, it is taken as 6 ; if the precision is zero and the \# flag is not specified, no decimal-point wide character appears. If a decimal-point wide character appears, at least one digit appears before it. The value is rounded to the appropriate number of digits.

A double argument representing an infinity is converted to one of the styles [-]inf or [-]infinity - which style is implementationdefined. A double argument representing a NaN is converted to one of the styles [-]nan or [-]nan ( \(n\)-wchar-sequence) - which style, and the meaning of any \(n\)-wchar-sequence, is implementation-defined. The \(\mathbf{F}\) conversion specifier produces INF, INFINITY, or NAN instead of inf, infinity, or nan, respectively. \({ }^{252}\)
\(\mathbf{e , E} \quad\) A double argument representing a floating-point number is converted in the style [-]d. \(d d d \mathbf{e} \pm d d\), where there is one digit before the decimal-point
252. When applied to infinite and NaN values, the \(\boldsymbol{-}, \boldsymbol{+}\), and space flag wide characters have their usual meaning; the \# and 0 flag wide characters have no effect.
wide character (which is nonzero if the argument is nonzero) and the number of digits after it is equal to the precision; if the precision is missing, it is taken as 6 ; if the precision is zero and the \# flag is not specified, no decimal-point wide character appears. The value is rounded to the appropriate number of digits. The \(\mathbf{E}\) conversion specifier produces a number with \(\mathbf{E}\) instead of \(\mathbf{e}\) introducing the exponent. The exponent always contains at least two digits, and only as many more digits as necessary to represent the exponent. If the value is zero, the exponent is zero.

A double argument representing an infinity or a NaN is converted in the style of an \(\mathbf{f}\) or \(\mathbf{F}\) conversion specifier.

G,G A double argument representing a floating-point number is converted in style \(\mathbf{f}\) or \(\mathbf{e}\) (or in style \(\mathbf{F}\) or \(\mathbf{E}\) in the case of a \(\mathbf{G}\) conversion specifier), with the precision specifying the number of significant digits. If the precision is zero, it is taken as 1 . The style used depends on the value converted; style e ( \(\boldsymbol{o r} \mathbf{E}\) ) is used only if the exponent resulting from such a conversion is less than -4 or greater than or equal to the precision. Trailing zeros are removed from the fractional portion of the result; a decimal-point wide character appears only if it is followed by a digit. A double argument representing an infinity or a NaN is converted in the style of an \(\mathbf{f}\) or \(\mathbf{F}\) conversion specifier.
a, A A double argument representing a floating-point number is converted in the style \([-] 0 \mathbf{x} h . h h h h \mathbf{p} \pm d\). The number of hexadecimal digits \(h\) after the decimal-point wide character is equal to the precision; if the precision is missing and FLT_RADIX is a power of 2, then the precision is sufficient for an exact representation of the value; if the precision is missing and FLT_RADIX is not a power of 2, then the precision is sufficient to distinguish \({ }^{253}\) values of type double, except that trailing zeros may be omitted. The hexadecimal digit to the left of the decimal-point wide character is nonzero for normalized floating-point numbers and is otherwise unspecified; \({ }^{254}\) if the precision is zero and the \# flag is not

\footnotetext{
253. The precision \(p\) is sufficient to distinguish values of the source type if \(16^{p-1}>b^{n}\)
where \(b\) is FLT_RADIX and \(n\) is the number of base- \(b\) digits in the significand of the source type. A smaller \(p\) might suffice depending on the implementation's scheme for determining the digit to the left of the decimal-point wide character.
254. Binary implementations can choose the hexadecimal digit to the left of the decimal-point wide character so that subsequent digits align to nibble (4-bit) boundaries.
}
specified, no decimal-point wide character appears. The letters abcdef are used for a conversion and the letters ABCDEF for \(\mathbf{A}\) conversion. The a conversion specifier will produce a number with \(\mathbf{x}\) and \(\mathbf{p}\) and the \(\mathbf{A}\) conversion specifier will produce a number with \(\mathbf{X}\) and \(\mathbf{P}\). The exponent always contains at least one digit, and only as many more digits as necessary to represent the decimal exponent of 2 . If the value is zero, the exponent is zero.

A double argument representing an infinity or a NaN is converted in the style of an \(\mathbf{f}\) or \(\mathbf{F}\) conversion specifier.

If no 1 qualifier is present, the int argument is converted to a wide character as if by calling btowc and the resulting wide character is written. Otherwise, the wint_t argument is converted to wchar_t and written.

If no 1 qualifier is present, the argument shall be a pointer to the initial element of a character array containing a multibyte sequence beginning in the initial shift state. Characters from the array are converted as if by repeated calls to the mbrtowc function, with the conversion state described by an mbstate_t object initialized to zero before the first multibyte character is converted, and written up to (but not including) the terminating null wide character. If the precision is specified, no more than that many wide characters are written. If the precision is not specified or is greater than the size of the converted array, the converted array shall contain a null wide character.

If an 1 qualifier is present, the argument shall be a pointer to the initial element of an array of wchar_t type. Wide characters from the array are written up to (but not including) a terminating null wide character. If the precision is specified, no more than that many wide characters are written. If the precision is not specified or is greater than the size of the array, the array shall contain a null wide character.

The argument shall be a pointer to void. The value of the pointer is converted to a sequence of printable wide characters, in an implementation-defined manner.

The argument shall be a pointer to signed integer into which is written the number of wide characters written to the output stream so far by this call to fwprintf. No argument is converted, but one is consumed. If the conversion specification with this conversion specifier is not one of \(\% n\), \(\% \ln , \% \ln , \% \mathrm{hn}\), or \(\% \mathrm{hhn}\), the behavior is undefined.

A \% wide character is written. No argument is converted. The complete conversion specification shall be \(\% \%\).

8 If a conversion specification is invalid, the behavior is undefined. \({ }^{255}\)
9 If any argument is, or points to, a union or an aggregate (except for an array of char type using \%s conversion, an array of wchar_t type using \%ls conversion, or a pointer using \(\% \mathrm{p}\) conversion), the behavior is undefined.

10 In no case does a nonexistent or small field width cause truncation of a field; if the result of a conversion is wider than the field width, the field is expanded to contain the conversion result.

11 For a and A conversions, if FLT_RADIX is a power of 2, the value is correctly rounded to a hexadecimal floating number with the given precision.

\section*{Recommended practice}

If FLT_RADIX is not a power of 2 , the result is one of the two adjacent numbers in hexadecimal floating style with the given precision, with the extra stipulation that the error have a correct sign for the current rounding direction.

12 For \(\mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}\), and \(\mathbf{G}\) conversions, if the number of significant decimal digits is at most DECIMAL_DIG, then the result is correctly rounded. \({ }^{256}\) If the number of significant decimal digits is more than DECIMAL_DIG but the source value is exactly representable with DECIMAL_DIG digits, then the result is an exact representation with trailing zeros. Otherwise, the source value is bounded by two adjacent decimal strings \(L<U\), both having DECIMAL_DIG significant digits; the value of the resultant decimal string \(D\) satisfies \(L \leq D \leq U\), with the extra stipulation that the error have a correct sign for the current rounding direction.

\section*{Returns}

13 The fwprintf function returns the number of wide characters transmitted, or a negative value if an output error occurred.

\footnotetext{
255. See 'future library directions'" (7.20).
256. For binary-to-decimal conversion, the result format's values are the numbers representable with the given format specifier. The number of significant digits is determined by the format specifier, and in the case of fixed-point conversion by the source value as well.
}

\section*{Environmental limit}

14 The minimum value for the maximum number of wide characters produced by any single conversion shall be 4095 .

Examples
15 To print a date and time in the form 'Sunday, July 3, 10:02', followed by \(\pi\) to five decimal places:
```

\#include <math.h>
\#include <stdio.h>
\#include <wchar.h>
/* ... */
wchar_t *weekday, *month; // pointers to wide strings
int day, hour, min;
fwprintf(stdout, L"%ls, %ls %d, %.2d:%.2d\n",
weekday, month, day, hour, min);
fwprintf(stdout, L"pi = %.5f\n", 4 * atan(1.0));

```

16 Forward References: the btowc function (7.19.7.1.1), the mbrtowc function (7.19.7.3.2).

\subsection*{7.19.2.2 The fwscanf function}

Synopsis
```

\#include <stdio.h>
\#include <wchar.h>
int fwscanf(FILE * restrict stream,
const wchar_t * restrict format, ...);

```

\section*{Description}

2 The fwscanf function reads input from the stream pointed to by stream, under control of the wide string pointed to by format that specifies the admissible input sequences and how they are to be converted for assignment, using subsequent arguments as pointers to the objects to receive the converted input. If there are insufficient arguments for the format, the behavior is undefined. If the format is exhausted while arguments remain, the excess arguments are evaluated (as always) but are otherwise ignored.

3 The format is composed of zero or more directives: one or more white-space wide characters; an ordinary wide character (neither \% nor a white-space wide character); or a conversion specification. Each conversion specification is introduced by the wide character \%. After the \%, the following appear in sequence:
- An optional assignment-suppressing wide character *.
- An optional nonzero decimal integer that specifies the maximum field width (in wide characters).
- An optional \(\mathbf{h h}, \mathbf{h}, \mathbf{l}\) (ell), or \(1 \mathbf{l}\) (ell-ell), or \(\mathbf{L}\) indicating the size of the receiving object. The conversion specifiers \(\mathbf{d}\), \(\mathbf{i}\), and \(\mathbf{n}\) shall be preceded by \(\mathbf{h h}\) if the corresponding argument is a pointer to signed char rather than a pointer to int, by \(h\) if it is a pointer to short int rather than a pointer to int, by 1 if it is a pointer to long int, or by \(l \mathrm{l}\) if it is a pointer to long long int. Similarly, the conversion specifiers \(\mathbf{o}, \mathbf{u}\), and \(\mathbf{x}\) shall be preceded by \(\mathbf{h h}\) if the corresponding argument is a pointer to unsigned char rather than a pointer to unsigned int, by \(h\) if it is a pointer to unsigned short int rather than a pointer to unsigned int, by \(l\) if it is a pointer to unsigned long int, or by 11 if it is a pointer to unsigned long long int. The conversion specifiers \(\mathbf{a}, \mathbf{e}, \mathbf{f}\), and \(\mathbf{g}\) shall be preceded by \(\mathbf{l}\) if the corresponding argument is a pointer to double rather than a pointer to float, or by \(L\) if it is a pointer to long double. Finally, the conversion specifiers \(\mathbf{c}, \mathbf{s}\), and [ shall be preceded by \(\mathbf{l}\) if the corresponding argument is a pointer to wchar_t rather than a pointer to a character type. If an \(\mathbf{h h}, \mathbf{h}, \mathbf{l}, \mathbf{l l}\), or \(\mathbf{L}\) appears with any other conversion specifier, the behavior is undefined.
- A wide character that specifies the type of conversion to be applied. The valid conversion specifiers are described below.

The fwscanf function executes each directive of the format in turn. If a directive fails, as detailed below, the fwscanf function returns. Failures are described as input failures (if an encoding error occurs or due to the unavailability of input characters), or matching failures (due to inappropriate input).

5 A directive composed of white-space wide character(s) is executed by reading input up to the first non-white-space wide character (which remains unread), or until no more wide characters can be read.

6 A directive that is an ordinary wide character is executed by reading the next wide character of the stream. If the wide character differs from the directive, the directive fails, and the differing and subsequent wide characters remain unread.

7 A directive that is a conversion specification defines a set of matching input sequences, as described below for each specifier. A conversion specification is executed in the following steps:
skipped, unless the specification includes a [, c, or \(\mathbf{n}\) specifier. \({ }^{257}\)
9 An input item is read from the stream, unless the specification includes an \(\mathbf{n}\) specifier. An input item is defined as the longest sequence of input wide characters which does not exceed any specified field width and which is, or is a prefix of, a matching input sequence. The first wide character, if any, after the input item remains unread. If the length of the input item is zero, the execution of the directive fails: this condition is a matching failure, unless end-of-file, an encoding error, or a read error prevented input from the stream, in which case it is an input failure.

10 Except in the case of a \% specifier, the input item (or, in the case of a \(\%\) n directive, the count of input wide characters) is converted to a type appropriate to the conversion specifier. If the input item is not a matching sequence, the execution of the directive fails: this condition is a matching failure. Unless assignment suppression was indicated by a *, the result of the conversion is placed in the object pointed to by the first argument following the format argument that has not already received a conversion result. If this object does not have an appropriate type, or if the result of the conversion cannot be represented in the object, the behavior is undefined.

11 The conversion specifiers and their meanings are:
d Matches an optionally signed decimal integer, whose format is the same as expected for the subject sequence of the wcstol function with the value 10 for the base argument. The corresponding argument shall be a pointer to signed integer.
i Matches an optionally signed integer, whose format is the same as expected for the subject sequence of the wcstol function with the value 0 for the base argument. The corresponding argument shall be a pointer to signed integer.
- Matches an optionally signed octal integer, whose format is the same as expected for the subject sequence of the wcstoul function with the value 8 for the base argument. The corresponding argument shall be a pointer to unsigned integer.
u Matches an optionally signed decimal integer, whose format is the same as expected for the subject sequence of the wcstoul function with the value 10 for the base argument. The corresponding argument shall be a pointer to unsigned integer.

\footnotetext{
257. These white-space wide characters are not counted against a specified field width.
}

Matches an optionally signed hexadecimal integer, whose format is the same as expected for the subject sequence of the wcstoul function with the value 16 for the base argument. The corresponding argument shall be a pointer to unsigned integer.
\(\mathbf{a}, \mathbf{e}, \mathbf{f}, \mathbf{g}\) Matches an optionally signed floating-point number, whose format is the same as expected for the subject sequence of the wcstod function. The corresponding argument shall be a pointer to floating.
\(\mathbf{s} \quad\) Matches a sequence of non-white-space wide characters. If no 1 qualifier is present, characters from the input field are converted as if by repeated calls to the wortomb function, with the conversion state described by an mbstate_t object initialized to zero before the first wide character is converted. The corresponding argument shall be a pointer to the initial element of a character array large enough to accept the sequence and a terminating null character, which will be added automatically.

Otherwise, the corresponding argument shall be a pointer to the initial element of an array of wchar_t type large enough to accept the sequence and a terminating null wide character, which will be added automatically.
[ Matches a nonempty sequence of wide characters from a set of expected characters (the scanset). f no 1 qualifier is present, characters from the input field are converted as if by repeated calls to the wcrtomb function, with the conversion state described by an mbstate_t object initialized to zero before the first wide character is converted. The corresponding argument shall be a pointer to the initial element of a character array large enough to accept the sequence and a terminating null character, which will be added automatically.

If an 1 qualifier is present, the corresponding argument shall be a pointer to the initial element of an array of wchar_t type large enough to accept the sequence and a terminating null wide character, which will be added automatically.

The conversion specifier includes all subsequent wide characters in the format string, up to and including the matching right bracket wide character (]). The wide characters between the brackets (the scanlist) comprise the scanset, unless the wide character after the left bracket is a circumflex ( \(\wedge\) ), in which case the scanset contains all wide characters that do not appear in the scanlist between the circumflex and the right bracket. If the conversion specifier begins with [] or [^], the right bracket wide character is in the scanlist and the next right bracket wide character is the matching right bracket that ends the specification; otherwise the first right bracket wide character is the one that ends the specification. If a - wide
character is in the scanlist and is not the first, nor the second where the first wide character is a ^, nor the last character, the behavior is implementation-defined.
c Matches a sequence of wide characters of exactly the number specified by the field width ( 1 if no field width is present in the directive). If no 1 qualifier is present, characters from the input field are converted as if by repeated calls to the wcrtomb function, with the conversion state described by an mbstate_t object initialized to zero before the first wide character is converted. The corresponding argument shall be a pointer to the initial element of a character array large enough to accept the sequence. No null character is added.

If an 1 qualifier is present, the corresponding argument shall be a pointer to the initial element of an array of wchar_t type large enough to accept the sequence. No null wide character is added.
\(\mathbf{p} \quad\) Matches an implementation-defined set of sequences, which should be the same as the set of sequences that may be produced by the \(\% \mathrm{p}\) conversion of the fwprintf function. The corresponding argument shall be a pointer to a pointer to void. The interpretation of the input item is implementation-defined. If the input item is a value converted earlier during the same program execution, the pointer that results shall compare equal to that value; otherwise the behavior of the \(\% \mathrm{p}\) conversion is undefined.
n No input is consumed. The corresponding argument shall be a pointer to signed integer into which is to be written the number of wide characters read from the input stream so far by this call to the fwscanf function. Execution of a \%n directive does not increment the assignment count returned at the completion of execution of the fwscanf function. No argument is converted, but one is consumed. If the conversion specification with this conversion specifier is not one of \(\% \mathbf{n}, \% \ln , \% \ln\), \(\% \mathrm{hn}\), or \(\% \mathrm{hhn}\), the behavior is undefined.
\%
Matches a single \%; no conversion or assignment occurs. The complete conversion specification shall be \(\% \%\)

12 If a conversion specification is invalid, the behavior is undefined. \({ }^{258}\)

\footnotetext{
258. See "future library directions" (7.20).
}

13 The conversion specifiers \(\mathbf{A}, \mathbf{E}, \mathbf{G}\), and \(\mathbf{x}\) are also valid and behave the same as, respectively, a, e, g, and \(\mathbf{x}\).

14 If end-of-file is encountered during input, conversion is terminated. If end-of-file occurs before any wide characters matching the current directive have been read (other than leading white space, where permitted), execution of the current directive terminates with an input failure; otherwise, unless execution of the current directive is terminated with a matching failure, execution of the following directive (other than \(\% \mathbf{n}\), if any) is terminated with an input failure.

15 Trailing white space (including new-line wide characters) is left unread unless matched by a directive. The success of literal matches and suppressed assignments is not directly determinable other than via the \%n directive.

16 If conversion terminates on a conflicting input wide character, the offending input wide character is left unread in the input stream. \({ }^{259}\)

\section*{Returns}

17 The fwscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the fwscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\section*{Examples}
1. The call:
```

\#include <stdio.h>
\#include <wchar.h>
/* ... */
int n, i; float x; wchar_t name[50];
n = fwscanf(stdin, L"%d%f%ls", \&i, \&x, name);

```

18 with the input line:

\section*{25 54.32E-1 thompson}

19 will assign to \(n\) the value 3 , to \(i\) the value 25 , to \(\mathbf{x}\) the value 5.432 , and to name the sequence thompson \(\backslash 0\).

\footnotetext{
259. fwscanf pushes back at most one input wide character onto the input stream. Therefore, some sequences that are acceptable to wcstod, wcstol, etc., are unacceptable to fwscanf.
}

Working Draft, 1997-11-21, WG14/N794 J11/97-158
2. The call:
```

\#include <stdio.h>
\#include <wchar.h>
/* ... */
int i; float x; double y;
fwscanf(stdin, L"%2d%f%*d %lf", \&i, \&x, \&y);

```

20 with input:
567890123 56a72
21 will assign to \(\mathbf{i}\) the value 56 and to \(\mathbf{x}\) the value 789.0 , will skip past 0123 , and will assign to \(\mathbf{y}\) the value \(\mathbf{5 6 . 0}\). The next wide character read from the input stream will be a.

22 Forward References: the wcstod function (7.19.4.1.1), the wcstol function (7.19.4.1.4), the westoul function (7.19.4.1.6), the wertomb function (7.19.7.3.3).
7.19.2.3 The wprintf function

Synopsis
1 \#include <wchar.h>
int wprintf(const wchar_t * restrict format, ...);
Description
2 The wprintf function is equivalent to fwprintf with the argument stdout interposed before the arguments to wprintf.

\section*{Returns}

3 The wprintf function returns the number of wide characters transmitted, or a negative value if an output error occurred.

\subsection*{7.19.2.4 The wscanf function}

Synopsis
```

\#include <wchar.h>
int wscanf(const wchar_t * restrict format, ...);

```

\section*{Description}

2 The wscanf function is equivalent to fwscanf with the argument stdin interposed before the arguments to wscanf.

\section*{Returns}

3 The wscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the wscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\subsection*{7.19.2.5 The swprintf function}

Synopsis
```

\#include <wchar.h>
int swprintf(wchar_t * restrict s,
size_t n,
const wchar_t * restrict format, ...);

```

\section*{Description}

2 The swprintf function is equivalent to fwprintf, except that the argument s specifies an array of wide characters into which the generated output is to be written, rather than written to a stream. No more than \(\mathbf{n}\) wide characters are written, including a terminating null wide character, which is always added (unless \(\mathbf{n}\) is zero).

\section*{Returns}

3 The swprintf function returns the number of wide characters written in the array, not counting the terminating null wide character, or a negative value if \(\mathbf{n}\) or more wide characters were requested to be written.

\subsection*{7.19.2.6 The swscanf function}

\section*{Synopsis}
```

\#include <wchar.h>
int swscanf(const wchar_t * restrict s,
const wchar_t * restrict format, ...);

```

\section*{Description}

2 The swscanf function is equivalent to fwscanf, except that the argument s specifies a wide string from which the input is to be obtained, rather than from a stream. Reaching the end of the wide string is equivalent to encountering end-of-file for the fwscanf function.

\section*{Returns}

3 The swscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the swscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\subsection*{7.19.2.7 The vfwprintf function}

Synopsis
```

\#include <stdarg.h>
\#include <stdio.h>
\#include <wchar.h>
int vfwprintf(FILE * restrict stream,
const wchar_t * restrict format,
va_list arg);

```

\section*{Description}

2 The vfwprintf function is equivalent to fwprintf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vfwprintf function does not invoke the va_end macro. \({ }^{260}\)

\section*{Returns}

3 The vfwprintf function returns the number of wide characters transmitted, or a negative value if an output error occurred.

\section*{Examples}

4 The following shows the use of the vfwprintf function in a general error-reporting routine.

\footnotetext{
260. As the functions vfwprintf, vswprintf, vfwscanf, vwprintf, vwscanf, and vswscanf invoke the va_arg macro, the value of arg after the return is indeterminate.
}
```

\#include <stdarg.h>
\#include <stdio.h>
\#include <wchar.h>
void error(char *function_name, wchar_t *format, ...)
{
va_list args;
va_start(args, format);
// print out name of function causing error
fwprintf(stderr, L"ERROR in %s: ", function_name);
// print out remainder of message
vfwprintf(stderr, format, args);
va_end(args);
}

```

\subsection*{7.19.2.8 The vwprintf function}

Synopsis
```

\#include <stdarg.h>
\#include <wchar.h>
int vwprintf(const wchar_t * restrict format,
va_list arg);

```

Description
2 The vwprintf function is equivalent to wprintf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vwprintf function does not invoke the va_end macro. \({ }^{260}\)

\section*{Returns}

3 The vwprintf function returns the number of wide characters transmitted, or a negative value if an output error occurred.

\subsection*{7.19.2.9 The vswprintf function}

Synopsis
1
```

\#include <stdarg.h>
\#include <wchar.h>
int vswprintf(wchar_t * restrict s,
size_t n,
const wchar_t * restrict format,
va_list arg);

```

\section*{Description}

2 The vswprintf function is equivalent to swprintf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vswprintf function does not invoke the va_end macro. \({ }^{260}\)

\section*{Returns}

3 The vswprintf function returns the number of wide characters written in the array, not counting the terminating null wide character, or a negative value if \(\mathbf{n}\) or more wide characters were requested to be generated.

\subsection*{7.19.2.10 The vfwscanf function}

Synopsis
```

\#include <stdarg.h>
\#include <stdio.h>
\#include <wchar.h>
int vfwscanf(FILE * restrict stream,
const wchar_t * restrict format,
va_list arg);

```

\section*{Description}

2 The vfwscanf function is equivalent to fwscanf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vfwscanf function does not invoke the va_end macro. \({ }^{260}\)

\section*{Returns}

3 The vfwscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the vfwscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\subsection*{7.19.2.11 The vwscanf function}

Synopsis
```

\#include <stdarg.h>
\#include <stdio.h>
\#include <wchar.h>
int vwscanf(FILE * restrict stream,
const wchar_t * restrict format,
va_list arg);

```

\section*{Description}

2 The vwscanf function is equivalent to wscanf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vwscanf function does not invoke the va_end macro. \({ }^{260}\)

\section*{Returns}

3 The vwscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the vwscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\subsection*{7.19.2.12 The vswscanf function}

Synopsis
```

\#include <stdarg.h>
\#include <stdio.h>
\#include <wchar.h>
int vswscanf(const wchar_t * restrict s,
const wchar_t * restrict format,
va_list arg);

```

Description
2 The vswscanf function is equivalent to swscanf, with the variable argument list replaced by arg, which shall have been initialized by the va_start macro (and possibly subsequent va_arg calls). The vswscanf function does not invoke the va_end macro. \({ }^{260}\)

\section*{Returns}

3 The vswscanf function returns the value of the macro EOF if an input failure occurs before any conversion. Otherwise, the vswscanf function returns the number of input items assigned, which can be fewer than provided for, or even zero, in the event of an early matching failure.

\subsection*{7.19.3 Wide-character input/output functions}

\subsection*{7.19.3.1 The fgetwc function}

Synopsis
```

\#include <stdio.h>
\#include <wchar.h>
wint_t fgetwc(FILE *stream);

```

\section*{Description}

2 If a next wide character is present from the input stream pointed to by stream, the fgetwc function obtains that wide character and advances the associated file position indicator for the stream (if defined).

\section*{Returns}

3 The fgetwc function returns the next wide character from the input stream pointed to by stream. If the stream is at end-of-file, the end-of-file indicator for the stream is set and fgetwc returns WEOF. If a read error occurs, the error indicator for the stream is set and fgetwc returns WEOF. If an encoding error occurs (including too few bytes), the value of the macro EILSEQ is stored in errno and fgetwc returns WEOF. \({ }^{261}\)

\subsection*{7.19.3.2 The fgetws function}

Synopsis
```

\#include <stdio.h>
\#include <wchar.h>
wchar_t *fgetws(wchar_t * restrict s,
int n, FILE * restrict stream);

```

\footnotetext{
261. An end-of-file and a read error can be distinguished by use of the feof and ferror functions. Also, errno will be set to EILSEQ by input/output functions only if an encoding error occurs.
}

\section*{Description}

2 The fgetws function reads at most one less than the number of wide characters specified by \(\mathbf{n}\) from the stream pointed to by stream into the array pointed to by \(\mathbf{s}\). No additional wide characters are read after a new-line wide character (which is retained) or after end-of-file. A null wide character is written immediately after the last wide character read into the array.

\section*{Returns}

3 The fgetws function returns \(\mathbf{s}\) if successful. If end-of-file is encountered and no characters have been read into the array, the contents of the array remain unchanged and a null pointer is returned. If a read or encoding error occurs during the operation, the array contents are indeterminate and a null pointer is returned.

\subsection*{7.19.3.3 The fputwc function}

Synopsis
```

\#include <stdio.h>
\#include <wchar.h>
wint_t fputwc(wchar_t c, FILE *stream);

```

\section*{Description}

2 The fputwc function writes the wide character specified by \(\mathbf{c}\) to the output stream pointed to by stream, at the position indicated by the associated file position indicator for the stream (if defined), and advances the indicator appropriately. If the file cannot support positioning requests, or if the stream was opened with append mode, the character is appended to the output stream.

\section*{Returns}

3 The fputwc function returns the wide character written. If a write error occurs, the error indicator for the stream is set and fputwc returns WEOF. If an encoding error occurs, the value of the macro EILSEQ is stored in errno and fputwc returns WEOF.

\subsection*{7.19.3.4 The fputws function}

\section*{Synopsis}
```

\#include <stdio.h>
\#include <wchar.h>
int fputws(const wchar_t * restrict s,
FILE * restrict stream);

```

\section*{Description}

2 The fputws function writes the wide string pointed to by \(\mathbf{s}\) to the stream pointed to by stream. The terminating null wide character is not written.

\section*{Returns}

3 The fputws function returns EOF if a write or encoding error occurs; otherwise, it returns a nonnegative value.

\subsection*{7.19.3.5 The getwc function}

Synopsis
```

\#include <stdio.h>
\#include <wchar.h>
wint_t getwc(FILE *stream);

```

\section*{Description}

2 The getwc function is equivalent to fgetwc, except that if it is implemented as a macro, it may evaluate stream more than once, so the argument should never be an expression with side effects.

\section*{Returns}

3 The getwc function returns the next wide character from the input stream pointed to by stream, or WEOF.

\subsection*{7.19.3.6 The getwchar function}

Synopsis
```

\#include <wchar.h>
wint_t getwchar(void);

```

\section*{Description}

2 The getwchar function is equivalent to getwc with the argument stdin.

\section*{Returns}

3 The getwchar function returns the next wide character from the input stream pointed to by stdin, or WEOF.

\subsection*{7.19.3.7 The putwc function}

Synopsis
```

\#include <stdio.h>
\#include <wchar.h>
wint_t putwc(wchar_t c, FILE *stream);

```

\section*{Description}

2 The putwc function is equivalent to fputwc, except that if it is implemented as a macro, it may evaluate stream more than once, so the argument should never be an expression with side effects.

\section*{Returns}

3 The putwc function returns the wide character written, or WEOF.

\subsection*{7.19.3.8 The putwchar function}

Synopsis
```

\#include <wchar.h>
wint_t putwchar(wchar_t c);

```

Description
2 The putwchar function is equivalent to putwc with the second argument stdout.

\section*{Returns}

3 The putwchar function returns the character written, or WEOF.

\subsection*{7.19.3.9 The ungetwc function}

Synopsis
```

    #include <stdio.h>
    #include <wchar.h>
    wint_t ungetwc(wint_t c, FILE *stream);
    ```

\section*{Description}

2 The ungetwc function pushes the wide character specified by cack onto the input stream pointed to by stream. The pushed-back wide characters will be returned by subsequent reads on that stream in the reverse order of their pushing. A successful intervening call (with the stream pointed to by stream to a file positioning function (fseek, fsetpos, or rewind) discards any pushed-back wide characters for the stream. The external storage corresponding to the stream is unchanged.

3 One wide character of pushback is guaranteed, even if the call to the ungetwc function follows just after a call to a formatted wide character input function fwscanf, vfwscanf, vwscanf, or wscanf. If the ungetwc function is called too many times on the same stream without an intervening read or file positioning operation on that stream, the operation may fail.

4 If the value of cequals that of the macro WEOF, the operation fails and the input stream is unchanged.

5 A successful call to the ungetwc function clears the end-of-file indicator for the stream. The value of the file position indicator for the stream after reading or discarding all pushed-back wide characters is the same as it was before the wide characters were pushed back. For a text or binary stream, the value of its file position indicator after a successful call to the ungetwc function is unspecified until all pushed-back wide characters are read or discarded.

\section*{Returns}

6 The ungetwc function returns the wide character pushed back, or WEOF if the operation fails.

\subsection*{7.19.3.10 The fwide function}

Synopsis
```

\#include <stdio.h>
\#include <wchar.h>
int fwide(FILE *stream, int mode);

```

\section*{Description}

2 The fwide function determines the orientation of the stream pointed to by stream. If mode is greater than zero, the function first attempts to make the stream wide oriented. If mode is less than zero, the function first attempts to make the stream byte oriented. \({ }^{262}\) Otherwise, mode is zero and the function does not alter the orientation of the stream.

\footnotetext{
262. If the orientation of the stream has already been determined, fwide does not change it.
}

\section*{Returns}

3 The fwide function returns a value greater than zero if, after the call, the stream has wide orientation, a value less than zero if the stream has byte orientation, or zero if the stream has no orientation.

\subsection*{7.19.4 General wide-string utilities}

1 The header <wchar.h> declares a number of functions useful for wide-string manipulation. Various methods are used for determining the lengths of the arrays, but in all cases a wchar_t * argument points to the initial (lowest addressed) element of the array. If an array is accessed beyond the end of an object, the behavior is undefined.

\subsection*{7.19.4.1 Wide-string numeric conversion functions}

\subsection*{7.19.4.1.1 The wcstod function}

\section*{Synopsis}
```

\#include <wchar.h>
double wcstod(const wchar_t * restrict nptr,
wchar_t ** restrict endptr);

```

\section*{Description}

2 The wcstod function converts the initial portion of the wide string pointed to by nptr to double representation. First, it decomposes the input string into three parts: an initial, possibly empty, sequence of white-space wide characters (as specified by the iswspace function), a subject sequence resembling a floating-point constant; and a final wide string of one or more unrecognized wide characters, including the terminating null wide character of the input wide string. Then, it attempts to convert the subject sequence to a floating-point number, and returns the result.

3 The expected form of the subject sequence is an optional plus or minus sign, then one of the following:
- a nonempty sequence of decimal digits optionally containing a decimal-point wide character, then an optional exponent part as defined for the corresponding singlebyte characters in subclause 6.1.3.1;
- a \(\mathbf{0 x}\) or \(\mathbf{0 x}\), then a nonempty sequence of hexadecimal digits optionally containing a decimal-point wide character, then an optional binary-exponent part, where either the decimal-point character or the binary-exponent part is present;
- one of INF or INFINITY, or any other wide string equivalent except for case
- one of NAN or NAN ( \(n\)-wchar-sequence \({ }_{o p t}\) ), or any other wide string equivalent except for case in the NAN part, where:

Library
```

n-wchar-sequence:
digit
nondigit
n-wchar-sequence digit
n-wchar-sequence nondigit

```
but no floating suffix. The subject sequence is defined as the longest initial subsequence of the input wide string, starting with the first non-white-space wide character, that is of the expected form. The subject sequence contains no wide characters if the input wide string is not of the expected form.

If the subject sequence has the expected form for a floating-point number, the sequence of wide characters starting with the first digit or the decimal-point wide character (whichever occurs first) is interpreted as a floating constant according to the rules of subclause 6.1.3.1, except that the decimal-point wide character is used in place of a period, and that if neither an exponent part, a binary-exponent part, nor a decimal-point wide character appears, a decimal point is assumed to follow the last digit in the wide string. A wide character sequence INF or INFINITY is interpreted as an infinity, if representable in the double type, else like a floating constant that is too large for the range of double. A wide character sequence NAN or NAN ( \(n\) -wchar-sequence \({ }_{o p t}\) ) is interpreted as a quiet NaN , if supported in the double type, else like a subject sequence part that does not have the expected form; the meaning of the n-wchar sequences is implementation-defined. \({ }^{263}\) If the subject sequence begins with a minus sign, the value resulting from the conversion is negated. \({ }^{264}\) A pointer to the final wide string is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

5 In other than the "C" locale, additional locale-specific subject sequence forms may be accepted.

6 If the subject sequence has the hexadecimal form and FLT_RADIX is a power of 2, then the value resulting from the conversion is correctly rounded.

7 If the subject sequence is empty or does not have the expected form, no conversion is performed; the value of nptr is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

\footnotetext{
263. An implementation may use the n -wchar sequence to determine extra information to be represented in the NaN's significand.
264. The wcstod function honors the sign of zero if the arithmetic supports signed zeros.
}

\section*{Recommended practice}

8 If the subject sequence has the hexadecimal form and FLT_RADIX is not a power of 2 , then the result is one of the two numbers in the appropriate internal format that are adjacent to the hexadecimal floating source value, with the extra stipulation that the error have a correct sign for the current rounding direction.

9 If the subject sequence has the decimal form and at most DECIMAL_DIG (defined in <math.h>) significant digits, then the value resulting from the conversion is correctly rounded. If the subject sequence \(D\) has the decimal form and more than DECIMAL_DIG significant digits, consider the two bounding, adjacent decimal strings \(L\) and \(U\), both having DECIMAL_DIG significant digits, such that the values of \(L, D\), and \(U\) satisfy \(L \leq D \leq U\). The result of conversion is one of the (equal or adjacent) values that would be obtained by correctly rounding \(L\) and \(U\) according to the current rounding direction, with the extra stipulation that the error with respect to \(D\) has a correct sign for the current rounding direction. \({ }^{265}\)

\section*{Returns}

10 The wcstod function returns the converted value, if any. If no conversion could be performed, zero is returned. If the correct value is outside the range of representable values, plus or minus HUGE_VAL is returned (according to the sign of the value), and the value of the macro ERANGE is stored in errno. If the result underflows (7.7.1), the function returns a value whose magnitude is no greater than the smallest normalized positive number in the result type; whether errno acquires the value ERANGE is implementation-defined.

\subsection*{7.19.4.1.2 The wcstof function}

\section*{Synopsis}
```

\#include <wchar.h>
float wcstof(
const wchar_t * restrict nptr,
wchar_t ** restrict endptr);

```

\footnotetext{
265. DECIMAL_DIG, defined in <math.h>, is recommended to be sufficiently large that \(L\) and \(U\) will usually round to the same internal floating value, but if not will round to adjacent values.
}

\section*{Description}

2 The wcstof function is similar to the wcstod function, except the returned value has type float and plus or minus HUGE_VALF is returned for values outside the range.

\subsection*{7.19.4.1.3 The wcstold function}

\section*{Synopsis}
```

\#include <wchar.h>
long double wcstold(
const wchar_t * restrict nptr,
wchar_t ** restrict endptr);

```

\section*{Description}

2 The wcstold function is similar to the wcstod function, except the returned value has type long double and plus or minus HUGE_VALL is returned for values outside the range.

\subsection*{7.19.4.1.4 The wastol function}

\section*{Synopsis}
```

\#include <wchar.h>
long int wcstol(
const wchar_t * restrict nptr,
wchar_t ** restrict endptr,
int base);

```

\section*{Description}

2 The wcstol function converts the initial portion of the wide string pointed to by nptr to long int representation. First, it decomposes the input string into three parts: an initial, possibly empty, sequence of white-space wide characters (as specified by the iswspace function), a subject sequence resembling an integer represented in some radix determined by the value of base, and a final wide string of one or more unrecognized wide characters, including the terminating null wide character of the input wide string. Then, it attempts to convert the subject sequence to an integer, and returns the result.

3 If the value of base is zero, the expected form of the subject sequence is that of an integer constant as described for the corresponding single-byte characters in subclause 6.1.3.2, optionally preceded by a plus or minus sign, but not including an integer suffix. If the value of base is between 2 and 36 (inclusive), the expected form of the subject sequence is a sequence of letters and digits representing an integer with the radix specified by base, optionally preceded by a plus or minus sign, but not
including an integer suffix. The letters from \(\mathbf{a}(\operatorname{or} \mathbf{A})\) through \(\mathbf{z}\) (or \(\mathbf{Z}\) ) are ascribed the values 10 through 35 ; only letters and digits whose ascribed values are less than that of base are permitted. If the value of base is 16 , the wide characters \(0 \mathbf{x}\) or \(0 \mathbf{x}\) may optionally precede the sequence of letters and digits, following the sign if present.

The subject sequence is defined as the longest initial subsequence of the input wide string, starting with the first non-white-space wide character, that is of the expected form. The subject sequence contains no wide characters if the input wide string is empty or consists entirely of white space, or if the first non-white-space wide character is other than a sign or a permissible letter or digit.

5 If the subject sequence has the expected form and the value of base is zero, the sequence of wide characters starting with the first digit is interpreted as an integer constant according to the rules of subclause 6.1.3.2. If the subject sequence has the expected form and the value of base is between 2 and 36 , it is used as the base for conversion, ascribing to each letter its value as given above. If the subject sequence begins with a minus sign, the value resulting from the conversion is negated. A pointer to the final wide string is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

6 In other than the "C" locale, additional locale-specific subject sequence forms may be accepted.

7 If the subject sequence is empty or does not have the expected form, no conversion is performed; the value of nptr is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

\section*{Returns}

The wcstol function returns the converted value, if any. If no conversion could be performed, zero is returned. If the correct value is outside the range of representable values, LONG_MAX or LONG_MIN is returned (according to the sign of the value), and the value of the macro ERANGE is stored in errno.

\subsection*{7.19.4.1.5 The wcstoll function}

Synopsis
```

\#include <wchar.h>
long long int wcstoll(
const wchar_t * restrict nptr,
wchar_t ** restrict endptr,
int base);

```

\section*{Description}

2 The wcstoll function is equivalent to the wcstol function, except that it converts the initial portion of the wide string pointed to by nptr to long long int representation.

\section*{Returns}

3 The wcstoll function returns the converted value, if any. If no conversion could be performed, zero is returned. If the correct value is outside the range of representable values, LLONG_MAX or LLONG_MIN is returned (according to the sign of the value), and the value of the macro ERANGE is stored in errno.

\subsection*{7.19.4.1.6 The wcstoul function}

\section*{Synopsis}
```

\#include <wchar.h>
unsigned long int wcstoul(
const wchar_t * restrict nptr,
wchar_t ** restrict endptr,
int base);

```

\section*{Description}

2 The wcstoul function converts the initial portion of the wide string pointed to by nptr to unsigned long int representation. First, it decomposes the input string into three parts: an initial, possibly empty, sequence of white-space wide characters (as specified by the iswspace function), a subject sequence resembling an unsigned integer represented in some radix determined by the value of base, and a final wide string of one or more unrecognized wide characters, including the terminating null wide character of the input wide string. Then, it attempts to convert the subject sequence to an unsigned integer, and returns the result.

3 If the value of base is zero, the expected form of the subject sequence is that of an integer constant as described for the corresponding single-byte characters in subclause 6.1.3.2, optionally preceded by a plus or minus sign, but not including an integer suffix. If the value of base is between 2 and 36 (inclusive), the expected form of the subject sequence is a sequence of letters and digits representing an integer with the radix specified by base, optionally preceded by a plus or minus sign, but not including an integer suffix. The letters from \(\mathbf{a}(\operatorname{or} \mathbf{A})\) through \(\mathbf{z}\) (or \(\mathbf{Z}\) ) are ascribed the values 10 through 35 ; only letters and digits whose ascribed values are less than that of base are permitted. If the value of base is 16 , the wide characters \(0 \mathbf{x}\) or \(0 \times\) may optionally precede the sequence of letters and digits, following the sign if present.

4 The subject sequence is defined as the longest initial subsequence of the input wide string, starting with the first non-white-space wide character, that is of the expected
form. The subject sequence contains no wide characters if the input wide string is empty or consists entirely of white space, or if the first non-white-space wide character is other than a sign or a permissible letter or digit.

5 If the subject sequence has the expected form and the value of base is zero, the sequence of wide characters starting with the first digit is interpreted as an integer constant according to the rules of subclause 6.1.3.2. If the subject sequence has the expected form and the value of base is between 2 and 36, it is used as the base for conversion, ascribing to each letter its value as given above. If the subject sequence begins with a minus sign, the value resulting from the conversion is negated. A pointer to the final wide string is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

6 In other than the "C" locale, additional locale-specific subject sequence forms may be accepted.

7 If the subject sequence is empty or does not have the expected form, no conversion is performed; the value of nptr is stored in the object pointed to by endptr, provided that endptr is not a null pointer.

\section*{Returns}

8 The westoul function returns the converted value, if any. If no conversion could be performed, zero is returned. If the correct value is outside the range of representable values, ULONG_MAX is returned, and the value of the macro ERANGE is stored in errno.

\subsection*{7.19.4.1.7 The wastoull function}

\section*{Synopsis}
```

\#include <wchar.h>
unsigned long long int wcstoull(
const wchar_t * restrict nptr,
wchar_t ** restrict endptr,
int base);

```

\section*{Description}

2 The wcstoull function is equivalent to the wcstoul function, except that it converts the initial portion of the wide string pointed to by nptr to unsigned long long int representation.

\section*{Returns}

3 The wcstoull function returns the converted value, if any. If no conversion could be performed, zero is returned. If the correct value is outside the range of representable values, ULLONG_MAX is returned, and the value of the macro ERANGE is stored in errno.

\subsection*{7.19.4.2 Wide-string copying functions}

\subsection*{7.19.4.2.1 The wcscpy function}

Synopsis
```

\#include <wchar.h>
wchar_t *wcscpy(wchar_t * restrict s1,
const wchar_t * restrict s2);

```

\section*{Description}

2 The wcscpy function copies the wide string pointed to by s2 (including the terminating null wide character) into the array pointed to by \(\mathbf{s} 1\).

\section*{Returns}

3 The wcscpy function returns the value of \(\mathbf{s} 1\).

\subsection*{7.19.4.2.2 The wosncpy function}

Synopsis
```

\#include <wchar.h>
wchar_t *wcsncpy(wchar_t * restrict s1,
const wchar_t * restrict s2,
size_t n);

```

\section*{Description}

2 The wesncpy function copies not more than \(\mathbf{n}\) wide characters (those that follow a null wide character are not copied) from the array pointed to by s2 to the array pointed to by \(\boldsymbol{s} 1 .{ }^{266}\)

3 If the array pointed to by \(\mathbf{s} \mathbf{2}\) is a wide string that is shorter than \(\mathbf{n}\) wide characters, null wide characters are appended to the copy in the array pointed to by \(\boldsymbol{s} \mathbf{1}\), until \(\mathbf{n}\) wide characters in all have been written.

\footnotetext{
266. Thus, if there is no null wide character in the first \(\mathbf{n}\) wide characters of the array pointed to by \(\mathbf{s} 2\), the result will not be null-terminated.
}

\section*{Returns}

4 The wcsncpy function returns the value of \(\boldsymbol{s} 1\).

\subsection*{7.19.4.3 Wide-string concatenation functions}

\subsection*{7.19.4.3.1 The wcscat function}

\section*{Synopsis}
```

\#include <wchar.h>
wchar_t *wcscat(wchar_t * restrict s1,
const wchar_t * restrict s2);

```

\section*{Description}

2 The wcscat function appends a copy of the wide string pointed to by \(\mathbf{s} 2\) (including the terminating null wide character) to the end of the wide string pointed to by \(\mathbf{s} 1\). The initial wide character of \(\mathbf{s} \mathbf{2}\) overwrites the null wide character at the end of \(\mathbf{s} \mathbf{1}\).

\section*{Returns}

3 The wcscat function returns the value of \(\boldsymbol{s} 1\).

\subsection*{7.19.4.3.2 The wesncat function}

\section*{Synopsis}
```

\#include <wchar.h>
wchar_t *wcsncat(wchar_t * restrict s1,
const wchar_t * restrict s2,
size_t n);

```

\section*{Description}

2 The wcsncat function appends not more than \(\mathbf{n}\) wide characters (a null wide character and those that follow it are not appended) from the array pointed to by \(\boldsymbol{s} 2\) to the end of the wide string pointed to by \(\mathbf{s 1}\). The initial wide character of s2 overwrites the null wide character at the end of s1 A terminating null wide character is always appended to the result. \({ }^{267}\)

\footnotetext{
267. Thus, the maximum number of wide characters that can end up in the array pointed to by \(\boldsymbol{s} 1\) is
} wcslen ( s 1 ) \(+\mathrm{n}+1\).

\section*{Returns}

3 The wcsncat function returns the value of s1.

\subsection*{7.19.4.4 Wide-string comparison functions}

1 Unless explicitly stated otherwise, the functions described in this subclause order two wide characters the same way as two integers of the underlying integer type designated by wchar_t.

\subsection*{7.19.4.4.1 The wcscmp function}

Synopsis
```

\#include <wchar.h>
int wcscmp(const wchar_t *s1, const wchar_t *s2);

```

\section*{Description}

2 The wcscmp function compares the wide string pointed to by \(\mathbf{s} 1\) to the wide string pointed to by s2.

\section*{Returns}

3 The wcscmp function returns an integer greater than, equal to, or less than zero, accordingly as the wide string pointed to by \(\mathbf{s} 1\) is greater than, equal to, or less than the wide string pointed to by s2.

\subsection*{7.19.4.4.2 The woscoll function}

\section*{Synopsis}
```

\#include <wchar.h>
int wcscoll(const wchar_t *s1, const wchar_t *s2);

```

\section*{Description}

2 The wcscoll function compares the wide string pointed to by \(\boldsymbol{s} 1\) to the wide string pointed to by s2, both interpreted as appropriate to the LC_COLLATE category of the current locale.

\section*{Returns}

3 The wcscoll function returns an integer greater than, equal to, or less than zero, accordingly as the wide string pointed to by \(\mathbf{s} 1\) is greater than, equal to, or less than the wide string pointed to by s2 when both are interpreted as appropriate to the current locale.

\subsection*{7.19.4.4.3 The wcsncmp function}

Synopsis
```

\#include <wchar.h>
int wcsncmp(const wchar_t *s1, const wchar_t *s2,
size_t n);

```

\section*{Description}

2 The wcsncmp function compares not more than \(\mathbf{n}\) wide characters (those that follow a null wide character are not compared) from the array pointed to by \(\boldsymbol{s} 1\) to the array pointed to by s2.

\section*{Returns}

3 The wcsncmp function returns an integer greater than, equal to, or less than zero, accordingly as the possibly null-terminated array pointed to by \(\mathbf{s 1}\) is greater than, equal to, or less than the possibly null-terminated array pointed to by \(\boldsymbol{s 2}\).

\subsection*{7.19.4.4.4 The wosxfrm function}

Synopsis
```

\#include <wchar.h>
size_t wcsxfrm(wchar_t * restrict s1,
const wchar_t * restrict s2,
size_t n);

```

\section*{Description}

2 The wcsxfrm function transforms the wide string pointed to by \(\mathbf{s} 2\) and places the resulting wide string into the array pointed to by \(\boldsymbol{s} 1\). The transformation is such that if the wcscmp function is applied to two transformed wide strings, it returns a value greater than, equal to, or less than zero, corresponding to the result of the wcscoll function applied to the same two original wide strings. No more than \(\mathbf{n}\) wide characters are placed into the resulting array pointed to by \(\boldsymbol{s} 1\), including the terminating null wide character. If \(\mathbf{n}\) is zero, \(\mathbf{s} \mathbf{1}\) is permitted to be a null pointer.

\section*{Returns}

3 The wcsxfrm function returns the length of the transformed wide string (not including the terminating null wide character). If the value returned is \(\mathbf{n}\) or greater, the contents of the array pointed to by s1 are indeterminate.

\section*{Examples}

4 The value of the following expression is the length of the array needed to hold the transformation of the wide string pointed to by \(\mathbf{s}\) :
```

1 + wCsxfrm(NULL, s, 0)

```

\subsection*{7.19.4.5 Wide-string search functions}

\subsection*{7.19.4.5.1 The wcschr function}

\section*{Synopsis}
```

\#include <wchar.h>
wchar_t *wcschr(const wchar_t *s, wchar_t c);

```

\section*{Description}

2 The wcschr function locates the first occurrence of \(\mathbf{c}\) in the wide string pointed to by \(\mathbf{s}\). The terminating null wide character is considered to be part of the wide string.

\section*{Returns}

3 The wcschr function returns a pointer to the located wide character, or a null pointer if the wide character does not occur in the wide string.

\subsection*{7.19.4.5.2 The woscspn function}

Synopsis
```

\#include <wchar.h>
size_t wcscspn(const wchar_t *s1, const wchar_t *s2);

```

\section*{Description}

2 The wcscspn function computes the length of the maximum initial segment of the wide string pointed to by \(\mathbf{s} 1\) which consists entirely of wide characters not from the wide string pointed to by \(\mathbf{s} 2\).

\section*{Returns}

3 The wcscspn function returns the length of the segment.

\subsection*{7.19.4.5.3 The wcspbrk function}

Synopsis
```

\#include <wchar.h>
wchar_t *wcspbrk(const wchar_t *s1, const wchar_t *s2);

```

\section*{Description}

2 The wcspbrk function locates the first occurrence in the wide string pointed to by \(\mathbf{s} 1\) of any wide character from the wide string pointed to by \(\mathbf{s} 2\).

\section*{Returns}

3 The wcspbrk function returns a pointer to the wide character in \(\mathbf{s} \mathbf{1}\), or a null pointer if no wide character from s2 occurs in \(\mathbf{s} 1\).

\subsection*{7.19.4.5.4 The wosrchr function}

Synopsis
```

\#include <wchar.h>
wchar_t *wcsrchr(const wchar_t *s, wchar_t c);

```

\section*{Description}

2 The wcsrchr function locates the last occurrence of \(\mathbf{c}\) in the wide string pointed to by \(\mathbf{s}\). The terminating null wide character is considered to be part of the wide string.

\section*{Returns}

3 The wcsrchr function returns a pointer to the wide character, or a null pointer if \(\mathbf{c}\) does not occur in the wide string.

\subsection*{7.19.4.5.5 The wasspn function}

Synopsis
```

\#include <wchar.h>
size_t wcsspn(const wchar_t *s1, const wchar_t *s2);

```

\section*{Description}

2 The wcsspn function computes the length of the maximum initial segment of the wide string pointed to by \(\mathbf{s} 1\) which consists entirely of wide characters from the wide string pointed to by \(\mathbf{s} 2\).

\section*{Returns}

3 The wcsspn function returns the length of the segment.

\subsection*{7.19.4.5.6 The wasstr function}

Synopsis
```

\#include <wchar.h>
wchar_t *wcsstr(const wchar_t *s1, const wchar_t *s2);

```

\section*{Description}

2 The wcsstr function locates the first occurrence in the wide string pointed to by \(\boldsymbol{s} 1\) of the sequence of wide characters (excluding the terminating null wide character) in the wide string pointed to by s2.

\section*{Returns}

3 The wcsstr function returns a pointer to the located wide string, or a null pointer if the wide string is not found. If s2 points to a wide string with zero length, the function returns s1.

\subsection*{7.19.4.5.7 The wcstok function}

Synopsis
```

\#include <wchar.h>
wchar_t *wcstok(wchar_t * restrict s1,
const wchar_t * restrict s2,
wchar_t ** restrict ptr);

```

\section*{Description}

2 A sequence of calls to the wcstok function breaks the wide string pointed to by \(\boldsymbol{s} 1\) into a sequence of tokens, each of which is delimited by a wide character from the wide string pointed to by s2. The third argument points to a caller-provided wchar_t pointer into which the wcstok function stores information necessary for it to continue scanning the same wide string.

3 For the first call in the sequence, s1 shall point to a wide string, while in subsequent calls for the same wide string, s1 shall be a null pointer. If \(\boldsymbol{s} 1\) is a null pointer, the value pointed to by ptr shall match that stored by the previous call for the same wide string; otherwise the value pointed to by ptr is ignored. The separator wide string pointed to by s2 may be different from call to call.

The first call in the sequence searches the wide string pointed to by \(\mathbf{s} 1\) for the first wide character that is not contained in the current separator wide string pointed to by s2. If no such wide character is found, then there are no tokens in the wide string pointed to by s1 and the wcstok function returns a null pointer. If such a wide character is found, it is the start of the first token.

5 The wcstok function then searches from there for a wide character that is contained in the current separator wide string. If no such wide character is found, the current token extends to the end of the wide string pointed to by \(\mathbf{s 1}\), and subsequent searches in the same wide string for a token return a null pointer. If such a wide character is found, it is overwritten by a null wide character, which terminates the current token.

6 In all cases, the wcstok function stores sufficient information in the pointer pointed to by ptr so that subsequent calls, with a null pointer for s1 and the unmodified pointer value for ptr, shall start searching just past the element overwritten by a null wide character (if any).

\section*{Returns}

7 The wcstok function returns a pointer to the first wide character of a token, or a null pointer if there is no token.

Examples
```

\#include <wchar.h>
static wchar_t str1[] = L"?a???b, ,,\#c";
static wchar_t str2[] = L"\t \t";
wchar_t *t, *ptr1, *ptr2;
// t points to the token L"a"
t = wcstok(str1, L"?", \&ptr1);
// t points to the token L"??b"
t = wcstok (NULL, L",", \&ptr1);
// t is a null pointer
t = wcstok(str2, L" \t", \&ptr2);
// t points to the token L"c"
t = wcstok(NULL, L"\#,", \&ptr1);
// t is a null pointer
t = wcstok(NULL, L"?", \&ptr1);

```

\subsection*{7.19.4.5.8 The wcslen function}

Synopsis
```

\#include <wchar.h>
size_t wcslen(const wchar_t *s);

```

\section*{Description}

2 The wcslen function computes the length of the wide string pointed to by \(\mathbf{s}\).

\section*{Returns}

3 The wcslen function returns the number of wide characters that precede the terminating null wide character.

\subsection*{7.19.4.6 Wide-character array functions}

1 These functions operate on arrays of type wchar_t whose size is specified by a separate count argument. These functions are not affected by locale and all wchar_t values are treated identically. The null wide character and wchar_t values not corresponding to valid multibyte characters are not treated specially.

2 Unless explicitly stated otherwise, the functions described in this subclause order two wide characters the same way as two integers of the underlying integer type designated by wchar_t.

3 Where an argument declared as size_t \(\mathbf{n}\) determines the length of the array for a function, \(\mathbf{n}\) can have the value zero on a call to that function. Unless stated explicitly otherwise in the description of a particular function in this subclause, pointer arguments on such a call must still have valid values, as described in subclause 7.1.8. On such a call, a function that locates a wide character finds no occurrence, a function that compares two wide character sequences returns zero, and a function that copies wide characters copies zero wide characters.

\subsection*{7.19.4.6.1 The wmemchr function}

\section*{Synopsis}
```

\#include <wchar.h>
wchar_t *wmemchr(const wchar_t *s, wchar_t c,
size_t n);

```

\section*{Description}

2 The wmemchr function locates the first occurrence of \(\mathbf{c}\) in the initial \(\mathbf{n}\) wide characters of the object pointed to by \(\mathbf{s}\).

\section*{Returns}

3 The wmemchr function returns a pointer to the located wide character, or a null pointer if the wide character does not occur in the object.

\subsection*{7.19.4.6.2 The wmememp function}

Synopsis
1 \#include <wchar.h>
int wmemcmp (wchar_t * restrict s1, const wchar_t * restrict s2, size_t n);

\section*{Description}

2 The wmemcmp function compares the first \(\mathbf{n}\) wide characters of the object pointed to by \(\mathbf{s} 1\) to the first \(\mathbf{n}\) wide characters of the object pointed to by \(\mathbf{s} 2\).

\section*{Returns}

3 The wmemcmp function returns an integer greater than, equal to, or less than zero, accordingly as the object pointed to by \(\mathbf{s} 1\) is greater than, equal to, or less than the object pointed to by s2.

\subsection*{7.19.4.6.3 The wmemcpy function}

Synopsis
```

\#include <wchar.h>
wchar_t *wmemcpy(wchar_t * restrict s1,
const wchar_t * restrict s2,
size_t n);

```

\section*{Description}

2 The wmemcpy function copies \(\mathbf{n}\) wide characters from the object pointed to by \(\mathbf{s} 2\) to the object pointed to by \(\mathbf{s} 1\).

\section*{Returns}

3 The wmemcpy function returns the value of \(\boldsymbol{s} 1\).

\subsection*{7.19.4.6.4 The wmemmove function}

Synopsis
```

\#include <wchar.h>
wchar_t *wmemmove(wchar_t *s1, const wchar_t *s2,
size_t n);

```

\section*{Description}

2 The wmemmove function copies \(\mathbf{n}\) wide characters from the object pointed to by \(\mathbf{s} \mathbf{2}\) to the object pointed to by \(\mathbf{s 1}\). Copying takes place as if the \(\mathbf{n}\) wide characters from the object pointed to by \(\mathbf{s} \mathbf{2}\) are first copied into a temporary array of \(\mathbf{n}\) wide characters that does not overlap the objects pointed to by \(\mathbf{s} \mathbf{1}\) or \(\mathbf{s} 2\), and then the \(\mathbf{n}\) wide characters from the temporary array are copied into the object pointed to by \(\mathbf{s} 1\).

\section*{Returns}

3 The wmemmove function returns the value of \(\boldsymbol{s} 1\).

\subsection*{7.19.4.6.5 The wmemset function}

Synopsis
1 \#include <wchar.h>
```

wchar_t *wmemset(wchar_t *s, wchar_t c, size_t n);

```

Description
2 The wmemset function copies the value of \(\mathbf{c}\) into each of the first \(\mathbf{n}\) wide characters of the object pointed to by s.

\section*{Returns}

3 The wmemset function returns the value of \(\mathbf{s}\).

\subsection*{7.19.5 The wcsftime function}

Synopsis
```

\#include <time.h>
\#include <wchar.h>
size_t wcsftime(wchar_t * restrict s,
size_t maxsize,
const wchar_t * restrict format,
const struct tm * restrict timeptr);

```

\section*{Description}

2 The wcsftime function is equivalent to the strftime function, except that:
- The argument s points to the initial element of an array of wide characters into which the generated output is to be placed.
- The argument maxsize indicates the limiting number of wide characters.
- The argument format is a wide string and the conversion specifiers are replaced by corresponding sequences of wide characters.
- The return value indicates the number of wide characters.

\section*{Returns}

3 If the total number of resulting wide characters including the terminating null wide character is not more than maxsize, the wcsftime function returns the number of wide characters placed into the array pointed to by \(\mathbf{s}\) not including the terminating null wide character. Otherwise, zero is returned and the contents of the array are indeterminate.

\subsection*{7.19.6 The wcsfxtime function}

Synopsis
```

\#include <time.h>
\#include <wchar.h>
size_t wcsfxtime(wchar_t * restrict s,
size_t maxsize,
const wchar_t * restrict format,
const struct tmx * restrict timeptr);

```

\section*{Description}

2 The wcsfxtime function is equivalent to the wcsftime function, except that the timeptr parameter has a different type, and the \(\% \mathbf{z}\) and \(\% \mathbf{Z}\) conversion specifiers depend on both the tm_zone and tm_isdst members.

\subsection*{7.19.7 Extended multibyte and wide-character conversion utilities}

1 The header <wchar. h > declares an extended set of functions useful for conversion between multibyte characters and wide characters.

2 Most of the following functions - those that are listed as "restartable", subclauses 7.19.7.3 and 7.19.7.4 - take as a last argument a pointer to an object of type mbstate_t that is used to describe the current conversion state from a particular multibyte character sequence to a wide-character sequence (or the reverse) under the rules of a particular setting for the LC_CTYPE category of the current locale.

3 The initial conversion state corresponds, for a conversion in either direction, to the beginning of a new multibyte character in the initial shift state. A zero-valued mbstate_t object is (at least) one way to describe an initial conversion state. A zero-valued mbstate_t object can be used to initiate conversion involving any multibyte character sequence, in any LC_CTYPE category setting. If an mbstate_t object has been altered by any of the functions described in this subclause, and is then used with a different multibyte character sequence, or in the other conversion direction, or with a different LC_CTYPE category setting than on earlier function calls, the behavior is undefined. \({ }^{268}\)

4 On entry, each function takes the described conversion state (either internal or pointed to by ps) as current. The conversion state described by the pointed-to object is altered as needed to track the shift state, and the position within a multibyte character, for the associated multibyte character sequence.

\subsection*{7.19.7.1 Single-byte wide-character conversion functions}

\subsection*{7.19.7.1.1 The btowc function}

\section*{Synopsis}
```

\#include <stdio.h>
\#include <wchar.h>
wint_t btowc(int c);

```

\footnotetext{
268. Thus a particular mbstate_t object can be used, for example, with both the mbrtowc and mbsrtowcs functions as long as they are used to step sequentially through the same multibyte character string.
}

\section*{Description}

2 The btowc function determines whether constitutes a valid (one-byte) multibyte character in the initial shift state.

\section*{Returns}

3 The btowc returns WEOF if \(\mathbf{c}\) has the value EOF or if (unsigned char) c does not constitute a valid (one-byte) multibyte character in the initial shift state. Otherwise, it returns the wide-character representation of that character.

\subsection*{7.19.7.1.2 The wctob function}

\section*{Synopsis}

1 \#include <stdio.h>
\#include <wchar.h>
int wctob(wint_t c);

\section*{Description}

2 The wctob function determines whether \(\mathbf{c}\) corresponds to a member of the extended character set whose multibyte character representation is a single byte when in the initial shift state.

\section*{Returns}

3 The wctob returns EOF if \(\mathbf{c}\) does not correspond to a multibyte character with length one in the initial shift state. Otherwise, it returns the single-byte representation of that character.

\subsection*{7.19.7.2 The mbsinit function}

Synopsis
```

\#include <wchar.h>
int mbsinit(const mbstate_t *ps);

```

\section*{Description}

2 If \(\mathbf{p s}\) is not a null pointer, the mbsinit function determines whether the pointed-to mbstate_t object describes an initial conversion state.

\section*{Returns}

3 The mbsinit function returns nonzero if ps is a null pointer or if the pointed-to object describes an initial conversion state; otherwise, it returns zero.

\subsection*{7.19.7.3 Restartable multibyte/wide-character conversion functions}

1 These functions differ from the corresponding multibyte character functions of subclause 7.14 .7 (mblen, mbtowc, and wctomb) in that they have an extra parameter, ps, of type pointer to mbstate_t that points to an object that can completely describe the current conversion state of the associated multibyte character sequence. If ps is a null pointer, each function uses its own internal mbstate_t object instead, which is initialized at program startup to the initial conversion state. The implementation behaves as if no library function calls these functions with a null pointer for \(\mathbf{p s}\).

2 Also unlike their corresponding functions, the return value does not represent whether the encoding is state-dependent.

\subsection*{7.19.7.3.1 The mbrlen function}

Synopsis
```

\#include <wchar.h>
size_t mbrlen(const char * restrict s,
size_t n,
mbstate_t * restrict ps);

```

\section*{Description}

2 The mbrlen function is equivalent to the call:
```

mbrtowc(NULL, s, n, ps != NULL ? ps : \&internal)

```
where internal is the mbstate_t object for the mbrlen function, except that the expression designated by ps is evaluated only once.

\section*{Returns}

3 The mbrlen function returns (size_t)-2, (size_t)-1, a value between zero and \(\mathbf{n}\), inclusive.

4 Forward References: the mbrtowc function (7.19.7.3.2).

\subsection*{7.19.7.3.2 The mbrtowc function}

Synopsis
```

\#include <wchar.h>
size_t mbrtowc(wchar_t * restrict pwc,
const char * restrict s,
size_t n,
mbstate_t * restrict ps);

```

\section*{Description}

2 If \(\mathbf{s}\) is a null pointer, the mbrtowc function is equivalent to the call:
```

mbrtowc(NULL, "", 1, ps)

```

In this case, the values of the parameters pwc and \(\mathbf{n}\) are ignored.
3 If \(\mathbf{s}\) is not a null pointer, the mbrtowc function inspects at most \(\mathbf{n}\) bytes beginning with the byte pointed to by \(\mathbf{s}\) to determine the number of bytes needed to complete the next multibyte character (including any shift sequences). If the function determines that the next multibyte character is completed, it determines the value of the corresponding wide character and then, if pwc is not a null pointer, stores that value in the object pointed to by pwc. If the corresponding wide character is the null wide character, the resulting state described is the initial conversion state.

\section*{Returns}

4 The mbrtowc function returns the first of the following that applies (given the current conversion state):
\(0 \quad\) if the next \(\mathbf{n}\) or fewer bytes complete the multibyte character that corresponds to the null wide character (which is the value stored).
positive if the next \(\mathbf{n}\) or fewer bytes complete a valid multibyte character (which is the value stored), the value returned is the number of bytes that complete the multibyte character.
(size_t)-2 if the next \(\mathbf{n}\) bytes contribute to an incomplete (but potentially valid) multibyte character, and all \(\mathbf{n}\) bytes have been processed (no value is stored). \({ }^{269}\)
(size_t)-1 if an encoding error occurs, in which case the next \(\mathbf{n}\) or fewer bytes do not contribute to a complete and valid multibyte character (no value is stored); the value of the macro EILSEQ is stored in errno, and the conversion state is undefined.

\footnotetext{
269. When \(\mathbf{n}\) has at least the value of the MB_CUR_MAX macro, this case can only occur if \(\mathbf{s}\) points at a sequence of redundant shift sequences (for implementations with state-dependent encodings).
}

\subsection*{7.19.7.3.3 The wcrtomb function}

\section*{Synopsis}
```

\#include <wchar.h>
size_t wcrtomb(char * restrict s,
wchar_t wc,
mbstate_t * restrict ps);

```

\section*{Description}

2 If \(\boldsymbol{s}\) is a null pointer, the wcrtomb function is equivalent to the call
```

wcrtomb(buf, L'\0', ps)

```
where buf is an internal buffer.
3 If \(\mathbf{s}\) is not a null pointer, the wcrtomb function determines the number of bytes needed to represent the multibyte character that corresponds to the wide character given by wc (including any shift sequences), and stores the resulting bytes in the array whose first element is pointed to by \(\mathbf{s}\). At most MB_CUR_MAX bytes are stored. If wc is a null wide character, a null byte is stored, preceded by any shift sequence needed to restore the initial shift state; the resulting state described is the initial conversion state.

\section*{Returns}

4 The wcrtomb function returns the number of bytes stored in the array object (including any shift sequences). When wc is not a valid wide character, an encoding error occurs: the function stores the value of the macro EILSEQ in errno and returns (size_t)-1; the conversion state is undefined.

\subsection*{7.19.7.4 Restartable multibyte/wide-string conversion functions}

1 These functions differ from the corresponding multibyte string functions of subclause 7.14 .8 (mbstowcs and wcstombs) in that they have an extra parameter, ps, of type pointer to mbstate_t that points to an object that can completely describe the current conversion state of the associated multibyte character sequence. If \(\mathbf{p s}\) is a null pointer, each function uses its own internal mbstate_t object instead, which is initialized at program startup to the initial conversion state. The implementation behaves as if no library function calls these functions with a null pointer for ps.

2 Also unlike their corresponding functions, the conversion source parameter, src, has a pointer-to-pointer type. When the function is storing the results of conversions (that is, when dst is not a null pointer), the pointer object pointed to by this parameter is updated to reflect the amount of the source processed by that invocation.

\subsection*{7.19.7.4.1 The mbsrtowcs function}

\section*{Synopsis}
```

\#include <wchar.h>
size_t mbsrtowcs (wchar_t * restrict dst,
const char ** restrict src,
size_t len,
mbstate_t * restrict ps);

```

\section*{Description}

2 The mbsrtowcs function converts a sequence of multibyte characters, beginning in the conversion state described by the object pointed to by ps, from the array indirectly pointed to by src into a sequence of corresponding wide characters. If dst is not a null pointer, the converted characters are stored into the array pointed to by dst. Conversion continues up to and including a terminating null character, which is also stored. Conversion stops earlier in two cases: when a sequence of bytes is encountered that does not form a valid multibyte character, or (if dst is not a null pointer) when len codes have been stored into the array pointed to by dst. \({ }^{270}\) Each conversion takes place as if by a call to the mbrtowc function.

3 If dst is not a null pointer, the pointer object pointed to by src is assigned either a null pointer (if conversion stopped due to reaching a terminating null character) or the address just past the last multibyte character converted (if any). If conversion stopped due to reaching a terminating null character and if dst is not a null pointer, the resulting state described is the initial conversion state.

\section*{Returns}

4 If the input conversion encounters a sequence of bytes that do not form a valid multibyte character, an encoding error occurs: the mbsrtowcs function stores the value of the macro EILSEQ in errno and returns (size_t)-1; the conversion state is undefined. Otherwise, it returns the number of multibyte characters successfully converted, not including the terminating null (if any).

\footnotetext{
270. Thus, the value of len is ignored if dst is a null pointer.
}

\subsection*{7.19.7.4.2 The wasrtombs function}

\section*{Synopsis}
```

\#include <wchar.h>
size_t wcsrtombs(char * restrict dst,
const wchar_t ** restrict src,
size_t len,
mbstate_t * restrict ps);

```

\section*{Description}

2 The wcsrtombs function converts a sequence of wide characters from the array indirectly pointed to by src into a sequence of corresponding multibyte characters, beginning in the conversion state described by the object pointed to by ps. If dst is not a null pointer, the converted characters are then stored into the array pointed to by dst. Conversion continues up to and including a terminating null wide character, which is also stored. Conversion stops earlier in two cases: when a code is reached that does not correspond to a valid multibyte character, or (if dst is not a null pointer) when the next multibyte character would exceed the limit of len total bytes to be stored into the array pointed to by dst. Each conversion takes place as if by a call to the wartomb function. \({ }^{271}\)

3 If dst is not a null pointer, the pointer object pointed to by src is assigned either a null pointer (if conversion stopped due to reaching a terminating null wide character) or the address just past the last wide character converted (if any). If conversion stopped due to reaching a terminating null wide character, the resulting state described is the initial conversion state.

\section*{Returns}

4 If conversion stops because a code is reached that does not correspond to a valid multibyte character, an encoding error occurs: the wcsrtombs function stores the value of the macro EILSEQ in errno and returns (size_t)-1; the conversion state is undefined. Otherwise, it returns the number of bytes in the resulting multibyte character sequence, not including the terminating null (if any).

\footnotetext{
271. If conversion stops because a terminating null wide character has been reached, the bytes stored include those necessary to reach the initial shift state immediately before the null byte.
}

\subsection*{7.20 Future library directions}

1 The following names are grouped under individual headers for convenience. All external names described below are reserved no matter what headers are included by the program.

\subsection*{7.20.1 Errors <errno.h>}

1 Macros that begin with \(\mathbf{E}\) and a digit or \(\mathbf{E}\) and an uppercase letter (possibly followed by any combination of digits, letters, and underscore) may be added to the declarations in the <errno. h> header.

\subsection*{7.20.2 Character handling <ctype. h >}

1 Function names that begin with either is or to, and a lowercase letter (possibly followed by any combination of digits, letters, and underscore) may be added to the declarations in the <ctype. \(\mathrm{h}>\) header.

\subsection*{7.20.3 Integer types <inttypes.h>}

1 Type names beginning with int or uint and ending with _t may be added to the types defined in the <inttypes.h> header. Macro names beginning with INT or UINT and ending with _MAX or _MIN, or macro names beginning with PRI or SCN followed by any lower case letter or \(\mathbf{x}\) may be added to the macros defined in the <inttypes.h> header.

\subsection*{7.20.4 Localization <locale. h >}

1 Macros that begin with LC_ and an uppercase letter (possibly followed by any combination of digits, letters, and underscore) may be added to the definitions in the <locale. \(\mathrm{h}>\) header.

\subsection*{7.20.5 Signal handling <signal.h>}

1 Macros that begin with either SIG and an uppercase letter or SIG_ and an uppercase letter (possibly followed by any combination of digits, letters, and underscore) may be added to the definitions in the <signal. \(\mathrm{h}>\) header.

\subsection*{7.20.6 Input/output <stdio.h>}

1 Lowercase letters may be added to the conversion specifiers in fprintf and fscanf. Other characters may be used in extensions.

2 The use of ungetc on a binary stream where the file position indicator is zero prior to the call is an obsolescent feature.

\subsection*{7.20.7 General utilities <stdlib.h>}

1 Function names that begin with str and a lowercase letter (possibly followed by any combination of digits, letters, and underscore) may be added to the declarations in the <stdlib.h> header.

\subsection*{7.20.8 Complex arithmetic <complex.h>}

1 The function names
```

cexp2
cexpm1
clog10
clog1p
clog2
cerf
cerfc
cgamma
clgamma

```
and the same names suffixed with \(\mathbf{f}\) or \(\mathbf{l}\) are reserved for the functions with complex arguments and return values.

\subsection*{7.20.9 String handling <string.h>}

1 Function names that begin with str, mem, or wcs and a lowercase letter (possibly followed by any combination of digits, letters, and underscore) may be added to the declarations in the <string. \(\boldsymbol{h}>\) header.
7.20.10 Wide-character classification and mapping utilities <wctype.h>

1 Function names that begin with is or to and a lowercase letter (possibly followed by any combination of digits, letters, and underscore) may be added to the declarations in the <wctype. \(\mathrm{h}>\) header.

\subsection*{7.20.11 Extended multibyte and wide-character utilities <wchar.h>}

Function names that begin with wcs and a lowercase letter (possibly followed by any combination of digits, letters, and underscore) may be added to the declarations in the <wchar.h> header.

2 Lowercase letters may be added to the conversion specifiers in fwprintf and fwscanf.

\section*{Annex A \\ (informative) \\ Bibliography}
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12. ISO/IEC 9899:1993, Programming languages - \(C\).
13. ISO/IEC TR 10176, Information technology - Guidelines for the preparation of programming language standards.
14. ISO/IEC 10646-1:1993, Information technology - Universal Multiple-Octet Coded Character Set (UCS) - Part 1: Architecture and Basic Multilingual Plane.
15. ISO/IEC 10967-1:1994, Information technology - Language independent arithmetic - Part 1: Integer and floating point arithmetic.

\section*{Annex B \\ (informative) \\ Language syntax summary}

1 Note - The notation is described in the introduction to clause 6 (Language).

\section*{B. 1 Lexical grammar}

\section*{B.1.1 Character sets}

1
(5.2.1) hex-quad:
hexadecimal-digit hexadecimal-digit
hexadecimal-digit hexadecimal-digit
(5.2.1) universal-character-name:
\u hex-quad
\U hex-quad hex-quad

\section*{B.1.2 Lexical elements}

1
(6.1) token:
keyword
identifier
constant
string-literal
operator
punctuator
(6.1) preprocessing-token:
header-name
identifier
pp-number
character-constant
string-literal
operator
punctuator
each non-white-space character that cannot be one of the above

\section*{B.1.3 Keywords}
(6.1.1) keyword: one of
\begin{tabular}{llll} 
auto & break & case & char \\
complex & const & continue & default \\
do & double & else & enum \\
extern & float & for & goto \\
if & imaginary & inline & int \\
long & register & restrict & return \\
short & signed & sizeof & static \\
struct & switch & typedef & union \\
unsigned & void & volatile & while
\end{tabular}

\section*{B.1.4 Identifiers}
(6.1.2) identifier:
nondigit
identifier nondigit
identifier digit
(6.1.2) nondigit: one of
universal-character-name

(6.1.2) digit: one of
\(\begin{array}{llllllllll}0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9\end{array}\)

\section*{B.1.5 Constants}
(6.1.3) constant:
floating-constant
integer-constant
enumeration-constant
character-constant
(6.1.3.1) floating-constant:
decimal-floating-constant
hexadecimal-floating-constant
(6.1.3.1) decimal-floating-constant:
fractional-constant exponent-part opt floating-suffix \({ }_{\text {opt }}\)
digit-sequence exponent-part floating-suffix \({ }_{\text {opt }}\)
(6.1.3.1) hexadecimal-floating-constant:

0x hexadecimal-fractional-constant binary-exponent-part floating-suffix
0x hexadecimal-fractional-constant binary-exponent-part floating-suffix \({ }_{\text {opt }}\) 0x hexadecimal-digit-sequence binary-exponent-part floating-suffix opt 0x hexadecimal-digit-sequence binary-exponent-part floating-suffix opt
(6.1.3.1) fractional-constant:
digit-sequence \({ }_{\text {opt }}\). digit-sequence
digit-sequence .
(6.1.3.1) exponent-part:
e sign \({ }_{\text {opt }}\) digit-sequence
\(\mathbf{E}\) sign opt digit-sequence
(6.1.3.1) sign: one of
\(+\quad-\)
(6.1.3.1) digit-sequence:
digit
digit-sequence digit
(6.1.3.1) hexadecimal-fractional-constant:
hexadecimal-digit-sequence \({ }_{\text {opt }} \cdot\) hexadecimal-digit-sequence hexadecimal-digit-sequence .
(6.1.3.1) binary-exponent-part:
\(\begin{array}{lll}\mathbf{P} & \text { sign } \\ \mathbf{P} & \text { sign } n_{\text {opt }} \text { digit-sequence } \\ \text { digit-sequence }\end{array}\)
(6.1.3.1) hexadecimal-digit-sequence:
hexadecimal-digit
hexadecimal-digit-sequence hexadecimal-digit
(6.1.3.1) hexadecimal-digit: one of
\(\begin{array}{llllllllll}0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9\end{array}\)
a b c d e f
A \(\quad\) B \(\quad\) C \(\quad\) D \(\quad\) E \(\quad\) F
(6.1.3.1) floating-suffix: one of
\(\mathbf{f} \mathbf{I} \quad \mathbf{L}\)
(6.1.3.2) integer-constant:
decimal-constant integer-suffix
octal-constant integer-suffix
hexadecimal-constant integer-suffix \({ }_{\text {opt }}\)
(6.1.3.2) decimal-constant:
nonzero-digit
decimal-constant digit
(6.1.3.2) octal-constant:

0
octal-constant octal-digit
(6.1.3.2) hexadecimal-constant:
\(\mathbf{0 x}\) hexadecimal-digit
0X hexadecimal-digit
hexadecimal-constant hexadecimal-digit
(6.1.3.2) nonzero-digit: one of
\(\begin{array}{lllllllll}1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9\end{array}\)
(6.1.3.2) octal-digit: one of
\(\begin{array}{llllllll}0 & 1 & 2 & 3 & 4 & 5 & 6 & 7\end{array}\)
(6.1.3.2) integer-suffix:
unsigned-suffix long-suffix opt
long-suffix unsigned-suffix opt
unsigned-suffix long-long-suffix opt
long-long-suffix unsigned-suffix opt
(6.1.3.2) unsigned-suffix: one of
u U
(6.1.3.2) long-suffix: one of

1 L
(6.1.3.2) long-long-suffix: one of

11 LL
(6.1.3.3) enumeration-constant:
identifier
(6.1.3.4) character-constant:
' c-char-sequence'
\(\mathrm{L}^{\prime}\) c-char-sequence'
(6.1.3.4) c-char-sequence:
c-char
c-char-sequence c-char
(6.1.3.4) c-char:
any member of the source character set except
the single-quote \({ }^{\prime}\), backslash \(\backslash\), or new-line character
escape-sequence
universal-character-name
(6.1.3.4) escape-sequence:
simple-escape-sequence
octal-escape-sequence
hexadecimal-escape-sequence
(6.1.3.4) simple-escape-sequence: one of
\' \" \? \\
\(\backslash a \quad \backslash b \quad \backslash f \quad \backslash n \quad \backslash r \quad \backslash t \quad \backslash v\)
(6.1.3.4) octal-escape-sequence:
\octal-digit
\ octal-digit octal-digit
\ octal-digit octal-digit octal-digit
(6.1.3.4) hexadecimal-escape-sequence:
\(\backslash \mathbf{x}\) hexadecimal-digit
hexadecimal-escape-sequence hexadecimal-digit

\section*{B.1.6 String literals}
(6.1.4) string-literal:
"s-char-sequence opt "
L "s-char-sequence \({ }_{\text {opt }} "\)
(6.1.4) \(s\)-char-sequence:
s-char
s-char-sequence s-char
(6.1.4) s-char:
any member of the source character set except
the double-quote ", backslash \\, or new-line character
escape-sequence
universal-character-name

\section*{B.1.7 Operators}
(6.1.5) operator: one of


\section*{B.1.8 Punctuators}
(6.1.6) punctuator: one of


\section*{B.1.9 Header names}
(6.1.7) header-name:
<h-char-sequence>
" \(q\)-char-sequence"
"q-char-sequence"
(6.1.7) h-char-sequence:
h-char
\(h\)-char-sequence \(h\)-char
(6.1.7) h-char:
any member of the source character set except the new-line character and >
(6.1.7) q-char-sequence:
q-char
q-char-sequence q-char
(6.1.7) q-char:
any member of the source character set except the new-line character and "

\section*{B.1.10 Preprocessing numbers}
(6.1.8) pp-number:
digit
. digit
pp-number digit
pp-number nondigit
pp-number e sign
pp-number E sign
pp-number \(\mathbf{p}\) sign
pp-number \(\mathbf{P}\) sign
pp-number .

\section*{B. 2 Phrase structure grammar}

\section*{B.2.1 Expressions}
(6.3.1) primary-expression:
identifier
constant
string-literal
( expression )
(6.3.2) postfix-expression:
primary-expression
postfix-expression [ expression]
postfix-expression ( argument-expression-list \({ }_{\text {opt }}\) )
postfix-expression . identifier
postfix-expression -> identifier
postfix-expression ++
postfix-expression --
( type-name ) \{ initializer-list \}
( type-name ) \{ initializer-list , \}
(6.3.2) argument-expression-list:
assignment-expression
argument-expression-list , assignment-expression
(6.3.3) unary-expression:
postfix-expression
++ unary-expression
-- unary-expression
unary-operator cast-expression
sizeof unary-expression
sizeof (type-name )
(6.3.3) unary-operator: one of
\& * + - ~ !
(6.3.4) cast-expression:
unary-expression
( type-name ) cast-expression
(6.3.5) multiplicative-expression:
cast-expression
multiplicative-expression * cast-expression
multiplicative-expression / cast-expression
multiplicative-expression \% cast-expression
(6.3.6) additive-expression:
multiplicative-expression
additive-expression + multiplicative-expression
additive-expression - multiplicative-expression
(6.3.7) shift-expression:
additive-expression
shift-expression << additive-expression
shift-expression >> additive-expression
(6.3.8) relational-expression:
shift-expression
relational-expression < shift-expression
relational-expression \(>\) shift-expression
relational-expression \(<=\) shift-expression
relational-expression \(>=\) shift-expression
(6.3.9) equality-expression:
relational-expression
equality-expression \(==\) relational-expression
equality-expression \(!=\) relational-expression
(6.3.10) AND-expression:
equality-expression
AND-expression \& equality-expression
(6.3.11) exclusive-OR-expression:

AND-expression
exclusive-OR-expression ^ AND-expression
(6.3.12) inclusive-OR-expression:
exclusive-OR-expression
inclusive-OR-expression | exclusive-OR-expression
(6.3.13) logical-AND-expression:
inclusive-OR-expression
logical-AND-expression \(\& \&\) inclusive-OR-expression
(6.3.14) logical-OR-expression:
logical-AND-expression
logical-OR-expression || logical-AND-expression
(6.3.15) conditional-expression:
logical-OR-expression
logical-OR-expression ? expression : conditional-expression
(6.3.16) assignment-expression:
conditional-expression
unary-expression assignment-operator assignment-expression
(6.3.16) assignment-operator: one of
\(=*=1=\%=+=-=\ll=\gg=\Leftrightarrow=1=\)
(6.3.17) expression:
assignment-expression
expression , assignment-expression
(6.4) constant-expression:
conditional-expression

\section*{B.2.2 Declarations}
(6.5) declaration:
declaration-specifiers init-declarator-list \({ }_{\text {opt }} ;\)
(6.5) declaration-specifiers:
storage-class-specifier declaration-specifiers \({ }_{\text {opt }}\)
type-specifier declaration-specifiers \({ }_{\text {opt }}\)
type-qualifier declaration-specifiers opt
function-specifiers
(6.5) init-declarator-list:
init-declarator
init-declarator-list , init-declarator
(6.5) init-declarator:
declarator
declarator \(=\) initializer
(6.5.1) storage-class-specifier:
typedef
extern
static
auto
register
(6.5.2) type-specifier:
void
char
short
int
long
float
double
complex
signed
unsigned
struct-or-union-specifier
enum-specifier
typedef-name
(6.5.2.1) struct-or-union-specifier:
struct-or-union identifier
struct-or-union identifier \{ struct-declaration-list \}
(6.5.2.1) struct-or-union:

> struct
> union
(6.5.2.1) struct-declaration-list:
struct-declaration
struct-declaration-list struct-declaration
(6.5.2.1) struct-declaration:
specifier-qualifier-list struct-declarator-list ;
(6.5.2.1) specifier-qualifier-list:
\(\begin{array}{ll}\text { type-specifier } & \text { specifier-qualifier-list } \\ \text { type-qualifier } & \text { specifier-qualifier-list }\end{array}\)
(6.5.2.1) struct-declarator-list:
struct-declarator
struct-declarator-list , struct-declarator
(6.5.2.1) struct-declarator:
declarator
declarator \(_{\text {opt }}\) : constant-expression
(6.5.2.2) enum-specifier:
enum identifier \({ }_{o p t}\{\) enumerator-list \}
enum identifier opt \(\{\) enumerator-list , \}
enum identifier
(6.5.2.2) enumerator-list:
enumerator
enumerator-list , enumerator
(6.5.2.2) enumerator:
enumeration-constant
enumeration-constant \(=\) constant-expression
(6.5.3) type-qualifier:
```

const
restrict

```
volatile
(6.5.4) function-specifier:

\section*{inline}
(6.5.5) declarator:
pointer \(_{\text {opt }}\) direct-declarator
(6.5.5) direct-declarator:
identifier
( declarator )
direct-declarator [ assignment-expression \(\left.{ }_{o p t}\right]\)
direct-declarator [ * ]
direct-declarator ( parameter-type-list )
direct-declarator ( identifier-list \({ }_{\text {opt }}\) )
(6.5.5) pointer:
* type-qualifier-list \({ }_{\text {opt }}\)
* type-qualifier-list \({ }_{\text {opt }}\) pointer
(6.5.5) type-qualifier-list:
type-qualifier
type-qualifier-list type-qualifier
(6.5.5) parameter-type-list:
parameter-list
parameter-list , ...
(6.5.5) parameter-list:
parameter-declaration
parameter-list , parameter-declaration
(6.5.5) parameter-declaration:
declaration-specifiers declarator
declaration-specifiers abstract-declarator \({ }_{o p t}\)
(6.5.5) identifier-list:
identifier
identifier-list , identifier
(6.5.6) type-name:
specifier-qualifier-list abstract-declarator \({ }_{\text {opt }}\)
(6.5.6) abstract-declarator:
pointer
pointer \(_{\text {opt }}\) direct-abstract-declarator
(6.5.6) direct-abstract-declarator:
```

( abstract-declarator )
direct-abstract-declarator ${ }_{\text {opt }}$ [ assignment-expression ${ }_{\text {opt }}$ ]
direct-abstract-declarator [ * ]
direct-abstract-declarator ${ }_{\text {opt }}$ ( parameter-type-list ${ }_{\text {opt }}$ )

```
(6.5.7) typedef-name:
identifier
(6.5.8) initializer:
assignment-expression
\{ initializer-list \}
\{ initializer-list , \}
(6.5.8) initializer-list:
designation \(_{\text {opt }}\) initializer
initializer-list , designation opt initializer
(6.5.8) designation:
designator-list \(=\)
(6.5.8) designator-list:
designator
designator-list designator
(6.5.8) designator:
[ constant-expression ]
. identifier

\section*{B.2.3 Statements}
(6.6) statement:
labeled-statement
compound-statement
expression-statement
selection-statement
iteration-statement
jump-statement
(6.6.1) labeled-statement:
identifier : statement
case constant-expression : statement
default : statement
(6.6.2) compound-statement:
\{ block-item-list \(\left.{ }_{\text {opt }}\right\}\)
(6.6.2) block-item-list:
block-item
block-item-list block-item
(6.6.2) block-item:
declaration
statement
(6.6.3) expression-statement:
expression \(_{o p t}\);
(6.6.4) selection-statement:
```

if ( expression ) statement
if ( expression ) statement else statement
switch ( expression ) statement

```
(6.6.5) iteration-statement:
while ( expression ) statement
do statement while ( expression ) ;
for ( expression \({ }_{o p t}\); expression \({ }_{o p t}\); expression \({ }_{o p t}\) ) statement
for ( declaration ; expression \({ }_{o p t}\); expression \({ }_{o p t}\) ) statement
(6.6.6) jump-statement:
goto identifier ;
continue ;
break ;
return expression \(_{\text {opt }}\);

\section*{B.2.4 External definitions}
(6.7) translation-unit:
external-declaration
translation-unit external-declaration
(6.7) external-declaration:
function-definition
declaration
(6.7.1) function-definition:
declaration-specifiers declarator declaration-list \({ }_{\text {opt }}\) compound-statement

\section*{B. 3 Preprocessing directives}
(6.8) preprocessing-file:
group \(_{\text {opt }}\)
(6.8) group:
group-part
group group-part
(6.8) group-part:
pp-tokens
if-section new-line
control-line
(6.8.1) if-section:
if-group elif-groups \({ }_{\text {opt }}\) else-group \(_{\text {opt }}\) endif-line
(6.8.1) if-group:
\# if constant-expression new-line group \({ }_{\text {opt }}\)
\# ifdef identifier new-line group \({ }_{\text {opt }}\)
\# ifndef identifier new-line group opt
(6.8.1) elif-groups:
elif-group
elif-groups elif-group
(6.8.1) elif-group:
\# elif constant-expression new-line group \({ }_{\text {opt }}\)
(6.8.1) else-group:
\# else new-line group \({ }_{\text {opt }}\)
(6.8.1) endif-line:
\# endif new-line
control-line:
\# include pp-tokens new-line
\# define identifier replacement-list new-line
\# define identifier lparen identifier-list \({ }_{\text {opt }}\) )
replacement-list new-line
(6.8.3) lparen:
the left-parenthesis character without preceding white space
(6.8.3) replacement-list:
pp-tokens \({ }_{\text {opt }}\)
(6.8) pp-tokens:
preprocessing-token
pp-tokens preprocessing-token
(6.8) new-line:
the new-line character

\section*{Annex C \\ (informative) \\ Sequence points}

1 The following are the sequence points described in 5.1.2.3.
2 - The call to a function, after the arguments have been evaluated (6.3.2.2).
- The end of the first operand of the following operators: logical AND \(\& \&\) (6.3.13); logical OR || (6.3.14); conditional ? (6.3.15); comma , (6.3.17).
- The end of a full declarator: declarators (6.5.5);
- The end of a full expression: an initializer (6.5.8); the expression in an expression statement (6.6.3); the controlling expression of a selection statement (if or switch) (6.6.4); the controlling expression of a while or do statement (6.6.5); each of the three expressions of a for statement (6.6.5.3); the expression in a return statement (6.6.6.4).
- Immediately before a library function return (7.1.8).

\title{
Annex D \\ (informative) \\ \\ Library summary
} \\ \\ Library summary
}

1 D. 1 Errors <errno. h>
EDOM
EILSEQ
ERANGE
errno
D. 2 Common definitions <stddef. h >

\section*{NULL}
offsetof (type, member-designator)
ptrdiff_t
size_t
wchar_t
D. 3 Boolean type and values <stdbool.h>
bool
true
false
__bool_true_false_are_defined
D. 4 Diagnostics <assert. h>

NDEBUG
void assert (int expression);
D. 5 Character handling <ctype. h>
```

int isalnum(int c);
int isalpha(int c);
int iscntrl(int c);
int isdigit(int c);
int isgraph(int c);
int islower(int c);
int isprint(int c);
int ispunct(int c);
int isspace(int c);
int isupper(int c);
int isxdigit(int c);
int tolower(int c);
int toupper(int c);

```

\section*{D. 6 Integer types <inttypes.h>}
```

int8_t
int16_t
int32_t
int64_t
uint8_t
uint16_t
uint32_t
uint64_t
int_least8_t
int_least16_t
int_least32_t
int_least64_t
uint_least8_t
uint_least16_t
uint_least32_t
uint_least64_t
int_fast8_t
int_fast16_t
int_fast32_t
int_fast64_t
uint_fast8_t
uint_fast16_t
uint_fast32_t
uint_fast64_t
intptr_t
uintptr_t
intmax_t

```
```

uintmax_t
INT8_MIN
INT16_MIN
INT32_MIN
INT64_MIN
INT8_MAX
INT16_MAX
INT32_MAX
INT64_MAX
UINT8_MAX
UINT16_MAX
UINT32_MAX
UINT64_MAX
INT_LEAST8_MIN
INT_LEAST16_MIN
INT_LEAST32_MIN
INT_LEAST64_MIN
INT_LEAST8_MAX
INT_LEAST16_MAX
INT_LEAST32_MAX
INT_LEAST64_MAX
UINT_LEAST8_MAX
UINT_LEAST16_MAX
UINT_LEAST32_MAX
UINT_LEAST64_MAX
INT_FAST8_MIN
INT_FAST16_MIN
INT_FAST32_MIN
INT_FAST64_MIN
INT_FAST8_MAX
INT_FAST16_MAX
INT_FAST32_MAX
INT_FAST64_MAX
UINT_FAST8_MAX
UINT_FAST16_MAX
UINT_FAST32_MAX
UINT_FAST64_MAX
INTPTR_MIN
INTPTR_MAX
UINTPTR_MAX
INTMAX_MIN
INTMAX_MAX

```
```

UINTMAX_MAX
INT8_C(value)
INT16_C(value)
INT32_C(value)
INT64_C(value)
UINT8_C(value)
UINT16_C(value)
UINT32_C (value)
UINT64_C (value)
INTMAX_C (value)
UINTMAX_C (value)
PRId8
PRId16
PRId32
PRId64
PRIdLEAST8
PRIdLEAST16
PRIdLEAST32
PRIdLEAST64
PRIdFAST8
PRIdFAST16
PRIdFAST32
PRIdFAST64
PRIdMAX
PRIdPTR
PRIi8
PRIi16
PRIi32
PRIi64
PRIiLEAST8
PRIiLEAST16
PRIiLEAST32
PRIiLEAST64
PRIiFAST8
PRIiFAST16
PRIiFAST32
PRIiFAST64
PRIiMAX
PRIiPTR
PRIO8
PRIo16
PRIO32

```

PRIO64
PRIOLEAST8
PRIOLEAST16
PRIOLEAST32
PRIOLEAST64
PRIOFAST8
PRIoFAST16
PRIOFAST32
PRIOFAST64
PRIOMAX
PRIOPTR
PRIu8
PRIu16
PRIu32
PRIU64
PRIuLEAST8
PRIuLEAST16
PRIULEAST32
PRIULEAST64
PRIUFAST8
PRIuFAST16
PRIUFAST32
PRIUFAST64
PRIUMAX
PRIUPTR
PRIx8
PRIx16
PRIx32
PRIx64
PRIxLEAST8
PRIxLEAST16
PRIxLEAST32
PRIxLEAST64
PRIxFAST8
PRIxFAST16
PRIxFAST32
PRIxFAST64
PRIxMAX
PRIxPTR
PRIX8
PRIX16
PRIX32

\section*{PRIX64}

PRIXLEAST8
PRIXLEAST16
PRIXLEAST32
PRIXIEAST64
PRIXFAST8
PRIXFAST16
PRIXFAST32
PRIXFAST64
PRIXMAX
PRIXPTR
SCNd8
SCNd1 6
SCNd32
SCNd64
SCNDLEAST8
SCNDLEAST16
SCNDLEAST32
SCNDLEAST64
SCNDFAST8
SCNDFAST16
SCNDFAST32
SCNDFAST64
SCNDMAX
SCNDPTR
SCNi8
SCNi16
SCNi 32
SCNi 64
SCNiLEAST8
SCNiLEAST16
SCNiLEAST32
SCNiLEAST64
SCNiFAST8
SCNiFAST16
SCNiFAST32
SCNiFAST64
SCNiMAX
SCNiPTR
SCNO8
SCNo16
SCNo32

SCNO64
SCNOLEAST8
SCNOLEAST16
SCNoLEAST32
SCNOLEAST64
SCNoFAST8
SCNOFAST16
SCNOFAST32
SCNOFAST64
SCNOMAX
SCNOPTR
SCNu 8
SCNu16
SCNu32
SCNu64
SCNuLEAST8
SCNuLEAST16
SCNuLEAST32
SCNuLEAST64
SCNuFAST8
SCNuFAST16
SCNuFAST32
SCNuFAST64
SCNUMAX
SCNuPTR
SCNx 8
SCNx1 6
SCNx 32
SCNx 64
SCNxLEAST8
SCNxLEAST16
SCNxLEAST32
SCNxLEAST64
SCNxFAST8
SCNxFAST16
SCNxFAST32
SCNxFAST64
SCNxMAX
SCNxPTR
PTRDIFF_MIN
PTRDIFF_MAX
SIG_ATOMIC_MIN
```

SIG_ATOMIC_MAX
SIZE_MAX
WCHAR_MIN
WCHAR_MAX
WINT_MIN
WINT_MAX
intmax_t strtoimax(const char * restrict nptr,
char ** restrict endptr, int base);
uintmax_t strtoumax(const char * restrict nptr,
char ** restrict endptr, int base);
intmax_t wcstoimax(const wchar_t * restrict nptr,
wchar_t ** restrict endptr, int base);
uintmax_t wcstoumax(const wchar_t * restrict nptr,
wchar_t ** restrict endptr, int base);

```
D. 7 Floating-point environment <fenv.h>
fenv_t
fexcept_t
FE_ALL_EXCEPT
FE_DFL_ENV
FE_DIVBYZERO
FE_INEXACT
FE_INVALID
FE_OVERFLOW
FE_UNDERFLOW
FE_DOWNWARD
FE_TONEAREST
FE_TOWARDZERO
FE_UPWARD
\#pragma STDC FENV_ACCESS on-off-switch
void feclearexcept (int excepts);
void fegetexceptflag(fexcept_t *flagp, int excepts);
void feraiseexcept (int excepts);
void fesetexceptflag(const fexcept_t *flagp, int excepts);
int fetestexcept (int excepts);
int fegetround (void);
int fesetround (int round);
void fegetenv (fenv_t *envp);
int feholdexcept (fenv_t *envp);
void fesetenv (const fenv_t *envp);
void feupdateenv(const fenv_t *envp);
D. 8 Localization <locale.h>

LC_ALL
LC_COLLATE
LC_CTYPE
LC_MONETARY
LC_NUMERIC
LC_TIME
NULL
struct lconv
char *setlocale(int category, const char *locale);
struct lconv *localeconv(void);

\section*{D. 9 Mathematics <math.h>}
```

float_t
double_t
INFINITY
NAN
FP_NAN
FP_INFINITE
FP_NORMAL
FP_SUBNORMAL
FP_ZERO
FP_FAST_FMA
FP_FAST_FMAF
FP_FAST_FMAL
FP_ILOGBO
FP_ILOGBNAN
DECIMAL_DIG
\#pragma STDC FP_CONTRACT on-off-switch
int signbit(real-floating x);
HUGE_VAL
HUGE_VALL
HUGE_VALF
int fpclassify(real-floating x);
int isfinite(real-floating x);
int isnormal(real-floating x);
int isnan(real-floating x);
int isinf(real-floating x);
double acos(double x);
float acosf(float x);
long double acosl(long double x);
double asin(double x);
float asinf(float x);
long double asinl(long double x);
double atan(double x);
float atanf(float x);
long double atanl(long double x);
double atan2(double y, double x);
float atan2f(float y, float x);
long double atan2l(long double y, long double x);
double cos(double x);
float cosf(float x);
long double cosl(long double x);

```
```

double sin(double x);
float sinf(float x);
long double sinl(long double x);
double tan(double x);
float tanf(float x);
long double tanl(long double x);
double cosh(double x);
float coshf(float x);
long double coshl(long double x);
double sinh(double x);
float sinhf(float x);
long double sinhl(long double x);
double tanh(double x);
float tanhf(float x);
long double tanhl(long double x);
double acosh(double x);
float acoshf(float x);
long double acoshl(long double x);
double asinh(double x);
float asinhf(float x);
long double asinhl(long double x);
double atanh(double x);
float atanhf(float x);
long double atanhl(long double x);
double exp(double x);
float expf(float x);
long double expl(long double x);
double frexp(double value, int *exp);
float frexpf(float value, int *exp);
long double frexpl(long double value, int *exp);
double ldexp(double x, int exp);
float ldexpf(float x, int exp);
long double ldexpl(long double x, int exp);
double log(double x);
float logf(float x);
long double logl(long double x);
double log10(double x);
float log10f(float x);
long double log101(long double x);
double modf(double value, double *iptr);
float modff(float value, float *iptr);
long double modfl(long double value, long double *iptr);

```
```

double exp2(double x);
float exp2f(float x);
long double exp2l(long double x);
double expm1 (double x);
float expm1f(float x);
long double expm1l(long double x);
double log1p(double x);
float log1pf(float x);
long double log1pl(long double x);
double log2(double x);
float log2f(float x);
long double log2l(long double x);
double logb(double x);
float logbf(float x);
long double logbl(long double x);
double scalbn(double x, int n);
float scalbnf(float x, int n);
long double scalbnl(long double x, int n);
double scalbln(double x, long int n);
float scalblnf(float x, long int n);
long double scalblnl(long double x, long int n);
int ilogb(double x);
int ilogbf(float x);
int ilogbl(long double x);
double fabs(double x);
float fabsf(float x);
long double fabsl(long double x);
double pow(double x, double y);
float powf(float x, float y);
long double powl(long double x, long double y);
double sqrt(double x);
float sqrtf(float x);
long double sqrtl(long double x);
double cbrt(double x);
float cbrtf(float x);
long double cbrtl(long double x);
double hypot(double x, double y);
float hypotf(float x, float y);
long double hypotl(long double x, long double y);
double erf(double x);
float erff(float x);
long double erfl(long double x);

```
```

double erfc(double x);
float erfcf(float x);
long double erfcl(long double x);
double gamma(double x);
float gammaf(float x);
long double gammal(long double x);
double lgamma(double x);
float lgammaf(float x);
long double lgammal(long double x);
double ceil(double x);
float ceilf(float x);
long double ceill(long double x);
double floor(double x);
float floorf(float x);
long double floorl(long double x);
double nearbyint(double x);
float nearbyintf(float x);
long double nearbyintl(long double x);
double rint(double x);
float rintf(float x);
long double rintl(long double x);
long int lrint(double x);
long int lrintf(float x);
long int lrintl(long double x);
long long llrint(double x);
long long llrintf(float x);
long long llrintl(long double x);
double round(double x);
float roundf(float x);
long double roundl(long double x);
long int lround(double x);
long int lroundf(float x);
long int lroundl(long double x);
long long llround(double x);
long long llroundf(float x);
long long llroundl(long double x);
double trunc(double x);
float truncf(float x);
long double truncl(long double x);
double fmod(double x, double y);
float fmodf(float x, float y);
long double fmodl(long double x, long double y);

```
double remainder(double x, double y);
float remainderf(float x, float y);
long double remainderl(long double x, long double y);
double remquo(double x, double y, int *quo);
float remquof(float x, float y, int *quo);
long double remquol(long double x, long double y, int *quo);
double copysign(double x, double y);
float copysignf(float x, float y);
long double copysignl(long double x, long double y);
double nan(const char *tagp);
double nextafter(double x, double y);
float nextafterf(float x, float y);
long double nextafterl(long double x, long double y);
double nextafterx(double x, long double y);
float nextafterxf(float x, long float y);
long double nextafterxl(long double x, long long double y);
double fdim(double x, double y);
float fdimf(float x, float y);
long double fdiml(long double x, long double y);
double fmax(double x, double y);
float fmaxf(float x, float y);
long double fmaxl(long double x, long double y);
double fmin(double x, double y);
float fminf(float x, float y);
long double fminl(long double x, long double y);
double fma(double x, double y, double z);
float fmaf(float x, float y, float z);
long double fmal(long double x, long double y,
    long double z);
int isgreater(real-floating \mathbf{x}, real-floating y);
int isgreaterequal(real-floating \mathbf{x}, real-floating y);
int isless(real-floating \mathbf{x}, real-floating y);
int islessequal(real-floating \mathbf{x, real-floating y);}
int islessgreater(real-floating \mathbf{x, real-floating y);}
int isunordered(real-floating \mathbf{x, real-floating y);}
```

D. 10 Complex <complex.h>

```
_Complex_I
_Imaginary_I
I
#pragma STDC CX_LIMITED_RANGE on-off-switch
double complex cacos(double complex z);
float complex cacosf(float complex z);
long double complex cacosl(long double complex z);
double complex casin(double complex z);
float complex casinf(float complex z);
long double complex casinl(long double complex z);
double complex catan(double complex z);
float complex catanf(float complex z);
long double complex catanl(long double complex z);
double complex ccos(double complex z);
float complex ccosf(float complex z);
long double complex ccosl(long double complex z);
double complex csin(double complex z);
float complex csinf(float complex z);
long double complex csinl(long double complex z);
double complex ctan(double complex z);
float complex ctanf(float complex z);
long double complex ctanl(long double complex z);
double complex cacosh(double complex z);
float complex cacoshf(float complex z);
long double complex cacoshl(long double complex z);
double complex casinh(double complex z);
float complex casinhf(float complex z);
long double complex casinhl(long double complex z);
double complex catanh(double complex z);
float complex catanhf(float complex z);
long double complex catanhl(long double complex z);
double complex ccosh(double complex z);
float complex ccoshf(float complex z);
long double complex ccoshl(long double complex z);
double complex csinh(double complex z);
float complex csinhf(float complex z);
long double complex csinhl(long double complex z);
double complex ctanh(double complex z);
float complex ctanhf(float complex z);
long double complex ctanhl(long double complex z);
```

```
double complex cexp(double complex z);
float complex cexpf(float complex z);
long double complex cexpl(long double complex z);
double complex clog(double complex z);
float complex clogf(float complex z);
long double complex clogl(long double complex z);
double complex csqrt(double complex z);
float complex csqrtf(float complex z);
long double complex csqrtl(long double complex z);
double cabs(double complex z);
float cabsf(float complex z);
long double cabsl(long double complex z);
double complex cpow(double complex x, double complex y);
float complex cpowf(float complex x, float complex y);
long double complex cpowl(long double complex x,
    long double complex y);
double carg(double complex z);
float cargf(float complex z);
long double cargl(long double complex z);
double complex conj(double complex z);
float complex conjf(float complex z);
long double complex conjl(long double complex z);
double cimag(double complex z);
float cimagf(float complex z);
long double cimagl(long double complex z);
double complex cproj(double complex z);
float complex cprojf(float complex z);
long double complex cprojl(long double complex z);
double creal(double complex z);
float crealf(float complex z);
long double creall(long double complex z);
```

```
D.11 Type-generic math <tgmath.h>
    acos
    asin
    atan
    acosh
    asinh
    atanh
    cos
    sin
    tan
    cosh
    sinh
    tanh
    exp
    log
    pow
    sqrt
    fabs
    atan2
    cbrt
    ceil
    copysign
    erf
    erfc
    exp2
    expm1
    fdim
    floor
    fmax
    fmin
    fma
    fmod
    frexp
    gamma
    hypot
    ilogb
    ldexp
    lgamma
    llrint
    llround
    log10
```

```
log1p
log2
logb
lrint
lround
nearbyint
nextafter
nextafterx
remainder
remquo
rint
round
scalbn
scalbln
trunc
fma
```

D. 12 Nonlocal jumps <set jmp.h>

```
jmp_buf
int setjmp(jmp_buf env);
void longjmp(jmp_buf env, int val);
```

D. 13 Signal handling <signal.h>

```
sig_atomic_t
SIG_DFL
SIG_ERR
SIG_IGN
SIGABRT
SIGFPE
SIGILL
SIGINT
SIGSEGV
SIGTERM
void (*signal(int sig, void (*func) (int))) (int);
int raise(int sig);
```

D. 14 Variable arguments <stdarg. h>
va_list
void va_start (va_list ap, parmN);
type va_arg(va_list ap, type);
void va_end(va_list ap);
void va_copy(va_list dest, va_list src);
D. 15 Input/output <stdio.h>
_IOFBF
_IOLBF
_IONBF
BUFSIZ
EOF
FILE
FILENAME_MAX
FOPEN_MAX
fpos_t
L_tmpnam
NULL
SEEK_CUR
SEEK_END
SEEK_SET
size_t
stderr
stdin
stdout
TMP_MAX
int remove (const char *filename);
int rename (const char *old, const char *new);
FILE *tmpfile(void);
char *tmpnam(char *s);
int fclose(FILE *stream);
int fflush (FILE *stream);
FILE *fopen (const char * restrict filename, const char * restrict mode);
FILE *freopen(const char * restrict filename, const char * restrict mode, FILE * restrict stream);
void setbuf (FILE * restrict stream, char * restrict buf);
int setvbuf(FILE * restrict stream,

```
    char * restrict buf,
    int mode, size_t size);
int fprintf(FILE * restrict stream,
    const char * restrict format, ...);
int fscanf(FILE * restrict stream,
    const char * restrict format, ...);
int printf(const char * restrict format, ...);
int scanf(const char * restrict format, ...);
int sprintf(char * restrict s,
    const char * restrict format, ...);
int snprintf(char * restrict s, size_t n,
    const char * restrict format, ...);
int sscanf(const char * restrict s,
    const char * restrict format, ...);
int vfprintf(FILE * restrict stream,
    const char * restrict format,
    va_list arg);
int vprintf(const char * restrict format,
    va_list arg);
int vsprintf(char * restrict s,
    const char * restrict format,
    va_list arg);
int vsnprintf(char * restrict s, size_t n,
        const char * restrict format,
        va_list arg);
int vfscanf(FILE * restrict stream,
        const char * restrict format,
        va_list arg);
int vscanf(const char * restrict format,
        va_list arg);
int vsscanf(const char * restrict s,
        const char * restrict format,
        va_list arg);
int fgetc(FILE *stream);
char *fgets(char * restrict s, int n,
    FILE * restrict stream);
int fputc(int c, FILE *stream);
int fputs(const char * restrict s,
    FILE * restrict stream);
int getc(FILE *stream);
int getchar(void);
char *gets(char *s);
```

```
int putc(int c, FILE *stream);
int putchar(int c);
int puts(const char *s);
int ungetc(int c, FILE *stream);
size_t fread(void * restrict ptr,
    size_t size, size_t nmemb,
    FILE * restrict stream);
size_t fwrite(const void * restrict ptr,
    size_t size, size_t nmemb,
    FILE * restrict stream);
int fgetpos(FILE * restrict stream,
    fpos_t * restrict pos);
int fseek(FILE *stream, long int offset, int whence);
int fsetpos(FILE *stream, const fpos_t *pos);
long int ftell(FILE *stream);
void rewind(FILE *stream);
void clearerr(FILE *stream);
int feof(FILE *stream);
int ferror(FILE *stream);
void perror(const char *s);
```

D. 16 General utilities <stdlib.h>

```
EXIT_FAILURE
EXIT_SUCCESS
MB_CUR_MAX
NULL
RAND_MAX
div_t
ldiv_t
lldiv_t
size_t
wchar_t
double atof(const char *nptr);
int atoi(const char *nptr);
long int atol(const char *nptr);
long long int atoll(const char *nptr);
double strtod(const char * restrict nptr,
    char ** restrict endptr);
float strtof(const char * restrict nptr,
    char ** restrict endptr);
long double strtold(const char * restrict nptr,
    char ** restrict endptr);
long int strtol(const char * restrict nptr,
    char ** restrict endptr, int base);
long long int strtoll(const char * restrict nptr,
    char ** restrict endptr, int base);
unsigned long int strtoul(
    const char * restrict nptr,
    char ** restrict endptr,
    int base);
unsigned long long int strtoull(
        const char * restrict nptr,
        char ** restrict endptr,
        int base);
int rand(void);
void srand(unsigned int seed);
void *calloc(size_t nmemb, size_t size);
void free(void *ptr);
void *malloc(size_t size);
void *realloc(void *ptr, size_t size);
void abort(void);
int atexit(void (*func) (void));
```

```
void exit(int status);
char *getenv(const char *name);
int system(const char *string);
void *bsearch(const void *key, const void *base,
        size_t nmemb, size_t size,
        int (*compar) (const void *, const void *));
void qsort(void *base, size_t nmemb, size_t size,
    int (*compar) (const void *, const void *));
int abs(int j);
div_t div(int numer, int denom);
long int labs(long int j);
long long int llabs(long long int j);
ldiv_t ldiv(long int numer, long int denom);
lldiv_t lldiv(long long int numer,
    long long int denom);
int mblen(const char *s, size_t n);
int mbtowc(wchar_t * restrict pwc,
    const char * restrict s,
    size_t n);
int wctomb(char *s, wchar_t wchar);
size_t mbstowcs(wchar_t * restrict pwcs,
    const char * restrict s,
    size_t n);
size_t wcstombs(char * restrict s,
    const wchar_t * restrict pwcs,
    size_t n);
```

D. 17 String handling <string. h>

NULL

```
size_t
void *memcpy(void * restrict s1,
        const void * restrict s2,
        size_t n);
void *memmove(void *s1, const void *s2, size_t n);
char *strcpy(char * restrict s1,
        const char * restrict s2);
char *strncpy(char * restrict s1,
        const char * restrict s2,
        size_t n);
char *strcat(char * restrict s1,
        const char * restrict s2);
char *strncat(char * restrict s1,
        const char * restrict s2,
        size_t n);
int memcmp(const void *s1, const void *s2, size_t n);
int strcmp(const char *s1, const char *s2);
int strcoll(const char *s1, const char *s2);
int strncmp(const char *s1, const char *s2, size_t n);
size_t strxfrm(char * restrict s1,
    const char * restrict s2,
    size_t n);
void *memchr(const void *s, int c, size_t n);
char *strchr(const char *s, int c);
size_t strcspn(const char *s1, const char *s2);
char *strpbrk(const char *s1, const char *s2);
char *strrchr(const char *s, int c);
size_t strspn(const char *s1, const char *s2);
char *strstr(const char *s1, const char *s2);
char *strtok(char * restrict s1,
    const char * restrict s2);
void *memset(void *s, int c, size_t n);
char *strerror(int errnum);
size_t strlen(const char *s);
```

D. 18 Date and time <time. $\mathrm{h}>$

```
CLOCKS_PER_SEC
NULL
_NO_LEAP_SECONDS
_LOCALTIME
clock_t
time_t
size_t
struct tm
struct tmx
clock_t clock(void);
double difftime(time_t time1, time_t time0);
time_t mktime(struct tm *timeptr);
time_t mkxtime(struct tmx *timeptr);
time_t time(time_t *timer);
char *asctime(const struct tm *timeptr);
char *ctime(const time_t *timer);
struct tm *gmtime(const time_t *timer);
struct tm *localtime(const time_t *timer);
struct tmx *zonetime(const time_t *timer);
size_t strftime(char * restrict s,
    size_t maxsize,
    const char * restrict format,
    const struct tm * restrict timeptr);
size_t strfxtime(char * restrict s,
    size_t maxsize,
    const char * restrict format,
    const struct tmx * restrict timeptr);
```

D. 19 Alternative spellings <iso646.h>

```
and
and_eq
bitand
bitor
compl
not
not_eq
or
or_eq
xor
xor_eq
```

D. 20 Wide-character classification and mapping utilities <wctype.h>

```
wctrans_t
wctype_t
WEOF
wint_t
```

D.20.1 Wide-character classification functions

```
int iswalnum(wint_t wc);
int iswalpha(wint_t wc);
int iswcntrl(wint_t wc);
int iswdigit(wint_t wc);
int iswgraph(wint_t wc);
int iswlower(wint_t wc);
int iswprint(wint_t wc);
int iswpunct(wint_t wc);
int iswspace(wint_t wc);
int iswupper(wint_t wc);
int iswxdigit(wint_t wc);
wctype_t wctype(const char *property);
int iswctype(wint_t wc, wctype_t desc);
```

D.20.2 Wide-character mapping functions

```
wint_t towlower(wint_t wc);
wint_t towupper(wint_t wc);
wctrans_t wctrans(const char *property);
wint_t towctrans(wint_t wc, wctrans_t desc);
```

D. 21 Extended multibyte and wide-character utilities <wchar.h>
mbstate_t
NULL
size_t
struct tm
wchar_t
WCHAR_MAX
WCHAR_MIN
WEOF
wint_t
D.21.1 Formatted wide-character input/output functions

```
int fwprintf(FILE * restrict stream,
    const wchar_t * restrict format, ...);
int fwscanf(FILE * restrict stream,
        const wchar_t * restrict format, ...);
int wprintf(const wchar_t * restrict format, ...);
int wscanf(const wchar_t * restrict format, ...);
int swprintf(wchar_t * restrict s,
        size_t n,
        const wchar_t * restrict format, ...);
int swscanf(const wchar_t * restrict s,
        const wchar_t * restrict format, ...);
int vfwprintf(FILE * restrict stream,
        const wchar_t * restrict format,
        va_list arg);
int vwprintf(const wchar_t * restrict format,
        va_list arg);
int vswprintf(wchar_t * restrict s,
        size_t n,
        const wchar_t * restrict format,
        va_list arg);
int vfwscanf(FILE * restrict stream,
        const wchar_t * restrict format,
        va_list arg);
int vwscanf(FILE * restrict stream,
        const wchar_t * restrict format,
        va_list arg);
int vswscanf(const wchar_t * restrict s,
        const wchar_t * restrict format,
        va_list arg);
```


## D.21.2 Wide-character input/output functions

```
wint_t fgetwc(FILE *stream);
wchar_t *fgetws(wchar_t * restrict s,
    int n, FILE * restrict stream);
wint_t fputwc(wchar_t c, FILE *stream);
int fputws(const wchar_t * restrict s,
    FILE * restrict stream);
wint_t getwc(FILE *stream);
wint_t getwchar(void);
wint_t putwc(wchar_t c, FILE *stream);
wint_t putwchar(wchar_t c);
wint_t ungetwc(wint_t c, FILE *stream);
int fwide(FILE *stream, int mode);
```

D.21.3 Wide-string numeric conversion functions

```
double wcstod (const wchar_t * restrict nptr,
    wchar_t ** restrict endptr);
float wcstof (const wchar_t * restrict nptr,
    wchar_t ** restrict endptr);
long double wcstold(const wchar_t * restrict nptr,
    wchar_t ** restrict endptr);
long int wcstol(const wchar_t * restrict nptr,
    wchar_t ** restrict endptr, int base);
long long int wcstoll(const wchar_t * restrict nptr,
    wchar_t ** restrict endptr, int base);
unsigned long int wcstoul(const wchar_t * restrict nptr,
    wchar_t ** restrict endptr, int base);
unsigned long long int wcstoull(
    const wchar_t * restrict nptr,
    wchar_t ** restrict endptr, int base);
```


## D.21.4 Wide-string functions

```
wchar_t *wcscpy (wchar_t * restrict s1,
    const wchar_t * restrict s2);
wchar_t *wcsncpy (wchar_t * restrict s1,
    const wchar_t * restrict s2, size_t n);
wchar_t *wcscat (wchar_t * restrict s1,
    const wchar_t * restrict s2);
wchar_t *wcsncat(wchar_t * restrict s1,
    const wchar_t * restrict s2, size_t n);
int wcscmp (const wchar_t *s1, const wchar_t *s2);
int wcscoll (const wchar_t *s1, const wchar_t *s2);
int wcsncmp (const wchar_t *s1, const wchar_t *s2,
        size_t n);
size_t wcsxfrm(wchar_t * restrict s1,
        const wchar_t * restrict s2, size_t n);
wchar_t *wcschr (const wchar_t *s, wchar_t c);
size_t wcscspn (const wchar_t *s1, const wchar_t *s2);
wchar_t *wcspbrk (const wchar_t *s1, const wchar_t *s2);
wchar_t *wcsrchr (const wchar_t *s, wchar_t c);
size_t wcsspn (const wchar_t *s1, const wchar_t *s2);
wchar_t *wcsstr(const wchar_t *s1, const wchar_t *s2);
wchar_t *wcstok (wchar_t * restrict s1,
    const wchar_t * restrict s2,
    wchar_t ** restrict ptr);
size_t wcslen(const wchar_t *s);
wchar_t *wmemchr (const wchar_t *s, wchar_t c, size_t n);
int wmemcmp (wchar_t * restrict s1,
    const wchar_t * restrict s2,
    size_t n);
wchar_t *wmemcpy (wchar_t * restrict s1,
    const wchar_t * restrict s2,
    size_t n);
wchar_t *wmemmove (wchar_t *s1, const wchar_t *s2,
    size_t n);
wchar_t *wmemset (wchar_t *s, wchar_t c, size_t n);
```

D.21.5 Wide-string time conversion functions

```
size_t wcsftime(wchar_t *s, size_t maxsize,
    const wchar_t *format, const struct tm *timeptr);
size_t wcsfxtime(wchar_t *s, size_t maxsize,
        const wchar_t *format, const struct tmx *timeptr);
```

D.21.6 Extended multibyte/wide-character conversion functions

```
wint_t btowc(int c);
int wctob(wint_t c);
int mbsinit(const mbstate_t *ps);
size_t mbrlen(const char * restrict s, size_t n,
    mbstate_t * restrict ps);
size_t mbrtowc(wchar_t * restrict pwc,
    const char * restrict s, size_t n,
    mbstate_t * restrict ps);
size_t wcrtomb(char * restrict s, wchar_t wc,
    mbstate_t * restrict ps);
size_t mbsrtowcs(wchar_t * restrict dst,
    const char ** restrict src, size_t len,
    mbstate_t * restrict ps);
size_t wcsrtombs(char * restrict dst,
    const wchar_t ** restrict src, size_t len,
    mbstate_t * restrict ps);
```


## Annex $\mathbf{E}$ (informative) <br> Implementation limits

1 The contents of a header <limits.h> are given below, in alphabetic order. The minimum magnitudes shown shall be replaced by implementation-defined magnitudes with the same sign. The values shall all be constant expressions suitable for use in \#if preprocessing directives. The components are described further in 5.2.4.2.1.

| \#define CHAR_BIT | 8 |
| :--- | ---: |
| \#define CHAR_MAX | UCHAR_MAX or SCHAR_MAX |
| \#define CHAR_MIN | 0 or SCHAR_MIN |
| \#define INT_MAX | +32767 |
| \#define INT_MIN | -32767 |
| \#define LONG_MAX | +2147483647 |
| \#define LONG_MIN | -2147483647 |
| \#define LLONG_MAX | +9223372036854775807 |
| \#define LLONG_MIN | -9223372036854775807 |
| \#define MB_LEN_MAX | 1 |
| \#define SCHAR_MAX | +127 |
| \#define SCHAR_MIN | -127 |
| \#define SHRT_MAX | +32767 |
| \#define SHRT_MIN | -32767 |
| \#define UCHAR_MAX | 255 |
| \#define USHRT_MAX | 65535 |
| \#define UINT_MAX | 65535 |
| \#define ULONG_MAX |  |
| \#define ULLONG_MAX | 18446744073709551615 |

2 The contents of a header <float.h> are given below. The value of FLT_RADIX shall be a constant expression suitable for use in \#if preprocessing directives. Values that need not be constant expressions shall be supplied for all other components. The components are described further in 5.2.4.2.2.

```
#define FLT_EVAL_METHOD
#define FLT_ROUNDS
```

3 The values given in the following list shall be replaced by implementation-defined expressions that shall be equal or greater in magnitude (absolute value) to those shown, with the same sign:

```
#define DBL_DIG 10
#define DBL_MANT_DIG
#define DBL_MAX_10_EXP +37
#define DBL_MAX_EXP
#define DBL_MIN_10_EXP -37
#define DBL_MIN_EXP
#define FLT_DIG 6
#define FLT_MANT_DIG
#define FLT_MAX_10_EXP +37
#define FLT_MAX_EXP
#define FLT_MIN_10_EXP -37
#define FLT_MIN_EXP
#define FLT_RADIX 2
#define LDBL_DIG 10
#define LDBL_MANT_DIG
#define LDBL_MAX_10_EXP +37
#define LDBL_MAX_EXP
#define LDBL_MIN_10_EXP -37
#define LDBL_MIN_EXP
```

4 The values given in the following list shall be replaced by implementation-defined expressions that shall be equal to or greater than those shown:

```
#define DBL_MAX
1E+37
```

\#define FLT_MAX 1E+37
\#define LDBL_MAX 1E+37

5 The values given in the following list shall be replaced by implementation-defined expressions that shall be equal to or less than those shown:

| \#define DBL_EPSILON | $1 \mathrm{E}-9$ |
| :--- | ---: |
| \#define DBL_MIN | $1 \mathrm{E}-37$ |
| \#define FLT_EPSILON | $1 \mathrm{E}-5$ |
| \#define FLT_MIN | $1 \mathrm{E}-37$ |
| \#define LDBL_EPSILON | $1 \mathrm{E}-9$ |
| \#define LDBL_MIN | $1 \mathrm{E}-37$ |

## Annex $\mathbf{F}$ (normative) <br> IEC 559 floating-point arithmetic

## F. 1 Introduction

1 This annex specifies C language support for the IEC 559 floating-point standard. The IEC 559 floating-point standard is specifically Binary floating-point arithmetic for microprocessor systems, second edition (IEC 559:1993), also known as IEEE Standard for Binary Floating-Point Arithmetic (ANSI/IEEE 754-1985). IEEE Standard for RadixIndependent Floating-Point Arithmetic (ANSI/IEEE 854-1987) generalizes the binary standard to remove dependencies on radix and word length. IEC 559 generally refers to the floating-point standard, as in IEC 559 operation, IEC 559 format, etc. An implementation that defines __STDC_IEC_559__ conforms to the specification in this annex. Where a binding between the C language and IEC 559 is indicated, the IEC 559 -specified behavior is adopted by reference, unless stated otherwise.

## F. 2 Types

1 The C floating types match the IEC 559 formats as follows:

- The float type matches the IEC 559 single format.
- The double type matches the IEC 559 double format.
- The long double type matches an IEC 559 extended format, ${ }^{272}$ else a non-IEC 559 extended format, else the IEC 559 double format.

Any non-IEC 559 extended format used for the long double type has more precision than IEC 559 double and at least the range of IEC 559 double. ${ }^{273}$

## Recommended practice

The long double type matches an IEC 559 extended format.
272. Extended is IEC 559's double-extended data format. Extended refers to both the common 80-bit and quadruple 128-bit IEC 559 formats.
273. A non-IEC 559 long double type must provide infinity and NaNs , as its values must include all double values.

## F.2.1 Infinities, signed zeros, and NaNs

1 This specification does not define the behavior of signaling NaNs. ${ }^{274}$ It generally uses the term NaN to denote quiet NaNs. The NAN and INFINITY macros and the nan function in <math. h > provide designations for IEC 559 NaNs and infinities.

## F. 3 Operators and functions

1 C operators and functions provide IEC 559 required and recommended facilities as listed below.

- The +, -, *, and / operators provide the IEC 559 add, subtract, multiply, and divide operations.
- The sqrt function in <math. h> provides the IEC 559 square root operation.
- The remainder function in <math. $\mathrm{h}>$ provides the IEC 559 remainder operation. The remquo function in <math. $\mathrm{h}>$ provides the same operation but with additional information.
- The rint function in <math.h> provides the IEC 559 operation that rounds a floating-point number to an integer value (in the same precision). The C nearbyint function in <math.h> provides the nearbyinteger function recommended in the Appendix to IEEE standard 854.
- The conversions for floating types provide the IEC 559 conversions between floatingpoint precisions.
- The conversions from integer to floating types provide the IEC 559 conversions from integer to floating point.
- The conversions from floating to integer types provide IEC 559-like conversions but always round toward zero.
- The lrint and llrint functions in <math.h> provide the IEC 559 conversions, which honor the directed rounding mode, from floating point to the long and long long integer formats. The lrint and llrint functions can be used to implement IEC 559 conversions from floating to other integer formats.
- The translation time conversion of floating constants and the strtod, fprintf, fscanf, and related library functions in <stdlib.h>, <stdio.h>, and <wchar.h> provide IEC 559 binary-decimal conversions. The strtold function in

274. Since NaNs created by IEC 559 operations are always quiet, quiet NaNs (along with infinities) are sufficient for closure of the arithmetic.
<stdlib. h> provides the conv function recommended in the Appendix to IEEE standard 854.

- The relational and equality operators provide IEC 559 comparisons. IEC 559 identifies a need for additional comparison predicates to facilitate writing code that accounts for NaNs. The comparison macros (isgreater, isgreaterequal, isless, islessequal, islessgreater, and isunordered) in <math.h> supplement the language operators to address this need. The islessgreater and isunordered macros provide respectively a quiet version of the <> predicate and the unordered predicate recommended in the Appendix to IEC 559.
- The feclearexcept, feraiseexcept, and fetestexcept functions in <fenv.h> provide the facility to test and alter the IEC 559 floating-point exception flags. The fegetexceptflag and fesetexceptflag functions in <fenv.h> provide the facility to save and restore all five status flags at one time. These functions are used in conjunction with the type fexcept_t and the exception macros (FE_INEXACT, FE_DIVBYZERO, FE_UNDERFLOW, FE_OVERFLOW, FE_INVALID) also in <fenv.h>.
- The fegetround and fesetround functions in <fenv.h> provide the facility to select among the IEC 559 directed rounding modes represented by the rounding direction macros (FE_TONEAREST, FE_UPWARD, FE_DOWNWARD, FE_TOWARDZERO) also in <fenv.h>.
- The fegetenv, feholdexcept, fesetenv, and feupdateenv functions in <fenv.h> provide a facility to manage the floating-point environment, comprising the IEC 559 status flags and control modes.
- The copysign function in <math.h> provides the copysign function recommended in the Appendix to IEC 559.
- The unary minus (-) operator provides the minus (-) operation recommended in the Appendix to IEC 559.
- The scalbn function in <math. h > provides the scalb function recommended in the Appendix to IEC 559.
- The logb function in <math. h> provides the logb function recommended in the Appendix to IEC 559, but following the newer specification in IEEE 854.
- The nextafter and nextafterx functions in <math.h> provide the nextafter function recommended in the Appendix to IEC 559 (but with a minor change to better handle signed zeros).
- The isfinite macro in <math. h> provides the finite function recommended in the Appendix to IEC 559.
- The isnan macro in <math.h> provides the isnan function recommended in the Appendix to IEC 559.
- The signbit macro and the fpclassify macro in <math.h>, used in conjunction with the number classification macros (FP_NAN, FP_INFINITE, FP_NORMAL, FP_SUBNORMAL, FP_ZERO), provide the facility of the class function recommended in the Appendix to IEC 559 (except that fpclassify does not distinguish signaling from quiet NaNs ).


## F. 4 Floating to integer conversion

1 If the floating value is infinite or NaN or if the integral part of the floating value exceeds the range of the integer type, then the invalid exception is raised and the resulting value is unspecified. Whether conversion of non-integer floating values whose integral part is within the range of the integer type raises the inexact exception is unspecified. ${ }^{275}$

## F. 5 Binary-decimal conversion

1 Conversion from the widest supported IEC 559 format to decimal with DECIMAL_DIG digits and back is the identity function. ${ }^{276}$

2 Conversions involving IEC 559 formats follow all pertinent recommended practice. In particular, conversion between any supported IEC 559 format and decimal with DECIMAL_DIG or fewer significant digits is correctly rounded.

[^97]
## F. 6 Contracted expressions

1 A contracted expression treats infinities, NaNs, signed zeros, subnormals, and the rounding directions in a manner consistent with the basic arithmetic operations covered by IEC 559.

## Recommended practice

2 A contracted expression raises exceptions in a manner generally consistent with the basic arithmetic operations. A contracted expression delivers the same value as its uncontracted counterpart, else is correctly rounded (once).

## F. 7 Environment

1 The floating-point environment defined in <fenv.h> includes the IEC 559 exception status flags and directed-rounding control modes. It includes also IEC 559 dynamic rounding precision and trap enablement modes, if the implementation supports them. ${ }^{277}$

## F.7.1 Environment management

1 IEC 559 requires that floating-point operations implicitly raise exception status flags, and that rounding control modes can be set explicitly to affect result values of floating-point operations. When the state for the FENV_ACCESS pragma (defined in <fenv.h>) is on, these changes to the floating-point state are treated as side effects which respect sequence points. ${ }^{278}$

## F.7.2 Translation

1 During translation the IEC 559 default modes are in effect:

- The rounding direction mode is rounding to nearest.
- The rounding precision mode (if supported) is set so that results are not shortened.
- Trapping or stopping (if supported) is disabled on all exceptions.


## Recommended practice

2 The implementation produces a diagnostic message for each translation-time floating-point exception, other than inexact; ${ }^{279}$ the implementation then proceeds with the translation of

[^98]the program.

## F.7.3 Execution

1 At program startup the floating-point environment is initialized as prescribed by IEC 559:

- All exception status flags are cleared.
- The rounding direction mode is rounding to nearest.
- The dynamic rounding precision mode (if supported) is set so that results are not shortened.
- Trapping or stopping (if supported) is disabled on all exceptions.


## F.7.4 Constant expressions

1 An arithmetic constant expression of floating type, other than one in an initializer for an object that has static storage duration, is evaluated (as if) during execution. As executiontime evaluation, it is affected by any operative modes and raises exceptions as required by IEC 559 (provided the state for the FENV_ACCESS pragma is on). ${ }^{280}$

## Examples

```
#include <fenv.h>
#pragma STDC FENV_ACCESS ON
void f(void)
{
    float w[] = { 0.0/0.0 }; // raises an exception
    static float x = 0.0/0.0; // does not raise an exception
    float y = 0.0/0.0; // raises an exception
    double z = 0.0/0.0; // raises an exception
    /* ... */
}
```

For the static initialization, the division is done at translation time, raising no (executiontime) exceptions. On the other hand, for the three automatic initializations the invalid division occurs at execution time.

[^99]const static double one_third = 1.0/3.0;

## F.7.5 Initialization

All computation for automatic initialization is done (as if) at execution time. As execution-time evaluation, it is affected by any operative modes and raises exceptions as required by IEC 559 (provided the state for the FENV_ACCESS pragma is on). All computation for initialization of objects that have static storage duration is done (as if) at translation time.

## Examples

```
#include <fenv.h>
#pragma STDC FENV_ACCESS ON
void f(void)
{
    float u[] = { 1.1e75 }; // raises exceptions
    static float v = 1.1e75; // does not raise exceptions
    float w = 1.1e75; // raises exceptions
    double x = 1.1e75; // may raise exceptions
    float y = 1.1e75f; // may raise exceptions
    long double z = 1.1e75; // does not raise exceptions
    /* ... */
}
```

The static initialization of $\mathbf{v}$ raises no (execution-time) exceptions because its computation is done at translation time. The automatic initialization of $\mathbf{u}$ and $\mathbf{w}$ require an executiontime conversion to float of the wider value $1.1 \mathrm{e75}$, which raises exceptions. The automatic initializations of $\mathbf{x}$ and $\mathbf{y}$ entail execution-time conversion; however, in some expression evaluation methods, the conversions is not to a narrower format, in which case no exception is raised. ${ }^{281}$ The automatic initialization of $\mathbf{z}$ entails execution-time conversion, but not to a narrower format, so no exception is raised. Note that the conversions of the floating constants 1.1 e 75 and $1.1 \mathrm{e75f}$ to their internal representations occur at translation time in all cases.

[^100]
## F.7.6 Changing the environment

1 Operations defined in 6.3 and functions and macros defined for the standard libraries change flags and modes just as indicated by their specification (including conformance to IEC 559). They do not change flags or modes (so as to be detectable by the user) in any other cases.

2 If the argument to the feraiseexcept function in <fenv.h> represents IEC 559 valid coincident exceptions for atomic operations (namely overflow and inexact, or underflow and inexact) then overflow or underflow is raised before inexact.

## F. 8 Optimization

1 This section identifies code transformations that might subvert IEC 559 -specified behavior, and others that do not.

## F.8.1 Global transformations

1 Floating-point arithmetic operations and external function calls may entail side effects which optimization must honor, at least where the state of the FENV_ACCESS pragma is on. The flags and modes in the floating-point environment may be regarded as global variables; floating-point operations ( $\boldsymbol{+}, \boldsymbol{*}$, etc.) implicitly read the modes and write the flags.

2 Concern about side effects may inhibit code motion and removal of seemingly useless code. For example, in

```
#include <fenv.h>
#pragma STDC FENV_ACCESS ON
void f(double x)
{
    /* ... */
    for (i = 0; i < n; i++) x + 1;
    /* ... */
}
```

$\mathbf{x}+1$ might raise exceptions, so cannot be removed. And since the loop body might not execute (maybe $0 \geq \mathbf{n}$ ), $\mathbf{x}+\mathbf{1}$ cannot be moved out of the loop. (Of course these optimizations are valid if the implementation can rule out the nettlesome cases.)

3 This specification does not require support for trap handlers that maintain information about the order or count of exceptions. Therefore, between function calls exceptions need not be precise: the actual order and number of occurrences of exceptions (>1) may vary from what the source code expresses. Thus the preceding loop could be treated as

```
if (0<n) x + 1;
```


## F.8.2 Expression transformations

| $\mathbf{x} / 2 \leftrightarrow x * 0.5$ | Although similar transformations involving inexact constants generally do not yield numerically equivalent expressions, if the constants are exact then such transformations can be made on IEC 559 machines and others that round perfectly. |
| :---: | :---: |
| 1 * $\mathbf{x}$ and $\mathbf{x} / 1 \rightarrow \mathbf{x}$ | The expressions $\mathbf{1} \boldsymbol{*} \mathbf{x}, \mathbf{x} / \mathbf{1}$, and $\mathbf{x}$ are equivalent (on IEC 559 machines, among others). ${ }^{282}$ |
| $\mathbf{x} / \mathrm{x} \rightarrow \mathbf{1 . 0}$ | The expressions $\mathbf{x} / \mathbf{x}$ and $\mathbf{1 . 0}$ are not equivalent if $\mathbf{x}$ can be zero, infinite, or NaN . |
| $\mathbf{x}-\mathbf{y} \leftrightarrow \mathbf{x}+(-\mathbf{y})$ | The expressions $\mathbf{x}-\mathbf{y}, \mathbf{x}+(-\mathbf{y})$, and $(-\mathbf{y})+\mathbf{x}$ are equivalent (on IEC 559 machines, among others). |
| $\mathbf{x}-\mathbf{y} \leftrightarrow-(\mathbf{y}-\mathbf{x})$ | The expressions $\mathbf{x}-\mathbf{y}$ and $-(\mathbf{y}-\mathbf{x})$ are not equivalent because $1-1$ is +0 but $-(1-1)$ is -0 (in the default rounding direction). ${ }^{283}$ |
| $\mathbf{x}-\mathbf{x} \rightarrow 0.0$ | The expressions $\mathbf{x}-\mathbf{x}$ and $\mathbf{0 . 0}$ are not equivalent if $\mathbf{x}$ is a NaN or infinite. |
| $0 * x \rightarrow 0.0$ | The expressions $\mathbf{0} \boldsymbol{*} \mathbf{x}$ and $\mathbf{0 . 0}$ are not equivalent if $\mathbf{x}$ is a NaN, infinite, or -0 . |
| $\mathbf{x}+\mathbf{0} \rightarrow \mathbf{x}$ | The expressions $\mathbf{x}+\mathbf{0}$ and $\mathbf{x}$ are not equivalent if $\mathbf{x}$ is -0 , because $(-0)+(+0)$ yields +0 (in the default rounding direction), not -0 . |
| $\mathbf{x}-\mathbf{0} \rightarrow \mathbf{x}$ | $(+0)-(+0)$ yields -0 when rounding is downward (toward $-\infty$ ), but +0 otherwise, and ( -0 ) - ( +0 ) always yields -0 ; so, if the state of the FENV_ACCESS pragma is off, promising default rounding, then the implementation can |

Although similar transformations involving inexact constants generally do not yield numerically equivalent expressions, if the constants are exact then such transformations can be made on IEC 559 machines and others that round perfectly.
$\mathbf{1} * \mathbf{x}$ and $\mathbf{x} / \mathbf{1} \rightarrow \mathbf{x} \quad$ The expressions $\mathbf{1} \mathbf{x} \mathbf{x} / \mathbf{1}$, and $\mathbf{x}$ are equivalent (on IEC 559 machines, among others). ${ }^{282}$

The expressions $\mathbf{x} / \mathbf{x}$ and $\mathbf{1 . 0}$ are not equivalent if $\mathbf{x}$ can be zero, infinite, or NaN .

The expressions $\mathbf{x}-\mathbf{y}, \mathbf{x}+(-\mathbf{y})$, and $(-\mathbf{y})+\mathbf{x}$ are equivalent (on IEC 559 machines, among others).

The expressions $\mathbf{x}-\mathbf{y}$ and $-(\mathbf{y}-\mathbf{x})$ are not equivalent because $1-1$ is +0 but $-(1-1)$ is -0 (in the default rounding direction). ${ }^{283}$

The expressions $\mathbf{x}-\mathbf{x}$ and 0.0 are not equivalent if $\mathbf{x}$ is a NaN or infinite.

The expressions $\mathbf{0} \boldsymbol{*} \mathbf{x}$ and $\mathbf{0 . 0}$ are not equivalent if $\mathbf{x}$ is a NaN, infinite, or -0 . -0 , because $(-0)+(+0)$ yields +0 (in the default rounding direction), not -0 .
$(+0)-(+0)$ yields -0 when rounding is downward (toward $-\infty$ ), but +0 otherwise, and ( -0 ) - ( +0 ) always yields -0 ; so, if the state of the FENV_ACCESS pragma is off, promising default rounding, then the implementation can
282. Strict support for signaling NaNs - not required by this specification - would invalidate these and other transformations that remove arithmetic operators.
283. IEC 559 prescribes a signed zero to preserve mathematical identities across certain discontinuities. Examples include:

and
$\operatorname{conj}(\operatorname{csqrt}(z))$ is $\operatorname{csqrt}(\operatorname{conj}(z))$, for complex $z$.
replace $\mathbf{x}-\mathbf{0}$ by $\mathbf{x}$, even if $\mathbf{x}$ might be zero.

```
-x}\leftrightarrow0-\mathbf{x
```

The expressions $-\mathbf{x}$ and $\mathbf{0}-\mathbf{x}$ are not equivalent if $\mathbf{x}$ is +0 , because $-(+0)$ yields -0 , but $0-(+0)$ yields +0 (unless rounding is downward).

## F.8.3 Relational operators

```
x != x }->\mathrm{ false
x == x }->\mathrm{ true
x < y }->\mathrm{ isless(x,y)
```

The statement $\mathbf{x} \quad \boldsymbol{!}=\mathbf{x}$ is true if $\mathbf{x}$ is a NaN.
The statement $\mathbf{x}=\mathbf{x}$ is false if $\mathbf{x}$ is a NaN.
(and similarly for $<=,>,>=$ ) Though numerically equal, these expressions are not equivalent because of side effects when $\mathbf{x}$ or $\mathbf{y}$ is a NaN and the state of the FENV_ACCESS pragma is on. This transformation, which would be desirable if extra code were required to cause the invalid exception for unordered cases, could be performed provided the state of the FENV_ACCESS pragma is off.

The sense of relational operators must be maintained. This includes handling unordered cases as expressed by the source code.

## Examples

```
    // calls \(\mathbf{g}\) and raises invalid
    // if \(\mathbf{a}\) and \(\mathbf{b}\) are unordered
    if (a < b)
        f();
    else
        g();
```

is not equivalent to

```
// calls £ and raises invalid
// if \mathbf{a}}\mathrm{ and }\mathbf{b}\mathrm{ are unordered
if (a >= b)
    g();
else
    f();
```

nor to

```
// calls f without raising invalid
// if a and b are unordered
if (isgreaterequal(a,b))
    g();
else
    f();
```

nor, unless the state of the FENV_ACCESS pragma is off, to

```
// calls g}\mathrm{ without raising invalid
```

// if $\mathbf{a}$ and $\mathbf{b}$ are unordered
if (isless (a,b))
f();
else
g();
but is equivalent to

```
if (!(a < b))
    g();
else
    f();
```


## F.8.4 Constant arithmetic

1 The implementation must honor exceptions raised by execution-time constant arithmetic wherever the state of the FENV_ACCESS pragma is on. (See F.7.4 and F.7.5.) An operation on constants that raises no exception can be folded during translation, except, if the state of the FENV_ACCESS pragma is on, a further check is required to assure that changing the rounding direction to downward does not alter the sign of the result, ${ }^{284}$ and implementations that support dynamic rounding precision modes must assure further that the result of the operation raises no exception when converted to the semantic type of the operation.

[^101]
## F. 9 <math.h>

1 This subclause contains specification of <math. $\mathrm{h}>$ facilities that is particularly suited for IEC 559 implementations.

2 The Standard C macro HUGE_VAL and its float and long double analogs, HUGE_VALF and HUGE_VALL, expand to expressions whose values are positive infinities.

3 Special cases for functions in <math.h> are covered directly or indirectly by IEC 559. F. 3 identifies the functions that IEC 559 specifies directly. The other functions in <math.h> treat infinities, NaNs, signed zeros, subnormals, and (provided the state of the FENV_ACCESS pragma is on) the exception flags in a manner consistent with the basic arithmetic operations covered by IEC 559.

4 The invalid and divide-by-zero exceptions are raised as specified in subsequent subclauses of this annex.

5 The overflow exception is raised whenever an infinity - or, because of rounding direction, a maximal-magnitude finite number - is returned in lieu of a value whose magnitude is too large.

6 The underflow exception is raised whenever a result is tiny (essentially subnormal or zero) and suffers loss of accuracy. ${ }^{285}$

7 Whether or when the trigonometric, hyperbolic, base-e exponential, base-e logarithmic, error, and $\log$ gamma functions raise the inexact exception is implementation-defined. For other functions, the inexact exception is raised whenever the rounded result is not identical to the mathematical result.

8 Whether the inexact exception may be raised when the rounded result actually does equal the mathematical result is implementation-defined. Whether the underflow (and inexact) exception may be raised when a result is tiny but not inexact is implementation-defined. ${ }^{286}$ Otherwise, as implied by F.7.6, the <math.h> functions do not raise spurious exceptions (detectable by the user).

9 Whether the functions honor the rounding direction mode is implementation-defined.
10 Generally, one-parameter functions of a NaN argument return that same NaN and raise no exception.

[^102]11 The specification in the following subclauses appends to the definitions in <math. h >.

## F.9.1 Trigonometric functions

## F.9.1.1 The acos function

1 - $\operatorname{acos}(1)$ returns +0 .

- acos ( $\mathbf{x}$ ) returns a NaN and raises the invalid exception for $|\mathbf{x}|>1$.


## F.9.1.2 The asin function

$1-\operatorname{asin}( \pm 0)$ returns $\pm 0$.
— asin (x) returns a NaN and raises the invalid exception for $|\mathbf{x}|>1$.

## F.9.1.3 The atan function

1 - atan ( $\pm 0$ ) returns $\pm 0$.

- atan ( $\pm \infty$ ) returns $\pm \pi / 2$.


## F.9.1.4 The atan2 function

1 - If one argument is a NaN then atan2 returns that same NaN ; if both arguments are NaNs then atan2 returns one of its arguments.

- atan2 $( \pm 0, \mathbf{x})$ returns $\pm 0$, for $\mathbf{x}>0$.
- atan2 $( \pm 0,+0)$ returns $\pm 0 .{ }^{287}$
- atan2 ( $\pm 0, \mathbf{x}$ ) returns $\pm \pi$, for $\mathbf{x}<0$.
- atan2 ( $\pm 0,-0$ ) returns $\pm \pi$.
- atan2 $(\mathbf{y}, \pm 0)$ returns $\pi / 2$ for $\mathrm{y}>0$.
- atan2 $(\mathbf{y}, \pm 0)$ returns $-\pi / 2$ for $\mathbf{y}<0$.
- atan2 $( \pm \mathbf{y}, \infty)$ returns $\pm 0$, for finite $\mathbf{y}>0$.
- atan2 ( $\pm \infty, \mathbf{x}$ ) returns $\pm \pi / 2$, for finite $\mathbf{x}$.
- atan2 $( \pm \mathbf{y},-\infty)$ returns $\pm \pi$, for finite $\mathbf{y}>0$.
- atan2 $( \pm \infty, \infty)$ returns $\pm \pi / 4$.

287. atan2 $(0,0)$ does not raise the invalid exception, nor does atan2 $(\mathbf{y}, 0)$ raise the divide-by-zero exception.

- atan2 $( \pm \infty,-\infty)$ returns $\pm 3 \pi / 4$.


## F.9.1.5 The cos function

$1-\cos ( \pm 0)$ returns 1 .

- $\boldsymbol{\operatorname { c o s }}( \pm \infty)$ returns a NaN and raises the invalid exception.
F.9.1.6 The sin function
$1-\sin ( \pm 0)$ returns $\pm 0$.
- $\boldsymbol{\operatorname { s i n }}( \pm \infty)$ returns a NaN and raises the invalid exception.
F.9.1.7 The $\tan$ function
$1-\tan ( \pm 0)$ returns $\pm 0$.
- $\tan ( \pm \infty)$ returns a NaN and raises the invalid exception.


## F.9.2 Hyperbolic functions

## F.9.2.1 The acosh function

1 - acosh (1) returns +0 .
$-\operatorname{acosh}(+\infty)$ returns $+\infty$.

- $\operatorname{acosh}(\mathbf{x})$ returns a NaN and raises the invalid exception if $\mathbf{x}<1$.
F.9.2.2 The asinh function
$1-\operatorname{asinh}( \pm 0)$ returns $\pm 0$.
- asinh ( $\pm \infty$ ) returns $\pm \infty$.


## F.9.2.3 The atanh function

1 - atanh ( $\pm 0$ ) returns $\pm 0$.

- atanh ( $\pm 1$ ) returns $\pm \infty$ and raises the divide-by-zero exception.
— $\operatorname{atanh}(\mathbf{x})$ returns a NaN and raises the invalid exception if $|\mathbf{x}|>1$.
F.9.2.4 The cosh function
$1-\cosh ( \pm 0)$ returns 1 .
$-\cosh ( \pm \infty)$ returns $+\infty$.


## F.9.2.5 The sinh function

$1-\sinh ( \pm 0)$ returns $\pm 0$.

- $\sinh ( \pm \infty)$ returns $\pm \infty$.


## F.9.2.6 The tanh function

$1-\tanh ( \pm 0)$ returns $\pm 0$.

- $\tanh ( \pm \infty)$ returns $\pm 1$.


## F.9.3 Exponential and logarithmic functions

F.9.3.1 The exp function
$1-\exp ( \pm 0)$ returns 1 .
$-\exp (+\infty)$ returns $+\infty$.
$-\exp (-\infty)$ returns +0 .

## F.9.3.2 The exp2 function

$1-\exp 2( \pm 0)$ returns 1 .

- $\exp 2(+\infty)$ returns $+\infty$.
$-\exp 2(-\infty)$ returns +0 .
F.9.3.3 The expm1 function

1 - expm1 ( $\pm 0$ ) returns $\pm 0$.

- expm1 ( $+\infty$ ) returns $+\infty$.
- expm1 ( $-\infty$ ) returns -1 .


## F.9.3.4 The frexp function

1 - frexp ( $\pm 0, \exp$ ) returns $\pm 0$, and returns 0 in *exp.

- frexp ( $\pm \infty$, exp) returns $\pm \infty$, and returns an unspecified value in *exp.
- frexp ( $\mathbf{x}, \mathbf{e x p}$ ) returns $\mathbf{x}$ if $\mathbf{x}$ is a NaN, and stores an unspecified value in *exp.
- frexp raises no exception.

2 On a binary system, frexp is equivalent to the comma expression

```
( (*exp = (value == 0) ? 0 :
    (int)(1 + logb(value))), scalbn(value,-(*exp)) )
```


## F.9.3.5 The 1dexp function

1 On a binary system, $1 \operatorname{dexp}(\mathbf{x}, \exp )$ is equivalent to

```
scalbn (x, exp)
```


## F.9.3.6 The log function

$1-\log ( \pm 0)$ returns $-\infty$ and raises the divide-by-zero exception.
$-\log (1)$ returns +0 .
$-\log (\mathbf{x})$ returns a NaN and raises the invalid exception if $\mathbf{x}<0$.

- $\log (+\infty)$ returns $+\infty$.


## F.9.3.7 The log10 function

$1-\log 10( \pm 0)$ returns $-\infty$ and raises the divide-by-zero exception.
$-\log 10(1)$ returns +0 .

- $\log 10(\mathbf{x})$ returns a NaN and raises the invalid exception if $\mathbf{x}<0$.
- $\log 10(+\infty)$ returns $+\infty$.


## F.9.3.8 The log1p function

$1-\log 1 \mathrm{p}( \pm 0)$ returns $\pm 0$.

- $\log 1 \mathrm{p}(-1)$ returns $-\infty$ and raises the divide-by-zero exception.
- $\log 1 \mathrm{p}(\mathbf{x})$ returns a NaN and raises the invalid exception if $\mathbf{x}<-1$.
- $\log 1 \mathrm{p}(+\infty)$ returns $+\infty$.


## F.9.3.9 The $\log 2$ function

$1-\log 2( \pm 0)$ returns $-\infty$ and raises the divide-by-zero exception.

- $\log 2(\mathbf{x})$ returns a NaN and raises the invalid exception if $\mathbf{x}<0$.
- $\log 2(+\infty)$ returns $+\infty$.


## F.9.3.10 The logb function

$1-\operatorname{logb}( \pm \infty)$ returns $+\infty$.

- logb ( $\pm 0$ ) returns $-\infty$ and raises the divide-by-zero exception.


## F.9.3.11 The modf functions

- modf (value,iptr) returns a result with the same sign as the argument value.
- modf ( $\pm \infty$,iptr) returns $\pm 0$ and stores $\pm \infty$ through iptr.
- modf of a NaN argument returns that same NaN and also stores it through iptr.
modf behaves as though implemented by

```
#include <math.h>
#include <fenv.h>
#pragma STDC FENV_ACCESS ON
double modf(double value, double *iptr)
{
    int save_round = fegetround();
    fesetround (FE_TOWARDZERO);
    *iptr = nearbyint(value);
    fesetround(save_round);
    return copysign(
        isinf(value) ? 0.0 :
                value - (*iptr), value);
}
```

F.9.3.12 The scalbn function

1 - scalbn $(\mathbf{x}, \mathbf{n})$ returns $\mathbf{x}$ if $\mathbf{x}$ is infinite, zero, or a NaN.

- scalbn ( $\mathbf{x}, 0$ ) returns $\mathbf{x}$.


## F.9.4 Power and absolute value functions

## F.9.4.1 The fabs function

1 - fabs ( $\pm 0$ ) returns +0 .

- fabs ( $\pm \infty$ ) returns $+\infty$.


## F.9.4.2 The hypot function

$1-\operatorname{hypot}(\mathbf{x}, \mathbf{y}), \operatorname{hypot}(\mathbf{y}, \mathbf{x})$, and $\operatorname{hypot}(\mathbf{x},-\mathbf{y})$ are equivalent.

- hypot ( $\mathbf{x}, \mathbf{y}$ ) returns $+\infty$ if $\mathbf{x}$ is infinite.
- If both arguments are NaNs then hypot returns one of its arguments; otherwise, if $\mathbf{x}$ is a NaN and y in not infinite then hypot returns that same NaN .
- hypot ( $\mathbf{x}, \pm 0$ ) is equivalent to $\mathbf{f a b s}(\mathbf{x})$.


## F.9.4.3 The pow function

- pow ( $\mathbf{x}, \pm 0$ ) returns 1 for any $\mathbf{x}$.
- pow $(\mathbf{x},+\infty)$ returns $+\infty$ for $|\mathbf{x}|>1$.
- pow $(\mathbf{x},+\infty)$ returns +0 for $|\mathbf{x}|<1$.
- pow $(\mathbf{x},-\infty)$ returns +0 for $|\mathbf{x}|>1$.
- pow $(\mathbf{x},-\infty)$ returns $+\infty$ for $|\mathbf{x}|<1$.
- pow $(+\infty, y)$ returns $+\infty$ for $\mathbf{y}>0$.
- pow $(+\infty, y)$ returns +0 for $\mathbf{y}<0$.
- pow $(-\infty, y)$ returns $-\infty$ for $\mathbf{y}$ an odd integer $>0$.
- pow $(-\infty, \mathbf{y})$ returns $+\infty$ for $\mathbf{y}>0$ and not an odd integer.
- pow $(-\infty, y)$ returns -0 for $\mathbf{y}$ an odd integer $<0$.
- pow $(-\infty, \mathbf{y})$ returns +0 for $\mathbf{y}<0$ and not an odd integer.
- pow ( $\mathbf{x}, \mathbf{y}$ ) returns one of its NaN arguments if $\mathbf{y}$ is a NaN , or if $\mathbf{x}$ is a NaN and $\mathbf{y}$ is nonzero.
- pow ( $\pm \mathbf{1}, \pm \infty$ ) returns a NaN and raises the invalid exception.
- pow ( $\mathbf{x}, \mathbf{y}$ ) returns a NaN and raises the invalid exception for finite $\mathbf{x}<0$ and finite non-integer $\mathbf{y}$.
- pow ( $\pm \mathbf{0}, \mathbf{y}$ ) returns $\pm \infty$ and raises the divide-by-zero exception for $\mathbf{y}$ an odd integer < 0 .
- pow $( \pm \mathbf{0}, \mathbf{y})$ returns $+\infty$ and raises the divide-by-zero exception for $\mathbf{y}<0$ and not an odd integer.
- pow ( $\pm \mathbf{0}, \mathbf{y}$ ) returns $\pm 0$ for $\mathbf{y}$ an odd integer $>0$.
- pow ( $\pm \mathbf{0}, \mathbf{y}$ ) returns +0 for $\mathbf{y}>0$ and not an odd integer.


## F.9.4.4 The sqrt function

sqrt is fully specified as a basic arithmetic operation in IEC 559.

## F.9.4.5 The cbrt function

- cbrt ( $\pm \infty$ ) returns $\pm \infty$ (for all rounding directions).
- cbrt ( $\pm 0$ ) returns $\pm 0$ (for all rounding directions).


## F.9.5 Error and gamma functions

F.9.5.1 The erf function

1 - erf $( \pm 0)$ returns $\pm 0$.

- erf $( \pm \infty)$ returns $\pm 1$.
F.9.5.2 The erfc function

1 - erfc ( $+\infty$ ) returns +0 .

- erfc $(-\infty)$ returns 2 .
F.9.5.3 The gamma function

1 - gamma ( $+\infty$ ) returns $+\infty$.

- gamma ( $\mathbf{x}$ ) returns a NaN and raises the invalid exception if $\mathbf{x}$ is a negative integer or zero.
- gamma ( $-\infty$ ) returns a NaN and raises the invalid exception.


## F.9.5.4 The lgamma function

1 - Igamma ( $+\infty$ ) returns $+\infty$.

- lgamma ( $\mathbf{x}$ ) returns $+\infty$ and raises the divide-by-zero exception if $\mathbf{x}$ is a negative integer or zero.
- lgamma ( $-\infty$ ) returns $+\infty$.


## F.9.6 Nearest integer functions

## F.9.6.1 The ceil function

1 - ceil (x) returns $\mathbf{x}$ if $\mathbf{x}$ is $\pm \infty$ or $\pm 0$.
The double version of ceil behaves as though implemented by

```
#include <math.h>
#include <fenv.h>
#pragma STDC FENV_ACCESS ON
double ceil(double x)
{
    double result;
    int save_round = fegetround();
    fesetround (FE_UPWARD);
    result = rint(x); // or nearbyint instead of rint
    fesetround(save_round);
    return result;
}
```


## F.9.6.2 The floor function

1 - $\mathbf{f l o o r}(\mathbf{x})$ returns $\mathbf{x}$ if $\mathbf{x}$ is $\pm \infty$ or $\pm 0$.
See the sample implementation for ceil in F.9.6.1.

## F.9.6.3 The nearbyint function

1 The nearbyint function differs from the rint function only in that the nearbyint function does not raise the inexact flag.

## F.9.6.4 The rint function

1 The rint function uses IEC 559 rounding according to the current rounding direction. It raises the inexact exception if its argument differs in value from its result.

- rint ( $\pm 0$ ) returns $\pm 0$ (for all rounding directions).
- rint $( \pm \infty)$ returns $\pm \infty$ (for all rounding directions).


## F.9.6.5 The lrint function

1 The lrint function provides floating-to-integer conversion as prescribed by IEC 559. It rounds according to the current rounding direction. If the rounded value is outside the range of long int, the numeric result is unspecified and the invalid exception is raised. When it raises no other exception and its argument differs from its result, lrint raises the inexact exception.

## F.9.6.6 The round function

1 The double version of round behaves as though implemented by

```
#include <math.h>
#include <fenv.h>
#pragma STDC FENV_ACCESS ON
double round(double x)
{
        double result;
        fenv_t save_env;
        feholdexcept (&save_env);
        result = rint(x);
        if (fetestexcept(FE_INEXACT)) {
            fesetround (FE_TOWARDZERO);
            result = rint(copysign(0.5 + fabs(x),x));
        }
        feupdateenv(&save_env);
        return result;
}
```

The round function may but is not required to raise the inexact exception for non-integer numeric arguments, as this implementation does.

## F.9.6.7 The lround function

1 The lround function differs from lrint with the default rounding direction just in that lround (1) rounds halfway cases away from zero and (2) may but need not raise the inexact exception for non-integer arguments that round to within the range of long int.

## F.9.6.8 The trunc function

1 The trunc function uses IEC 559 rounding toward zero (regardless of the current rounding direction).

## F.9.7 Remainder functions

## F.9.7.1 The fmod function

1 - If one argument is a NaN then fmod returns that same NaN ; if both arguments are NaNs then fmod returns one of its arguments.

- $\operatorname{fmod}( \pm \mathbf{0}, \mathbf{y})$ returns $\pm 0$ if $\mathbf{y}$ is not zero.
- $\mathbf{f m o d}(\mathbf{x}, \mathbf{y})$ returns a NaN and raises the invalid exception if $\mathbf{x}$ is infinite or $\mathbf{y}$ is zero.
- $\operatorname{fmod}(\mathbf{x}, \pm \infty)$ returns $\mathbf{x}$ if $\mathbf{x}$ is not infinite.

The double version of fmod behaves as though implemented by

```
#include <math.h>
#include <fenv.h>
#pragma STDC FENV_ACCESS ON
double fmod(double x, double y)
{
double result;
result = remainder(fabs(x),(y = fabs(y)));
if (signbit(result)) result += y;
return copysign(result,x);
}
```


## F.9.7.2 The remainder function

1 The remainder function is fully specified as a basic arithmetic operation in IEC 559.

## F.9.7.3 The remquo function

1 The remquo function follows the specification for remainder. It has no further specification special to IEC 559 implementations.

## F.9. 8 Manipulation functions

F.9.8.1 The copysign function

1 copysign is specified in the Appendix to IEC 559.

## F.9.8.2 The nan function

1 All IEC 559 implementations support quiet NaNs , in all floating formats.

## F.9.8.3 The nextafter function

- If one argument is a NaN then nextafter returns that same NaN ; if both arguments are NaNs then nextafter returns one of its arguments.
- nextafter $(\mathbf{x}, \mathbf{y})$ raises the overflow and inexact exceptions if $\mathbf{x}$ is finite and the function value is infinite.
- nextafter ( $\mathbf{x}, \mathbf{y}$ ) raises the underflow and inexact exceptions if the function value is subnormal or zero and $\mathbf{x} \neq \mathbf{y}$.


## F.9.9 Maximum, minimum, and positive difference functions

## F.9.9.1 The fdim function

1 - If one argument is a NaN then fdim returns that same NaN ; if both arguments are NaNs then fdim returns one of its arguments.

## F.9.9.2 The fmax function

1 - If just one argument is a NaN then fmax returns the other argument; if both arguments are NaNs then fmax returns one of its arguments.

The body of the fmax function might be ${ }^{288}$
\{ return (isgreaterequal ( $\mathbf{x}, \mathrm{y}$ ) ||

$$
\text { isnan(y)) ? x : y; \} }
$$

## F.9.9.3 The fmin function

$1 \quad \mathrm{fmin}$ is analogous to fmax. See F.9.9.2.

## F.9.10 Floating point multiply-add

## F.9.10.1 The fma function

1 - fma $(\mathbf{x}, \mathbf{y}, \mathbf{z})$ computes the sum $\mathbf{z}$ plus the product $\mathbf{x}$ times $\mathbf{y}$, correctly rounded once.

- $\mathbf{f m a}(\mathbf{x}, \mathbf{y}, \mathbf{z})$ returns one of its NaN arguments and raises no exception if $\mathbf{x}$ or $\mathbf{y}$ is a NaN .
- $\mathbf{f m a}(\mathbf{x}, \mathbf{y}, \mathbf{z})$ returns a NaN and optionally raises the invalid exception if one of $\mathbf{x}$ and $\mathbf{y}$ is infinite, the other is zero, and $\mathbf{z}$ is a NaN.
- fma $(\mathbf{x}, \mathbf{y}, \mathbf{z})$ returns a NaN and raises the invalid exception if one of $\mathbf{x}$ and $\mathbf{y}$ is infinite, the other is zero, and $\mathbf{z}$ is not a NaN.
- $\mathbf{f m a}(\mathbf{x}, \mathbf{y}, \mathbf{z})$ returns a NaN and raises the invalid exception if $\mathbf{x}$ times $\mathbf{y}$ is an exact infinity and $\mathbf{z}$ is also an infinity but with the opposite sign.

[^103]
## Annex G <br> (informative) <br> IEC 559-compatible complex arithmetic

## G. 1 Introduction

1 This annex supplements Annex F to specify complex arithmetic for compatibility with IEC 559 real floating-point arithmetic. An implementation supports this specification if and only if it defines the macro __STD_IEC_559_COMPLEX__.

## G. 2 Types

1 There are three imaginary types, designated as float imaginary, double imaginary, and long double imaginary. The imaginary types (along with the real floating and complex types) are floating types.

2 For imaginary types, the corresponding real type is given by deleting the keyword imaginary from the type name.

3 Each imaginary type has the same representation and alignment requirements as the corresponding real type. The value of an object of imaginary type is the value of the real representation times the imaginary unit.

4 The imaginary type-domain comprises the imaginary types.

## G. 3 Conversions

## G.3.1 Imaginary types

1 Conversions among imaginary types follow rules analogous to those for real floating types.

## G.3.2 Real and imaginary

1 When a value of imaginary type is converted to a real type, the result is a positive zero.
2 When a value of real type is converted to an imaginary type, the result is a positive imaginary zero.

## G.3.3 Imaginary and complex

1 When a value of imaginary type is converted to a complex type, the real part of the complex result value is a positive zero and the imaginary part of the complex result value is determined by the conversion rules for the corresponding real types.

2 When a value of complex type is converted to an imaginary type, the real part of the complex value is discarded and the value of the imaginary part is converted according to the conversion rules for the corresponding real types.

## G. 4 Binary operators

1 The following subclauses supplement 6.3 in order to specify the type of the result for an operation with an imaginary operand.

2 For most operand types, the value of the result of a binary operator with an imaginary or complex operand is completely determined, with reference to real arithmetic, by the usual mathematical formula. For some operand types, the usual mathematical formula is problematic because of its treatment of infinities and because of undue overflow or underflow; in these cases the result satisfies certain properties (specified in G.4.1), but is not completely determined.

## G.4.1 Multiplicative operators

## Semantics

1 If one operand has real type and the other operand has imaginary type, then the result has imaginary type. If both operands have imaginary type, then the result has real type. (If either operand has complex type, then the result has complex type.)

2 If the operands are not both complex, then the result and exception behavior of the * operator is defined by the usual mathematical formula:

| $*$ | real $x$ | imaginary $y * I$ | complex $x+y * I$ |
| :--- | :--- | :--- | :--- |
| real $u$ | $x * u$ | $(y * u) * I$ | $(x * u)+(y * u) * I$ |
| imaginary $v * I$ | $(x * v) * I$ | $-y * v$ | $(-y * v)+(x * v) * I$ |
| complex $u+v * I$ | $(x * u)+$ <br> $(x * v) * I$ | $(-y * v)+$ <br> $(y * u) * I$ |  |

If the second operand is not complex, then the result and exception behavior of the / operator is defined by the usual mathematical formula:

| $/$ | real $x$ | imaginary $y * I$ | complex $x+y$ *I |
| :--- | :--- | :--- | :--- |
| real $u$ | $x / u$ | $(y / u) * I$ | $(x / u)+(y / u) * I$ |
| imaginary $v * I$ | $(-x / v) * I$ | $y / v$ | $(y / v)+(-x / v) * I$ |

4 A complex or imaginary value with at least one infinite part is regarded as an infinity (even if its other part is a NaN ). A complex or imaginary value is a finite number if each of its parts is a finite number (neither infinite nor NaN ). A complex or imaginary value is a zero if each of its parts is a zero. The * and / operators satisfy the following infinity properties for all real, imaginary, and complex operands: ${ }^{289}$

- if one operand is an infinity and the other operand is a nonzero finite number or an infinity, then the result of the * operator is an infinity;
- if the first operand is an infinity and the second operand is a finite number, then the result of the / operator is an infinity;
- if the first operand is a finite number and the second operand is an infinity, then the result of the / operator is a zero;
- if the first operand is a nonzero finite number or an infinity and the second operand is a zero, then the result of the / operator is an infinity.

If both operands of the * operator are complex or if the second operand of the / operator is complex, the operator raises exceptions if appropriate for the calculation of the parts of the result, and may raise spurious exceptions.

## Examples

1. Multiplication of double complex operands could be implemented as follows. Note that the imaginary unit I has imaginary type (see G.5).
[^104]```
#include <math.h>
#include <complex.h>
/* Multiply z * w ... */
double complex _Cmultd(double complex z, double complex w)
{
    #pragma STDC FP_CONTRACT OFF
    double a, b, c, d, ac, bd, ad, bc, x, y;
    a = creal(z); b = cimag(z)
    c = creal(w); d = cimag(w);
    ac = a * c; bd = b * d;
    ad = a * d; bc = b * c;
    x = ac - bd;
    y = ad + bc;
    /* Recover infinities that computed as NaN+iNaN ... */
    if (isnan(x) && isnan(y)) {
    int recalc = 0;
    if (isinf(a) || isinf(b) ) { /* z is infinite */
        /* "Box" the infinity ... */
            a = copysign(isinf(a) ? 1.0 : 0.0, a);
            b = copysign(isinf(b) ? 1.0 : 0.0, b);
            /* Change NaNs in the other factor to 0 ... */
            if (isnan(c)) c = copysign(0.0, c);
            if (isnan(d)) d = copysign(0.0, d);
            recalc = 1;
            }
            if (isinf(c) || isinf(d) ) { /* w is infinite */
            /* "Box" the infinity ... */
            c = copysign(isinf(c) ? 1.0 : 0.0, c);
            d = copysign(isinf(d) ? 1.0 : 0.0, d);
            /* Change NaNs in the other factor to 0 ... */
            if (isnan(a)) a = copysign(0.0, a);
            if (isnan(b)) b = copysign(0.0, b);
            recalc = 1;
            }
            if (!recalc) {
            /* *Recover infinities from overflow cases ... */
            if (isinf(ac) || isinf(bd) ||
                isinf(ad) || isinf(bc)) {
                /* Change all NaNs to 0 ... */
                if (isnan(a)) a = copysign(0.0, a);
                if (isnan(b)) b = copysign(0.0, b);
```

```
                if (isnan(c)) c = copysign(0.0, c);
                if (isnan(d)) d = copysign(0.0, d);
                recalc = 1;
        }
        }
        if (recalc) {
            x = INFINITY * ( a * c - b * d );
        y = INFINITY * ( a * d + b * c );
    }
    }
    return x + I * y;
}
```

In ordinary (finite) cases, the cost to satisfy the infinity property for the * operator is only one isnan test. This implementation opts for performance over guarding against undue overflow and underflow.
2. Division of two double complex operands could be implemented as follows.

```
#include <math.h>
#include <complex.h>
/* Divide z / w ... */
double complex _Cdivd(double complex z, double complex w)
{
    #pragma STDC FP_CONTRACT OFF
    double a, b, c, d, logbw, denom, x, y;
    int ilogbw = 0;
    a = creal(z); b = cimag(z);
    c = creal(w); d = cimag(w);
    logbw = logb(fmax(fabs(c), fabs(d)));
    if (isfinite(logbw)) {
        ilogbw = (int)logbw;
        c = scalbn(c, -ilogbw);
        d = scalbn(d, -ilogbw);
    }
denom = c * c + d * d;
x = scalbn((a * c + b * d) / denom, -ilogbw);
y = scalbn((b * c - a * d) / denom, -ilogbw);
/*
    * Recover infinities and zeros that computed
    * as NaN+iNaN; the only cases are non-zero/zero,
    * infinite/finite, and finite/infinite, ...
```

```
        */
        if (isnan(x) && isnan(y)) {
        if ((denom == 0.0) &&
            (!isnan(a) || !isnan(b))) {
            x = copysign(INFINITY, c) * a;
            y = copysign(INFINITY, c) * b;
        }
        else if ((isinf(a) || isinf(b)) &&
            isfinite(c) && isfinite(d)) {
            a = copysign(isinf(a) ? 1.0 : 0.0, a);
            b = copysign(isinf(b) ? 1.0 : 0.0, b);
            x = INFINITY * ( a * c + b * d );
            y = INFINITY * ( b * c - a * d );
        }
        else if (isinf(logbw) &&
            isfinite(a) && isfinite(b)) {
            c = copysign(isinf(c) ? 1.0 : 0.0, c);
            d = copysign(isinf(d) ? 1.0 : 0.0, d);
            x = 0.0 * ( a * c + b * d );
            y = 0.0 * ( b * c - a * d );
    }
    }
    return x + I * y;
}
```

7 Scaling the denominator alleviates the main overflow and underflow problem, which is more serious than for multiplication. In the spirit of the multiplication example above, this code does not defend against overflow and underflow in the calculation of the numerator. Scaling with the scalbn function, instead of with division, provides better roundoff characteristics.

## G.4.2 Additive operators

## Semantics

1 If one operand has real type and the other operand has imaginary type, then the result has complex type. If both operands have imaginary type, then the result has imaginary type. (If either operand has complex type, then the result has complex type.)

2 In all cases the result and exception behavior of a + or - operator is defined by the usual mathematical formula:

| $\pm$ | $\mathbf{x}$ | Y* 1 | $\mathbf{x}+\mathrm{y}^{*} \mathrm{I}$ |
| :---: | :---: | :---: | :---: |
| u | $\mathbf{x} \pm \mathbf{u}$ | $\pm \mathbf{~ + ~} \mathbf{Y}^{*} \mathbf{I}$ | $\begin{aligned} & (x \pm u)+ \\ & y^{*} I \end{aligned}$ |
| v*I | $\mathbf{x} \pm \mathrm{v}^{*} \mathrm{I}$ | $(\mathrm{y} \pm \mathrm{v}) * \mathrm{I}$ | $\begin{aligned} & x+ \\ & (y \pm v) * I \end{aligned}$ |
| $\mathbf{u}+\mathrm{v}^{*} \mathrm{I}$ | $\begin{aligned} & (x \pm u) \pm \\ & v * I \end{aligned}$ | $\begin{aligned} & \pm u+ \\ & (\mathrm{y} \pm \mathrm{v}) * \mathrm{I} \end{aligned}$ | $\begin{aligned} & (\mathrm{x} \pm \mathrm{u})+ \\ & (\mathrm{y} \pm \mathrm{v}) * \mathrm{I} \end{aligned}$ |

## G. 5 <complex.h>

1 The macro

```
_Imaginary_I
```

is defined, and the macro
I
is defined to be _Imaginary_I (7.8).
2 This subclause contains specification for the <complex.h> functions that is particularly suited to IEC 559 implementations.

3 The functions are continuous onto both sides of their branch cuts, taking into account the sign of zero. For example, csqrt ( $-2 \pm 0 * I$ ) $== \pm$ sqrt (2)*I.

4 Since complex and imaginary values are composed of real values, each function may be regarded as computing real values from real values. Except as noted, the functions treat real infinities, NaNs, signed zeros, subnormals, and the exception flags in a manner consistent with the specification for real functions in F.9.

5 The functions conj, cimag, cproj, and creal are fully specified for all implementations, including IEC 559 ones, in 7.9.2. These functions raise no exceptions.

6 Each of the functions cabs and carg is specified by a formula in terms of a real function (whose special cases are covered in annex F ):

```
cabs(x + i*y) = hypot(x, y)
carg(x + i*y) = atan2 (y, x)
```

7 Each of the functions casin, catan, ccos, csin, ctan, and cpow is specified implicitly by a formula in terms of other complex functions (whose special cases are specified below):

```
casin(z) = -i*casinh(i*z)
catan(z) = -i*catanh(i*z)
ccos(z) = ccosh(i*z)
csin(z) = -i*csinh(i*z)
ctan(z) = -i*ctanh(i*z)
cpow(z, c) = cexp(c * clog(z))
```

For the other functions, the following subclauses specify behavior for special cases, including treatment of the invalid and divide-by-zero exceptions. For a function $f$ satisfying $f(\operatorname{conj}(z))=\operatorname{conj}(f(z))$, the specification for the upper half-plane implies the specification for the lower half-plane; if also the function $f$ is either even, $f(-z)=f(z)$, or odd, $f(-z)=-f(z)$, then the specification for the first quadrant implies the specification for the other three quadrants.

## G.5.1 The cacos function

$1-\operatorname{cacos}(\operatorname{conj}(z))=\operatorname{conj}(\operatorname{cacos}(z))$.

- cacos ( $\pm 0+\mathrm{i} 0)$ returns $\pi / 2-i 0$.
- cacos ( $-\infty+\mathrm{i} \infty$ ) returns $3 \pi / 4-i \infty$.
- cacos ( $+\infty+\mathrm{i} \infty$ ) returns $\pi / 4-i \infty$.
- cacos ( $\mathbf{x}+\mathbf{i} \infty$ ) returns $\pi / 2-i \infty$, for finite $\mathbf{x}$.
- cacos ( $-\infty+\mathrm{iy}$ ) returns $\pi-i \infty$, for positive-signed finite $\mathbf{y}$.
- cacos ( $+\infty+\mathrm{iy}$ ) returns $+0-i \infty$, for positive-signed finite $\mathbf{y}$.
- cacos ( $\pm \infty+\mathrm{iNAN}$ ) returns $N a N \pm i \infty$ (where the sign of the imaginary part of the result is unspecified).
$-\operatorname{cacos}( \pm 0+\mathrm{iNAN})$ returns $\pi / 2+i N a N$.
- cacos (NAN+im) returns $N a N-i \infty$.
- cacos (x+iNAN) returns $N a N+i N a N$ and optionally raises the invalid exception, for nonzero finite $\mathbf{x}$.
- cacos (NAN+iy) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{y}$.
- cacos (NAN+iNAN) returns $N a N+i N a N$.


## G.5.2 The cacosh function

$-\operatorname{cacosh}(\operatorname{conj}(z))=\operatorname{conj}(\operatorname{cacosh}(z))$.

- cacosh ( $\pm 0+\mathrm{i} 0$ ) returns $+0+i \pi / 2$.
$-\operatorname{cacosh}(-\infty+\mathrm{i} \infty)$ returns $+\infty+i 3 \pi / 4$.
- cacosh ( $+\infty+\mathrm{i} \infty$ ) returns $+\infty+i \pi / 4$.
- $\mathbf{c a c o s h}(\mathbf{x}+\mathrm{i} \infty)$ returns $+\infty+i \pi 2$, for finite $\mathbf{x}$.
- cacosh ( $-\infty+\mathrm{i} \mathbf{y}$ ) returns $+\infty+i \pi$, for positive-signed finite $\mathbf{y}$.
- cacosh ( $+\infty+\mathrm{iy}$ ) returns $+\infty+i 0$, for positive-signed finite $\mathbf{y}$.
- cacosh (NAN $+\mathrm{i} \infty$ ) returns $+\infty+i N a N$.
- cacosh ( $\pm \infty+$ inAN $)$ returns $+\infty+i N a N$.
- cacosh (x+iNAN) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{x}$.
- cacosh (NAN+iy) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{y}$.
- cacosh (NAN+iNAN) returns $N a N+i N a N$.


## G.5.3 The casinh function

- casinh (conj(z)) = conj(casinh(z)) and casinh is odd.
- casinh (+0+i0) returns $0+i 0$.
- casinh ( $\infty+\mathrm{i} \infty$ ) returns $+\infty+i \pi / 4$.
- casinh $(\mathbf{x}+\mathrm{i} \infty)$ returns $+\infty+i \pi / 2$ for positive-signed finite $\mathbf{x}$.
- casinh ( $+\infty+\mathrm{iy}$ ) returns $+\infty+i 0$ for positive-signed finite $\mathbf{y}$.
- casinh (NAN $+\mathrm{i} \infty$ ) returns $\pm \infty+i N a N$ (where the sign of the real part of the result is unspecified).
- casinh ( $+\infty+$ inAN) returns $+\infty+i N a N$.
- casinh (NAN+iO) returns $N a N+i 0$.
- casinh (NAN+iy) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite nonzero $\mathbf{y}$.
- casinh (x+iNAN) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{x}$.
- casinh (NAN+iNAN) returns $N a N+i N a N$.


## G.5.4 The catanh function

- catanh $(\operatorname{conj}(z))=\operatorname{conj}(\operatorname{catanh}(z))$ and catanh is odd.
- catanh ( $+0+\mathrm{i} 0$ ) returns $+0+i 0$.
- catanh ( $+\infty+\mathrm{i} \infty$ ) returns $+0+i \pi /$.
- catanh ( $+\infty+\mathrm{iy}$ ) returns $+0+i \pi / 2$, for finite positive-signed $\mathbf{y}$.
- catanh $(\mathbf{x}+\mathrm{i} \infty)$ returns $+0+i \pi / 2$, for finite positive-signed $\mathbf{x}$.
- catanh (+0+iNAN) returns $+0+i N a N$.
- catanh (NAN $+\mathrm{i} \infty$ ) returns $\pm 0+i \pi / 2$ (where the sign of the real part of the result is unspecified).
- catanh ( $+\infty+$ iNAN ) returns $+0+i N a N$.
- catanh (NAN+iy) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{y}$.
- catanh (x+iNAN) returns NAN+iNAN and optionally raises the invalid exception for nonzero finite $\mathbf{x}$.
- catanh (NAN+iNAN) returns $N a N+i N a N$.


## G.5.5 The ccosh function

$1-\operatorname{ccosh}(\operatorname{conj}(z))=\operatorname{conj}(\operatorname{ccosh}(z))$ and $\operatorname{ccosh}$ is even.
$-\operatorname{ccosh}(+0+i 0)$ returns $1+i 0$.

- ccosh ( $+0+\mathrm{i} \infty$ ) returns $N a N \pm i 0$ (where the sign of the imaginary part of the result is unspecified) and raises the invalid exception.
$-\operatorname{ccosh}(+\infty+i 0)$ returns $+\infty+i 0$.
- $\operatorname{ccosh}(+\infty+\mathrm{i} \infty)$ returns $+\infty+i N a N$ and raises the invalid exception.
- $\operatorname{ccosh}(\mathbf{x}+\mathrm{i} \infty)$ returns $N a N+i N a N$ and raises the invalid exception, for finite nonzero $\mathbf{x}$.
$-\operatorname{ccosh}(+\infty+\mathbf{i y})$ returns $(+\infty) *$ cis $(\mathbf{y})$, for finite nonzero $\mathbf{y} .{ }^{290}$

290. $\operatorname{cis}(y)$ is defined by $\cos (y)+i * \sin (y)$.

- ccosh (+0+iNaN) returns $N a N \pm i 0$ (where the sign of the imaginary part of the result is unspecified).
- $\operatorname{ccosh}(+\infty+i N a N)$ returns $+\infty+i N a N$.
- ccosh (x+iNAN) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite nonzero $\mathbf{x}$.
- ccosh (NAN+i0) returns $N a N \pm i 0$ (where the sign of the imaginary part of the result is unspecified).
- $\operatorname{ccosh}(\mathbf{N A N}+\mathrm{iy})$ returns NAN+iNAN and optionally raises the invalid exception, for all nonzero numbers $\mathbf{y}$.
- ccosh (NAN+iNAN) returns $N a N+i N a N$.


## G.5.6 The csinh function

$1-\operatorname{csinh}(\operatorname{conj}(z))=\operatorname{conj}(\operatorname{csinh}(z))$ and $\operatorname{csinh}$ is odd.
$-\operatorname{csinh}(+0+\mathrm{iO})$ returns $+0+i 0$.

- csinh ( $+0+\mathrm{i} \infty$ ) returns $\pm 0+i N a N$ (where the sign of the real part of the result is unspecified) and raises the invalid exception.
- csinh ( $+\infty+\mathrm{iO}$ ) returns $+\infty+i 0$.
- csinh $(+\infty+\mathrm{i} \infty$ ) returns $\pm \infty+i N a N$ (where the sign of the real part of the result is unspecified) and raises the invalid exception.
— csinh ( $+\infty+\mathbf{i y}$ ) returns $(+\infty) *$ cis $(\mathbf{y})$, for positive finite $\mathbf{y}$.
- csinh ( $\mathbf{x}+\mathrm{i} \infty$ ) returns $N a N+i N a N$ and raises the invalid exception, for positive finite $\mathbf{x}$.
- csinh (+0+iNAN) returns $\pm 0+i N a N$ (where the sign of the real part of the result is unspecified).
- csinh ( $+\infty+$ inAN) returns $\pm \infty+i N a N$ (where the sign of the real part of the result is unspecified).
- csinh (x+iNAN) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite nonzero $\mathbf{x}$.
- csinh (NAN+iO) returns $N a N+i 0$.
- csinh (NAN+iy) returns $N a N+i N a N$ and optionally raises the invalid exception, for all nonzero numbers $\mathbf{y}$.
- csinh (NAN+iNAN) returns $N a N+i N a N$.


## G.5.7 The ctanh function

- ctanh $(\operatorname{conj}(z))=\operatorname{conj}(\operatorname{ctanh}(z))$ and ctanh is odd.
$-c \tanh (+0+i 0)$ returns $+0+i 0$.
- ctanh ( $+\infty+\mathrm{i} \mathbf{y}$ ) returns $1+i 0$, for all positive-signed numbers $\mathbf{y}$.
- ctanh ( $\mathbf{x}+\mathrm{i} \infty$ ) returns $N a N+i N a N$ and raises the invalid exception, for finite $\mathbf{x}$.
- ctanh ( $+\infty+\mathrm{iNAN}$ ) returns $1 \pm i 0$ (where the sign of the imaginary part of the result is unspecified).
- ctanh (NAN+iO) returns $N a N+i 0$.
- ctanh (NAN+iy) returns $N a N+i N a N$ and optionally raises the invalid exception, for all nonzero numbers $\mathbf{y}$.
- ctanh (x+iNAN) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{x}$.
- ctanh (NAN+iNAN) returns $N a N+i N a N$.


## G.5.8 The cexp function

$1-\operatorname{cexp}(\operatorname{conj}(z))=\operatorname{conj}(\operatorname{cexp}(z))$.
$-\operatorname{cexp}( \pm 0+\mathrm{i} 0)$ returns $1+i 0$.
$-\operatorname{cexp}(+\infty+i 0)$ returns $+\infty+i 0$.

- $\mathbf{c e x p}(-\infty+\mathrm{i} \infty)$ returns $\pm 0 \pm i 0$ (where the signs of the real and imaginary parts of the result are unspecified).
- cexp ( $+\infty+\mathrm{i} \infty$ ) returns $\pm \infty+i N a N$ and raises the invalid exception (where the sign of the real part of the result is unspecified).
$-\operatorname{cexp}(\mathbf{x}+\mathrm{i} \infty)$ returns $N a N+i N a N$ and raises the invalid exception, for finite $\mathbf{x}$.
$-\operatorname{cexp}(-\infty+\mathbf{i y})$ returns +0 *is $(\mathbf{y})$, for finite $\mathbf{y}$.
- cexp ( $+\infty+\mathrm{iy}$ ) returns $+\infty *$ cis $(\mathbf{y})$, for finite nonzero $\mathbf{y}$.
- cexp ( $-\infty+\mathbf{i N A N}$ ) returns $\pm 0 \pm i 0$ (where the signs of the real and imaginary parts of the result are unspecified).
- cexp ( $+\infty+$ iNAN) returns $\pm \infty+i N a N$ (where the sign of the real part of the result is unspecified).
- $\mathbf{c e x p}($ NAN+iO) returns $N a N+i 0$.
- cexp (NAN+iy) returns $N a N+i N a N$ and optionally raises the invalid exception, for all non-zero numbers $\mathbf{y}$.
IEC 559-compatible complex arithmetic
- cexp (x+iNAN) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{x}$.
- Cexp (NAN+iNAN) returns $N a N+i N a N$.


## G.5.9 The clog function

$-\operatorname{clog}(\operatorname{conj}(z))=\operatorname{conj}(c l o g(z))$.

- clog ( $-0+\mathrm{i} 0$ ) returns $-\infty+i \pi$ and raises the divide-by-zero exception.
- clog (+0+i0) returns $-\infty+i 0$ and raises the divide-by-zero exception.
$-\mathrm{clog}(-\infty+\mathrm{i} \infty)$ returns $+\infty+i 3 \pi / 4$.
- clog ( $+\infty+\mathrm{i} \infty$ ) returns $+\infty+i \pi / 4$.
- $\mathbf{c l o g}(\mathbf{x}+\mathrm{i} \infty)$ returns $+\infty+i \pi / 2$, for finite $\mathbf{x}$.
- clog ( $-\infty+\mathbf{i y}$ ) returns $+\infty+i \pi$, for finite positive-signed $\mathbf{y}$.
- clog ( $+\infty+\mathrm{i} \mathbf{y}$ ) returns $+\infty+i 0$, for finite positive-signed $\mathbf{y}$.
$-\mathrm{clog}( \pm \infty+\mathrm{iNAN})$ returns $+\infty+i N a N$.
- clog (NAN+i $\infty$ ) returns $+\infty+i N a N$.
- clog ( $\mathbf{x}+\mathbf{i N A N}$ ) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{x}$.
- clog (NAN+iy) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{y}$.
- clog(NAN+iNAN) returns $N a N+i N a N$.
G.5.10 The csqre function
$1-\operatorname{csqrt}(\operatorname{conj}(z))=\operatorname{conj}(\operatorname{csqrt}(z))$.
- csqrt ( $\pm 0+\mathrm{i} 0)$ returns $+0+i 0$.
- csqrt ( $-\infty+\mathrm{iy}$ ) returns $+0+i \infty$, for finite positive-signed $\mathbf{y}$.
- csqrt ( $+\infty+\mathrm{iy}$ ) returns $+\infty+i 0$, for finite positive-signed $\mathbf{y}$.
- csqrt ( $\mathbf{x}+\mathrm{i} \infty$ ) returns $+\infty+i \infty$, for all $\mathbf{x}$ (including NaN ).
- csqrt ( $-\infty+\mathrm{iNAN}$ ) returns $N a N \pm i \infty$ (where the sign of the imaginary part of the result is unspecified).
- csqrt ( $+\infty+$ iNAN) returns $+\infty+i N a N$.
- csqrt ( $\mathbf{x}+\mathbf{i N A N}$ ) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{x}$.
- csqrt (NAN+iy) returns $N a N+i N a N$ and optionally raises the invalid exception, for finite $\mathbf{y}$.
- csqrt (NAN+iNAN) returns $N a N+i N a N$.


## G. 6 <tgmath.h>

1 Type-generic macros that accept complex arguments also accept imaginary arguments. If an argument is imaginary, the macro expands to an expression whose type is real, imaginary, or complex, as appropriate for the particular function: if the argument is imaginary, then the types of cos, cosh, fabs, carg, cimag, and creal are real; the types of sin, tan, sinh, tanh, asin, atan, asinh, and atanh are imaginary; and the types of the others are complex.

2 Given an imaginary argument, each of the type-generic macros cos, sin, tan, cosh, sinh, tanh, asin, atan, asinh, atanh is specified by a formula in terms of real functions:

```
cos(i*y) = cosh(y)
sin(i*y) = i*sinh(y)
tan(i*y) = i*tanh(y)
cosh(i*y) = cos(y)
sinh(i*y) = i*sin(y)
tanh(i*y) = i*tan(y)
asin(i*y) = i*asinh(y)
atan(i*y) = i*atanh(y)
asinh(i*y) = i*asin(y)
atanh(i*y) = i*atan(y)
```


## Annex H <br> (informative) <br> Language independent arithmetic

## H. 1 Introduction

1 This annex documents the extent to which the C language supports the standard: ISO/IEC 10967-1 Language independent arithmetic, part 1, integer and floating-point arithmetic (LIA-1). LIA-1 is more general than IEC 559 (annex F) in that it covers integer and diverse floating-point arithmetics.

## H. 2 Types

1 The relevant C arithmetic types meet the requirements of LIA-1 types if an implementation adds notification of exceptional arithmetic operations and meets the 1-ULP accuracy requirement.

## H.2.1 Boolean type

1 The LIA-1 data type Boolean is implemented by the C data type bool with values of true and false, all from <stdbool.h>.

## H.2.2 Integer types

1 The signed C integer types int, long, long long and the corresponding unsigned types are compatible with LIA-1. If an implementation adds support for the LIA-1 exceptional values integer_overflow and undefined, then those types are LIA-1 conformant types. C's unsigned integer types are "modulo" in the LIA-1 sense in that overflows or out-of-bounds results silently wrap. An implementation that defines signed integer types as also being modulo need not detect integer overflow, in which case, only integer divide-by-zero need be detected.

2 The parameters for the integer data types can be accessed by the following:

```
maxint INT_MAX, LONG_MAX, LLONG_MAX, UINT_MAX, ULONG_MAX,
    ULLONG_MAX
minint INT_MIN, LONG_MIN, LLONG_MIN
```

3 The parameter "bounded" is always true, and is not provided. The parameter "minint" is always 0 for the unsigned types, and is not provided for those types.

## H.2.2.1 Integer operations

1 The integer operations on integer types are the following:

| addI | $\mathbf{x}+\mathbf{y}$ |
| :--- | :--- |
| subI | $\mathbf{x}-\mathbf{y}$ |
| mulI | $\mathbf{x} * \mathbf{y}$ |
| divI, divtI | $\mathbf{x} / \mathbf{y}$ |
| remI, remtI | $\mathbf{x} \% \mathbf{y}$ |
| negI | $-\mathbf{x}$ |
| absI | abs (x), labs (x), llabs (x) |
| eqI | $\mathbf{x}=\mathbf{y}$ |
| neqI | $\mathbf{x} \mathbf{!}=\mathbf{y}$ |
| lssI | $\mathbf{x}<\mathbf{y}$ |
| leqI | $\mathbf{x}<=\mathbf{y}$ |
| gtrI | $\mathbf{x}>\mathbf{y}$ |
| geqI | $\mathbf{x}>=\mathbf{y}$ |

where $\mathbf{x}$ and $\mathbf{y}$ are expressions of the same integer type.

## H.2.3 Floating-point types

1 The C floating-point types float, double, and long double are compatible with LIA-1. If an implementation adds support for the LIA-1 exceptional values underflow, floating_overflow, and undefined, then those types are conformant with LIA-1. An implementation that uses IEC 559 floating-point formats and operations (see Annex F) along with IEC 559 status flags and traps has LIA-1 conformant types.

## H.2.3.1 Floating-point parameters

1 The parameters for a floating point data type can be accessed by the following:

| $r$ | FLT_RADIX |
| :--- | :--- |
| $p$ | FLT_MANT_DIG, DBL_MANT_DIG, LDBL_MANT_DIG |
| emax | FLT_MAX_EXP, DBL_MAX_EXP, LDBL_MAX_EXP |
| emin | FLT_MIN_EXP, DBL_MIN_EXP, LDBL_MIN_EXP |

2 The derived constants for the floating point types are accessed by the following:

```
fmax FLT_MAX, DBL_MAX, LDBL_MAX
fminN FLT_MIN, DBL_MIN, LDBL_MIN
epsilon FLT_EPSILON, DBL_EPSILON, LDBL_EPSILON
rnd_style FLT_ROUNDS
```


## H.2.3.2 Floating-point operations

1 The floating-point operations on floating-point types are the following:

| addF | $\mathbf{x}+\mathbf{y}$ |
| :---: | :---: |
| subF | $\mathbf{x}-\mathbf{y}$ |
| mulF | x * $\mathbf{y}$ |
| $\operatorname{div} F$ | $\mathbf{x} / \mathbf{y}$ |
| negF | - x |
| absF | fabsf(x), fabs (x), fabsl (x) |
| exponent $F$ | 1.f+logbf (x), 1.0+logb (x), 1.L+logbl (x) |
| scaleF | ```scalbnf(x, n), scalbn(x, n), scalbnl (x, n), scalblnf(x, li), scalbln(x, li), scalblnl(x, li)``` |
| intpartF | $\operatorname{modff}(x, \& y), \operatorname{modf}(x, \& y), \operatorname{modfl}(x, \& y)$ |
| fractpartF | $\operatorname{modff}(\mathrm{x}, ~ \& y), \operatorname{modf}(\mathrm{x}, ~ \& y), \operatorname{modfl}(\mathrm{x}, \& y)$ |
| $e q F$ | $\mathrm{x}==\mathrm{y}$ |
| neqF | $\mathbf{x}!=y$ |
| $l s s F$ | $\mathbf{x}<\mathbf{y}$ |
| leqF | $x<=y$ |
| gtrF | $\mathbf{x}>\mathrm{y}$ |
| geqF | $\mathbf{x}>=\mathrm{y}$ |

where $\mathbf{x}$ and $\mathbf{y}$ are expressions of the same floating point type, $\mathbf{n}$ is of type int, and $\mathbf{l i}$ is of type long int.

## H.2.3.3 Rounding styles

1 The C Standard requires all floating types use the same radix and rounding style, so that only one identifier for each is provided to map to LIA-1.

2 The FLT_ROUNDS parameter can be used to indicate the LIA-1 rounding styles:

```
truncate FLT_ROUNDS == 0
nearest FLT_ROUNDS == 1
other FLT_ROUNDS != 0 && FLT_ROUNDS != 1
```

provided that an implementation extends FLT_ROUNDS to cover the rounding style used in all relevant LIA-1 operations, not just addition as in C.

## H.2.4 Type conversions

1 The LIA-1 type conversions are the following type casts:

```
cvtl'->I (int)i, (long)i, (long long)i, (unsigned int)i, (unsigned
                long)i, (unsigned long long)i
cvtF->I (int) x, (long) x, (long long) x, (unsigned int) x,
    (unsigned long) x, (unsigned long long) x
cvtI->F (float) i, (double) i, (long double) i
cvtF'->F (float) x, (double) x, (long double) x
```

2 In the above conversions from floating to integer, the use of (cast) $\mathbf{x}$ can be replaced with (cast) round ( $\mathbf{x}$ ), (cast) rint ( $\mathbf{x}$ ), (cast) nearbyint ( $\mathbf{x}$ ), (cast)trunc ( $\mathbf{x}$ ), (cast) ceil ( $\mathbf{x}$ ), or (cast) $\mathbf{f l o o r}(\mathbf{x})$ as they all meet LIA's requirements on floating to integer rounding. The remainder() function is useful for doing silent wrapping to unsigned integer types.

3 C's floating-point to integer conversion functions, lrint(), llrint(), lround (), and llround (), can meet LIA-1's requirements if an implementation handles out of range as per LIA-1.
4 C's conversions (type casts) from floating-point to floating-point can meet LIA-1 if an implementation uses round to nearest.

5 C's conversions (type casts) from integer to floating-point can meet LIA-1 if an implementation uses round to nearest.

## H. 3 Notification

1 Notification is the process by which a user or program is informed that an exceptional arithmetic operation has occurred. C's operations are compatible with LIA-1 in that C allows an implementation to cause a notification to occur when any arithmetic operation returns an exceptional value as defined in LIA-1 clause 5.

## H.3.1 Notification alternatives

1 LIA-1 requires at least the following two alternatives for handling of notifications: setting indicators or trap-and-terminate. LIA-1 allows a third alternative: trap-and-resume.

2 An implementation need only support a given notification alternative for the entire program. An implementation may support the ability to switch between notification alternatives during execution, but is not required to do so. An implementation can provide separate selection for each kind of notification, but this is not required.

3 C allows an implementation to provide notification. C's SIGFPE (for traps) and FE_INVALID, FE_DIVBYZERO, FE_OVERFLOW, FE_UNDERFLOW (for indicators) can provide LIA notification.

4 C's signal handlers are compatible with LIA-1. Default handling of SIGFPE can provide trap-and-terminate behavior. User provided signal handlers for SIGFPE allow for trap-and-resume behavior.

## H.3.1.1 Indicators

1 C's <fenv. h > status flags are compatible with LIA's indicators.
2 The following mapping is for floating-point types:

```
undefined
floating_overflow FE_OVERFLOW
underflow FE_UNDERFLOW
```

3 The floating-point indicator interrogation and manipulation operations are:

```
set_indicators feraiseexcept (i)
clear_indicators feclearexcept (i)
test_indicators fetestexcept(i)
current_indicators fetestexcept (FE_ALL_EXCEPT)
```

where $\mathbf{i}$ is an expression of type int representing a LIA-1 indicator subset.
4 C allows an implementation to provide the following LIA-1 required behavior: at program termination if any indicator is set the implementation shall send an unambiguous and "hard
to ignore" message (see LIA-1 subclause 6.1.2)
5 LIA-1 does not make the distinction between floating-point and integer for undefined. This documentation is making that distinction because <fenv.h> covers only the floating-point indicators.

## H.3.1.2 Traps

1 C is compatible with LIA's trap requirements. An implementation can provide an alternative of notification through termination with a 'hard-to-ignore'" message (see LIA-1 subclause 6.1.3).

2 LIA-1 does not require that traps be precise.
3 C does require that SIGFPE be the signal corresponding to arithmetic exceptions, if there is any signal raised for them.

4 C has signal handlers for SIGFPE and allows trapping arithmetic exceptions. When arithmetic exceptions do trap, C's signal-handler mechanism allows trap-and-terminate (either default implementation behavior or user replacement for it) and trap-and-resume, at the programmer's option.

## Annex I (normative) <br> Universal character names for identifiers

1 This clause lists the hexadecimal code values that are valid in universal character names in identifiers.

2 This table is reproduced unchanged from ISO/IEC PDTR 10176, producted by ISO/IEC JTC1/SC22/WG20, except that the range 0041-005a and 0061-007a designate the upper and lower case Latin alphabets, which are part of the basic source character set, and are not repeated in the table below.

- Latin: 00c0-00d6, 00d8-00f6, 00f8-01f5, 01fa-0217, 0250-02a8, 1e00-1e9a, 1ea0-1ef9
- Greek: 0384, 0388-038a, 038c, 038e-03a1, 03a3-03ce, 03d0-03d6, 03da, 03dc, 03de, 03 e 0 , 03e2-03f3, $1 \mathrm{f} 00-1 \mathrm{f} 15$, $1 \mathrm{f} 18-1 \mathrm{f} 1 \mathrm{~d}$, $1 \mathrm{f} 20-1 \mathrm{f} 45$, $1 \mathrm{f} 48-1 \mathrm{f} 4 \mathrm{~d}, 1 \mathrm{f} 50-1 \mathrm{f} 57$, 1 f 59 , 1 f 5 b , $1 f 5 d, 1 f 5 f-1 f 7 d, 1 f 80-1 \mathrm{fb} 4,1 \mathrm{fb} 6-1 \mathrm{fbc}, 1 \mathrm{fc} 2-1 \mathrm{fc} 4$, $1 \mathrm{fc} 6-1 \mathrm{fcc}, 1 \mathrm{fd} 0-1 \mathrm{fd} 3$, $1 \mathrm{fd} 6-1 \mathrm{fdb}, 1 \mathrm{fe} 0-$ 1fec, 1ff2-1ff4, 1ff6-1ffc,
- Cyrilic: 0401-040d, 040f-044f, 0451-045c, 045e-0481, 0490-04c4, 04c7-04c8, 04cb04cc, 04d0-04eb, 04ee-04f5, 04f8-04f9
- Armenian: 0531-0556, 0561-0587
- Hebrew: 05d0-05ea, 05f0-05f4
- Arabic: 0621-063a, 0640-0652, 0670-06b7, 06ba-06be, 06c0-06ce, 06e5-06e7,
- Devanagari: 0905-0939, 0958-0962
- Bengali: 0985-098c, 098f-0990, 0993-09a8, 09aa-09b0, 09b2, 09b6-09b9, 09dc-09dd, 09df-09e1, 09f0-09f1
- Gurmukhi: 0a05-0a0a, 0a0f-0a10, 0a13-0a28, 0a2a-0a30, 0a32-0a33, 0a35-0a36, 0a380a39, 0a59-0a5c, 0a5e
- Gujarati: 0a85-0a8b, 0a8d, 0a8f-0a91, 0a93-0aa8, 0aaa-0ab0, 0ab2-0ab3, 0ab5-0ab9, 0ae0,
- Oriya: 0b05-0b0c, 0b0f-0b10, 0b13-0b28, 0b2a-0b30, 0b32-0b33, 0b36-0b39, 0b5c0b5d, 0b5f-0b61,
- Tamil: 0b85-0b8a, 0b8e-0b90, 0b92-0b95, 0b99-0b9a, 0b9c, 0b9e-0b9f, 0ba3-0ba4, 0ba8-0baa, 0bae-0bb5, 0bb7-0bb9,
— Telugu: 0c05-0c0c, 0c0e-0c10, 0c12-0c28, 0c2a-0c33, 0c35-0c39, 0c60-0c61,
- Kannada: 0c85-0c8c, 0c8e-0c90, 0c92-0ca8, 0caa-0cb3, 0cb5-0cb9, 0ce0-0ce1,
— Malayalam: 0d05-0d0c, 0d0e-0d10, 0d12-0d28, 0d2a-0d39, 0d60-0d61,
- Thai: 0e01-0e30, 0e32-0e33, 0e40-0e46, 0e4f-0e5b,
- Lao: 0e81-0e82, 0e84, 0e87, 0e88, 0e8a, 0e0d, 0e94-0e97, 0e99-0e9f, 0ea1-0ea3, 0ea5, 0ea7, 0eaa, 0eab, 0ead-0eb0, 0eb2, 0eb3, 0ebd, 0ec0-0ec4, 0ec6,
- Georgian: 10a0-10c5, 10d0-10f6,
- Hiragana: 3041-3094, 309b-309e
- Katakana: 30a1-30fe,
- Bopmofo: 3105-312c,
- Hangul: 1100-1159, 1161-11a2, 11a8-11f9, ac00-d7af
- CJK Unified Ideographs: f900-fa2d, fb1f-fb36, fb38-fb3c, fb3e, fb40-fb41, fb42-fb44, fb46-fbb1, fbd3-fd3f, fd50-fd8f, fd92-fdc7, fdf0-fdfb, fe70-fe72, fe74, 5e76-fefc, ff21ff3a, ff41-ff5a, ff66-ffbe, ffc2-ffc7, ffca-ffcf, ffd2-ffd7, ffda-ffdc, 4e00-9fa5


## Annex J (informative) <br> Common warnings

1 An implementation may generate warnings in many situations, none of which is specified as part of this International Standard. The following are a few of the more common situations.

- A block with initialization of an object that has automatic storage duration is jumped into (6.1.2.4).
- An integer character constant includes more than one character or a wide character constant includes more than one multibyte character (6.1.3.4).
- The characters /* are found in a comment (6.1.7).
- An implicit narrowing conversion is encountered, such as the assignment of a long int or a double to an int, or a pointer to void to a pointer to any type other than a character type (6.2).
- An 'unordered'" binary operator (not comma, $\& \&$ or $\|$ ) contains a side-effect to an lvalue in one operand, and a side-effect to, or an access to the value of, the identical lvalue in the other operand (6.3).
- A function is called but no prototype has been supplied (6.3.2.2).
- The arguments in a function call do not agree in number and type with those of the parameters in a function definition that is not a prototype (6.3.2.2).
- An object is defined but not used (6.5).
- A value is given to an object of an enumeration type other than by assignment of an enumeration constant that is a member of that type, or an enumeration variable that has the same type, or the value of a function that returns the same enumeration type (6.5.2.2).
- An aggregate has a partly bracketed initialization (6.5.7).
- A statement cannot be reached (6.6).
- A statement with no apparent effect is encountered (6.6).
- A constant expression is used as the controlling expression of a selection statement (6.6.4).
- A function has return statements with and without expressions (6.6.6.4).
- An incorrectly formed preprocessing group is encountered while skipping a preprocessing group (6.8.1).
- An unrecognized \#pragma directive is encountered (6.8.6).


## Annex K (informative) <br> Portability issues

1 This annex collects some information about portability that appears in this International Standard.

## K. 1 Unspecified behavior

1 The following are unspecified:

- The termination status returned to the hosted environment if the \} that terminates the main function is reached (5.1.2.2.3).
- The manner and timing of static initialization (5.1.2).
- The behavior if a printable character is written when the active position is at the final position of a line (5.2.2).
- The behavior if a backspace character is written when the active position is at the initial position of a line (5.2.2).
- The behavior if a horizontal tab character is written when the active position is at or past the last defined horizontal tabulation position (5.2.2).
- The behavior if a vertical tab character is written when the active position is at or past the last defined vertical tabulation position (5.2.2).
- Many aspects of the representations of types (6.1.2.8).
- The value of padding bytes when storing values in structures or unions (6.1.2.8.1).
- The value of a union member other than the last one stored into (6.1.2.8.1).
- The representation used when storing a value in an object that has more than one object representation for that value (6.1.2.8.1).
- The value of padding bits in integer representations (6.1.2.8.2).
- The order in which subexpressions are evaluated and the order in which side effects take place, except as indicated by the syntax or otherwise specified for the function-call (), $\& \&, \|, ?:$, and comma operators (6.3).
- The order in which the function designator, arguments, and subexpressions within the arguments are evaluated in a function call (6.3.2.2).
- The order of side effects among compound literal initialization list expressions (6.3.2.5).
- The order in which the operands of an assignment operator are evaluated (6.3.16).
- The alignment of the addressable storage unit allocated to hold a bit-field (6.5.2.1).
- The choice of using an inline definition or external definition of a function when both definitions are in scope (6.5.4).
- The size of an array when * is written instead of a size expression (6.5.5.2).
- Whether side effects are produced when evaluating the size expression in a declaration of an array of variable length (6.5.5.2).
- The layout of storage for function parameters (6.7.1).
- The order in which \# and \#\# operations are evaluated during macro substitution (6.8.3.2, 6.8.3.3).
- Whether errno is a macro or an external identifier (7.1.4).
- The order of raising floating-point exceptions, except as stated in F.7.6 (7.6.2.3).
— The result of rounding when the value is out of range (7.7.9.5, 7.7.9.8, F.9.6.5).
- The result of the nan function when its argument does not point to an $n$-char-sequence string, or when the implementation does not support quiet NaNs for the double type (7.7.11.2).
- Whether set jmp is a macro or an external identifier (7.10).
- Whether va_end is a macro or an external identifier (7.12.1).
- The hexadecimal digit left of the decimal point when a non-normalized floating-point number is printed with an a or $\mathbf{A}$ conversion specifier (7.13.6.1, 7.19.2.1).
- The value of the file position indicator after a successful call to the ungetc function for a text stream, or the ungetwc function for any stream, until all pushed-back characters are read or discarded (7.13.7.11, 7.19.3.9).
- The details of the value stored by the fgetpos function (7.13.9.1).
- The details of the value returned by the ftell function for a text stream (7.13.9.4).
- The order and contiguity of storage allocated by successive calls to the calloc, malloc, and realloc functions (7.14.3).
- The amount of storage allocated by a successful call to the calloc, malloc, or realloc function when 0 bytes was requested (7.14.3).
- Which of two elements that compare as equal is matched by the bsearch function (7.14.5.1).
- The order of two elements that compare as equal in an array sorted by the qsort function (7.14.5.2).
- The value of the tm_extlen member of the $\operatorname{tmx}$ structure to which a pointer is returned by the zonetime function when it has set the tm_ext member to a null pointer (7.16.1).
- The encoding of the calendar time returned by the time function (7.16.2.5).
- The resulting value when the invalid exception is raised during IEC 559 floating to integer conversion (F.4).
- Whether conversion of non-integer IEC 559 floating values raises the inexact exception (F.4).
- The value stored by frexp for a NaN or infinity (F.9.3.4).
- The sign of one part of the complex result of several math functions for certain exceptional values in IEC-559 compatible implementations (G.5.1, G.5.3, G.5.4, G.5.5, G.5.6, G.5.7, G.5.8, G.5.10).


## K. 2 Undefined behavior

1 The behavior is undefined in the following circumstances:

- A nonempty source file does not end in a new-line character, ends in new-line character immediately preceded by a backslash character, or ends in a partial preprocessing token or comment (5.1.1.2).
- Token concatenation produces a character sequence matching the syntax of a universal character name (5.1.1.2).
- A universal character name specifies a character identifier in the range 0000 through 0020 or 007 F through 009 F , or designates a character in the basic source character set (5.1.1.2).
- A program in a hosted environment does not define a function named main using one of the two specified forms (5.1.2.2.1).
- A character not in the required basic source character set is encountered in a source file, except in a character constant, a string literal, a header name, a comment, or a preprocessing token that is never converted to a token (5.2.1).
- A comment, string literal, character constant, or header name contains an invalid multibyte character or does not begin and end in the initial shift state (5.2.1.2).
- An unmatched ' or " character is encountered on a logical source line during tokenization (6.1).
- A reserved keyword token is used in translation phase 7 or 8 for some purpose other than as a keyword (6.1.1).
- The reserved token complex or imaginary is used before <complex.h> is included (6.1.1).
- A universal character name does not designate a code value in one of the specified ranges (6.1.2).
- The first character of an identifier is a digit (6.1.2).
- Two identifiers differ only in nonsignificant characters (6.1.2).
- The same identifier is used more than once as a label in the same function (6.1.2.1).
- The same identifier has both internal and external linkage in the same translation unit (6.1.2.2).
- A block containing a variably modified object having automatic storage duration is entered by a jump to a labeled statement (6.1.2.4).
- The value of a pointer that referred to an object with automatic storage duration is used after the storage is no longer guaranteed to be reserved (6.1.2.4).
- Two declarations of the same object or function specify types that are not compatible (6.1.2.6).
- A trap representation is accessed by an lvalue expression that does not have character type (6.1.2.8.1).
- A trap representation is produced by a side effect that modifies any part of the object by an lvalue expression that does not have character type. (6.1.2.8.1).
- The whole-number and fraction parts of a floating constant are both omitted (6.1.3.1).
- The period and the exponent part of a decimal floating constant are both omitted (6.1.3.1).
- An unspecified escape sequence is encountered in a character constant or a string literal (6.1.3.4).
- An attempt is made to modify a string literal of either form (6.1.4).
- The characters ', <br>, ", //, or /* are encountered between the $<$ and $>$ delimiters, or the characters ', <br>, //, or /* are encountered between the " delimiters, in a header name preprocessing token (6.1.7).
- Conversion to or from an integer type produces a value outside the range that can be represented (6.2.1.3).
- Conversion between two pointer types produces a result that is incorrectly aligned (6.2.1.3).
- Demotion of one real floating type to another produces a value outside the range that can be represented (6.2.1.4).
- A non-array lvalue with an incomplete type is used in a context that requires the value of the designated object (6.2.2.1).
- An lvalue having array type is converted to a pointer to the initial element of the array, and the array object has register storage class (6.2.2.1).
- An attempt is made to use the value of a void expression, or an implicit or explicit conversion (except to void) is applied to a void expression (6.2.2.2).
- Conversion of a pointer to an integer type produces a value outside the range that can be represented (6.2.2.3).
- Conversion between two pointer types produces a result that is incorrectly aligned (6.2.2.3).
- A pointer to a function is converted to point to a function of a different type and used to call a function of a type not compatible with the type of the called function (6.2.2.3).
- Between two sequence points, an object is modified more than once, or is modified and the prior value is accessed other than to determine the value to be stored (6.3).
- Non-integer operands are used with a bitwise operator (6.3).
- An exception occurs during the evaluation of an expression (6.3).
- An object has its stored value accessed other than by an lvalue expression having one of the following types: a type compatible with the effective type of the object, a qualified version of a type compatible with the effective type of the object, a type that is the signed or unsigned type corresponding to the effective type of the object, a type that is the signed or unsigned type corresponding to a qualified version of the effective type of the object, an aggregate or union type that includes one of the aforementioned types among its members (including, recursively, a member of a subaggregate or contained union), or a character type (6.3).
- For a function call without a function prototype, the number of arguments does not agree with the number of parameters (6.3.2.2).
- For a function call without a function prototype, the function is defined without a function prototype, and the types of the arguments after promotion are not compatible with those of the parameters after promotion (6.3.2.2).
- For a function call without a function prototype, the function is defined with a function prototype, and the types of the arguments after promotion are not compatible with the
types of the parameters, or the prototype ends with an ellipsis (6.3.2.2).
- A function is defined with a type that is not compatible with the type pointed to by the expression that denotes the called function (6.3.2.2).
— The operand of the unary * operator has an invalid value (6.3.3.2).
- A pointer is converted to other than an integer or pointer type (6.3.4).
- The value of the second operand of the / or \% operator is zero (6.3.5).
- Addition or subtraction of a pointer into, or just beyond, an array object and an integer type produces a result that does not point into, or just beyond, the same array object (6.3.6).
- Addition or subtraction of a pointer into, or just beyond, an array object and an integer type produces a result that does not point into the same array object, and the result is used as an operand of a unary * operator that is actually evaluated. (6.3.6).
- An array subscript is out of range, even if an object is apparently accessible with the given subscript (as in the lvalue expression a[1][7] given the declaration int a[4][5]) (6.3.6).
- The result of subtracting two pointers is not representable in an object of type ptrdiff_t (6.3.6).
- Pointers that do not point into, or just beyond, the same array object are subtracted (6.3.6).
- An expression is shifted by a negative number or by an amount greater than or equal to the width of the promoted expression (6.3.7).
- An expression having signed promoted type is left-shifted and either the value of the expression is negative or the result of shifting would be not be representable in the promoted type. (6.3.7).
- Pointers are compared that do not point to the same aggregate or union (nor just beyond the same array object) (6.3.8).
- An object is assigned to an inexactly overlapping object or to an exactly overlapping object with incompatible type (6.3.16.1).
- An expression that is required to be an integer constant expression does not have an integer type, contains casts (outside operands to sizeof operators) other than conversions of arithmetic types to integer types, or has operands that are not integer constants, enumeration constants, character constants, fixed-length sizeof expressions, or immediately-cast floating constants (6.4).
- A constant expression in an initializer does not evaluate to one of the following: an arithmetic constant expression, a null pointer constant, an address constant, or an address constant for an object type plus or minus an integer constant expression (6.4).
- An arithmetic constant expression does not have arithmetic type, contains casts (outside operands to sizeof operators) other than conversions of arithmetic types to arithmetic types, or has operands that are not integer constants, floating constants, enumeration constants, character constants, or sizeof expressions (6.4).
- An address constant is created neither explicitly using the unary $\&$ operator or an integer constant cast to pointer type, nor implicitly by the use of an expression of array or function type (6.4).
- The value of an object is accessed by an array-subscript [], member-access . or $\rightarrow$, address $\&$, or indirection * operator or a pointer cast in creating an address constant (6.4).
- An identifier for an object is declared with no linkage and the type of the object is incomplete after its declarator, or after its init-declarator if it has an initializer (6.5).
- A function is declared at block scope with an explicit storage-class specifier other than extern (6.5.1).
- A structure or union is defined as containing no named members (6.5.2.1).
- A bit-field is declared with a type other than a qualified or unqualified version of signed int or unsigned int (6.5.2.1).
- An attempt is made to access a flexible array member of a structure when the the referenced object provides no elements for that array (6.5.2.1).
- A tag is declared with the bracketed list twice within the same scope (6.5.2.3).
- When the complete type is needed, an incomplete structure or union type is not completed in the same scope by another declaration of the tag that defines the content (6.5.2.3).
- An attempt is made to modify an object defined with a const-qualified type through use of an lvalue with non-const-qualified type (6.5.3).
- An attempt is made to refer to an object defined with a volatile-qualified type through use of an lvalue with non-volatile-qualified type (6.5.3).
- An attempt is made to access an object through two different restrict-qualified pointers (6.5.3, 6.5.3.1).
- The specification of a function type includes a type qualifier (6.5.3).
- Two qualified types that are required to be compatible do not have the identically qualified version of a compatible type (6.5.3).
- A function with external linkage is declared with an inline function specifier, but is not also defined in the same translation unit (6.5.4).
- A storage-class specifier or type qualifier modifies the keyword void as a function parameter type list (6.5.4.3).
- Two pointer types that are required to be compatible are not identically qualified, or are not pointers to compatible types (6.5.5.1).
- The size expression in an array declaration is not a constant expression and evaluates at program execution time to a nonpositive value (6.5.5.2).
- In a context requiring two array types to be compatible, they do not have compatible element types, or their size specifiers evaluate to unequal values (6.5.5.2).
- In a context requiring two function types to be compatible, they do not have compatible return types, or their parameters disagree in use of the ellipsis terminator or the number and type of parameters (after default argument promotion, when there is no parameter type list or when one type is specified by a function definition with identifier list) (6.5.5.3).
- An identifier used as a typedef name is redeclared in an inner scope, or is declared as a member of a structure or union in the same or an inner scope, and the type specifiers are omitted in this (re)declaration (6.5.7).
- The value of an unnamed member of a structure or union is used (6.5.8).
- The value of an uninitialized object that has automatic storage duration, or of an uninitialized variable-length array that has block scope, is used before a value is assigned (6.5.8, 6.6.2).
- The initializer for a scalar is neither a single expression nor a single expression enclosed in braces (6.5.8).
- The initializer for a structure or union object is neither an initializer list nor a single expression that has compatible structure or union type (6.5.8).
- The initializer for an aggregate, other than an array initialized by a character string literal or wide string literal, is not a brace-enclosed list of initializers for its members, or the initializer for a union object is not a brace-enclosed initializer for its first member (6.5.8).
- The \} that terminates a function is reached, and the value of the function call is used by the caller (6.6.6.4).
- An identifier with external linkage is used, but in the program there does not exist exactly one external definition for the identifier, or the identifier is not used and there exist multiple external definitions for the identifier. (6.7).
- A function definition includes an identifier list, but the types of the parameters are not declared in a following declaration list. (6.7.1).
- A function that accepts a variable number of arguments is defined without a parameter type list that ends with the ellipsis notation (6.7.1).
- An adjusted parameter type in a function definition is not an object type (6.7.1).
- An identifier for an object with internal linkage and an incomplete type is declared with a tentative definition (6.7.2).
- The token defined is generated during the expansion of a \#if or \#elif preprocessing directive, or the use of the defined unary operator does not match one of the two specified forms prior to macro replacement (6.8.1).
- The \#include preprocessing directive that results after expansion does not match one of the two header name forms (6.8.2).
- The character sequence in an \#include preprocessing directive does not start with a letter (6.8.2).
- There are sequences of preprocessing tokens within the list of macro arguments that would otherwise act as preprocessing directive lines (6.8.3).
- A fully expanded macro replacement list contains a function-like macro name as its last preprocessing token (6.8.3).
- The result of the preprocessing operator \# is not a valid character string literal (6.8.3.2).
- The result of the preprocessing operator \#\# is not a valid preprocessing token (6.8.3.3).
- The \#line preprocessing directive that results after expansion does not match one of the two well-defined forms, or its digit sequence specifies zero or a number greater than 2147483647 (6.8.4).
- A non-STDC \#pragma preprocessing directive that is documented as causing translation failure or some other form of undefined behavior is encountered (6.8.6).
- A \#pragma STDC preprocessing directive does not match one of the nine well-defined forms (6.8.6).
- One of the following identifiers is the subject of a \#define or \#undef preprocessing directive: __LINE__, __FILE__, __DATE__, __TIME__, __STDC__, STDC_VERSION__, __STDC_IEC_559__, __STDC_IEC_559_COMPLEX__, or defined (6.8.8).
- An attempt is made to copy an object to an overlapping object by use of a library function other than memmove (7).
- A file with the same name as one of the standard headers, not provided as part of the implementation, is placed in any of the standard places for a source file to be included (7.1.2).
- A header is included within an external declaration or definition (7.1.2).
- A function, object, type, or macro that is specified as being declared or defined by some standard header is used before any header that declares or defines it is included (7.1.2).
- A standard header is included while a macro is defined with the same name as a keyword (7.1.2).
- The program attempts to declare a library function itself, rather than via a standard header, but the declaration does not have external linkage (7.1.2).
- The program declares or defines a reserved identifier (other than as allowed by 7.1.8) (7.1.3).
- The program removes the definition of a macro whose name begins with an underscore and either an uppercase letter or another underscore (7.1.3).
- A macro definition of errno is suppressed in order to access an actual object, or the program defines an identifier with the name errno (7.1.4).
- The member-designator parameter of an offsetof macro is an invalid right operand of the . operator for the type parameter, or the member-designator parameter designates a bit-field (7.1.6).
- An argument to a library function has an invalid value or a type not expected by a function with variable number of arguments (7.1.8).
- The macro definition of assert is suppressed in order to access an actual function (7.2).
- The value of an argument to a character handling function is neither equal to the value of EOF nor representable as an unsigned char (7.3).
- The argument in an instance of one of the integer-constant macros is not a decimal, octal, or hexadecimal constant, or it has a value that exceeds the limits for the corresponding type (7.4.3).
- The program modifies the string pointed to by the value returned by the setlocale function (7.5.1.1).
- The program modifies the structure pointed to by the value returned by the localeconv function (7.5.2.1).
- The FENV_ACCESS pragma is used in any context other than outside all external declarations or preceding all explicit declarations and statements inside a compound statement (7.6.1).
- Part of the program tests flags or runs under non-default mode settings, but was translated with the state for the FENV_ACCESS pragma off (7.6.1).
- The exception-mask argument for one of the functions that provide access to the exception flags has a value not obtained by bitwise OR of the exception macros (7.6.2).
- The fesetexceptflag function is used to set the status for exception flags not specified in the call to the fegetexceptflag function that provided the value of the corresponding fexcept_t object (7.6.2.4).
- The FP_CONTRACT pragma is used in any context other than outside all external declarations or preceding all explicit declarations and statements inside a compound statement (7.7.2).
- An argument to a floating-point classification macro is not of real floating type (7.7.3).
- The CX_LIMITED_RANGE pragma is used in any context other than outside all external declarations or preceding all explicit declarations and statements inside a compound statement (7.8.1).
- A non-real argument is supplied for a generic parameter of a type-generic macro (7.9.1).
- A macro definition of set jmp is suppressed in order to access an actual function, or the program defines an identifier with the name set jmp (7.10.1).
- An invocation of the setjmp macro occurs in a context other than as the entire controlling expression in a selection or iteration statement, or in a comparison with an integer constant expression (possibly as implied by the unary ! operator) as the entire controlling expression of a selection or iteration statement, or as the entire expression of an expression statement (possibly cast to void) (7.10.2.1).
- The longjmp function is invoked to restore a nonexistent environment (7.10.2.1).
- After a longjmp, there is an attempt to access the value of an object of automatic storage class with non-volatile-qualified type, local to the function containing the invocation of the corresponding set jmp macro, that was changed between the setjmp invocation and longjmp call (7.10.2.1).
- The longjmp function is invoked from a nested signal handler (7.10.2.1).
- The program uses a nonpositive value for a signal number (7.11).
- The program specifies an invalid pointer to a signal handler function (7.11.1.1).
- A signal handler returns when the signal corresponded to a computational exception (7.11.1.1).
- A signal occurs other than as the result of calling the abort or raise function, and the signal handler calls a function in the standard library other than the signal function (for the same signal number) or refers to an object with static storage duration other than by assigning a value to an object declared as volatile sig_atomic_t (7.11.1.1).
- The value of errno is referred to after a signal occurred other than as the result of calling the abort or raise function and the corresponding signal handler obtained a SIG_ERR return from a call to the signal function (7.11.1.1).
- A function with a variable number of arguments attempts to access its varying arguments other than through a properly declared and initialized va_list object, or before the va_start macro is invoked (7.12, 7.12.1.1, 7.12.1.2).
- The macro va_arg is invoked using the parameter ap that was passed to a function that invoked the macro va_arg with the same parameter (7.12).
- A macro definition of va_start, va_arg, va_copy, or va_end is suppressed in order to access an actual function, or the program defines an external identifier with the name va_end (7.12.1).
- The va_end macro is invoked without a corresponding invocation of the va_start or va_copy macro, or vice versa. (7.12.1, 7.12.1.1, 7.12.1.4).
- The parameter parmN of a va_start macro is declared with the register storage class, with a function or array type, or with a type that is not compatible with the type that results after application of the default argument promotions (7.12.1.1).
- The va_arg macro is invoked when there is no actual next argument, or with a specified type that is not compatible with the promoted type of the actual next argument (7.12.1.2).
- A byte input/output function is applied to a wide-oriented stream, or a wide-character input/output function is applied to a byte-oriented stream (7.13.2).
- Use if made of any portion of a file beyond the most recent wide character written to a wide-oriented stream (7.13.2).
- The value of a pointer to a FILE object is used after the associated file is closed (7.13.3).
- The stream for the $\mathbf{f f l u s h}$ function points to an input stream or to an update stream in which the most recent operation was input (7.13.5.2).
- An output operation on an update stream is followed by an input operation without an intervening call to the fflush function or a file positioning function, or an input operation on an update stream is followed by an output operation with an intervening call to a file positioning function (7.13.5.3).
- The string pointed to by the mode argument in a call to the fopen function does not exactly match one of the specified character sequences (7.13.5.3).
- An attempt is made to use the contents of the array that was supplied in a call to the setvbuf function (7.13.5.6).
- There are insufficient arguments for the format in a call to the fprintf, fscanf, fwprintf, or fwscanf function, or an argument does not have an appropriate type (7.13.6.1, 7.13.6.2, 7.19.2.1, 7.19.2.2).
- The format in a call to the fprintf, fscanf, fwprintf, fwscanf, or strftime function is not a valid multibyte character sequence that begins and ends in its initial shift state (7.13.6.1, 7.13.6.2, 7.19.2.1, 7.19.2.2, 7.16.3.6).
- In a call to the fprintf or fwprintf function, a precision appears with a conversion specifier other than $\mathbf{a}, \mathbf{A}, \mathbf{d}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}, \mathbf{G}, \mathbf{i}, \mathbf{o}, \mathbf{s}, \mathbf{u}, \mathbf{x}$, or $\mathbf{x}$ (7.13.6.1, 7.19.2.1).
- A conversion specification for the fprintf, fscanf, fwprintf, or fwscanf function uses $\mathbf{h h}, \mathbf{h}, \mathbf{l}, \mathbf{l l}$, or $\mathbf{L}$ with a conversion specifier in a combination not specified in this International Standard (7.13.6.1, 7.13.6.2, 7.19.2.1, 7.19.2.2).
- An asterisk is used to denote an argument-supplied field width or precision, but the corresponding argument is not provided (7.13.6.1, 7.19.2.1).
- A conversion specification for the fprintf or fwprintf function uses a \# flag with a conversion specifier other than $\mathbf{a}, \mathbf{A}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}, \mathbf{G}, \mathbf{o}, \mathbf{x}$, or $\mathbf{x}(7.13 .6 .1,7.19 .2$.1).
- A conversion specification for the fprintf or fwprintf function uses a $\mathbf{0}$ flag with a conversion specifier other than $\mathbf{a}, \mathbf{A}, \mathbf{d}, \mathbf{e}, \mathbf{E}, \mathbf{f}, \mathbf{F}, \mathbf{g}, \mathbf{G}, \mathbf{i}, \mathbf{o}, \mathbf{u}, \mathbf{x}$, or $\mathbf{x}$ (7.13.6.1, 7.19.2.1).
- An s conversion specifier is encountered by the fprintf or fwprintf function, and the argument is missing the null terminator (unless a precision is specified that does not require null termination) (7.13.6.1, 7.19.2.1).
- An $n$ conversion specifier is encountered by the fprintf, fscanf, fwprintf, or fwscanf function, but the complete conversion specification is not exactly $\% \mathrm{hhn}$, \%hn, \%n, \% ln, or \% lln (7.13.6.1, 7.13.6.2, 7.19.2.1, 7.19.2.2).
- A \% conversion specifier is encountered by the fprintf, fscanf, fwprintf, or fwscanf function, but the complete conversion specification is not exactly $\% \%$ (7.13.6.1, 7.13.6.2, 7.19.2.1, 7.19.2.2).
- An invalid conversion specification is found in the format for the fprintf, fscanf, strftime, fwprintf, or fwscanf function (7.13.6.1, 7.16.3.6, 7.13.6.2, 7.19.2.1, 7.19.2.2).
- An argument to the fprintf or fwprintf function is, or points to, a union or an aggregate other than a pointer using \%p conversion, or an array of character type using $\% \mathbf{s}$ conversion, or an array of wchar_t type using \%ls conversion (7.13.6.1, 7.19.2.1).
- The result of a conversion by the fscanf or fwscanf function cannot be represented in the corresponding object, or the receiving object does not have an appropriate type (7.13.6.2, 7.19.2.2).
- A c, s, or [ conversion specifier is encountered by the fscanf or fwscanf function, and the character array pointed to by the corresponding argument is not large enough to accept the input sequence (and a null terminator if the conversion specifier is $\mathbf{s}$ or [) (7.13.6.2, 7.19.2.2).
- Anc, s, or [ conversion specifier with an 1 qualifier is encountered by the fscanf or fwscanf function, but the input is not a valid multibyte character sequence that begins in the initial shift state (7.13.6.2, 7.19.2.2).
- The input item for \%p conversion by the fscanf or fwscanf function is not a value converted earlier during the same program execution (7.13.6.2, 7.19.2.2).
- The sprintf, snprintf, sscanf, vsprintf, vsnprintf, mbstowcs, wcstombs, memcpy, strcpy, strncpy, strcat, strncat, strxfrm, or strftime function, or any of the functions declared by <wchar.h> (except where otherwise specified), is used to copy between overlapping objects (7.13.6.5, 7.13.6.6, 7.13.6.7, 7.13.6.10, 7.13.6.11, 7.14.8.1, 7.14.8.2, 7.15.2.1, 7.15.2.3, 7.15.2.4, 7.15.3.1, 7.15.3.2, 7.15.4.5, 7.16.3.6, 7.19).
- The vfprintf, vprintf, vsprintf, vsnprintf, vfscanf, vscanf, vsscanf, vfwprintf, vwprintf, vswprintf, vfwscanf, vwscanf, or vswscanf function is called with an improperly initialized va_list argument (7.13.6.8, 7.13.6.9, 7.13.6.10, 7.13.6.11, 7.13.6.12, 7.13.6.13, 7.13.6.14, 7.19.2.7, $7.19 .2 .8,7.19 .2 .9,7.19 .2 .10,7.19 .2 .11,7.19 .2 .12)$.
- The contents of the array supplied in a call to the fgets, gets, or fgetws function are used after a read error occurred (7.13.7.2, 7.13.7.7, 7.19.3.2).
- The file position indicator for a binary stream is used after a call to the ungetc function where its value was zero before the call (7.13.7.11).
- The file position indicator for a stream is used after an error occurred during a call to the fread or fwrite function (7.13.8.1, 7.13.8.2).
- A partial element read by a call to the fread function is used (7.13.8.1).
- The fseek function is called for a text stream with other than SEEK_SET, or with a non-zero offset that was not returned by a previous successful call to the ftell function for the same file (7.13.9.2).
- The fsetpos function is called to set a position that was not returned by a previous successful call to the fgetpos function for the same file (7.13.9.3).
- The value of the result of converting a string to a number by the atof, atoi, atol, or atoll function cannot be represented (7.14.1).
- A non-null pointer returned by a call to the calloc, malloc, or realloc function with a zero requested size is used to access an object (7.14.3).
- The value of a pointer that refers to space deallocated by a call to the free or realloc function is used (7.14.3).
- The pointer argument to the free or realloc function does not match a pointer earlier returned by calloc, malloc, or realloc, or the space has been deallocated by a call to free or realloc (7.14.3.2, 7.14.3.4).
- The value of the object allocated by the malloc function is used (7.14.3.3).
- The value of the newly allocated portion of an object expanded by the realloc function is used (7.14.3.4).
- The program executes more than one call to the exit function (7.14.4.3).
- The string set up by the getenv or strerror function is modified by the program (7.14.4.4, 7.15.6.2).
- A command is executed through the system function in a way that is documented as causing termination or some other form of undefined behavior (7.14.4.5).
- The comparison function called by the bsearch or qsort function returns ordering values inconsistently (7.14.5.1, 7.14.5.2).
- The array being searched by the bsearch function does not have its elements in proper order (7.14.5.1).
- The result of an integer arithmetic function (abs, div, labs, llabs, ldiv, or lldiv) cannot be represented (7.14.6.1, 7.14.6.2).
- The current shift state is used with a multibyte character function after the LC_CTYPE category was changed (7.14.7).
- A string or wide-string utility function is instructed to access an array beyond the end of an object (7.15.1, 7.19.4).
- The contents of the destination array are used after a call to the strxfrm, strftime, wcsxfrm, or wcsftime function in which the specified length was too small to hold the entire null-terminated result (7.15.4.5, 7.16.3.6, 7.19.4.4.4, 7.19.5).
- The tmx structure whose address is passed as an argument to the mkxtime or strfxtime function does not provide an appropriate extension block as required by the implementation (7.16.1).
- The value of the tm_version member of the tmx structure pointed to by the argument in a call to the mkxtime function is not 1 (7.16.2.3, 7.16.2.6).
- The value of one of the tm_year, tm_mon, tm_mday, tm_hour, tm_min, tm_sec, tm_leapsecs, tm_zone, or tm_isdst members of the tm or tmx structure pointed to by the argument in a call to the mktime or mkxtime function is drastically out of the normal range (7.16.2.6).
- The value of an argument of type wint_t to a wide-character classification or mapping function is neither equal to the value of WEOF nor representable as a wchar_t (7.18.1).
- The iswctype function is called using a different LC_CTYPE category from the one in effect for the call to the wctype function that returned the description (7.18.2.2.2).
- The towctrans function is called using a different LC_CTYPE category from the one in effect for the call to the wctrans function that returned the description (7.18.3.2.2).
- The argument corresponding to an specifier without an 1 qualifier in a call to the fwprintf function does not point to a valid multibyte character sequence that begins in the initial shift state (7.19.2.3).
- The first argument in a call to the wcstok function does not point to a wide string on the first call or is not a null pointer for subsequent calls to continue parsing the same wide string, or when continuing parsing, the saved pointer value pointed to by the third argument does not match that stored by the previous call for the same wide string (7.19.4.5.7).
- An mbstate_t object is used inappropriately (7.19.7).
- The conversion state is used after the mbrtowc, wcrtomb, mbsrtowcs, or wcsrtombs function reports an encoding error (7.19.7.3.2, 7.19.7.3.3, 7.19.7.4.1, 7.19.7.4.2).


## K. 3 Implementation-defined behavior

1 A conforming implementation shall document its choice of behavior in each of the areas listed in this subclause. The following are implementation-defined:

## K.3.1 Translation

1 - How a diagnostic is identified (3.4, 5.1.1.3).

- Whether each nonempty sequence of white-space characters other than new-line is retained or replaced by one space character in translation phase 3 (5.1.1.2).


## K.3.2 Environment

- The name and type of the function called at program startup in a freestanding environment (5.1.2.1).
- The effect of program termination in a freestanding environment (5.1.2.1).
- An alternate manner in which the main function may be defined (5.1.2.2.1).
- The values given to the strings pointed to by the argv argument to main (5.1.2.2.1).
- What constitutes an interactive device (5.1.2.3).
- Signals for which the equivalent of signal(sig, SIG_IGN); is executed at program startup (7.11.1.1).
- The form of the status returned to the host environment to indicate unsuccessful termination when the SIGABRT signal is raised and not caught (7.14.4.1).
- The forms of the status returned to the host environment by the exit function to report successful and unsuccessful termination (7.14.4.3).
- The status returned to the host environment by the exit function if the value of its argument is other than zero, EXIT_SUCCESS, or EXIT_FAILURE (7.14.4.3).
- The set of environment names and the method for altering the environment list used by the getenv function (7.14.4.4).
- The manner of execution of the string by the system function (7.14.4.5).


## K.3.3 Identifiers

1 - The number of significant initial characters (beyond 63) in an identifier without external linkage (6.1.2).

- The number of significant initial characters (beyond 31) in an identifier with external linkage (6.1.2).


## K.3.4 Characters

1 - The number of bits in a byte (3.4).

- The values of the members of the execution character set (5.2.1).
- The unique value of the member of the execution character set produced for each of the standard alphabetic escape sequences (5.2.2).
- The value of a char object into which has been stored any character other than a member of the required source character set (6.1.2.5).
- Which of signed char or unsigned char has the same range, representation, and behavior as 'plain'" char (6.1.2.5, 6.2.1.1).
- The mapping of members of the source character set (in character constants and string literals) to members of the execution character set (6.1.3.4).
- The value of an integer character constant that contains more than one character or a wide character constant that contains more than one multibyte character (6.1.3.4).
- The value of an integer character constant that contains a character or escape sequence not represented in the basic execution character set or a wide character constant that contains a multibyte character or escape sequence not represented in the extended execution character set (6.1.3.4).
- The current locale used to convert a wide character constant consisting of a single multibyte character that maps to a member of the extended execution character set into a corresponding wide-character code (6.1.3.4).


## K.3.5 Integers

- Any extended integer types that exist in the implementation (6.1.2.5).
- The rank of any extended integer type relative to another extended integer type with the same precision (6.2.1.1).
- The result of converting an integer to a signed integer type when the value cannot be represented in an object of that type (6.2.1.2).
— The results of some bit-wise operations on signed integers (6.3).


## K.3.6 Floating point

1 - Rounding behavior for values of FLT_ROUNDS less than -1 (5.2.4.2.2).

- Rounding behavior for values of FLT_EVAL_METHOD less than - 1 or greater than 2 (5.2.4.2.2).
- How the nearest representable value or the larger or smaller representable value immediately adjacent to the nearest representable value is chosen for certain floating constants (6.1.3.1).
- The direction of rounding when an integer is converted to a floating-point number that cannot exactly represent the original value (6.2.1.3).
- The direction of rounding when a floating-point number is converted to a narrower floating-point number (6.2.1.4).
- Whether and how floating expressions are contracted when not disallowed by the FP_CONTRACT pragma (6.3).


## K.3.7 Arrays and pointers

1 - The result of converting a pointer to an integer or vice versa (6.2.2.3).

- The size of the result of subtracting two pointers to elements of the same array (6.3.6).


## K.3.8 Hints

1 - The extent to which suggestions made by using the register storage-class specifier are effective (6.5.1).

- The extent to which suggestions made by using the inline function specifier are effective (6.5.4).


## K.3.9 Structures, unions, enumerations, and bit-fields

1 - The behavior when a member of a union object is accessed using a member of a different type (6.3.2.3).

- Whether a "plain" int bit-field is treated as a signed int bit-field or as an unsigned int bit-field (6.5.2).
- Whether a bit-field can straddle a storage-unit boundary (6.5.2.1).
- The order of allocation of bit-fields within a unit (6.5.2.1).
- The alignment of non-bit-field members of structures (6.5.2.1). This should present no problem unless binary data written by one implementation is read by another.
- The integer type compatible with each enumerated type (6.5.2.2).


## K.3.10 Qualifiers

1 - What constitutes an access to an object that has volatile-qualified type (6.5.3).

## K.3.11 Preprocessing directives

1 - How sequences in both forms of header names are mapped to headers or external source file names (6.1.7).

- Whether the value of a character constant in a constant expression that controls conditional inclusion matches the value of the same character constant in the execution character set (6.8.1).
- Whether the value of a single-character character constant in a constant expression that controls conditional inclusion may have a negative value (6.8.1).
- The places that are searched for an included < > delimited header, and how the places are specified or the header is identified (6.8.2).
- How the named source file is searched for in an included " " delimited header (6.8.2).
- The method by which preprocessing tokens are combined into a header name (6.8.2).
- The nesting limit for \#include processing (6.8.2).
- The behavior on each recognized non-STDC \#pragma directive (6.8.6).
_ The definitions for __DATE__ and __TIME__ when respectively, the date and time of translation are not available (6.8.8).


## K.3.12 Library functions

- All library facilities available to a freestanding program (5.1.2.1).
- The null pointer constant to which the macro NULL expands (7.1.6).
- The format of the diagnostic printed by the assert macro (7.2.1.1).
— Strings other than " C " and "" that may be passed as the second argument to the setlocale function (7.5.1.1).
- The default state for the FENV_ACCESS pragma (7.6.1)
- The representation of floating point exception flags stored by the fegetexceptflag function (7.6.2.2).
- Whether the feraiseexcept function raises the inexact exception in addition to the overflow or underflow exception (7.6.2.3).
- Floating environment macros other than FE_DFL_ENV that can be used as the argument to the fesetenv or feupdateenv function (7.6.4.3, 7.6.4.4).
- The types defined for float_t and double_t when the value of the FLT_EVAL_METHOD macro is less than 0 or greater than 2 (7.7).
- The infinity to which the INFINITY macro expands, if any (7.7).
- The quiet NaN to which the NAN macro expands, when it is defined (7.7).
- The value to which the DECIMAL_DIG macro expands (7.7).
- The default state for the FP_CONTRACT pragma (7.7.2)
- Domain errors for the mathematics functions, other than those required by this International Standard (7.7.1).
- The values returned by the mathematics functions, and whether errno is set to the value of the macro EDOM, on domain errors (7.7.1).
- Whether the mathematics functions set errno to the value of the macro ERANGE on overflow and/or underflow range errors (7.7.1).
- Whether a domain error occurs or zero is returned when the fmod function has a second argument of zero (7.7.10.1).
- The base-2 logarithm of the modulus used by the remquo function in reducing the quotient (7.7.10.3).
- The set of signals, their semantics, and their default handling (7.11).
- If the equivalent of signal (sig, SIG_DFL) ; is not executed prior to the call of a signal handler, the blocking of the signal that is performed (7.11.1.1).
- Whether the equivalent of signal (sig, SIG_DFL) ; is executed prior to the call of a signal handler for the signal SIGILL (7.11.1.1).
- Signal values other than SIGFPE, SIGILL, and SIGSEGV that correspond to a computational exception (7.11.1.1).
- Whether the last line of a text stream requires a terminating new-line character (7.13.2).
- Whether space characters that are written out to a text stream immediately before a new-line character appear when read in (7.13.2).
- The number of null characters that may be appended to data written to a binary stream (7.13.2).
- Whether the file position indicator of an append-mode stream is initially positioned at the beginning or end of the file (7.13.3).
- Whether a write on a text stream causes the associated file to be truncated beyond that point (7.13.3).
- The characteristics of file buffering (7.13.3).
- Whether a zero-length file actually exists (7.13.3).
- The rules for composing valid file names (7.13.3).
- Whether the same file can be open multiple times (7.13.3).
- The nature and choice of encodings used for multibyte characters in files (7.13.3).
- The effect of the remove function on an open file (7.13.4.1).
- The effect if a file with the new name exists prior to a call to the rename function (7.13.4.2).
- Whether an open temporary file is removed upon abnormal program termination (7.13.4.3).
- What happens when the tmpnam function is called more than TMP_MAX times (7.13.4.4).
- The style used to print an infinity or NaN , and the meaning of the $n$-char-sequence if that style is printed for a $\mathrm{NaN}(7.13 .6 .1,7.19 .2 .1)$.
- The output for $\%$ p conversion in the fprintf or fwprintf function (7.13.6.1, 7.19.2.1).
- The interpretation of a - character that is neither the first nor the last character, nor the second where a $\wedge$ character is the first, in the scanlist for $\%$ [ conversion in the fscanf or fwscanf function (7.13.6.2, 7.19.2.1).
- The set of sequences matched by the \%p conversion in the fscanf or fwscanf function (7.13.6.2, 7.19.2.2).
- The interpretation of the input item corresponding to a \%p conversion in the fscanf or fwscanf function (7.13.6.2, 7.19.2.2).
- The value to which the macro errno is set by the fgetpos, fsetpos, or ftell functions on failure (7.13.9.1, 7.13.9.3, 7.13.9.4).
- The meaning of the n-char-sequence in a string converted by the strtod or wcstod function (7.14.1.5, 7.19.4.1.1).
- Whether or not the strtod or wcstod function sets errno to ERANGE when underflow occurs (7.14.1.5).
- Whether the calloc, malloc, and realloc functions return a null pointer or a pointer to an allocated object when the size requested is zero (7.14.3).
- Whether open output streams are flushed, open streams are closed, or temporary files are removed when the abort function is called (7.14.4.1).
Portability issues
- The termination status returned to the host environment by the abort function (7.14.4.1).
- The value returned by the system function when its argument is not a null pointer (7.14.4.5).
- The local time zone and Daylight Saving Time (7.16.1).
- The era for the clock function (7.16.2.1).
- The positive value for $\mathbf{t m}$ _isdst in a normalized $\boldsymbol{t m x}$ structure (7.16.2.6).
- The replacement string for the $\% \mathbf{Z}$ specifier to the strftime function in the "C" locale (7.16.3.6).
- Whether or when the trigonometric, hyperbolic, base-e exponential, base-e logarithmic, error, and $\log$ gamma functions raise the inexact exception in an IEC 559 conformant implementation (F.9).
- Whether the inexact exception may be raised when the rounded result actually does equal the mathematical result in an IEC 559 conformant implementation (F.9).
- Whether the underflow (and inexact) exception may be raised when a result is tiny but not inexact in an IEC 559 conformant implementation (F.9).
- Whether the functions honor the rounding direction mode (F.9).


## K.3.13 Architecture

1 - The number, order, and encoding of bytes in any object (when not explicitly specified in this International Standard) (3.15).

- The values or expressions assigned to the macros specified in the headers <limits.h>, <float.h>, and <inttypes.h> (5.2.4.2, 5.2.4.2.1, 5.2.4.2.2, 7.4.2, 7.4.5).
- The value of the result of the sizeof operator (6.3.3.4).


## K. 4 Locale-specific behavior

1 The following characteristics of a hosted environment are locale-specific and shall be documented by the implementation:

- Additional members of the execution character set beyond the required members (5.2.1).
- The presence, meaning, and representation of additional multibyte characters in the execution character set beyond the required single-byte characters (5.2.1.2).
- The shift states used for the encoding of multibyte characters (5.2.1.2).
- The direction of writing of successive printable characters (5.2.2).
- The decimal-point character (7.1.1).
- The set of printing characters (7.3).
- The set of control characters (7.3).
- The sets of characters tested for by the isalpha, isblank, islower, ispunct, isspace, or isupper functions (7.3.1.2, 7.3.1.3, 7.3.1.7, 7.3.1.9, 7.3.1.10, 7.3.1.11).
- The native environment (7.5.1.1).
- Additional subject sequences accepted by the string conversion functions (7.14.1) and the wide string numeric conversion function (7.19.4.1).
- The collation sequence of the execution character set (7.15.4.3).
- The contents of the error message strings set up by the strerror function (7.15.6.2).
- The formats for time and date (7.16.3.6).
- Character mappings that are supported by the towctrans function (7.18.1).
- Character classifications that are supported by the iswctype function (7.18.1).
- The set of printing wide characters (7.18.2).
- The set of control wide characters (7.18.2).
- The sets of wide characters tested for by the iswalpha, iswblank, iswlower, iswpunct, iswspace, or iswupper functions (7.18.2.1.2, 7.18.2.1.3, 7.18.2.1.7, 7.18.2.1.9, 7.18.2.1.10, 7.18.2.1.11).


## K. 5 Common extensions

1 The following extensions are widely used in many systems, but are not portable to all implementations. The inclusion of any extension that may cause a strictly conforming program to become invalid renders an implementation nonconforming. Examples of such extensions are new keywords, extra library functions declared in standard headers, or predefined macros with names that do not begin with an underscore.

## K.5.1 Environment arguments

1 In a hosted environment, the main function receives a third argument, char *envp [], that points to a null-terminated array of pointers to char, each of which points to a string that provides information about the environment for this execution of the program (5.1.2.2.1).

## K.5.2 Specialized identifiers

1 Characters other than the underscore _, letters, and digits, that are not defined in the required source character set (such as the dollar sign \$, or characters in national character sets) may appear in an identifier (6.1.2).

## K.5.3 Lengths and cases of identifiers

1 All characters in identifiers (with or without external linkage) are significant (6.1.2).

## K.5.4 Scopes of identifiers

1 A function identifier, or the identifier of an object the declaration of which contains the keyword extern, has file scope (6.1.2.1).

## K.5.5 Writable string literals

1 String literals are modifiable (in which case, identical string literals should denote distinct objects) (6.1.4).

## K.5.6 Other arithmetic types

1 Additional arithmetic types, such as __int128, and their appropriate conversions are defined (6.1.2.5, 6.2.1.1).

## K.5.7 Function pointer casts

1 A pointer to an object or to void may be cast to a pointer to a function, allowing data to be invoked as a function (6.3.4).

2 A pointer to a function may be cast to a pointer to an object or to void, allowing a function to be inspected or modified (for example, by a debugger) (6.3.4).

## K.5.8 Non-int bit-field types

1 Types other than unsigned int or signed int can be declared as bit-fields, with appropriate maximum widths (6.5.2.1).

## K.5.9 The fortran keyword

1 The fortran function specifier may be used in a function declaration to indicate that calls suitable for FORTRAN should be generated, or that a different representation for the external name is to be generated (6.5.4).

## K.5.10 The asm keyword

1 The asm keyword may be used to insert assembly language directly into the translator output; the most common implementation is via a statement of the form
asm ( character-string-literal );

## K.5.11 Multiple external definitions

1 There may be more than one external definition for the identifier of an object, with or without the explicit use of the keyword extern; if the definitions disagree, or more than one is initialized, the behavior is undefined (6.7.2).

## K.5.12 Predefined macro names

1 Macro names that do not begin with an underscore, describing the translation and execution environments, are defined by the implementation before translation begins (6.8.8).

## K.5.13 Extra arguments for signal handlers

1 Handlers for specific signals may be called with extra arguments in addition to the signal number (7.11.1.1).

## K.5.14 Additional stream types and file-opening modes

 Additional mappings from files to streams may be supported (7.13.2).Additional file-opening modes may be specified by characters appended to the mode argument of the fopen function (7.13.5.3).

## K.5.15 Defined file position indicator

1 The file position indicator is decremented by each successful call to the ungetc or ungetwc function for a text stream, except if its value was zero before a call (7.13.7.11, 7.19.3.9).

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[^0]:    1. This International Standard is designed to promote the portability of $C$ programs among a variety of data-processing systems. It is intended for use by implementors and programmers.
[^1]:    5. Implementations must behave as if these separate phases occur, even though many are typically folded together in practice.
[^2]:    8. The intent is that an implementation should identify the nature of, and where possible localize, each violation. Of course, an implementation is free to produce any number of diagnostics as long as a valid program is still correctly translated. It may also successfully translate an invalid program.
[^3]:    9. Thus, int can be replaced by a typedef name defined as int, or the type of argv can be written as char ** argv, and so on.
[^4]:    10. In accordance with subclause 6.1.2.4, objects with automatic storage duration declared in main will no longer have storage guaranteed to be reserved in the former case even where they would in the latter.
[^5]:    12. The trigraph sequences enable the input of characters that are not defined in the Invariant Code Set as described in ISO/IEC 646:1991, which is a subset of the seven-bit ASCII code set.
[^6]:    14. See 6.1.2.5.
    15. The floating-point model is intended to clarify the description of each floating-point characteristic and does not require the floating-point arithmetic of the implementation to be identical.
    16. IEC 559:1993 specifies quiet and signaling NaNs. For implementations that do not support IEC 559:1993, the terms quiet NaN and signaling NaN are intended to apply to encodings with similar behavior.
    17. Evaluation of FLT_ROUNDS correctly reflects any execution-time change of rounding mode through the function fesetround in <fenv.h>.
[^7]:    20. On systems in which linkers cannot accept extended characters, an encoding of the universal-character-name may be used in forming valid external identifiers. For example, some otherwise unused character or sequence of characters may be used to encode the $\backslash \mathbf{u}$ in a universal-charactername. Extended characters may produce a long external identifier.
[^8]:    21. A function declaration can contain the storage-class specifier static only if it is at file scope; see 6.5.1.
    22. As specified in 6.1.2.1, the later declaration might hide the prior declaration.
[^9]:    23. There is only one name space for tags even though three are possible.
    24. The term constant address means that two pointers to the object constructed at possibly different times will compare equal. The address may be different during two different executions of the same program.
    In the case of a volatile object, the last store may not be explicit in the program.
[^10]:    25. Implementation-defined keywords must have the form of an identifier reserved for any use as described in 7.1.3.
    26. Therefore, any statement in this Standard about signed integer types also applies to the extended signed integer types.
[^11]:    27. The same representation and alignment requirements are meant to imply interchangeability as arguments to functions, return values from functions, and members of unions.
    28. A specification for imaginary types is in informative Annex G.
[^12]:    32. Note that aggregate type does not include union type because an object with union type can only contain one member at a time.
[^13]:    36. A positional representation for integers that uses the binary digits 0 and 1 , in which the values represented by successive bits are additive, begin with 1, and are multiplied by successive integral powers of 2, except perhaps the bit with the highest position. (Adapted from the American National Dictionary for Information Processing Systems.) A byte contains CHAR_BIT bits, and the values of type unsigned char range from 0 to $2^{\text {CHAR_BIT }}-1$.
    37. Thus an automatic variable can be initialized to a trap representation without causing undefined behavior, but the value of the variable cannot be used until a proper value is stored in it.
[^14]:    38. Thus, for example, structure assignment may be implemented element-at-a-time or via memcpy.
    39. Some combinations of padding bits might generate trap representations, for example, if one padding bit is a parity bit. Regardless, no arithmetic operation on valid values can generate a trap representation other than as part of an exception such as an overflow, and this cannot occur with unsigned types. All other combinations of padding bits are alternative object representations of the value specified by the value bits.
[^15]:    40. The specification for the library functions recommends more accurate conversion than required for floating constants. See strtod (7.14.1.5).
[^16]:    41. The semantics of these characters were discussed in 5.2.2.
    42. See 'future language directions" (6.9.1).
[^17]:    43. A character string literal need not be a string (see 7.1.1), because a null character may be embedded in it by a \o escape sequence.
[^18]:    45. Thus, sequences of characters that resemble escape sequences cause undefined behavior.
[^19]:    46. Thus, /* ... */ comments do not nest.
[^20]:    47. The integer promotions are applied only as part of the usual arithmetic conversions, to certain argument expressions, to the operands of the unary $\boldsymbol{+},-$, and $\sim$ operators, and to both operands of the shift operators, as specified by their respective subclauses.
    48. The remaindering operation performed when a value of integer type is converted to unsigned type need not be performed when a value of real floating type is converted to unsigned type. Thus, the range of portable real floating values is ( $-1, \mathbf{U}$ type_MAX+1).
[^21]:    49. For example, addition of a double complex and a float entails just the conversion of the float operand to double (and yields a double complex result).
    50. The cast and assignment operators still must perform their specified conversions, as described in 6.2.1.3 and 6.2.1.4.
    51. The name 'lvalue" comes originally from the assignment expression $\mathbf{E} \mathbf{1}=\mathbf{E} \mathbf{2}$, in which the left operand E1 must be a (modifiable) lvalue. It is perhaps better considered as representing an object "locator value." What is sometimes called "rvalue"' is in this International Standard described as the "value of an expression."
    An obvious example of an lvalue is an identifier of an object. As a further example, if $\mathbf{E}$ is a unary expression that is a pointer to an object, $\boldsymbol{*} \mathbf{E}$ is an lvalue that designates the object to which $\mathbf{E}$ points.
[^22]:    52. Because this conversion does not occur, the operand of the sizeof operator remains a function designator and violates the constraint in 6.3.3.4.
[^23]:    59. The intent of this list is to specify those circumstances in which an object may or may not be aliased.
    60. A contracted expression might also omit the raising of floating-point exception flags.
    61. This license is specifically intended to allow implementations to exploit fast machine instructions that combine multiple C operators. As contractions potentially undermine predictability, and can even decrease accuracy for containing expressions, their use must be well-defined and clearly documented.
[^24]:    64. Most often, this is the result of converting an identifier that is a function designator.
[^25]:    65. A function may change the values of its parameters, but these changes cannot affect the values of the arguments. On the other hand, it is possible to pass a pointer to an object, and the function may change the value of the object pointed to. A parameter declared to have array or function type is converted to a parameter with a pointer type as described in 6.7.1.
[^26]:    66. If $\boldsymbol{\&} E$ is a valid pointer expression (where $\boldsymbol{\&}$ is the "address-of" operator, which generates a pointer to its operand), the expression (\&E) $->$ MOS is the same as $\mathbf{E}$. MOS.
    67. The "byte orders" for scalar types are invisible to isolated programs that do not indulge in type punning (for example, by assigning to one member of a union and inspecting the storage by accessing another member that is an appropriately sized array of character type), but must be accounted for when conforming to externally imposed storage layouts.
[^27]:    68. Note that this differs from a cast expression. For example, a cast specifies a conversion to scalar
[^28]:    71. Thus $\& \mathbf{E E}^{2}$ is equivalent to $\mathbf{E}$ (even if $\mathbf{E}$ is a null pointer), and $\&(\mathbf{E} 1[\mathbf{E} 2]]$ ) to ( $\mathbf{E} 1+(\mathbf{E} 2)$ ). It is always true that if $\mathbf{E}$ is a function designator or an lvalue that is a valid operand of the unary $\&$ operator, $\boldsymbol{*} \boldsymbol{E}$ is a function designator or an lvalue equal to $\mathbf{E}$. If $\boldsymbol{*} \mathbf{P}$ is an lvalue and $\mathbf{T}$ is the name of an object pointer type, $\boldsymbol{*}(\mathbf{T}) \mathbf{P}$ is an lvalue that has a type compatible with that to which $\mathbf{T}$ points.
    Among the invalid values for dereferencing a pointer by the unary * operator are a null pointer, an address inappropriately aligned for the type of object pointed to, and the address of an automatic storage duration object when execution of the block with which the object is associated has terminated.
[^29]:    72. When applied to a parameter declared to have array or function type, the sizeof operator yields the size of the pointer obtained by converting as in 6.2.2.1; see 6.7.1.
[^30]:    75. This is often called "truncation toward zero".
[^31]:    76. Another way to approach pointer arithmetic is first to convert the pointer(s) to character pointer(s): In this scheme the integer expression added to or subtracted from the converted pointer is first multiplied by the size of the object originally pointed to, and the resulting pointer is converted back to the original type. For pointer subtraction, the result of the difference between the character pointers is similarly divided by the size of the object originally pointed to.

    When viewed in this way, an implementation need only provide one extra byte (which may overlap another object in the program) just after the end of the object in order to satisfy the "one past the last element'" requirements.

[^32]:    77. The expression $\mathbf{a}<\mathbf{b}<\mathbf{c}$ is not interpreted as in ordinary mathematics. As the syntax indicates, it means $(\mathbf{a}<\mathbf{b})<\mathbf{c}$; in other words, 'if $\mathbf{a}$ is less than $\mathbf{b}$ compare 1 to $\mathbf{c}$; otherwise, compare 0 to $\mathbf{c}$.',
    78. Because of the precedences, $\mathbf{a}<\mathbf{b}=\mathbf{c}<\mathbf{d}$ is 1 whenever $\mathbf{a}<\mathbf{b}$ and $\mathbf{c}<\mathbf{d}$ have the same truth-value.
    79. If invalid prior pointer operations, such as accesses outside array bounds, produced undefined behavior, the effect of subsequent comparisons is undefined.
[^33]:    80. A conditional expression does not yield an lvalue.
[^34]:    81. The asymmetric appearance of these constraints with respect to type qualifiers is due to the conversion (specified in 6.2.2.1) that changes lvalues to "the value of the expression' which removes any type qualifiers from the type category of the expression.
[^35]:    82. A comma operator does not yield an lvalue.
[^36]:    85. Thus, in the following initialization,
    static int $i=2| | 1 / 0 ;$
    the expression is a valid integer constant expression with value one.
[^37]:    89. A structure or union can not contain a member with a variably modified type because member names are not ordinary identifiers as defined in 6.1.2.3.
    90. The unary \& (address-of) operator may not be applied to a bit-field object; thus, there are no pointers to or arrays of bit-field objects.
    91. As specified in 6.5.2 above, if the actual type specifier used is int or there is no type specifier, or is a typedef-name defined using either of these, then it is implementation-defined whether the bit-field is signed or unsigned.
    92. An unnamed bit-field structure member is useful for padding to conform to externally imposed layouts.
[^38]:    93. Thus, the identifiers of enumeration constants declared in the same scope shall all be distinct from each other and from other identifiers declared in ordinary declarators.
[^39]:    94. An incomplete type may only by used when the size of an object of that type is not needed. It is not needed, for example, when a typedef name is declared to be a specifier for a structure or union, or when a pointer to or a function returning a structure or union is being declared. (See incomplete types in 6.1.2.5.) The specification shall be complete before such a function is called or defined.
[^40]:    95. If there is no identifier, the type can, within the translation unit, only be referred to by the declaration of which it is a part. Of course, when the declaration is of a typedef name, subsequent declarations can make use of that typedef name to declare objects having the specified structure, union, or enumerated type.
    96. A similar construction with enum does not exist.
[^41]:    97. The implementation may place a const object that is not volatile in a read-only region of storage. Moreover, the implementation need not allocate storage for such an object if its address is never used.
    98. This applies to those objects that behave as if they were defined with qualified types, even if they are never actually defined as objects in the program (such as an object at a memory-mapped input/output address).
[^42]:    99. A volatile declaration may be used to describe an object corresponding to a memory-mapped input/output port or an object accessed by an asynchronously interrupting function. Actions on objects so declared shall not be 'optimized out'" by an implementation or reordered except as permitted by the rules for evaluating expressions.
    100. Both of these can occur through the use of typedefs.
[^43]:    101. Inline substitution is not textual substitution, nor does it create a new function. Therefore, for example, the expansion of a macro used within the body of the function uses the definition it had at the point the function body appears, and not where the function is called; and identifiers refer to the declarations in scope where the body occurs. Similarly, the address of the function is not affected by the function's being inlined.
    102. For example, an implementation might never perform inline substitution, or might only perform inline substitutions to calls in the scope of an inline declaration.
[^44]:    103. When several 'array of'" specifications are adjacent, a multidimensional array is declared.
[^45]:    104. Arrays and functions are rewritten as pointers.
[^46]:    105. The macros defined in the <stdarg. h> header (7.12) may be used to access arguments that correspond to the ellipsis.
    106. See 'future language directions'" (6.9.3).
    107. If both function types are 'old style', parameter types are not compared.
[^47]:    108. As indicated by the syntax, empty parentheses in a type name are interpreted as "function with no parameter specification", rather than redundant parentheses around the omitted identifier.
[^48]:    109. If the initializer list for a subaggregate or contained union does not begin with a left brace, its subobjects are initialized as usual, but the subaggregate or contained union does not become the current object: current objects are associated only with brace-enclosed initializer lists.
    110. After a union member is initialized, the next object is not the next member of the union; instead, it is the next subobject of an object containing the union.
[^49]:    111. Thus, a designator can only specify a strict subobject of the aggregate or union that is associated with the surrounding brace pair. Note, too, that each separate designator list is independent.
[^50]:    112. In particular, the evaluation order need not be the same as the order of subobject initialization.
[^51]:    113. Such as assignments, and function calls which have side effects.
[^52]:    114. Thus, clause- 1 specifies initialization for the loop, possibly declaring one or more variables for use in the loop; expression-2, the controlling expression, specifies an evaluation made before each iteration, such that execution of the loop continues until the expression compares equal to 0 ; expression-3 specifies an operation (such as incrementing) that is performed after each iteration. If clause-1 is a declaration, then the scope of any variable it declares is the remainder of the declaration and the entire loop, including the other two expressions.
[^53]:    115. Following the contin: label is a null statement.
    116. The return statement is not an assignment. The overlap restriction of subclause 6.3.16.1 does not apply to the case of function return.
[^54]:    117. Thus, if an identifier declared with external linkage is not used in an expression, there need be no external definition for $i t$.
[^55]:    118. The intent is that the type category in a function definition cannot be inherited from a typedef:
    ```
    typedef int F(void); /* type F is "function of no arguments returning int" */
    F f, g; /* f and g both have type compatible with F */
    F f { /* ... */ } /* WRONG: syntax/constraint error */
    F g() { /* ... */ } /*WRONG: declares that g returns a function */
    int f(void) { /* ... */ } /* RIGHT: f has type compatible with F */
    int g() {/* .. */ } /* RIGHT: g has type compatible with F */
    F *e (void) { /* ... */ } /* e returns a pointer to a function */
    F *((e)) (void) { /* ... */ } /* same: parentheses irrelevant */
    int (*fp) (void);
    F *Fp;
    /* fp points to a function that has type F */
    /* Fp points to a function that has type F */
    ```

[^56]:    119. See "future language directions" (6.9.4).
    120. A parameter is in effect declared at the head of the compound statement that constitutes the function body, and therefore may not be redeclared in the function body (except in an enclosed block).
[^57]:    121. Thus, preprocessing directives are commonly called "lines." These "lines" have no other syntactic significance, as all white space is equivalent except in certain situations during preprocessing (see the \# character string literal creation operator in 6.8.3.2, for example).
[^58]:    122. This allows an implementation to choose to interpret the directive:
    \#define THIS\$AND\$THAT (a, b) ( a ) + (b))
    as defining a function-like macro THIS\$AND\$THAT, rather than an object-like macro THIS.
    Whichever choice it makes, it must also issue a diagnostic.
[^59]:    123. Because the controlling constant expression is evaluated during translation phase 4 , all identifiers either are or are not macro names - there simply are no keywords, enumeration constants, etc.
[^60]:    126. Note that adjacent string literals are not concatenated into a single string literal (see the translation phases in 5.1.1.2); thus, an expansion that results in two string literals is an invalid directive.
[^61]:    128. The value in ISO/IEC 9899:1994 was 199409L.
[^62]:    129. The functions that make use of the decimal-point character are atof, fprintf, fscanf, fwprintf, fwscanf, localeconv, printf, scanf, sprintf, sscanf, strtod, swprintf, swscanf, vfprintf, vfscanf, vfwprintf, vfwscanf, vprintf, vscanf, vsprintf, vsscanf, vswprintf, vswscanf, vwprintf, vwscanf, wprintf, and wscanf.
[^63]:    131. For state-dependent encodings, the values for MB_CUR_MAX and MB_LEN_MAX must thus be large enough to count all the bytes in any complete multibyte character plus at least one adjacent shift sequence of maximum length. Whether these counts provide for more than one shift sequence is the implementation's choice.
    132. A header is not necessarily a source file, nor are the < and > delimited sequences in header names necessarily valid source file names.
[^64]:    133. The list of reserved identifiers with external linkage includes errno, set jmp, and va_end.
    134. Since macro names are replaced whenever found, independent of scope and name space, macro names matching any of the reserved identifier names must not be defined if an associated header, if any, is included.
[^65]:    135. The macro errno need not be the identifier of an object. It might expand to a modifiable lvalue resulting from a function call (for example, *errno ()).
    136. Thus, a program that uses errno for error checking should set it to zero before a library function call, then inspect it before a subsequent library function call. Of course, a library function can save the value of errno on entry and then set it to zero, as long as the original value is restored if errno's value is still zero just before the return.
    137. See 'future library directions'" (7.20.1).
[^66]:    138. The traditional choice for type bool has been int, but this is not a requirement of this International Standard. Other available choices include, but are not limited to, char, unsigned int, and an enumeration type.
    If an enumeration type is chosen, the names of its true and false members are "masked" by the macros true and false, but the member names might be available to the debugger:
    ```
    typedef enum { false=0, true=1 } bool;
    #define false 0
    #define true 1
    ```

    The type is suitable for bit-fields if it is int, unsigned int, signed int, or some type allowed by an implementation extension. It is required that a bool bit-field of width 1 be unsigned. Thus, bool cannot be signed int, nor can it be plain int if width 1 plain int bit-fields are signed.

[^67]:    143. The message written might be of the form

    Assertion failed: expression, function $a b c$, file $x y z$, line $n n n$

[^68]:    144. See "future library directions" (7.20.2).
    145. In an implementation that uses the seven-bit ASCII character set, the printing characters are those whose values lie from $0 \times 20$ (space) through $0 \times 7 \mathrm{E}$ (tilde); the control characters are those whose values lie from 0 (NUL) through 0x1F (US), and the character 0x7F (DEL).
[^69]:    146. The functions islower and isupper test true or false separately for each of these additional characters; all four combinations are possible.
[^70]:    147. Some of these typedef names may denote implementation-defined extended integer types.
[^71]:    148. The designated type is not guaranteed to be fastest for all purposes; if the implementation has no clear grounds for choosing one type over another, it will simply pick some integer type satisfying the signedness and width requirements.
[^72]:    149. C++ implementations should define these macros only when $\qquad$ STDC_LIMIT_MACROS is defined before <inttypes.h> is included.
[^73]:    150. C++ implementations should define these macros only when $\qquad$ STDC_CONSTANT_MACROS is defined before <inttypes. h > is included.
[^74]:    151. C++ implementations should define these macros only when __STDC_FORMAT_MACROS is defined before <inttypes.h> is included.
    152. Separate macros are given for use with fprintf and fscanf functions because, typically, different format specifiers are required for fprintf and fscanf even when the type is the same.
[^75]:    153. ISO/IEC 9945-2, Information technology - Portable operating system interface (POSIX) - Part 2: shell and utilities specifies locale and charmap formats that may be used to specify locales for C.
[^76]:    154. See "future library directions" (7.20.4).
    155. The only functions in 7.3 whose behavior is not affected by the current locale are isdigit and isxdigit.
[^77]:    156. The implementation must arrange to encode in a string the various categories due to a heterogeneous locale when category has the value LC_ALL.
[^78]:    159. Even though the rounding direction macros may expand to constants corresponding to the values of FLT_ROUNDS, they are not required to do so.
    160. The purpose of the FENV_ACCESS pragma is to allow certain optimizations, for example global common subexpression elimination, code motion, and constant folding, that could subvert flag tests and mode changes. In general, if the state of FENV_ACCESS is off then the translator can assume that default modes are in effect and the flags are not tested.
[^79]:    161. The side effects impose a temporal ordering that requires two evaluations of $\mathbf{x}+1$. On the other hand, without the \#pragma STDC FENV_ACCESS ON pragma, and assuming the default state is off, just one evaluation of $\mathbf{x}+1$ would suffice.
    162. The functions fetestexcept, feraiseexcept, and feclearexcept support the basic abstraction of flags that are either set or clear. An implementation may endow exception flags with more information - for example, the address of the code which first raised the exception; the functions fegetexceptflag and fesetexceptflag deal with the full content of flags.
[^80]:    163. The effect is intended to be similar to that of exceptions raised by arithmetic operations. Hence, enabled traps for exceptions raised by this function are taken. The specification in F.7.6 is in the same spirit.
[^81]:    164. This mechanism allows testing several exceptions with just one function call.
[^82]:    165. IEC 559 systems have a default non-stop mode, and typically at least one other mode for trap handling or aborting; if the system provides only the non-stop mode then installing it is trivial. For such systems, the feholdexcept function can be used in conjunction with the feupdateenv function to write routines that hide spurious exceptions from their callers.
[^83]:    169. Typically, the FP_FAST_FMA macro is defined if and only if the fma function is implemented directly with a hardware multiply-add instruction. Software implementations are expected to be substantially slower.
[^84]:    170. DECIMAL_DIG is intended to give an appropriate number of digits to carry in canonical decimal representations.
    171. In order that correctly rounded conversion from an internal floating-point format with precision $m$ to decimal with DECIMAL_DIG digits and back be the identity function, DECIMAL_DIG should be a positive integer $n$ satisfying the inequality:
    $n \geq m$ if FLT_RADIX is 10
    $10^{n-1}>$ FLT_RADIX ${ }^{m}$ otherwise
    172. In an implementation that supports infinities, this allows an infinity as an argument to be a domain error if the mathematical domain of the function does not include the infinity.
[^85]:    173. The term underflow here is intended to encompass both gradual underflow as in IEC 559 and also flush-to-zero underflow.
[^86]:    174. Since an expression can be evaluated with more range and precision than its type has, it is important to know the type that classification is based on. For example, a normal long double value might become subnormal when converted to double, and zero when converted to float.
[^87]:    176. For the isnan macro, the type for determination doesn't matter unless the implementation supports NaNs in the evaluation type but not in the semantic type.
[^88]:    178. For small magnitude $\mathbf{x}, \log 1 \mathbf{p}(\mathbf{x})$ is expected to be more accurate than $\log (\mathbf{1}+\mathbf{x})$.
[^89]:    180. The argument values are converted to the type of the function, even by a macro implementation of the function.
    181. The result of the nextafterx function is determined in the type of the function, without loss of range or precision in a floating second argument.
[^90]:    184. IEC 559 requires that the built-in relational operators raise the invalid exception if the operands compare unordered, as an error indicator for programs written without consideration of NaNs; the result in these cases is false.
[^91]:    185. The imaginary unit is a number $i$ such that $i * i=-1$.
    186. A specification for imaginary types is in informative Annex G.
[^92]:    188. For a variable $\mathbf{z}$ of complex type, $\mathbf{z}==\operatorname{creal}(\mathbf{z})+\operatorname{cimag}(\mathbf{z}) * I$.
[^93]:    189. Like other function-like macros in Standard libraries, each type-generic macro can be suppressed to make available the corresponding ordinary function.
    190. If the type of the argument is incompatible with the type of the parameter for the selected function, the behavior is undefined.
[^94]:    191. These functions are useful for dealing with unusual conditions encountered in a low-level function of a program.
[^95]:    192. For example, by executing a return statement or because another longjmp call has caused a transfer to a set jmp invocation in a function earlier in the set of nested calls.
[^96]:    195. It is permitted to create a pointer to a va_list and pass that pointer to another function, in which case the original function may make further use of the original list after the other function returns.
[^97]:    275. IEEE 854, but not IEC 559 (IEEE 754), directly specifies that floating-to-integer conversions raise the inexact exception for non-integer in-range values. In those cases where it matters, library functions can be used to effect such conversions with or without raising the inexact exception. See rint, lrint, llrint, and nearbyint in <math. h>.
    276. If the minimum-width IEC 559 extended format ( 64 bits of precision) is supported, DECIMAL_DIG must be at least 21. If IEC 559 double ( 53 bits of precision) is the widest IEC 559 format supported, then DECIMAL_DIG must be at least 17. (By contrast, LDBL_DIG and DBL_DIG are 19 and 15, respectively, for these formats.)
[^98]:    277. This specification does not require dynamic rounding precision nor trap enablement modes.
    278. If the state for the FENV_ACCESS pragma is off, the implementation is free to assume the modes will be the default ones and the flags will not be tested, which allows certain optimizations (see F.8).
    279. As floating constants are converted to appropriate internal representations at translation time, their conversion is subject to default rounding modes and raises no execution-time exceptions (even where the state of the FENV_ACCESS pragma is on). Library functions, for example strtod, provide executiontime conversion of numeric strings.
[^99]:    280. Where the state for the FENV_ACCESS pragma is on, results of inexact expressions like $\mathbf{1 . 0 / 3 . 0}$ are affected by rounding modes set at execution time, and expressions such as $0.0 / 0.0$ and $1.0 / 0.0$ generate execution-time exceptions. The programmer can achieve the efficiency of translation-time evaluation through static initialization, such as
[^100]:    281. Use of float_t and double_t variables increases the likelihood of translation-time computation. For example, the automatic initialization
    double_t x = 1.1e75;
    could be done at translation time, regardless of the expression evaluation method.
[^101]:    284. $0-0$ yields -0 instead of +0 just when the rounding direction is downward.
[^102]:    285. IEC 559 allows different definitions of underflow. They all result in the same values, but differ on when the exception is raised.
    286. It is intended that undeserved underflow and inexact exceptions are raised only if determining inexactness would be too costly.
[^103]:    288. Ideally, fmax would be sensitive to the sign of zero, for example $\mathbf{f m a x}(-\mathbf{0 . 0},+\mathbf{0} \mathbf{0})$ would return +0 ; however, implementation in software might be impractical.
[^104]:    289. These properties are already implied for those cases covered in the tables, but are required for all cases (at least where the state for CX_LIMITED_RANGE is off).
