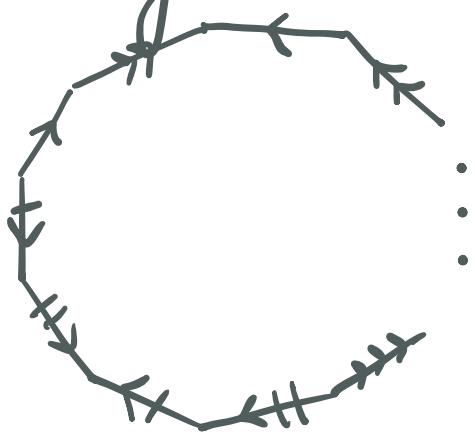


## I. Entry to surfaces:

1. Use the fact that every surface of genus  $g$  can be obtained from a polygon w/  $4g$  sides and alternating sides identified



to prove that  $\chi(S_g) = 2 - 2g$ .

Note that given a polygonal decomposition of a surface with  $F$  faces,  $E$  edges, and  $V$  vertices

$$\chi(S_g) = V - E + F.$$

## 2) Geometric intersection #:

a) Compute the geometric intersection number between two nonseparating curves, between two separating curves, and between a separating curve and a nonseparating curve.

(Choose pairs of curves that intersect nontrivially!)

b) Suppose  $\alpha$  is a separating curve. Prove that  $i(\alpha, \beta) \in 2\mathbb{Z}$  for any s.c.c.  $\beta$ .

## 3) True or False?

"Any two s.c.c. that intersect exactly once are in minimal position."

## II. Entry to the curve complex:

- 1) Prove the curve complex is quasi-isometric to its 1-skeleton
- 2) Prove that  $\delta$ -hyperbolicity is a quasi-isometry invariant.

Hint: Quasi-geodesic stability!

Every quasi-geodesic is bounded distance from a geodesic in a  $\delta$ -hyperbolic space.

Def: A quasi-geodesic is a geodesic up to additive and multiplicative constants. For example, the image of a geodesic under a

quasi-isometry is a quasi-geodesic.

3) Prove that the curve complex  
is connected for any closed surface  
with genus at least 2.

4) Show  $\hat{N}(S_{g,n})$  is connected.

5) Are  $\text{Map}(S_{1,0})$  and  $\text{Map}(S_{1,1})$   
finitely generated? If so,  
what are their respective finite  
generating sets? Can you prove/  
justify your answers?

### III. Intro to MCG's

$Z(G)$

1) The center<sup>v</sup> of a group,  $G$ , is the collection of elements that commute with every element in the group:

$$Z(G) = \{g \in G \mid ag = ga \ \forall a \in G\}$$

In this question you will compute the center of  $\text{Map}(S)$  for  $S$  finite type.

Some needed ingredients:

Vhm (Alexander Method)

If  $f \in \text{Map}(S)$  is nontrivial, then  
 $\exists$  a s.c.c.  $c$  on  $S$  s.t.  $f(c) \neq c$ .

This was extended to infinite type surfaces by  
(Fernández Fernández)-Morales-Maldonado

Two facts about Dehn twists:

Let  $f \in \text{Map}(S)$  and let  $a, b$  be s.c.c.

$$\left. \begin{array}{l} 1) T_a = T_b \Leftrightarrow a = b \\ 2) f T_a f^{-1} = T_{f(a)} \end{array} \right\} \text{Exercise within an exercise: prove these facts!}$$

Equipped with these ingredients, compute the center of  $\text{Map}(S)$ !

Does your proof rely on  $S$  being finite type? Why or why not?

## IV. Intro to big MCG's

1) Given an infinite-type surface  $S$  we can define its curve graph,  $C(S)$ , in the same way as we do for a finite-type surface.

- a) Is  $C(S)$  connected?
  - b) Infinite diameter?
  - c)  $\delta$ -hyperbolic?
- 2) Check out the survey of Aramayona-Ulamis and the open problem list compiled by Chandran-Patel-Ulamis. Pick out two or three problems

you are interested in and discuss them with at least one other person. Some discussion prompts if you're nervous:

- Share what about these problems piqued your interest.
- What types of tools/topics would you want to learn if you were tackling this problem?
- How close (mathematically) is this to other things you've studied or brand new?

Sketch of finite generation proof:

So what can we show about  $\text{Map}(S)$  using the action on  $\mathcal{C}(S)$ ?

Generating  $\text{Map}(S)$ :

Vhm (Dehn-Lickarish)

$\text{Map}(S_g)$  is generated by finitely many Dehn twists about non-separating curves.

Key Lemma:  $G$  a group  $\curvearrowright X$  a  $\overset{\vee}{\text{1-dim}}$  simplicial complex such that:

- $G \curvearrowright X$  by simplicial automorphisms
- $G \curvearrowright V(X)$  transitively
- $G \curvearrowright E(X)$  transitively

Let  $(v, w) \in E(X)$  and choose  $h \in G$

so that  $h(w) = v$ . Then  $G = \langle h, \text{Stab}(v) \rangle$ .

Key Complex:  $\hat{N}(S)$  - modified non-sep.  
curve complex

vertices - non-separating curves  
edges - correspond to non-sep.  
curves intersecting exactly  
once

Exercise: Show  $\hat{N}(S_{g,n})$  is connected.

Key assumption:  $\text{PMap}(S_{0,n})$  is finitely generated by Dehn twists for  $n \geq 0$ .

Key tool: The Birman exact sequence:

$$1 \rightarrow \pi_1(S_{g,n}) \rightarrow \text{PMap}(S_{g,n+1}) \rightarrow \text{PMap}(S_{g,n}) \rightarrow 1.$$

## Sketch of Proof: Proof by induction

Base case:  $\text{Map}(S_{1,0})$  and  $\text{Map}(S_{1,1})$  are generated by Dehn twists about the meridian and longitudinal curves.

(Exercise)

Inductive step: (On # of punctures)

Let  $g \geq 1$  and  $n \geq 0$ . Note that  $\pi_1(S_{g,n})$  is f.g. by non-sep loops. Each of these loops  $\alpha$  gives us  $P_\alpha = T_a T_b^{-1} \in \text{PMap}(S_{g,n+1})$ . By assumption,  $\text{PMap}(S_{g,n})$  is f.g. by Dehn twists about non-sep. curves.

Lift this generating set to  $\text{PMap}(S_{g,n+1})$ . So  $\text{PMap}(S_{g,n+1})$  is f.g. by Dehn twists about non-sep curves.

Together w/ base case we have completed the proof for  $\text{PMap}(S_{1,n})$  for  $n \geq 0$ .

Inductive step: (On genus)

Since  $\hat{N}(S_g)$  is connected and  $\text{Map}(S_g)$  acts transitively on ordered pairs of non-sep curves w/ geom. intersection number 1 (by change of coordinates). So apply the key lemma to  $\text{Map}(S_g) \curvearrowright \hat{N}(S_g)$ .

Let  $(v, w) \in E(\hat{N}(S_g))$ .

Exercise: Show  $T_w T_v(w) = v$ .

By the key lemma,  $\text{Map}(S_g)$  is generated by  $T_w T_v$  and  $\text{Stab}(v)$ .

All that remains is showing  $\text{Stab}(v) = \text{Map}(S_g, a)$  is f.g. by Dehn twists about non-sep. curves.

Exercise: Use the following SES's + the inductive hypothesis to complete the proof.

$$1 \rightarrow \underbrace{\text{Map}(S_g, \bar{a})}_{\substack{\text{preserve orientation of } \bar{a}}} \rightarrow \text{Map}(S_g, a) \rightarrow \mathbb{Z}/2\mathbb{Z} \rightarrow 1$$

and

$$1 \xrightarrow{\quad} \langle T_a \rangle \rightarrow \text{Map}(S_g, \bar{a}) \rightarrow \text{PMap}(S_g, -a) \xrightarrow{\quad} 1$$

□